

OPTIMIZATION AND PARALLELIZATION  
METHODS FOR THE DESIGN  
OF NEXT-GENERATION  
RADIO NETWORKS

Lucas Benedičić

**Doctoral Dissertation**

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**Evaluation Board:**

*Assist. Prof. Dr. Jurij Šilc, Chairman, Jožef Stefan Institute, Ljubljana, Slovenia*

*Assist. Prof. Dr. Aleš Švigelj, Member, Jožef Stefan Institute, Ljubljana, Slovenia*

*Prof. Dr. Marian Vajtersic, Member, Department of Computer Science, University of Salzburg, Austria*

MEDNARODNA PODIPLOMSKA ŠOLA JOŽEFA STEFANA  
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**OPTIMIZACIJSKE IN PARALELIZACIJSKE  
METODE ZA NAČRTOVANJE  
RADIJSKIH OMREŽIJ  
NASLEDNJE GENERACIJE**

**Doktorska disertacija**

*Supervisor:* Assoc. Prof. Dr. Peter Korošec

*Co-Supervisor:* Assist. Prof. Dr. Tomaž Javornik

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## Abstract

The complexity of the design of radio networks has grown with the adoption of modern standards. Therefore, the role of the computer for the faster delivery of accurate results has become increasingly important. In this thesis, novel methods for the planning and automatic optimization of radio networks are developed and discussed.

The state-of-the-art metaheuristic algorithms, which compare a large number of different network configurations, rely on model-based simulations for the evaluation of the solution quality and the exploration of the search space. However, current radio-network solutions, based on snapshot simulations, have major weaknesses with respect to the simulation time and flexibility provided. In particular, the size of networks that can be analyzed in a feasible time is typically very limited.

The new unified framework developed in this thesis significantly outperforms the currently available solutions for snapshot-based, radio-network simulations. It brings together novel and state-of-the-art parallelization methods, in order to allow for a detailed analysis of very large networks within an acceptable amount of time for everyday planning. This is achieved by the parallel features of the framework, which are exploitable on a single multi-core CPU, as well as on a network of standard PCs with GPU devices. Clearly, the significant speedup achieved at the simulation stage allows for an increased level of detail of the simulations, which improves the accuracy of the results.

Increasing the performance of the simulations involved during the objective-function evaluation is only the first step towards a practical running-time reduction for radio-network optimization. In addition to this, also the optimization algorithms have to be improved in terms of speed, but not at the expense of the quality of results. In this sense, a novel agent-based algorithm is presented and tailored to a classic optimization problem in radio networks. The algorithm, which is based on techniques of cellular automata and population-based metaheuristics, shows considerable gains with respect to the size of problem instances it may handle, as well as regarding its speed performance and solution quality.

The proposed unified framework is tested on complex optimization problems, namely (i) the problem of soft-handover balancing in third-generation systems, and (ii) the parameter optimization of empirical radio-propagation models.

Most mobile operators are aware of the soft-handover balancing problem, but so far, and due to its complexity, it has not yet been tackled by any modern optimization approach. This thesis identifies and formally defines the mentioned problem. Using a black-box approach, different metaheuristic algorithms are employed for solving the problem, the solutions of which show a substantial improvement of downlink and uplink balance.

Another use case of the presented methods is to optimize the parameters of empirical radio-propagation models. Until now, only one parameter set was used to adapt a radio-propagation model to a complete radio network. Using the proposed design automation, the model parameters can be adjusted locally, e.g., for each cell or region in the network, and thus greatly improve the accuracy of the calculated predictions.

On the one hand, the proposed approaches allow a more detailed analysis of radio networks within a reasonable time. On the other hand, the optimization of much larger radio networks is also possible.

## Povzetek

Zahtevnost načrtovanja radijskih omrežij se je povečala z večanjem števila celic v sodobnih celičnih standardih. Zato je vloga računalnika kot pomožno orodje za izvajanje hitrejših in natančnih izračunov vse bolj pomembna. Pričujoča disertacija predstavlja nove metode za načrtovanje in samodejno optimizacijo radijskih omrežij.

Klasični pristopi za optimizacijo in načrtovanje sistemov v velikih sistemih odpovedo, zato sodobni pristopi pri optimizaciji velikih sistemov uporabljajo metahevristične algoritme. Delovanje metahevrističnih algoritmov temelji na velikem številu poizkusov in oceni kakovosti rešitev ter usmerjanju raziskovanja algoritma v iskalnem prostoru. Poizkusi v radijskih omrežjih se izvajajo z modeliranjem radijskih omrežij in simulacijami, vendar imajo več pomanjkljivosti v zvezi s ponujeno fleksibilnostjo in s časom izvajanja simulacije. Velikost omrežij, ki jih je mogoče analizirati na izvedljivem času, je običajno zelo omejena.

Pričujoča disertacija zato predstavlja nove pristope za povečanje hitrosti optimizacije radijskih omrežij z metahevrističnimi algoritmi in njihovo uporabo v konkretnih optimizacijskih problemih radijskih omrežij.

Prvi korak, ki ga predлага disertacija, zajema pohitritev izvajanja poizkusov oz. simulacij z metodami vzporednega računanja. Disertacija podaja in primerja metode vzporednega računanja tako za izvajanje simulacij na večjedrnih procesorjih kot na gruči računalnikov z grafičnimi procesorji, ki so povezani preko lokalnega omrežja. Povečana hitrost izvajanja omogoča povečano stopnjo podrobnosti simulacij, kar posledično izboljšuje natančnost rezultatov.

Povečanje hitrosti izvajanja simulacij je le prvi korak k praktičnemu zmanjševanju računskega časa, potrebnega za evaluacijo cenovne funkcije in posledično optimizacijo radijskih omrežij. Poleg tega je potrebno tudi zmanjšati čas izvajanja optimizacijskih algoritmov, vendar ne na račun kakovosti rezultatov. Doktorska disertacija v ta namen predstavlja optimizacijski algoritem, katerega delovanje je prilagojeno reševanju klasičnih optimizacijskih problemov v radijskih omrežjih. Novi algoritem, ki temelji na tehnikah, uporabljenih v celičnih avtomatih ter metahevrističnih algoritmih, doseže boljše rezultate z vidika velikosti preučenih primerkov, kot tudi z vidika kakovosti rešitev.

Predlagani pristopi so preizkušeni na kompleksnih optimizacijskih problemih, in sicer (i) problem optimizacije ravnovesja povezave navzdol-navzgor pri mehkem izročanju v sistemih tretje generacije in (ii) optimizacija parametrov empiričnih modelov razširjanja radijskega valovanja.

Problem optimizacije ravnovesja povezave navzdol-navzgor pri mehkem izročanju v sistemih tretje generacije je problem, ki se ga zaveda večina mobilnih operaterjev, vendar do sedaj zaradi svoje kompleksnosti še ni bil rešen. Pričujoča disertacija formalno opredeljuje ta optimizacijski problem ter prikaže rešitev problema po principu črne škatle (angl. black box) z uporabo različnih metahevrističnih algoritmov. Rezultati uporabe predlaganih metod pokažejo precejšnje izboljšane razporeditve področij mehkega izročanja.

Drugi primer uporabe metod je optimizacija parametrov empiričnih modelov razširjanja radijskega valovanja. Do sedaj, so se parametri empiričnih modelov razširjanja radijskega valovanja uporabljali za celo radijsko omrežje. Z uporabo predlagane avtomatizacije načrtovanja, pa lahko parametre empiričnega radijskega kanala lokalno prilagodimo, npr. za vsako celico ali regijo v omrežju, s tem pa močno izboljšamo natančnost izračunane predikcije.

Predlagani pristopi na eni strani omogočajo bolj podrobno analizo omrežij v sprejemljivem času, po drugi strani pa z njimi lahko optimiziramo večja radijska omrežja.



<b>Abstract</b>	<b>V</b>
<b>Povzetek</b>	<b>VI</b>
<b>Abbreviations and symbols</b>	<b>XII</b>
<b>1 Introduction</b>	<b>1</b>
1.1 Problem statement . . . . .	2
1.2 Hypotheses and approaches . . . . .	3
1.3 Scientific contributions . . . . .	4
1.4 Organization . . . . .	5
1.5 Publications . . . . .	5
<b>2 Overview of Optimization Methods</b>	<b>7</b>
2.1 Gradient-based methods . . . . .	8
2.2 Linear and non-linear programming . . . . .	9
2.3 Metaheuristics . . . . .	10
2.3.1 Differential evolution . . . . .	12
2.3.2 Differential ant-stigmergy algorithm . . . . .	13
2.3.3 Simulated annealing . . . . .	14
2.4 Black-box optimization . . . . .	14
<b>3 Principles of Mobile Radio Networks</b>	<b>17</b>
3.1 Handover . . . . .	17
3.1.1 Hard handover . . . . .	18
3.1.2 Soft handover . . . . .	18
<b>4 Overview of Radio-Network Optimization</b>	<b>21</b>
4.1 Optimizing base-station locations . . . . .	22
4.2 Optimizing antenna parameters . . . . .	23
4.3 Optimizing coverage . . . . .	25
4.4 Summary . . . . .	26
<b>5 Principles of GPU Programming</b>	<b>27</b>
5.1 CUDA . . . . .	27
5.2 OpenCL . . . . .	29
5.3 Summary . . . . .	29
<b>6 A Parallel Framework for Radio-Network Planning and Optimization</b>	<b>31</b>
6.1 Motivation . . . . .	31
6.2 Related work . . . . .	33
6.3 Radio-coverage prediction for mobile networks . . . . .	34
6.3.1 Background . . . . .	34

6.3.2	Radio-propagation model . . . . .	35
6.4	Design and implementation . . . . .	36
6.4.1	Geographic resources analysis support system . . . . .	36
6.4.2	Multi-paradigm parallel programming . . . . .	36
6.4.3	Design of the serial version . . . . .	37
6.4.3.1	Input parameters . . . . .	37
6.4.3.2	Isotropic path-loss calculation . . . . .	38
6.4.3.3	Antenna diagram influence . . . . .	38
6.4.3.4	Transmitter path-loss prediction . . . . .	38
6.4.3.5	Coverage prediction . . . . .	39
6.4.4	Computational complexity . . . . .	39
6.4.5	Design of the parallel version . . . . .	41
6.4.5.1	Master process . . . . .	42
6.4.5.2	Worker processes on CPU . . . . .	44
6.4.5.3	Worker processes on GPU . . . . .	45
6.4.5.4	Master-worker communication . . . . .	46
6.5	Simulations . . . . .	46
6.5.1	Test networks . . . . .	47
6.5.2	Weak scalability . . . . .	48
6.5.3	Strong scalability . . . . .	50
6.5.3.1	Speedup . . . . .	51
6.5.3.2	Efficiency . . . . .	52
6.6	Summary . . . . .	52
<b>7</b>	<b>Service-Coverage Optimization</b>	<b>55</b>
7.1	Motivation . . . . .	55
7.2	Related work . . . . .	56
7.3	Radio-network model . . . . .	56
7.3.1	Basic elements . . . . .	57
7.3.2	Coverage . . . . .	57
7.4	Problem definition . . . . .	58
7.5	Optimization approaches . . . . .	58
7.5.1	Attenuation-based approach . . . . .	58
7.5.2	Parallel-agent approach . . . . .	59
7.5.2.1	Objective-function evaluation . . . . .	59
7.5.2.2	Autonomous agents . . . . .	60
7.6	Simulations . . . . .	62
7.6.1	Test networks . . . . .	62
7.6.2	Parameter settings of the parallel-agent approach . . . . .	62
7.6.3	Results . . . . .	64
7.6.4	Performance analysis . . . . .	64
7.7	Summary . . . . .	67
<b>8</b>	<b>Soft-Handover Optimization</b>	<b>69</b>
8.1	Motivation . . . . .	69
8.2	Related work . . . . .	71
8.3	Radio-network model . . . . .	72
8.3.1	Soft-handover areas . . . . .	72
8.4	Problem definition . . . . .	73
8.5	Optimization approaches . . . . .	74

