reponse times for such applications have to be guaranteed within a pre defined deadline in the worst case SC2005 Operating Systems EDF has poor deadline predictability 1b) True. It behaves unnaturally under overload. 1d) False. The under overload, because once a process hypervisor **TUTORIAL SEVEN** instance overloads, all subsequent creates and process deadlines are in risk. This is manages virtual because priorities are not fixed across **Real-Time OS and Virtualization** machine, helps in instances, unlike in RM/DM. communication Indicate whether the following statements are true or false. Justify your answers. between virtual machines and 1c) False. The Real-time CPS applications usually require a response within a very short time. virtual machine hypervisor migrations interacts with Managing the host OS which in EDF scheduler has poor deadline predictability under overload conditions. interactions between turn interacts with guest OS and the H/W hardware. A guest OS abstracts H/W and provides an interface between H/W and is the key function of a done under type-2 hypervisor. hypervisor. It does so virtualization. The through its virtual quest OS interacts Managing the interactions between a guest OS and its applications is one of the with the hypervisor machine management framework. The key functions of a hypervisor. and provides interactions between services to the guest OS and its Under a Type-1 para-virtualized hypervisor, the hypervisor directly interacts with applications applications are running on it. the H/W and provides a virtual call API to interact with the guest OS. completely independent of the hypervisor. 1e) True. For type Consider the following set of periodic real-time processes: 1 Virtualization, the hypervisor P₂<10, 3, 10> P₁<5, 2, 5> $P_3 < 20, 5, 20 >$ interacts directly with hardware. For Construct all possible schedules for this process set under the earliest deadline first para-virtualization. 1.Tie break (EDF) scheduling strategy. Is this process set schedulable under EDF2 not all of H/W there is an API between P1(R=15) between guest OS instructions are and P3 (R=0) Consider the following set of periodic real-time processes: virtualized. Guest OS and Virtual 2. Tie break needs to be modified Machine. between P2(R=10) $P_1 < 10, 4, 7 >$ $P_2 < 5, 2, 7 >$ P₃<10, 2, 10> to interact with the and P3(R=0) Construct all possible schedules for this process set under the deadline monotonic (DM) 3.Tie break between P1(R=15) scheduling strategy. Is this process set **schedulable** under DM? and P2(r=10) P3 <20,5,20> 2<10,3,10> P1<5,2,5> 1 7 Tie break: can choose schedulable between two or more process as deadline is met under all the above schedules as they tie in comparison P3<10,2,10> 10 *should draw all not schedulable arrows, process deadline is missed releases and deadlines

1a) False. Real time means predictable even in the worst case instead of being fast