

### 3.1 Stack Manipulation Operations

Fig. 3.1 shows the contents on the system stack and the current position of the stack pointer (SP) of an ARM processor. With reference to Fig. 3.1, answer the following questions.

- (1) What is the current content (in hexadecimal) of the register SP?
- (2) Give the instruction to push the content in R1 onto the system stack.
- (3) Give a single instruction to pop off four word-sized items from the stack to registers R6-R9.
- (4) Give the instruction to retrieve the 3<sup>rd</sup> word-sized item (i.e. 0x00000000) on the stack and place it into R0. You are not to pop any items from the stack.

Address	Contents
(SP) → 0xFEC	0x00123FFE
0xFF0	0x01AD4563
0xFF4	0x00000000
0xFF8	0xEBC19820
0xFFC	0xDE91235A

Fig. 3.1 – The system stack and its contents

### 3.2 Modular Programming – Subroutine Call and Parameter Passing

Fig. 3.2 shows a code segment for subroutine **MySub** and the calling program that makes use of it. With reference to Fig. 3.2, answer the following questions:

- (1) How is each of the parameters **NumX**, **NumY** and **Ans** passed into the **MySub** subroutine? Is it by value or by reference? Give reasons for each of your answers.
- (2) With the aid of the code address shown in Fig. 3.2, give the hexadecimal content of the **PC** and **SP** immediately after the execution of each of the instructions at (b1), (c1) and (s6). Assume execution begins at instruction (a1).
- (3) What is the hexadecimal content in **R0** after instruction (d1) is executed?
- (4) Give a single ARM instruction at (e1) that would remove all the parameters pushed to the stack?
- (5) Replace the **BL MySub** instruction with a branch and any other supporting instruction that supports branching to subroutine.
- (6) With the help of the comments given, complete the subroutine **MySub**. This subroutine implements the function **Ans = NumX \* NumY**. The unsigned word-sized values of **NumX** and **NumY** are multiplied and the resulting word-sized multiplication is stored directly to memory variable **Ans**. For example, given the values of 0x004 and 0x003 for **NumX** and **NumY** respectively, the subroutine **MySub** should compute the result 0x00C and place this result into the memory variable **Ans** as shown in Fig. 3.2.

Code Address					Address	Contents
0x000	Start	MOV	SP, #0xFFFFF0	; (a1) Initialize stack pointer		
0x004		MOV	R0, #0x100		NumX	0x100
0x008		LDR	R1, [R0]		NumY	0x104
0x00C		STR	R1, [SP, #-4]!	; (b1) NumX to stack	Ans	0x108
0x010		ADD	R0, R0, #4			0x00C
0x014		LDR	R1, [R0]			0x10C
0x018		STR	R1, [SP, #-4]!	; (b2) NumY to stack		0x000
0x01C		ADD	R0, R0, #4			0x110
0x020		STR	R0, [SP, #-4]!	; (b3) Ans to stack		
0x024		BL	MySub	; (c1)		
		LDR	R0, [SP, #-4]	; (d1)		
		?	?	; (e1) Remove stack parameters		
		:				
		:				
0x040	MySub	?	?	; (s1) Save registers R0,R1,R2,R3		
		?	?	; (s2) Retrieve NumX from stack		
		?	?	; (s3) Retrieve NumY from stack		
		:		; Complete the segment of code to compute the		
		:		; value of NumX*NumY using successive addition		
		?	?	; (s4a) Move the result directly to ...		
		?	?	; (s4b) ... the memory variable Ans		
		?	?	; (s5) Restore saved registers		
		MOV	PC, LR	; (s6) Return to calling program		

Fig. 3.2 – A partially completed ARM assembly language program

(Question 3.3 and 3.4 need not be covered during the tutorial)

### 3.3 Generating Time Delays using Software

Fig. 3.3 shows a subroutine for generating a short time delay.

- (1) Calculate the delay produced by the subroutine in terms of number of execution cycles.
- (2) Given that the ARM processor system clock is 10MHz, what is the duration of the time delay produced by the routine? You may assume that each memory access cycle uses one cycle of the system clock.
- (3) How can a time delay of 1ms (millisecond) be achieved?

Delay	MOV	R0, #10
Loop	SUBS	R0, R0, #1
	BNE	Loop
	MOV	PC, LR

Fig 3.3 – Delay subroutine

### 3.4 Multi-precision Arithmetic

An ARM assembly language program and the contents of several memory variables are given in Fig. 3.4. Each multi-precision integer **M1**, **M2** and **M3** is stored using two words in memory and in Little Endian format. The memory variable **N1** stores the number of two-word integers in the array. On completion of the multi-precision addition routine, the last two-word integer in the array is replaced by the total sum of all the two-word integers in the array.

- (1) With reference to Fig. 3.4, for each “?”, give the single ARM mnemonic that will implement the corresponding functionality described by each of the comments shown. (note: Some ‘?’ require more than 1 instruction)

<b>Start</b>	?	; push the value in <b>N1</b>	(m1)
	?	; and <b>M1</b>	(m2)
	?	;	(m3)
	?	; on to the stack	(m4)
	?	; make a call to subroutine <b>SubA</b>	(m5)
	:		
	:		
<b>SubA</b>	?	; save value registers <b>R4–R7</b> on to the stack	(s1)
	?	; create a 3-word stack frame on the stack	(s2)
	?	; get the contents of <b>N1</b> on the stack into <b>R4</b>	(s3)
	<b>STR R4, [SP]</b>	; use 1 <sup>st</sup> word (SF1) in stack frame to save value of <b>N</b>	(s4)
	?	; get the address of <b>M1</b> on the stack into <b>R4</b>	(s5)
	?	; clear 2 <sup>nd</sup> word (SF2) in stack frame	(s6)
	?	; clear 3 <sup>rd</sup> word (SF3) in stack frame	(s7)
<b>Loop</b>	?	; add lower word of integer and save result in (SF2)	(s8)
	?	; add upper word of integer and save result in (SF2)	(s9)
	?	; decrement value of <b>N1</b> in (SF1)	
	<b>BEQ Done</b>	; jump if <b>N1</b> =0	(s11)
	<b>ADD R4, R4, #8</b>	; point to next integer	(s12)
	<b>B Loop</b>	; jump back to add again	(s13)
<b>Done</b>	<b>LDR R5, [SP, #4]</b>	; copy result to	(s14)
	<b>STR R5, [R4]</b>		(s15)
	<b>LDR R5, [SP, #8]</b>		(s16)
	<b>STR R5, [R4, #4]</b>	; last array item	(s17)
	?	; collapse stack frame	(s18)
	?	; restore original registers	(s19)
	<b>BX LR</b>		

  

Address	Contents
<b>N</b> 0x100	0x003
<b>M1</b> 0x104	0xFFFFFFFF
	0x108 0x00000111
<b>M2</b> 0x10C	0x00000002
	0x110 0x00000600
<b>M3</b> 0x114	0x00000222
	0x118 0x00000100

**Memory Map**

Fig 3.4 – An incomplete ARM calling program and subroutine

- (2) Describe what changes to the program in Fig. 3.4 you would make if the multi-precision integers are stored using the Big Endian (word wise) format instead.
- (3) Give the two 32-bit hexadecimal values in memory addresses 0x114 and 0x118 at the end of the execution of the code segment shown in Fig 3.4.