1a) False. A ready process waiting to get access to CPU is

in the ready state SC2005 Operating Systems

A process in the waiting state is waiting for an TUTORIAL TWO event or the I/O operation

2. Data contains

and no local

contains local

variables in a

Size of data

while stack

statically fixed

can grow and shrink as

region is

process

executes

function

global parameters

variables while Stack

parameters and local

## completion of an **Processes and Threads**

b)True. Each PCB has a pointer to other PCBs in a queue and processes in the ready queue are in the ready state The short term scheduler uses this queue to schedule processes.

The head of the queue contains PCB of the process currently in running state

c) False. The wait() system call causes the parent to wait until all child processes. terminate.

Also called join(), reference to fork join processing model. The parent forks several child processes for independent computations and combines their results using join()

Indicate whether the following statements are true or false. Justify your answers. 1.

A ready process waiting to get access to the CPU is in the "waiting" state.

A ready queue is a queue of Process Control Blocks (PCBs) of all processes in the "ready" state.

The "wait()" system call is generally used by a child process to wait for instructions from a parent process.

Message passing based Inter-Process Communication (IPC) consumes less memory than shared memory based IPC.

What are two main differences between the data and stack regions of a process memory?

Explain the difference between a single-threaded and a multi-threaded process.

The figure below shows the execution of processes P0 and P1 in a multiprogramming system.

a) Identify state transitions of each process.

Describe operations A, B, C and D performed by the operating system kernel.

∂ Same as a process

3. Single threaded process is one unit of CPU utilization, using the register set and stack space. Multi threaded process share with its peer threads the code and data sections as well as operating system resources in the same process. While one thread is blocked and waiting, another thread can run, enabling higher throughput

2.

Individual threads can execute concurrently

Threads in a process share code, data and heap regions of memory, whereas stack space is unique to each thread. Each thread has its own Thread Control Block similar to PCB.

Process P0 OS Kernel Process P1 interrupt executing ready idle running -> ready-> running executing running-> I/O system call idle waiting C ready-> running idle executing

4. In the running state, instructions are executed to run the process P0 and P1

An interrupt system call triggers the process to be in the ready state as it waits to be assigned to the CPU

An I/O system call causes the process to wait for the I/O event

A: Change mode to kernel Save P1 state into PCB1 Invoke ISR B:

Run scheduler.select P0 Load P0 state from PCB Change mode to user

C: Change mode to kernel Save P0 state into PCB0 Invoke device driver D: Run scheduler, select P1 Load P1 state from PCB Change mode to user

d)False,