8.5.1	Differential evolution . . . . .	74
8.5.2	Differential ant-stigmergy algorithm . . . . .	75
8.5.3	Simulated annealing . . . . .	75
8.6	Simulations . . . . .	75
8.6.1	Test network . . . . .	76
8.6.2	Algorithm parameters . . . . .	76
8.6.3	Results . . . . .	77
8.6.4	Performance analysis . . . . .	78
8.7	Summary . . . . .	80
<b>9</b>	<b>Framework Automated Tuning</b>	<b>83</b>
9.1	Motivation . . . . .	83
9.2	Related work . . . . .	84
9.3	Parameter tuning of the radio-propagation model . . . . .	84
9.3.1	Field measurements . . . . .	85
9.3.2	Linear least squares . . . . .	85
9.3.3	Optimization objective . . . . .	86
9.3.4	Differential ant-stigmergy algorithm . . . . .	86
9.3.5	Simulations . . . . .	87
9.3.6	Results . . . . .	89
9.3.7	Performance analysis . . . . .	89
9.4	Clutter optimization . . . . .	93
9.4.1	Optimization objective . . . . .	93
9.4.2	Simulations . . . . .	94
9.4.3	Results . . . . .	96
9.4.4	Performance analysis . . . . .	98
9.5	Summary . . . . .	100
<b>10</b>	<b>Framework Verification</b>	<b>101</b>
10.1	Motivation . . . . .	101
10.2	Radio-environment setup . . . . .	102
10.3	Performance analysis . . . . .	102
10.4	Summary . . . . .	107
<b>11</b>	<b>Conclusion and Further Work</b>	<b>109</b>
11.1	Scientific contributions . . . . .	110
<b>12</b>	<b>Bibliography</b>	<b>113</b>
<b>Index of Figures</b>		<b>127</b>
<b>Index of Tables</b>		<b>131</b>
<b>Index of Algorithms</b>		<b>133</b>
<b>Appendix A: Bibliography</b>		<b>135</b>
<b>Appendix B: Biography</b>		<b>137</b>
<b>Appendix C: Acknowledgments</b>		<b>139</b>



## Abbreviations and symbols

2G	= Second generation of mobile networks
3G	= Third generation of mobile networks
4G	= Fourth generation of mobile networks
ACO	= Ant-colony optimization
AMC	= Adaptive modulation and coding
BS	= Base station
CDMA	= Code division multiple access
CUDA	= Compute unified device architecture
DASA	= Differential ant-stigmergy algorithm
DB	= Database system
DE	= Differential evolution
DEM	= Digital elevation model
GIS	= Geographical information system
GPS	= Global positioning system
GPU	= Graphics processing unit
GRASS	= Geographic resources analysis support system
GSM	= Global system for mobile communications
HPC	= High-performance computing
HSDPA	= High speed downlink packet access
HSPA	= High speed packet access
HSUPA	= High speed uplink packet access
I/O	= Input/output
INI	= Initialization file
KPI	= Key performance indicator
LAN	= Local area network
LOS	= Line of sight

LTE	= Long term evolution
MIMO	= Multiple input multiple output
MPI	= Message passing interface
MW	= Master-worker parallel paradigm
MWD	= Master-worker-database parallel paradigm
NACC	= Nagasaki advanced computing center
NFS	= Network file system
NLOS	= Non line of sight
OA	= Open area
PDF	= Probability distribution function
PDF	= Probability distribution function
PRATO	= Parallel radio-prediction tool
QoS	= Quality of service
RRM	= Radio-resource management
RSG	= Regular square grid
RSRP	= Reference signal received power
RSRQ	= Reference signal received quality
SA	= Simulated annealing
SHO	= Soft handover
SIMT	= Single instruction multiple thread
SM	= Streaming multiprocessor
SP	= Streaming processor
SPMD	= Single-program multiple-data
SQL	= Structured query language
TDMA	= Time division multiple access
UE	= User equipment
UMTS	= Universal mobile telecommunications system
WCDMA	= Wideband code division multiple access
$\alpha$	= Knife-edge diffraction control parameter
$\epsilon$	= Maximum parameter precision in DASA

$\gamma^{\text{cov}}$	= Signal-to-interference ratio coverage threshold
$\gamma^{\text{sho}}$	= SHO window
$\text{cov}(c, m)$	= Binary function to assert the coverage of a mobile $m \in M$ from a cell $c \in C$
$L(d_{(x,y)})$	= Path loss at the topography point with coordinates $(x, y)$
$L_{\text{CLUT}}$	= Signal loss due to clutter
$L_{\text{NLOS}}$	= Path loss in non-line-of-sight conditions
$L_{\text{OA}}$	= Path loss in an open-area environment
$N_0$	= Thermal noise
$\text{SINR}(c, m)$	= Signal-to-interference ratio from cell $c \in C$ to mobile $m \in M$
$\text{uniform}[0, 1)$	= Function representing a uniformly-distributed random number from the interval $[0, 1)$
$\rho$	= Pheromone dispersion factor in DASA
$\vec{\beta}$	= Control-parameter vector of the path-loss model
$A_{\text{covered}}$	= Area under service coverage of the mobile network
$A_{\text{total}}$	= Complete geographical area under optimization
$as^{\max}$	= Maximum number of neighbouring cells in the active set
$b$	= Discrete base in DASA
$C$	= Set of antenna installations (cells) in a mobile network
$C_m$	= Subset of cells, $C_m \subset C$ , that cover a mobile $m \in M$
$c_m^*$	= Best-serving cell of mobile $m \in M$
$CR$	= Crossover constant in DE
$d_{(x,y)}$	= Distance from the cell to the topography point with coordinates $(x, y)$
$E(d_{(x,y)})$	= Loss correction due to the Earth sphere
$F$	= Mutation scaling factor in DE
$f_{\text{param}}^*(c)$	= Objective function of the parameter-tuning problem for cell $c \in C$
$f_{\text{sho}}$	= Objective function of the SHO-balancing problem
$FM_c$	= Set of RSRP field measurements for cell $c \in C$
$g$	= Generation index in DE
$g_{\max}$	= Number of generations for DE to run
$H_A$	= Antenna height above the receiver

$H_R$	= Height of the receiver antenna above ground level
$it$	= Maximum number of iterations in SA
$K(d_{(x,y)})$	= Knife-edge diffraction loss
$l_{cm}^{\downarrow}$	= Downlink attenuation factor between cell $c \in C$ and mobile $m \in M$
$l_{mc}^{\uparrow}$	= Uplink attenuation factor between mobile $m \in M$ and cell $c \in C$
$M$	= Set of mobile devices or users of a mobile network
$m$	= Number of ants in DASA
$NP$	= Population size in DE
$P_c$	= Set of candidate pilot power settings for cell $c \in C$
$p_c$	= Pilot-power setting of cell $c \in C$
$p_m^{\uparrow}$	= Uplink transmit power of mobile $m \in M$
$s_+$	= Global scale-increasing factor in DASA
$s_-$	= Global scale-decreasing factor in DASA
$t_{\text{initial}}$	= Initial temperature in SA

# 1 Introduction

Many researchers believe the computer has become the third method to do research, in addition to theory and experimentation, for both science and engineering. Although there is no complete agreement on the position intended for scientific computing with respect to the other two methods, it is undeniable that computational methods are an essential tool in most disciplines, particularly in those related to decision making.

Nowadays, decision making is present practically everywhere. As scientists, engineers and managers have to make decisions in more complex and competitive circumstances every day, decision making involves dealing with rational and optimal approaches. According to Talbi [156], decision making consists of the following steps:

- formulating the problem,
- modeling the problem,
- solving the problem, and
- implementing a solution for the problem.

Formulating a decision problem means making an initial statement about it. Although this first formulation may be imprecise, the objectives of the problem are outlined, together with the internal and external factors that have some degree of influence over it. During the modeling of the problem, an abstract mathematical model is built for it. Sometimes this model is inspired by similar models in the literature, making it possible to tackle the problem with well-studied methods. After a model of the problem is available, one may start solving it. In the context of this thesis, optimization means generating “good” solutions for the problem. It is important to note that the resulting solutions are given for the abstract model, and not for the original problem itself. Therefore, the performance of the obtained solution is indicative when the model is an accurate one [156]. In the last step, the obtained solution is practically tested by the decision maker and implemented if it is an “acceptable” or “good” one. In case of “bad” or “unacceptable” solutions, the decision-making process is repeated, possibly improving the model and/or the optimization algorithm. The process, as described here, is depicted in Figure 1.1.

Scientific computing, by means of computer-science methodology, enables the study of problems that are too complex to be treated analytically, or those that are very expensive or dangerous to be studied by direct experimentation. Real-world problems are typically very complex systems to be directly assessed by analytical models, and require a numerical

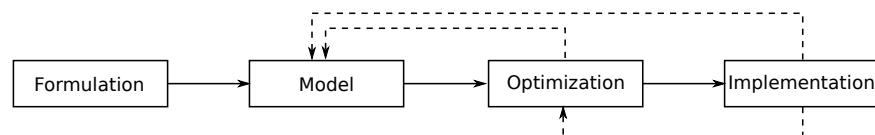


Figure 1.1: A diagram of the classic decision-making process. Adapted from [156].

simulation for their study. Computer simulations provide a resource to mimic the behavior of complex systems, by numerically evaluating a model and gathering its data to estimate their true characteristics [96].

A model is a simplified representation of a studied problem, and one of its purposes is to predict the effects of the variations within the system. A good model is a balance between realism and simplicity. The system simulation, on the other hand, is the operation of the model. Its configuration can be changed, thus allowing multiple experimental executions, something that might not be possible with the real system it represents [106]. However, it is important to understand that the models used in scientific simulations and engineering never offer a perfect representation of the system they resemble, but only a subset of its composition and dynamics. For this reason, experimentation and expert observation will always be essential as reference points for understanding the studied phenomena. Consequently, problems categorized as of large size and of considerable complexity represent a challenge, because of the different involved disciplines for their study and the degree of difficulty of their modeling. Cellular radio networks in general, and those of the latest generations in particular, fall under this categorization.

## 1.1 Problem statement

Radio networks represent one of the most fast-growing technology markets since the introduction of the Global System for Mobile communications (GSM) [1]. As an implementation of the second generation of mobile networks (2G), GSM appeared more than twenty years ago. Its successor, the Universal Mobile Telecommunications System (UMTS) [2] marks an evolution from 2G, representing a milestone for the third generation of mobile radio networks (3G). In recent years, the first commercial networks implementing the Long Term Evolution (LTE), also known as fourth generation (4G) LTE, have also appeared [61]. The always increasing demand for more bandwidth has been one of the main forces behind the standardization and later implementation of systems delivering higher-speed data services in order to improve the user's experience.

This evolution, first from 2G to 3G and later from 3G to 4G, has introduced not only the technology needed to increase data-transfer capacity and voice quality, but also a greater complexity in terms of radio-network planning, deployment, and configuration. This fact has attracted the attention of the research community into areas such as the design and optimization of radio-networks.

During the design phase of a radio network, a traditional or manual approach comprises the software tool, i.e., the model in Figure 1.1, that executes the analysis, and the human that makes the configuration changes, i.e., the optimization in Figure 1.1. Therefore, a radio engineer manually adjusts the network parameters and the software tool analyzes the given configuration. If the obtained results are not acceptable, the analysis process has to be repeated several times, until the goal is achieved and the changes are implemented, i.e., the implementation in Figure 1.1. In the context of this thesis, this process is referred to as manual radio-network optimization.

