FROM A BRANCH-AND-BOUND TO A STATE-OF-THE-ART HEURISTIC

AN EXAMPLE ON THE PERMUTATION FLOWSHOP PROBLEM

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THE PERMUTATION FLOWSHOP

Famous and fundamental problem

	I											_
m_1	j_1	j_2	j_2 j_3		j_4							
m_2		j_1		j_2			j_3		j_4			
m_3				j_1			j_2			j_3	j_4	
	0				8			11			16	T ₁₈

We study 2 objectives:

The makespan: completion time of the last job on the last machine

obj: 18

total completion time (flowtime): sum of completion times

obj: 8+11+16+18 = 53

THE BENCHMARKS

Makespan:

Taillard: sizes between 20 x 5 and 500 x 20

VFR small: sizes between 10 x 5 and 60 x 20

VFR large: sizes between 100 x 20 and 800 x 60

Total completion time:

Taillard: sizes between 20 x 5 and 500 x 20

RESOLUTION METHODS

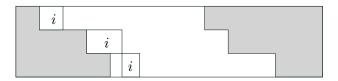
exact methods: mostly **branch-and-bounds**Gmys et al. [2020], Tomazella and Nagano [2020]

meta-heuristics: mostly iterated-greedy algorithms
Fernandez-Viagas and Framinan [2019]
Pagnozzi and Stützle [2019]
Kizilay et al. [2019]

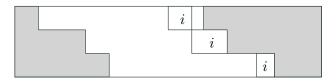
- 1. choose the insertion direction (forward or backward)
- 2. branch on jobs



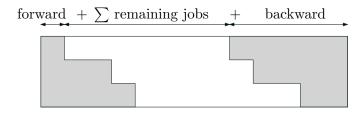
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HOW TO TRANSFORM THIS BRANCH-AND-BOUND INTO AN HEURISTIC?

Key idea: a branch-and-bound explores a search-space

can we focus on "a-priori" good solutions first? (spoiler: yes)

HOW TO TRANSFORM THIS BRANCH-AND-BOUND INTO AN HEURISTIC?

- I. The search strategy
- II. The guidance strategy

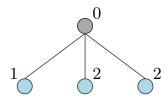
THE "CLASSICAL" SEARCH STRATEGY: DEPTH FIRST SEARCH

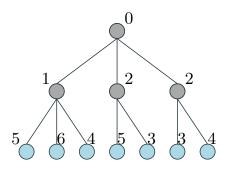
- small memory requirements
- simple

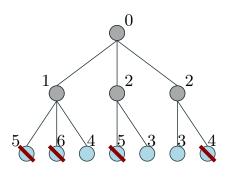
BUT

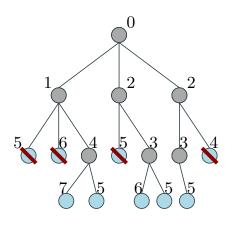
• gets usually stuck in "early bad decisions"

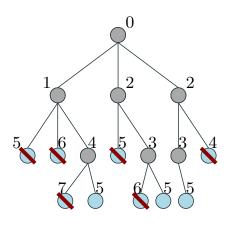












ITERATIVE BEAM SEARCH

- Runs a beam of size 1 (greedy)
- Then runs a beam of size 2, then 4, then 8 ...

Stops when no heuristic fathoming is done (proves optimality)

HOW TO TRANSFORM THIS BRANCH-AND-BOUND INTO AN HEURISTIC?

- I. The search strategy
- II. The guidance strategy

GUIDE FUNCTIONS: WHICH NODE TO CHOOSE?

which criterion?

the bound

the idle time

a bit of both $(\alpha \cdot bound + (1 - \alpha) \cdot idle)$

with weighted idle time (similar to the Liu & Reeves heuristic) $(\alpha \text{ . bound } + (1 - \alpha) \text{ . weighted idle})$

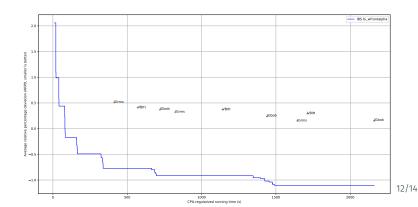
RESULTS

MAKESPAN VARIANT

101/240 new-best-known solutions

- excellent on large instances (500+ jobs, 40+ machines)
- less efficient on "small" instances (100 jobs, 60 machines)

800 jobs, 60 machines

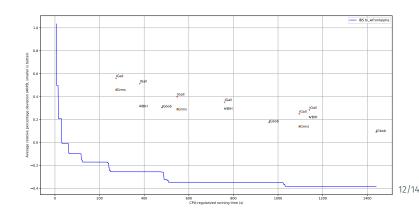


MAKESPAN VARIANT

101/240 new-best-known solutions

- excellent on large instances (500+ jobs, 40+ machines)
- less efficient on "small" instances (100 jobs, 60 machines)

800 jobs, 40 machines

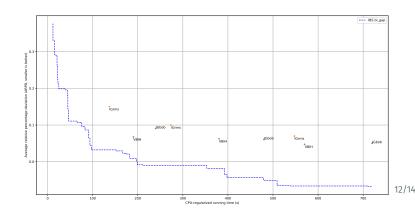


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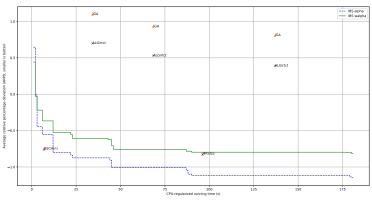
800 jobs, 20 machines



TOTAL COMPLETION TIME VARIANT

only the forward direction 51/120 new-best-known solutions

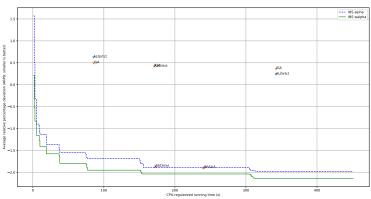
200 jobs, 20 machines



TOTAL COMPLETION TIME VARIANT

only the forward direction 51/120 new-best-known solutions

500 jobs, 20 machines



CONCLUSIONS

This methodology leads to interesting algorithms sometimes, improving upon the state-of-the-art used on other problems:

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This methodology leads to interesting algorithms sometimes, improving upon the state-of-the-art used on other problems:

- · sequential ordering problem (ECAI2020, Q1)
- EURO/ROADEF 2018 challenge (EJOR, Q1)
- · general 2D cutting & packing problems
- longest common subsequence
- pre-print: http://librallu.gitlab.io/pdfs/2020_pfsp_ibs.pdf
- code: https://github.com/librallu/dogs-pfsp

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