Advances in the past few years have improved the manual optimization process by introducing different problem-solving approaches that increase the role of the computer during the optimization of radio networks, consequently enlarging the scope of problems and instance sizes that may be subject to optimization. Still, there are some important aspects that restrict the automation level of these methods, not only in real-world environments, but also when doing research in the area of radio networks:

- A selected optimization method is typically a compromise between solution quality and computational-time complexity. The proposed state-of-the-art approach for the evaluation of radio networks is the Monte-Carlo snapshot analysis [159, 167]. However, real-world environments, where radio-network design is carried out, require the evaluation of networks with thousands of base stations in a reasonable amount of time. Moreover, for applications involving radio-network optimization, usually millions of evaluations are required to find a good solution, in which case also snapshot simulations are too time-consuming for practical use. Therefore, for such applications and environments, methods with improved time efficiency are required.
- A considerable number of publications in the field of radio-network optimization, a few references of which are given for illustration purposes [7, 28, 30, 66, 144, 143], base their simulations on platforms for which it is not possible to reproduce the experiments, either because they have used proprietary software or because the data is not available. This fact reduces the possibilities for comparing different approaches among each other, and significantly contrasts with other research areas, such as evolutionary computing or numerical optimization, where several sets of open and well-known benchmarks are available for the community to use. Consequently, an open and unified framework should allow researchers to compare different methods and results in a simplified and objective manner.
- Available commercial tools for radio-network evaluation present several drawbacks. In particular, regarding the size of networks that can be considered and their computational-time complexity. Yet, even if precise and fast methods were commercially available, they would lack the level of flexibility required by the scientific community. Consequently, an essential attribute of the framework is to be open source, so that anyone can extend it to meet some specific requirements. In the long term, this process should also extend the set of built-in functionality.
- Particularly for path-loss predictions, estimated with empirical or deterministic mathematical models, the inaccuracy of the input data directly deteriorates the precision of the calculated results. Moreover, since the physical properties that influence the propagation of radio signals are not constant and every environment introduces its own deviations, the calculated prediction may be considered as not more than a rough approximation. Therefore, there is a need for a technique that improves the accuracy of the state-of-the-art mathematical models for radio propagations, despite the various sources of error noted before.

This thesis introduces methods and tools to mitigate the above-mentioned drawbacks from a radio-planning perspective.

## 1.2 Hypotheses and approaches

The work presented in this thesis is based on the following hypotheses and their related approaches:

- Applying parallelization techniques should reduce the computational-time requirement of the radio-propagation prediction, thus making it possible to process larger, real-world radio networks.
  - Unlike several examples in the literature of radio-network simulators, which do not include support for parallel methods in order to reduce the computational-time complexity, the unified framework includes this feature in its basic structure.

In this sense, parallel techniques are used even when targeting a single computing host.

- Using hardware, specialized for parallel execution, e.g., a graphics-processing unit (GPU), should improve the execution time of parallel algorithms using threads or message-passing mechanisms.
  - From its arrival a few years ago, general GPU programming is continuously gaining the attention of the research community and it is extending to several areas of science. In particular, the combination of parallel-programming techniques with GPU hardware should reduce the computational-time requirement of the radio-propagation prediction even further.
- Using the parallel framework for the objective-function evaluation of different optimization problems should improve the quality of solutions and enable tackling larger problem instances.
  - On the one hand, a larger number of evaluations in the same amount of time translates into a better exploration of the search space by the optimization algorithm used. On the other hand, more complex or exact models may be evaluated in the same amount of time, which implies more accurate solutions.
- Using the parallel framework to evaluate the objective function of an automatic-optimization system should provide the means for solving new, previously inaccessible optimization problems.
  - The improved performance at the evaluation level opens new possibilities for formally defining and tackling optimization problems that were previously out-of-reach of state-of-the-art methods, either because of their complexity or size.

### 1.3 Scientific contributions

The contributions of this thesis to the fields of telecommunications and computer sciences include the following:

1. Design and development of a unified framework for radio-network planning and optimization that supports computer clusters and GPUs.
2. Proposal of a new approach for parallel programming that enhances the performance of a classic master-worker method.
3. Quality improvement of radio-propagation predictions by applying metaheuristic-optimization techniques.
4. A new algorithm, based on autonomous agents, to tackle the service-coverage problem in radio networks.
5. Identification and formalization of a new optimization problem in radio networks.

## 1.4 Organization

So far, a context has been provided for the following content of this thesis. An introduction of the three research areas addressed by this work, i.e., optimization, radio networks, and parallel methods, is given in Chapters 2, 3 and 5, respectively. They provide an overview of the theoretical elements needed for a better understanding of the rest of the thesis. Additionally, Chapter 4 provides a short literature survey about the optimization of radio networks.

The unified framework for radio-network planning and optimization, incorporating elements from the above-mentioned fields of study, is presented in Chapter 6.

Chapters 7 and 8 demonstrate how the framework is applied to tackle two optimization problems. In Chapter 7, the problem of minimizing the total amount of pilot power subject to a full coverage constraint is addressed with a novel approach. Next, Chapter 8 formally introduces a new optimization problem that deals with the balancing of downlink and uplink soft-handover areas in UMTS networks.

The applicability and performance of the framework for radio-network planning is validated in Chapters 9 and 10. Chapter 9 further extends the framework with automated-tuning capabilities that enable its adaptation to different environments, thus improving the accuracy of the radio-propagation predictions. In Chapter 10, the framework is tested in real network-planning conditions and compared to a commercial, enterprise-level tool in terms of solution quality and speed performance.

Finally, Chapter 11 gives a short summary of the thesis, outlines its main contributions and discusses further work.

## 1.5 Publications

The work presented in this thesis is supported by a number of previous publications. The initial work on the unified framework for radio-network planning was published by Benedičić et al. in [22]. To address the challenges of the computational-time complexity during optimization and its objective-function evaluations, Benedičić et al. published the GPU extensions for the framework in [21]. The agent-based algorithm for solving the service-coverage problem was initially published at the International multi-conference Information Society [19]. This work was further extended and later published in [21]. The primary ideas for the formal definition of the soft-handover balancing problem were published at the World Congress of Computational Intelligence [20].

A comprehensive list of the author's publications is given in Appendix A.



## 2 Overview of Optimization Methods

Optimization may be informally defined as the procedure of finding better solutions to a given problem that usually models some physical phenomenon. In our every day life, we are constantly solving small optimization problems, like choosing the shortest route to a friend's house, or organizing the appointments in our agenda. In general, these problems are small enough for us to find a good solution without extra help, but as they become larger and more complex, the aid of computers for their resolution is unavoidable.

Complex multidimensional optimization problems are popular in engineering, economics, physics and other scientific fields. When solving an optimization problem, the objective is to find a "good" solution in a "reasonable" computational time. In this respect, the field of mathematical optimization has received a lot of attention by the scientific community during the last decades. However, both "good" and "reasonable" are problem, application and context-specific concepts, in which the biggest challenge of selecting an appropriate optimization approach usually lays.

Mathematical optimization involves the process of finding solutions from a group of possible decisions, which may be defined as:

$$\min f(\vec{x}) \quad \vec{x} \in \Omega \subseteq \mathbb{R}^n, \quad (2.1)$$

where  $\vec{x} = (x_1, \dots, x_n)$  is a vector representing the decision variables,  $f(\vec{x})$  is the objective function measuring the quality of the decisions and  $\Omega$  is the set of possible solutions of the problem, also known as search space. Note that the objective function  $f$  makes it possible to define a total order relation between any pair of solutions in  $\Omega$ .

The search space  $\Omega$  may also be expressed as a solution to a system of equations or inequalities, e.g.:

$$\begin{aligned} g(x_1, \dots, x_n) &\leq 0 \\ h(x_1, \dots, x_n) &= 0. \end{aligned} \quad (2.2)$$

Optimization problems involving the maximization of the objective function also fall into this category, since:

$$\max f(\vec{x}) = -\min(-f(\vec{x})). \quad (2.3)$$

A point  $\vec{x}^*$  is considered to be an unrestricted local minimum of a function  $f$  if it holds a better value than all its neighbors, i.e., there exists  $\epsilon > 0$  so that:

$$f(\vec{x}^*) \leq f(\vec{x}) \quad \forall \vec{x} \in \mathbb{R}^n \quad |\vec{x} - \vec{x}^*| < \epsilon. \quad (2.4)$$

Similarly, a point  $\vec{x}^*$  is considered to be an unrestricted global minimum of a function  $f$  if it holds a better value than all others, i.e.:

$$f(\vec{x}^*) \leq f(\vec{x}) \quad \forall \vec{x} \in \mathbb{R}^n. \quad (2.5)$$

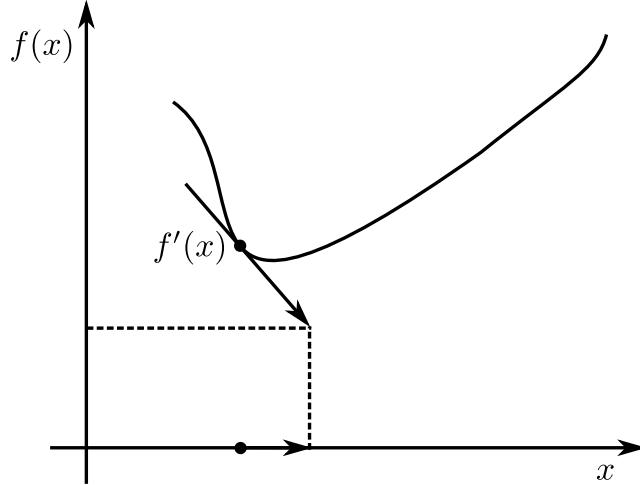


Figure 2.1: Gradient descent with a negative slope, i.e.  $x$  is increasing.

The concepts of local and global minimum are considered strict if the inequalities of Equations 2.4 and 2.5 are strict. Likewise, the definition of local and global maximum is given by the existing relation between a minimization and a maximization problem, as specified in Equation 2.3, i.e., a point  $\vec{x}^*$  is a local or global maximum of a function  $f$  if and only if  $\vec{x}^*$  is a local or global minimum of function  $-f$ , respectively.

## 2.1 Gradient-based methods

Gradient-based methods are among the oldest and most studied optimization approaches. They are based on the derivative of the optimized function, using the first and even the second derivative of a function  $f$ . The name gradient follows from the derivative of multidimensional functions,  $\nabla f(\vec{x})$ , which is simply a vector where each element is the slope of  $\vec{x}$  in that dimension, i.e.,  $\langle \frac{\partial f}{\partial x_1}, \dots, \frac{\partial f}{\partial x_n} \rangle$  [104].

The principle behind gradient-based methods is rather simple. Starting from an arbitrary value for  $x$ , a subtraction (or addition) of a small positive value is iteratively applied to it, e.g., for gradient descent:

$$x \leftarrow x - \alpha f'(x), \quad (2.6)$$

where  $\alpha$  is a small positive value, and  $f'(x)$  is the first derivative of  $f(x)$ . Consequently, a positive slope will make  $x$  decrease, whereas a negative slope will make it increase. Figure 2.1 shows an example of this behavior. Therefore,  $x$  will gradually move down the function until it finds its minimum, where  $f'(x)$  is zero, causing it to stop.

However, gradient methods have certain drawbacks that make them unsuitable for tackling a wide range of optimization problems. Take, for example, the time they take to converge. As gradient descent approaches a function minimum, it will skip this point and land on the other side. In the next step, something similar will happen, but this time from the other side of the minimum point, thus slowly approaching to the target in a “zig-zag” way. This behavior is directly related to the slope of the function at the given point, i.e., a steepest slope translates into a larger jump, and may be alleviated by adjusting the value of  $\alpha$ . However, some functions (or regions of functions) may require smaller values, while for

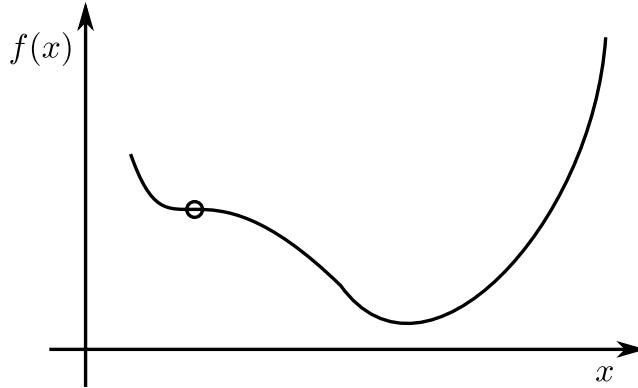


Figure 2.2: A saddle point or point of inflection, where the derivative is zero.

others, a bigger value would be more appropriate. Newton's method improves this by taking the second derivative of the function,  $f''(x)$ , into account, i.e.:

$$x \leftarrow x - \alpha \frac{f'(x)}{f''(x)}, \quad (2.7)$$

thus adjusting the value of  $\alpha$  as it converges towards a point with zero slope [104].

Another issue is how other points are handled. Beside maxima and minima points, some functions also contain saddle points (known as inflection points in one-dimensional functions). Clearly, the first derivative of a saddle point is zero, meaning gradient descent will stop looking for the minimum, even though it has not found it (see Figure 2.2). Newton's method, on the other hand, does not help either, even trying to divide by zero in this case. These observations clearly show how gradient methods get caught in local optima. Local optima of a function are defined as the optima (or minima in this case) of a local region. Similarly, global optima are defined as the optima of the whole domain of a function. It follows that gradient methods, as gradient descent or Newton's method, are local optimization algorithms [104].

But maybe the biggest concern with gradient-based methods is that they assume the function under optimization is derivable. This assumption holds only when optimizing a well-formed mathematical function. Unfortunately, this is generally not true, since in most cases, the gradient is not computable because the function is not known. The only available approach in such situations is creating inputs to the function in order to assess their quality. Metaheuristics (see Section 2.3) are good candidates for this class of problems, for solving both moderate and large instances.

## 2.2 Linear and non-linear programming

It was in the early 40s of the twentieth century, through the work of teams formed by mathematicians, economists and physicists, that the bases were established for the resolution of problems with a set of techniques known as linear and non-linear programming. Their initial goal was to solve different kinds of logistic problems during the second world war.

In a linear programming optimization problem, both the objective function  $f$  and a given set of constraints are linear functions. The constraints impose restrictions over  $x$ , i.e., they must meet certain requirements as, for example, fulfill a limited availability of resources. A problem may be formulated as follows:

$$\min f(x) = c \cdot x \quad (2.8)$$

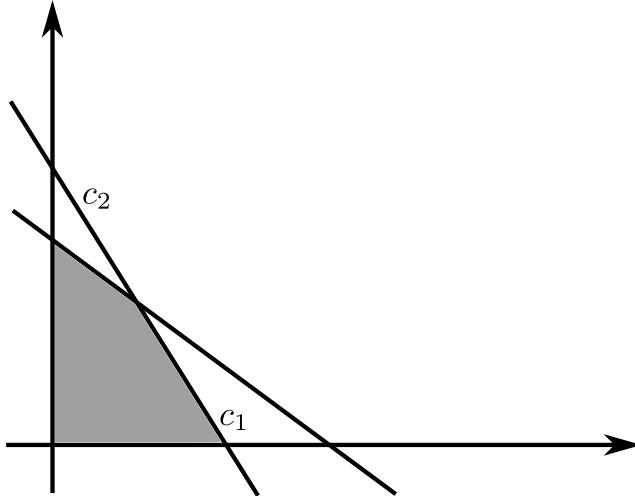


Figure 2.3: Graphical representation of a linear-programming example with two constraints,  $c_1$  and  $c_2$ . The grayed area is the polytope representing the region of feasible solutions.

subject to

$$\begin{aligned} A \cdot x &\leq b \\ x &\geq 0. \end{aligned} \tag{2.9}$$

In the example above, the inequalities defined in Equation (2.9) are the constraints to the linear program defined in Equation (2.8).

For solving continuous, linear-optimization problems, efficient exact algorithms exist, such as the simplex method [43] or the interior-points method [89]. Indeed, linear programming is one of the most satisfactory models for solving optimization problems, since the feasible region of the problem is a convex set and the objective function is a convex function. It follows that the global optimum is a node of the polytope representing the feasible region [156]. See Figure 2.3 for a linear-programming example with several constraints.

Non-linear programming models, on the other hand, consider problems where the objective function  $f$  and/or the constraints are non-linear [14]. However, non-linear continuous problems are more difficult to solve. Despite several existing techniques to linearize such models, they often introduce extra variables and constraints, and also some degree of approximation [63]. Moreover, some problem properties such as high dimensionality, parameter interaction, and multi-modality make these approaches ineffective.

Generally speaking, when dealing with real-world problems, the availability of analytical optimization models, such as those required by gradient methods or (non-)linear programming, is not guaranteed. Indeed, for some applications, only simulations or physical models are the available means for objective-function evaluation [54]. Once again, metaheuristics appear as good candidates to solve different instance sizes of this class of problems.

## 2.3 Metaheuristics

Metaheuristics, a term proposed by Glover in [64], represent a group of approximation algorithms designed to combine basic, heuristic principles with advanced high-level guidance methods, targeted at improving the efficiency of a search process. These techniques are

meant to find good solutions to a given problem, for which the mathematical function is not available or its search space is big enough for an exhaustive search to be unfeasible [58].

From the theoretical point of view, metaheuristics represent a subset of stochastic optimization, since they use some degree of randomness to find optimal (or as good as possible) solutions to hard problems. They are the most general of this kind of algorithms, and are applied to a wide range of problems [104]. Unlike the exact optimization methods introduced in the previous sections, metaheuristics do not generally guarantee the optimality of the obtained solutions [156]. Moreover, they do not define how close the obtained solutions are from the optimal ones, as approximation algorithms do.

The characterization given by Blum and Roli [25] provides a clear overview of the fundamental properties associated with metaheuristics:

- metaheuristics are strategies that “guide” the search process;
- their goal is to efficiently explore the search space in order to find optimal or near-optimal solutions;
- they build upon techniques which range from simple local search procedures to complex learning processes;
- they are approximate and usually non-deterministic;
- they may incorporate mechanisms to avoid getting trapped in confined areas of the search space;
- their basic concepts permit an abstract-level description, which is not problem-specific;
- they may make use of domain-specific knowledge in the form of heuristics that are controlled by the upper level strategy;
- advanced metaheuristics use search experience (implemented as some form of memory) to guide the search process.

The strategies used by metaheuristics should provide a dynamic balance between the exploitation of the accumulated search experience (commonly called intensification) and the exploration of the search space (commonly called diversification) [25]. This balance provides the necessary means to quickly identify promising regions, and early discarding those which have already been explored or do not provide solutions of better quality. Promising regions within the search space, which are identified by the obtained “good” solutions, are thoughtfully explored during the intensification phase, hoping to find better solutions. On the other hand, during the diversification phase, not-yet-visited regions are explored, making sure the search space as a whole is evenly explored, thus confining the search to a reduced number of regions is avoided. In this context, the ultimate search algorithm in terms of diversification is random search. Random search generates a random solution in the search space at each iteration, without using memory [156]. In terms of intensification, iterated local search [102] is a representative algorithm. The steepest local search algorithm selects, at each iteration, the best neighboring solution that improves the current one [156].

Metaheuristics are applicable where state-of-the-art exact algorithms cannot tackle the given instances within the required time, either because of their size or structure. The meaning of “required time” within this context directly depends on the target optimization problem itself. A feasible or acceptable time may vary from some seconds to several months, again, depending on the target optimization problem, e.g., real-time decisions against structural-design problems.

Based on the characterization given by Talbi [156], a summary of the essential properties of optimization problems that justify the use of metaheuristics follows:

- Very large problem instances. Even though exact polynomial-time algorithms might be known for solving the target problem, they are too expensive due to the size of the instances.
- Problems with hard real-time constraints, where a “good solution” has to be found online. Metaheuristics appear as an alternative to exact algorithms in order to reduce the search time.
- A difficult problem of moderate size, which input instances have an intricate structure.
- Optimization problems with time-consuming objective function(s) and/or constraints. Indeed, various real-world optimization problems are characterized by the huge computational cost of the objective functions. Several radio-network design problems fall into this category.
- Problems that cannot be solved with exhaustive search due to the non-analytical models on which they are based. These problems are defined by a black-box evaluation of the objective function (see Section 2.4).

The influence of these conditions may increase in the presence of non-deterministic optimization models, e.g., problems with complex Monte Carlo simulations [38].

Undoubtedly, metaheuristics are rapidly gaining popularity as optimization problems are increasing in both size and complexity. Indeed, as the computing power of commodity hardware increases, the possibility of building models of greater complexity is available for developing more accurate models of real-world problems in engineering and science.

Related literature groups metaheuristic algorithms due to their behavior, for example:

- differential evolution, from the family of evolutionary algorithms;
- differential ant-stigmergy algorithm, from the family of swarm-intelligence algorithms; and
- simulated annealing, from the group of classic metaheuristic algorithms.

Each of these algorithms minimize an objective-function value by adopting essentially distinct approaches. In the following sections, a short description of their functioning is given.

### 2.3.1 Differential evolution

Differential Evolution (DE) [149] is a simple and powerful evolutionary algorithm proposed for numerical optimization. The version of DE used for this thesis is known as DE/rand/1/bin, or "classic DE" [125]. At the beginning, the classic DE initializes a population of  $N$ ,  $D$ -dimensional vectors with parameter values that are distributed with random uniformity between pre-specified lower and upper initial parameter bounds,  $x_{i,\text{low}}$  and  $x_{i,\text{high}}$ , respectively [128], i.e.:

$$\begin{aligned} x_{i,j,g} &= x_{i,\text{low}} + \text{uniform}[0, 1] \cdot (x_{i,\text{high}} - x_{i,\text{low}}), \\ i &= (1, 2, \dots, D), j = (1, 2, \dots, N), g = 0. \end{aligned} \quad (2.10)$$

In Equation (2.10),  $g$  denotes the generation index, while  $i$  and  $j$  indicate the parameter and population indices, respectively. Consequently,  $x_{i,j,g}$  is the  $i^{\text{th}}$  parameter of the  $j^{\text{th}}$  population vector in generation  $g$ . The (pseudo)-random number generator is represented by a function “uniform[0, 1]” that selects a uniformly-distributed random number from the interval [0, 1].

DE generates solutions by applying the mutation and cross-over operations. A vector from the current population is mutated by adding the scaled difference of two other vectors from the same population, combined with the mutation-scale factor. A proportion of the parameters of the mutated vector are crossed with those belonging to the  $j^{\text{th}}$  population vector, resulting in a new solution vector. The cross-over constant controls the proportion of parameters that the mutant vector contributes to the new vector.

The generated solution vector is retained if it yields a lower objective-function value than the  $j^{\text{th}}$  population vector, otherwise the  $j^{\text{th}}$  vector is kept for at least one more generation.

The four parameters to control the search process of DE are: the population size ( $NP$ ), the number of generations for the algorithm to run ( $g_{\max}$ ), the crossover constant ( $CR$ ), and the mutation scaling factor ( $F$ ).

An extensive description of DE and its variants may be found in [125]. A wide range of optimization problems have been solved by applying DE [44], also in the area of radio-network optimization [20, 41, 110].

### 2.3.2 Differential ant-stigmergy algorithm

The Differential Ant-Stigmergy Algorithm (DASA) [92] is based on the metaheuristic Ant-Colony Optimization (ACO) [46]. ACO was pioneered by Dorigo in 1992 and is based on the foraging behavior of social ants, which use pheromone as a chemical messenger. When looking for a food source, a swarm of ants interacts in their local environment. An ant is able to follow the route marked with pheromone laid by other ants. When an ant finds a food source, it will mark the trail to and from it with pheromone.

The pheromone concentration,  $p$ , evaporates at a constant rate,  $\rho$ , over time,  $t$ , i.e.:

$$p(t) = p_0 e^{-\rho t}, \quad (2.11)$$

where  $p_0$  is the initial pheromone concentration at  $t = 0$ . The evaporation is important, as it provides the means for algorithm convergence and self-organization.

The pheromone concentration varies over time and the ants follow the route with higher concentration. In turn, the pheromone concentration increases with the number of ants. That is, as more and more ants follow the same route, it becomes the favored path, producing a positive feedback mechanism. Consequently, as the system evolves in time, it converges to a self-organized state, which is the essence of any ant-based algorithm.

As it was mentioned before, DASA is an ant-based algorithm that provides a framework to successfully cope with high-dimensional numerical optimization problems. It creates a fine-grained discrete form of the search space, representing it as a graph. This graph is then used as the walking paths for the ants, which iteratively improve the temporary best solution.

At the end of every iteration, and after all the ants have created solutions, they are evaluated to establish if any of them is better than the best solution found so far.

There are six parameters that control the way DASA explores the search space: the number of ants ( $m$ ), the discrete base ( $b$ ), the pheromone dispersion factor ( $\rho$ ), the global scale-increasing factor ( $s_+$ ), the global scale-decreasing factor ( $s_-$ ), and the maximum parameter precision ( $\epsilon$ ).

---

**Algorithm 2.1** Pseudo-code of the simulated-annealing metaheuristic. Adapted from [23].

---

**Require:**  $x \leftarrow$  initial solution  
**Require:**  $T_1 \leftarrow$  temperature parameter  
**for**  $k=1,2,\dots$  **do**  
     $y \leftarrow$  random solution from  $\Omega$   
    **if**  $f(y) \leq f(x)$  **then**  
         $x \leftarrow y$   
    **else if**  $e^{(\frac{f(x)-f(y)}{T_k})} \leq \text{uniform}[0, 1]$  **then**  
         $x \leftarrow y$   
    **end if**  
     $T_{k+1} \leftarrow$  update temperature  $T_k$   
**end for**

---

A more in-depth explanation about these parameters and the DASA algorithm itself can be found in [92].

### 2.3.3 Simulated annealing

This metaheuristic was introduced in the area of combinatorial optimization by Kirkpatrick et al. [90]. The algorithm resembles the physical annealing process, where particles of a solid material arrange themselves into a thermal equilibrium. Simulated Annealing (SA) is based on the principle of local-search heuristics, and uses a pre-defined neighborhood structure on the search space  $\Omega$ . The temperature parameter controls the search behavior. Its name follows the analogy to the physical annealing process. In each iteration, a solution  $y$ , which is a neighbor of the current solution  $x$ , is computed. If  $y$  has a better objective-function value than  $x$ , it becomes the new solution, i.e., the current solution  $x$  is replaced by  $y$ . If, on the other hand,  $y$  has a worse objective-function value than  $x$ , it is only accepted with a certain probability depending on two factors:

- the difference of the objective-function values of  $x$  and  $y$ , and
- the temperature parameter.

Intensification in SA is provided by the local-search nature of the algorithm, while diversification is produced with non-zero temperatures, by “uphill” moves. The pseudo-code of the SA metaheuristic can be represented as in Algorithm 2.1. In this pseudo-code,  $T_1, T_2, \dots$  represent an usually decreasing sequence of values for the temperature parameter, and “uniform[0, 1]” is a function for selecting a uniformly-distributed random number from the given interval, whereas the meaning of the rest of the symbols is as defined above.

SA has two parameters to control the search process: the initial temperature ( $t_{\text{initial}}$ ) and the total number of iterations or evaluations ( $it$ ). SA also allows to define the way the temperature is lowered during the annealing process.

SA has proved to be a solid optimization algorithm, capable of giving high-quality solutions to a wide scope of optimization problems [151]. Additionally, several works employ the SA metaheuristic for solving optimization problems in the area of radio network [20, 48, 56, 108, 142, 176]. A detailed description of SA and its applications can be found in [162].

## 2.4 Black-box optimization

The complexity of a problem is equivalent to the complexity of the best known algorithm solving that problem [156]. If there exists a polynomial-time algorithm to solve a problem,

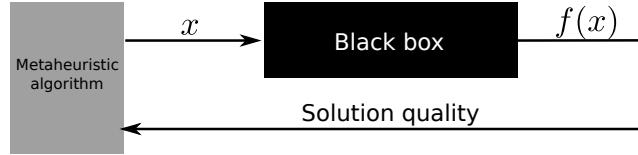


Figure 2.4: A metaheuristic algorithm process using a black box for the objective-function evaluation,  $f(x)$ , of a solution,  $x$ .

the problem is said to be easy or tractable. Similarly, if a problem is difficult or intractable, there is no known polynomial-time algorithm to solve it.

Many optimization problems cannot be formulated with a clear analytical mathematical notation. In such cases, the objective function may become a black box [88]. This is one of the main advantages when using metaheuristics, i.e., there is no need of a complete knowledge of the targeted model. Indeed, in a black-box optimization, no analytical formulation of the objective exists [156], as Figure 2.4 shows.

More specifically, a function  $f(\vec{x})$ ,  $\vec{x} \in \mathbb{R}^n$ , is a black-box function if and only if [156]:

- the domain  $\vec{x}$  is known,
- it is possible to get the value of  $f$  for each  $\vec{x}$  based on simulation, and
- there is no other information available for function  $f$ .

Typically, the experiments associated with this kind of problems are very expensive in terms of time and cost, since a simulation must be forced to evaluate the solution. Generally speaking, the most time-consuming part of a metaheuristic optimization process is the evaluation of the objective function [156]. This is especially true when dealing with real-world problems of areas such as structural design [13], molecular docking [157] and, the field on which this thesis focuses, radio-network design [19, 20, 21, 22? ]. A possible substitution for lengthy evaluations is to reduce their complexity by approximating the objective function, thus replacing it with an approximation during the optimization process. This approach is known as meta-modeling [156]. However, when dealing with approximations, some degree of solution quality is inevitably sacrificed. As it will be shown in the following chapters, there is a very fine balance between the number of evaluations and the quality of the achieved solutions. Consequently, reducing the time spent in objective-function evaluation should favorably influence the solution quality achieved by a preferred metaheuristic algorithm. A major portion of this thesis is dedicated to improve this specific aspect in the area of radio-network optimization, starting with a high-performance, unified framework for radio-network planning, which is presented in Chapter 6.

In practice, however, the black-box evaluation of the objective function presents a problem. In the context of research about radio-network optimization there is an inherent difficulty of providing the black-box used to evaluate a given approach. A quick review of the state-of-the-art in radio-network optimization indicates that this fact has become increasingly popular in several published works [7, 28, 30, 66, 144, 143]. Clearly, this fact creates a barrier to one of the most important phases of scientific methodology: experimental reproducibility [57].

There are several reasons behind this situation. For example, it is a known fact that proprietary software, providing good computational models for radio-network simulation, is a very expensive tool for science. Even neglecting the economical aspect, but considering the great variety of software packages and license combinations, it is practically impossible

for a research laboratory to have the whole palette of commercially-available solutions at its disposal. Moreover, genuine users of these applications are generally not allowed to mention the formats, protocols, or algorithms used by the proprietary software, since their disclosure is often explicitly forbidden by restrictive licenses.

Some non-commercial and open-source packages for radio-network simulation [82, 86, 109, 123, 124, 174] present two main drawbacks: poor documentation and/or low scalability. The scalability issues shown by some projects prevent these packages to be used in larger, real-world environments, where big problem instances are the rule. Despite this, the merit and acknowledgment go to the authors of these frameworks, not only for providing them to the scientific community, but fundamentally because of providing an environment in which different kinds of simulations are completely reproducible. Additionally, the lack of documentation represents a big hurdle for extending the base code, which becomes a difficult task without the help of the original authors of the package. In such cases, author's knowledge is required for effectively expanding the functionality of an open-source tool. This is especially true when dealing with complex simulation frameworks, as the ones used for radio networks.

### 3 Principles of Mobile Radio Networks

A cellular mobile radio network is a collection of individual cells that are served by several transmitters, called Base Stations (BSs). Each BS gives radio coverage to a small geographical area. The integration of the coverage of various BSs provides radio coverage over a much larger geographical area, thus defining a cellular radio network. The two basic functions of a radio-network system are:

- to locate and track both active and idle mobile devices, called User Equipments (UEs); and
- to attempt to connect each UE to the best available BS.

The first task involves a location-update procedure, which allows a UE to inform the network about its movement from one location area to the next. This process is called mobility management. The second task requires the constant evaluation of the radio-link quality with the serving BS, and the radio-link qualities of alternate BSs. This process is called radio-resource management, and is performed by the network using knowledge about the link-quality evaluations of the reference channels, e.g., the pilot channels.

The radio communications between a UE and a grid of BSs use low power. However, the movement of the UE causes highly irregular radio-link conditions, thus consistent monitoring and precise control are required to keep the radio-link quality at an acceptable level. At the core of the evaluation of radio-link quality is a statistical-measurement process based on a previous knowledge of the expected characteristics of the pilot channel. On the one hand, the link quality, and the size and distribution of the cells of a modern radio cellular system are limited by the speed of the link-quality measurement and network control. On the other hand, the spectral efficiency of a radio network is bounded by the cell sizes, the ability of radio links to withstand interference, and the ability of the system to react to variations in traffic [150].

Cellular radio systems partition the available spectrum among the BSs, and a given frequency is reused at the closest possible distance that the radio link will allow. Consequently, smaller cells have a shorter distance between reused frequencies, and this results in an increased spectral efficiency and traffic-carrying capacity. The radio links of a high-capacity radio network interfere with each other due to frequency reuse. For this reason, it is always desirable to use the lowest possible transmit power while maintaining each radio link above a given Quality-of-Service (QoS) threshold. Therefore, radio links should not significantly exceed their target QoS, since doing so will cause unnecessary interference to other radio links [150]. This particular situation is further discussed in Chapter 7, where an optimization approach that minimizes the total transmit power used in a radio network is presented.

#### 3.1 Handover

In radio networks, handover (or handoff) is one of the main features that allows the mobility of UEs [79]. The concept behind the handover operation is simple: when a UE moves from

the coverage area of a cell to the coverage area of a neighboring cell, the system creates a new connection with the latter cell and disconnects the user from the former one, so that an acceptable link quality can be maintained. Otherwise, the increase in transmit power that is required to compensate for path loss results in excessive interference. The handover procedure consists of two processes [150]:

- the link-quality evaluation before handover initiation, and
- the allocation of radio and network resources.

Generally speaking, radio networks with smaller cell sizes require faster and more reliable handover algorithms. Indeed, it has been shown that the number of cell-boundary crossings is inversely proportional to the cell size [93]. Since there is certain probability of dropping a connection whenever a handover is attempted, it is clear that the role of handover configuration becomes more important as the cell sizes decrease. Therefore, if the radio network does not detect poor signal quality fast enough, or makes too many handovers, the capacity is diminished due to increased interference and/or excessive control traffic [150].

### 3.1.1 Hard handover

During a hard handover, a UE can connect to only one BS at a time. A unique decision initiates and executes a handover without making a number of simultaneous connections among candidate BSs. Based on link measurements, the target BS is selected prior to executing the handover, and the active connection is instantly transferred to it. Moreover, the connection even experiences a brief interruption during the actual transfer, because the UE can only connect to one BS at a time. In contrast to soft handover (see next section), hard handovers do not take advantage of the diversity gain, where the signals from two or more BSs arrive at comparable strengths. Hard handover is a simple and inexpensive way to support UE mobility. It is used in Time-Division, Multiple-Access (TDMA) cellular systems such as GSM [150].

### 3.1.2 Soft handover

Soft Handover (SHO) enhances handover functionality by allowing a UE to potentially operate on multiple radio links at a time (see Figure 3.1). During SHO, the target BS is selected as the best candidate from among the available BSs. The UE performs the necessary link-quality measurements by monitoring the signals from the surrounding BSs. Simultaneously keeping multiple connections means that SHO enhances the system performance through diversity reception.

Despite the advantages it provides, SHO is complex and expensive to implement. Additionally, interference actually increases with SHO, since several BSs can connect to the same UE. This increase in forward interference can become a problem. If the handover region is large, such that there are many UEs in SHO mode, the increased interference due to several BSs connected to the same UE can become a problem [150].

The SHO procedure is important in systems using the channel-access method known as Code Division Multiple Access (CDMA), and especially in the Wideband CDMA (WCDMA), which is employed by the UMTS. CDMA systems are interference-limited meaning their capacities are closely related to the level of interference they can tolerate. Specifically, a CDMA system cell is affected by the interference within its own cell, and also interference from its neighboring cells. In order to mitigate the level of interference, and thus increase the capacity and quality, CDMA systems use power control. The main idea behind power control is to prevent the UEs and the BSs from transmitting more power than is strictly

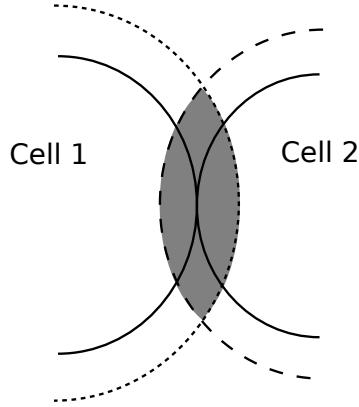


Figure 3.1: An example of the soft-handover region, showed as a grayed area. The area boundaries are configurable relative to the pilot-signal level of each cell. Adapted from [150].

necessary to meet the target QoS level. For the power control to work properly, the system must ensure that each UE is connected to the BS having the least attenuation at all times. If this is not the case, a positive-feedback problem appears, and it can potentially destabilize the entire system [170]. The SHO procedure helps to prevent such situations by ensuring that each UE is served by the best BS most of the time, i.e., by allowing connections to multiple BSs.

The SHO condition depends on the relative, received-signal quality from different cells and the SHO window, which triggers the addition of a cell to the active set of the UE. Depending on radio propagation characteristics, the radio transmission can gain more than 3 dB out of a SHO situation [79]. From this point of view, SHO is a method to reduce interference and improve radio quality, particularly at the cell border where the radio coverage is of inferior quality. In UMTS Release 99 [79], SHO is specified to work from the BS towards the UE (downlink), and from the UE towards the BS (uplink).

With the introduction of the High Speed Packet Access (HSPA) as an improvement of the performance existing in WCDMA protocols, the role SHO plays in mobile network configuration and functioning slightly changed. The key difference is that the High Speed Downlink Packet Access (HSDPA) does not support SHO, whereas the High Speed Uplink Packet Access (HSUPA) does. This particular distinction is further discussed in Chapter 8, because it has some important implications in the balanced distribution of SHO areas, and thus in the quality and capacity of HSPA services [80].



## 4 Overview of Radio-Network Optimization

Once a radio network is launched, an important part of its operation and maintenance deals with monitoring its quality characteristics. With the evolution of mobile communications, the complexity of network planning has grown along with its throughput capacity, thus making it practically impossible to plan modern radio networks with traditional methods. In this sense, an examination of colored coverage maps in conjunction with some statistical analysis are no longer appropriate tools for troubleshooting a network. Moreover, since real-world radio networks are large and many of their configuration parameters are interdependent, an engineer is not able to cope with the level of complexity present in these systems. For this reason, the computer, along with specialized software, guides the engineer to the most appropriate configuration for the network. In the context of this thesis, this process is referred to as radio-network optimization.

Radio-network optimization may be divided into two fundamental phases: analysis and decision [112]. The analysis phase consists of the examination of the network performance, which mainly focuses on the definition and collection of Key-Performance Indicators (KPIs). KPIs are quantifiable measurements that reflect different network-quality factors. The second phase deals with the decision making, based on the analytical results collected in the previous phase, about changing a particular configuration or parameter setting. The process, a representation of which is depicted in Figure 4.1, is repeated until the achieved results are acceptable. Notice the similarity with the general, decision-making process that was presented in Chapter 1, Figure 1.1.

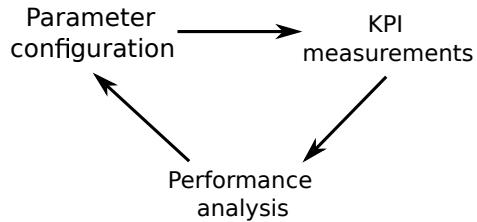


Figure 4.1: A typical optimization cycle for radio networks. This sequence is repeated until the achieved results are acceptable.

Since radio networks are increasingly more sophisticated, the need for optimization methods that are capable of dealing with greater complexity is far from declining. Indeed, several radio-network optimization problems were shown to be NP-hard, since the computational time grows non-polynomially with the problem size [6, 7, 67, 75, 97, 127, 143]. As described in [112], there are other reasons directly related with the growth of already deployed networks that also increase the need for optimization methods:

**Network performance improvement** more users receive service coverage with the same physical infrastructure, making parameter optimization the less expensive and only viable short-term approach.

**Changes in user profile** the introduction of new, faster services puts additional stress on the infrastructure, requiring additional optimization efforts.

**Changes in the propagation conditions** the allocation of different frequency bands for different systems, e.g., GSM, UMTS or LTE, requires the deployment of new BSs, the radio propagations of which behave differently, especially in urban areas.

Depending on the optimization problem being addressed, network operators define an optimization target that is represented by an objective function that maps possible configurations into a real value. Unfortunately, there is no universal objective function in the field of radio-network optimization [112]. However, it is possible to optimize for a different target at a time, such as service coverage, SHO balance or signal propagation. Particular optimization algorithms for solving these problems are presented in Chapters 7, 8 and 9, respectively. In all three cases, the introduced optimization approaches are performed “offline”, meaning that the optimization software is not an active part of the target radio network. As the feedback information of each optimization target, the statistical data about the network functioning is used.

Bellow, an overview of some well-known optimization problems for radio-mobile networks is given. Each section describes an optimization problem, and presents a short survey of recently proposed optimization methods for solving them.

## 4.1 Optimizing base-station locations

Some authors [76, 108] formulate the problem of locating BSs in terms of the minimum set-covering problem (see Figure 4.2). The set-covering problem is defined by considering the signal level in every test point from all BSs and requiring that at least one level is above a fixed threshold.

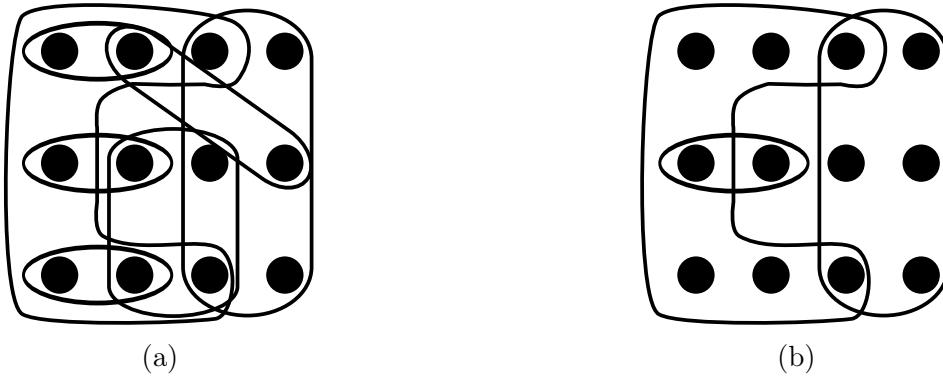


Figure 4.2: A graphical representation of the set-covering problem: (a) the problem input and (b) the solution.

A different formulation considers the BS-site location problem as a  $p$ -median problem [172], in which the BS location is the only decision variable considered. To each of the candidate solutions, an installation cost is also associated. The  $p$ -median problem constitutes seeking  $p$  different locations each time, regardless of how distant the sites are. The problem involves selecting one installation-candidate site from each region such that the traffic capacity and the size of the covered area of the network are maximized with the lowest installation cost.

## Related work

Aydin et al. [12] proposed a solution to the  $p$ -median problem based on three metaheuristic algorithms: a genetic algorithm, SA, and tabu search. Their experimental study focused on the performance comparison between the three approaches.

In [172], the authors also used a simplified  $p$ -median problem as the model. They presented the results of extensive simulations to compare the performance of three different metaheuristic algorithms.

A solution to the set covering problem is proposed by Hao et al. [76]. An implementation of SA was developed to solve the formulated combinatorial problem. The presented results showed the feasibility of the proposed approach.

Mathar and Niessen [108] proposed a hybrid method that combines a linear-programming approach with SA. The SA algorithm substituted linear programming whenever an exact solution was out of reach because of the complexity of the problem instance.

Amaldi et al. [6] presented a discussion about the computational results of two different heuristics: greedy search and tabu search. The problem formulation was based on a set of candidate sites where the BSs could be installed, an estimation of the traffic distribution and a propagation description of the area to be covered. Some years later, the same authors [7] extended the problem formulation by adding the BS configuration and the hardware characteristics as additional constraints of the integer program. In both works, they proposed a mixed, integer-programming model to maximize the trade-off between covered area and installation costs.

## 4.2 Optimizing antenna parameters

Since an antenna shapes the emitted energy, its configuration plays an important role in the coverage and interference of a radio network. The two most important parameters in this sense are the azimuth angle and the tilt (or elevation angle) of the antenna. The antenna azimuth, an example pattern of which is depicted in Figure 4.3, is the horizontal direction of the main antenna beam. The antenna tilt (see Figure 4.4) is defined as the angle of the main antenna beam relative to the horizontal plane. Both of these parameters have a great influence on network quality, although the antenna tilt usually requires less effort to implement, since most modern radio networks already support remote electrical tilt [11]. The adjustment of these two parameters optimizes some important aspects of the network, e.g.:

- the path loss between the BS and the UE,
- the interference between neighboring cells, which leads to an overall capacity increase of the network.

## Related work

One of the approaches proposed by Gerdenitsch et al. [60] involves an “ad-hoc” strategy for adjusting antenna azimuth and downtilt by analyzing the structure of the network. The objective of this optimization was to improve the number of served users in the target area.

Siomina and Yuan [144] proposed an approach for automated optimization of antenna azimuth and tilt, including support for both mechanical and electrical tilt. The implementation introduced a SA-based algorithm that searches the solution space of feasible antenna configurations. The goal of the optimization was to improve power sharing among different cell channels and ultimately improve the throughput of the network.

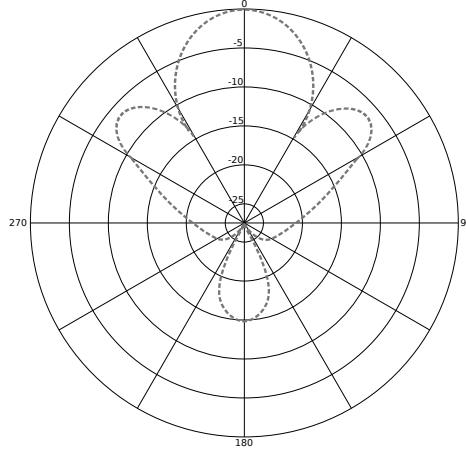


Figure 4.3: An example of an antenna-azimuth pattern, showing the gain in each direction.

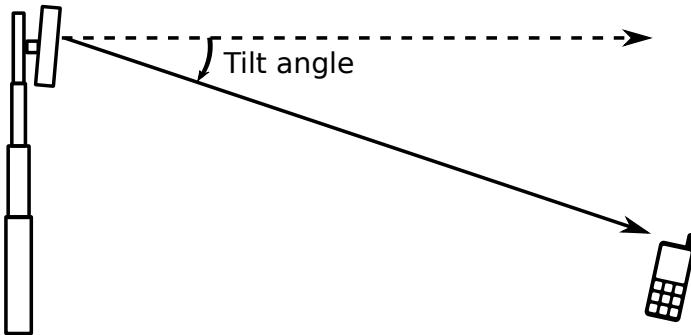


Figure 4.4: The antenna tilt, showing its angle with the horizontal plane. Adapted from [79].

In [66], the authors presented a compound optimization method containing two loops: the inner one and the outer one. The inner loop concentrated on frequency planning while the outer loop focused on finding the optimal settings of antenna azimuth and tilt for the current solution delivered by the inner loop. This approach is interesting because of its flexibility, e.g., the inner loop could be replaced with some other optimization objective, like service coverage.

In [49], the authors proposed an autonomous optimization approach for the antenna tilts. Based on the gradient-ascent method, the presented heuristic showed the fast convergence needed for an online-optimization method to be effective.

Combining the power setting of the pilot signal with the antenna configuration is also a common practice. For example, Siomina et al. [142] present an optimization approach for maximizing the service coverage that combines both antenna parameters with the power of the pilot signal. Their SA-based algorithm searches the solution space of possible configurations in order to improve the performance of the target radio network by reducing the total interference. The simulation results show the algorithm is capable of tackling large network instances.

Two optimization algorithms for finding an optimal setting of antenna tilt and pilot-signal power were also introduced in the previously cited work by Gerdenitsch et al. [60]. The first algorithm is based on a rule-based approach, while the second one extends it by incorporating SA. The evaluation of both techniques showed that the second algorithm achieves better solutions.

In a different work, Gerdenitsch et al. [59] proposed a genetic algorithm to tackle the optimization of antenna tilt and pilot-signal power, the goal of which is to increase network capacity. The implementation involves a deterministic fitness selection scheme, a problem-specific recombination operator and an improved mutation operator. After the initial identification of the best promising individuals, a local-optimization technique is applied to improve their fitness.

### 4.3 Optimizing coverage

Coverage is arguably the most common optimization objective considered in radio-network optimization. A general objective function for coverage optimization can be defined as follows:

$$f_{\text{cov}} = \frac{A_{\text{covered}}}{A_{\text{total}}},$$

where  $A_{\text{covered}}$  represents the area covered by the network and  $A_{\text{total}}$  represents the total area under optimization. Thus, the expression  $f_{\text{cov}}$  represents the proportion of the total area that is actually under network coverage, the value of which ranges from 0 (no coverage) to 1 (total coverage).

The area under optimization is usually divided into squares (or pixels) of a certain size, creating a Regular Square Grid (RSG) of a certain resolution. A pixel is considered covered if a given QoS measure is above a defined threshold. It is also common to use a binary function to check the coverage of a given pixel. The function returns 1 if the pixel located at  $(x, y)$  is covered, and 0 otherwise.

### Related work

Siomina and Yuan [143] considered the problem of minimizing the pilot-signal power subject to a coverage constraint. Their approach consists of an iterative mathematical program, based on a linear-integer formulation of the problem. The simulation results showed the trade-off between service coverage and power consumption for different test networks.

Almaldi et al. [7] investigated several mathematical-programming models for supporting decisions on where to install new BSs and how to select their configuration in order to find a trade-off between coverage and costs. The overall model takes into account signal-quality constraints, in both uplink and downlink directions, as well as the pilot-signal power.

Connecting UMTS and LTE-network optimization from the coverage point of view, Sallent et al. [132] presented a framework to automatically identify a cell with sub-optimal coverage. The framework input is a collection of KPIs, including usage statistics and UE field measurements. The authors used the experimental results to provide a projection for the LTE self-optimization functionality.

In [158], the authors introduced a sparse-sampling algorithm to optimize the coverage of LTE networks by means of antenna tilt. By applying a reinforcement-learning technique, their approach optimizes coverage without prior knowledge about the target network.

Parkinnen et al. [161] presented a gradient-descent method to minimize an objective function that combines some KPIs with the service coverage. The power of the pilot signal of the cells is periodically updated in order to improve the aforementioned function.

## 4.4 Summary

The variety of optimization problems described in the previous sections differ in many aspects, like implementation, running time and solution quality. Picking the right method for a given situation depends on the optimization task and the desired results. Since computational time is usually an important restriction, simpler and faster methods may be preferable.

Beside the convenience of a literature overview as presented here, it is very important to develop a feeling for the properties, advantages and drawbacks of the respective methods. Moreover, the recommendations of experts regarding the interpretation of the solutions and the feedback from everyday network operation are an essential input for establishing high-quality optimization methods.

Additionally, as noted for some of the cited works, hybrid methods, i.e., approaches that combine different optimization and/or search techniques, usually give competitive solutions to complex optimization problems. Therefore, a simple optimization algorithm can quickly find a subset of reasonable solutions, whereas a more complex method could be applied afterwards, to refine the search. Sometimes it may also be useful to apply a local search method at the end to find better solutions in the proximity of a current one.

## 5 Principles of GPU Programming

During the past few years, the computing industry has been delivering extra processing power in the form of parallel computing, i.e., more cores instead of higher-frequency rates. Concurrently to this situation, new hardware architectures have appeared. Among them is the programmable Graphics Processing Unit (GPU), the computing capacity of which is already showing a faster progress compared to CPUs. The peak performance of the latest GPU is several times higher than the latest CPU, but even more important is the trend [39].

The GPU market was almost exclusively driven by the games industry. This fact makes this commodity hardware cheaper than other alternatives for High-Performance Computing (HPC), e.g., the classic ‘mainframe’ servers. Moreover, the performance-per-watt of GPUs is an extra benefit over their raw performance [39].

In February 2007, the first real opportunity for using GPUs for general programming and scientific computing came about with the release of the Compute Unified Device Architecture (CUDA) [117].

In the following, the focus on the CUDA architecture. In particular, similar hardware produced by other vendors, e.g., AMD and Intel, as well as other programming frameworks like OpenCL [147], are based on the same technology paradigm of massively-parallel processors. For this reason, the principles presented in this section can be applied to current GPU technologies, in a vendor-independent manner.

### 5.1 CUDA

A GPU is effectively a large set of processor cores with the ability to directly address into a global memory. This makes it easier for developers to implement data-parallel kernels. The CUDA programming model [117] was created for developing applications for the GPU platform. A system within this model consists of a host, i.e., a traditional CPU, and one or more compute devices that are massively data-parallel co-processors, i.e., a GPU. The host code transfers data to and from the global memory of the GPU using CUDA-function calls. Each processor of a CUDA device supports the Single-Program Multiple-Data (SPMD) model [9], in which all concurrent threads are based on the same code, although they may not follow exactly the same path of execution.

A CUDA program consists of multiple parts that are executed on either the CPU or the GPU. For this reason, the CUDA programming model may be viewed as an environment to isolate program parts that are rich in data parallelism and thus should be executed on the GPU. The parallel parts of a CUDA program are implemented as kernels, i.e., functions written in the C language that describe the work of a single thread of execution. Several restrictions apply on kernel functions: there must be no recursion, no static variable declarations, and a non-variable number of arguments [130]. At run time, the kernel invocation is typically done on thousands of threads, which are organized within developer-defined groups called thread blocks. The threads can share data and synchronize their execution only within the block they belong to. Thread blocks are, in turn, organized in a higher-hierarchy level called grid. All threads within a given grid execute the same kernel function. During ex-

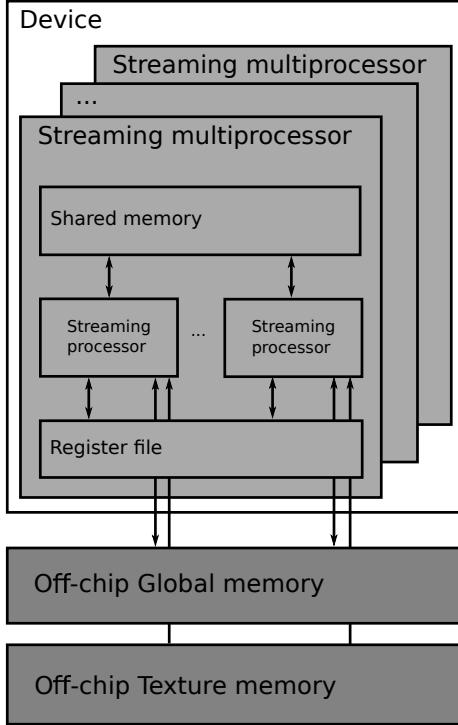


Figure 5.1: A simplified diagram of the architecture of a modern GPU. Adapted from [130].

ecution, continuous sections of threads within a block are grouped into warps of several parallel threads (typically 32 or 64, depending on the target architecture), which represent the multi-threading scheduling unit [130]. A GPU contains several Streaming Multiprocessors (SMs), each of which executes one instruction at a time for all the 32 threads in the warp. Consequently, if a thread block is not evenly divisible by the warp size, any remaining thread slots are wasted. Another undesirable situation appears when threads in a warp take different control paths, since the execution performance of the entire warp is reduced.

The architecture of a typical GPU is depicted in Figure 5.1. It consists of  $n$  SMs, each containing  $m$  Streaming Processors (SPs). Each SP executes a single instruction of a thread in a Single-Instruction, Multiple-Thread (SIMT) manner [39]. The registers of each SM are dynamically partitioned among the threads running on it.

Table 5.1 lists the different types of memory on a GPU, along with some of their properties. Specifically, the location of the memory, its hit-latency in terms of clock cycles, whether it is read-only or not, and the program scope it may be accessed from. Due to the different memory system a GPU has, the actual throughput an application can achieve depends on issues related to the access to memory, in particular the slow accesses to global memory from the GPU chip, and the use of shared memory in the SPs to mitigate the high-latency effects [39]. Since variables in the source code can be declared to reside in global, shared, local, or constant memory, a programmer has the means to organize the code in such a way that the application throughput can be maximized.

Clearly, there are hard limits to the memories, threads, and total bandwidth available to an application running on a GPU. Managing these limits is critical when optimizing applications, but applying strategies for avoiding one limit can easily cause other limits to be hit. Additionally, managing the behavior of threads so that those in the same warp follow the same control paths and load contiguous values from global memory can improve the execution performance [130].

A detailed discussion of the CUDA programming model can be found in [52].

Table 5.1: Properties of different memory types on a GeForce 8800 GPU. Adapted from [130].

Memory	Location	Latency	Read-only	Program scope
Global	off-chip	200-300 cycles	no	global
Texture	on-chip cache	>100 cycles	yes	global
Shared	on-chip	>1 cycle	no	function
Register	on-chip	~1 cycle	no	function

Table 5.2: Naming-convention translation between OpenCL and CUDA. Adapted from [91].

OpenCL	CUDA
Grid	Grid
Work group	Block
Work item	Thread
__kernel	__global__
__global	__device__
__local	__shared__
__private	__local__
imagend_t	texture<type,n,...>
barrier(L M F)	__syncthreads( )
get_local_id(0 1 2)	threadIdx.x y z
get_group_id(0 1 2)	blockIdx.x y z
get_global_id(0 1 2)	(not implemented)

## 5.2 OpenCL

The Open computing language (OpenCL) [147] is an open parallel computing API designed to enable GPUs and other co-processors to work together with the CPU, providing additional computing power. As a standard, OpenCL 1.0 was released in 2008, by The Khronos Group, an independent standards consortium [70].

The main advantage of OpenCL over CUDA relies on the fact that its source code can be compiled to run on a variety of hardware, including multicore CPUs and GPUs from different vendors. This provides a complete framework that is capable of exploiting the parallel features of different hardware without the need of changing the implementation. Moreover, being an open standard, its users are not tight to the decision of only one vendor. As it was mentioned before, the details described in the previous sections may be equally applied on CUDA and OpenCL.

One unfortunate consequence of the vendor variety is that CUDA and OpenCL have each introduced its own naming conventions. For the sake of consistency, in Table 5.2, a short “translation dictionary” between both platforms is presented.

In the remainder of this thesis, the naming convention introduced by the CUDA platform is used.

## 5.3 Summary

This chapter gave an overview of the basic concepts, the potential and the limitations of the GPUs when used as parallel processors for general programming. In particular, emphasis has been given to the details that differentiate this platform from CPUs. In this sense, it is important to recognize which applications can benefit from using a GPU and which not, also taking into consideration the effort required at the implementation time.



## 6 A Parallel Framework for Radio-Network Planning and Optimization

In this chapter, a parallel framework for radio-coverage simulation is presented. The objective of the framework is to provide an environment for the radio-coverage prediction of large radio networks. Due to its high performance, the framework also enables the evaluation of more complex optimization problems for radio networks.

The framework is implemented as a module of a geographical information system, since the prediction calculation employs digital elevation models and land-usage data. Following a master-worker parallel paradigm over a message-passing communication model proved to be a bottleneck for the performance of the parallel module. A new approach to overcome this performance constraint is introduced in this chapter. The efficiency improvement is based on overlapping process execution and communication. This minimizes the idle time of the worker processes and thus improves the overall efficiency of the system. To this end, the intermediate calculation results are saved into an external database (DB) instead of sending them back to the master process. This approach is implemented as part of a parallel radio-prediction tool (PRATO) for the open-source Geographic Resources Analysis Support System (GRASS) [114]. An extended analysis of the experimental results is provided, which are based on real data from an LTE network currently deployed in Slovenia. Based on these experiments, which were performed on a computer cluster, the new technique exhibits better scalability than the traditional master-worker approach. Some real-world data sets are presented, the coverage predictions of which are calculated in a shorter time while saturating the hardware utilization.

The content of this chapter extends the research work published by the author in [22]. The rest of this chapter is organized as follows. Section 6.1 describes the motivation behind the presented research, followed by an overview of the relevant publications in Section 6.2, describing how they relate to this work. Section 6.3 gives a description of the radio-coverage prediction problem, including the radio-propagation model. Section 6.4 concentrates on the design and implementation of the radio-propagation tool, for both the serial and parallel versions. Section 6.5 discusses the experimental results and their analysis.

### 6.1 Motivation

Although Gordon Moore's well-known and often cited prediction still holds [111], the fact is that for the past few years, CPU speeds have hardly been improving. Instead, the number of cores within a single CPU is increasing. This situation poses a challenge for software development in general and research in particular: a hardware upgrade will, most of the time, fail to double the serial-execution speed of its predecessor. However, since this commodity hardware is present in practically all modern desktop computers, it creates an opportunity for the parallel exploitation of these computing resources in order to enhance the performance of complex algorithms over large data sets. The challenge is thus to deliver the computing power of multi-core systems in order to tackle a computationally time-consuming problem, the completion of which is unfeasible using traditional serial approaches. Indeed, a performance

improvement opens new possibilities regarding the data sizes and accuracy a model may handle.

A traditional approach, when dealing with computationally expensive problem solving, is to simplify the models, thus reducing the time needed for their calculation. Clearly, this method increases the introduced error level, which is not an option for a certain group of simulations, e.g., those dealing with disaster-contingency planning and decision support [85, 175]. The conducted simulations during the planning phase of a radio network also belong to this group. The simulation results are the basis for the decision making prior to physically installing the BSs and antennas that will cover a certain geographical area. A larger deviation of these results increases the probability of making the wrong decisions at installation time, which may considerably increase the costs or even cause mobile-network operators to incur losses.

Various researchers have successfully deployed HPC systems and techniques to solve different problems dealing with spatial data [5, 8, 72, 85, 99, 119, 153, 154, 155, 169, 175]. Their work confirms that a parallel paradigm such as master-worker, and techniques like work pool (or task farming) and spatial-block partitioning are applicable when dealing with parallel implementations over large spatial data sets. However, it is well known that parallel programming and HPC often call for area experts in order to integrate them into a given environment [34]. Moreover, the wide range of options currently available creates even more barriers for general users wanting to benefit from HPC.

In this chapter, a high-performance, radio-planning framework for GSM (2G), UMTS (3G) and LTE (4G) networks is presented. Using a snapshot-based approach (see Section 7.3), the performance estimation of a radio network is divided into two parts, i.e., the coverage prediction and the performance analysis. During the coverage prediction, which is the subject of this chapter, path-loss matrices are created based on radio-propagation models, a network configuration, and digital maps of the target geographical area. Hence, in addition to a reliable radio-propagation model, also the resolution of the digital map should be high enough.

During the performance-analysis part, which is discussed in the following chapters, the predicted path losses are used for analyzing different phenomena, e.g., service availability and SHO performance. Therefore, the largest influence on the result performance of the framework comes from the calculated coverage predictions [36].

The framework is also suitable as a support tool for maintenance activities related to network troubleshooting in general and optimization in particular. Specifically, automatic radio-coverage optimization requires the evaluation of millions of radio-propagation predictions in order to find a good solution set. This is often unfeasible using other serial implementations of academic or commercial tools [82, 109, 124].

As a reference implementation, the publicly available radio-coverage prediction tool, presented in [82], was used. The authors developed a modular radio-coverage tool that performs separate calculations for radio-signal path loss and antenna-radiation patterns, also taking into account different configuration parameters, such as antenna tilting, azimuth and height. The output result, saved as a raster map, is the maximum signal level over the target area, in which each point represents the received signal from the best-serving transmitter (or cell). This work implements some well-known radio-propagation models, e.g., Okumura-Hata [77], the description of which is later presented in Section 6.3.2. Regarding the accuracy of the predicted values, the authors reported comparable results to those of an industrial tool [82]. To ensure that the presented implementation is completely compliant with this reference, a comparison test was designed that consists of running both tools with the same set of input

parameters. The test results from PRATO and the reference implementation were identical in all the tested cases.

## 6.2 Related work

There are few examples of radio-network simulators in the literature [82, 109, 123, 124, 133, 174]. Most of these tools were developed for academic research, thus not targeting industrial-sized environments.

A Matlab-based LTE simulator was proposed in [109]. It implements a standard LTE downlink physical layer, including Adaptive Modulation and Coding (AMC), multiple users, Multiple Input Multiple Output (MIMO) transmission and scheduler. Despite being open source and freely available, the fact of being implemented in Matlab makes it restrictive in terms of tackling big problem instances of real networks.

A promising tool from the performance point of view was presented in [124], where the authors implemented a full-stack LTE system in C++. Although the tool has no graphical capabilities for displaying the simulation progress or outcome, they might be included, since the source code is available. In our opinion, the main drawback of this tool is the lack of documentation, which makes it very difficult to continue extending this work without the direct help of some of the original authors.

As an extension to the well-known NS-2 network simulator, Filiposka and Trajanov [53] introduced a module for radio-propagation predictions, which takes the terrain profile into account. In this case, the authors focus on the relief, leaving out signal loss due to land-use, which is an important factor when targeting realistic radio-propagation scenarios.

The task-parallelization problem within the GRASS environment was addressed by several authors in a variety of studies. For example, in [27], the authors presented a collection of GRASS modules for a watershed analysis. Their work concentrates on different ways of slicing raster maps to take advantage of a Message Passing Interface (MPI) implementation.

In the field of HPC, the authors of [5] presented MPI and Ninf-G implementation examples of a GRASS raster module that processes vegetation indexes from satellite images. The authors acknowledge a limitation in the performance of their MPI implementation for big processing jobs. The restriction appears due to the computing nodes being fixed to a specific spatial range, since the input data are equally distributed among worker processes, creating an obstacle for load balancing in heterogeneous environments.

Using a master-worker technique, the work presented in [83] abstracts the GRASS data types into its own *struct* and MPI data types, thus not requiring the GRASS in the worker nodes. The data are evenly distributed by rows among the workers, with each one receiving an exclusive column extent to work on. Their test cluster contained heterogeneous hardware configurations. The authors noted that data-set size is bounded by the amount of memory on each of the nodes, since they allocate the memory for the whole map as part of the set-up stage, before starting the calculation. Regarding the data sets during the simulations, the largest one contained 3,265,110 points. They concluded that the data-set size should be large enough for the communication overhead to be hidden by the calculation time, so that the parallelization pays off.

In [153], the authors employed a master-worker approach, using one worker process per worker node. The complete exploitation of the computing resources of a single computing node is achieved with OpenMP. The experimental environment featured one host. The horizon-composition algorithm presents no calculation dependency among the spatial blocks. Consequently, the digital elevation model (DEM) was divided into separate blocks to be independently calculated by each worker process. The authors presented an improved algorithm

that can also be used to accelerate other applications like visibility maps. The tasks are dynamically assigned to idle processes using a task-farming paradigm over the MPI.

Similar to [153], in [154] there was no calculation dependency among the spatial blocks. The experimental evaluation was made over multiple cores of one CPU and a GPU, and a master-worker setup was used for process communication.

In [175], the authors presented a parallel framework for GIS integration. Based on the principle of spatial dependency, they lowered the calculation-processing time using a knowledge DB, delivering the heavy calculation load to the parallel back-end if a specific problem instance is not found in the DB. It is worth pointing out that the role of the external DB as a support structure to store knowledge models is essentially different than the one used in this chapter. In this work, the intermediate calculation results are effectively stored into an external DB, as it will be described later in Section 6.4.

An agent-based approach for simulating spatial interactions was presented in [65]. This technique decomposes the entire landscape into equally-sized regions, i.e., a spatial-block division as in [153], which are in turn processed by a different core of a multi-core CPU. This work used multi-core CPUs instead of a computing cluster.

Guan and Clarke [71] introduced a parallel programming library. Providing generic support for raster processing, the library is based on a cellular automaton model, thus distributing the parallel processing over different neighborhood-based regions. Similar to this work, PRATO also allows communication-computation overlap between the master and worker processes, but without the use of built-in, non-blocking operations, such as MPI\_ISend or MPI\_IRecv, as it will be explained later in Section 6.4.

Some years ago, grid computing received a lot of attention from the research community. It appeared to be a good alternative for accessing the extra computational power needed for the spatial analysis of large data sets [8, 166, 168]. However, several obstacles are still preventing this technology from being widely used. In particular, its adoption requires not only hardware and software compromises with respect to the involved parts, but also a behavioral change at the human level [8].

Parallel methods for radio-propagation predictions were presented by several authors [10, 101, 135]. Ray-tracing methods prevail among these works, since they exhibit a higher computational-time complexity. However, such methods are still inefficient in terms of the computational effort required to achieve satisfying results, particularly if applied to real-world networks containing thousands of transmitters.

## 6.3 Radio-coverage prediction for mobile networks

### 6.3.1 Background

As it was mentioned in Chapter 3, radio communications in a mobile network take place between a BS (or fixed transmitter) and a number of UEs (or mobile receivers). The effects of signal propagation limit the performance of a mobile-radio system. For this reason, the characterization and modeling of radio propagation is considered a fundamental aspect in radio-network planning [4]. Consequently, understanding the mathematical modeling of a frequency channel is necessary to accurately predict the system performance and to provide a mechanism to analyze the effects caused by signal propagation [122].

The coverage planning of radio networks is a key problem that all mobile operators have to deal with. Moreover, it has proven to be a fundamental issue, not only in LTE networks, but also in other standards for mobile communications [131, 136, 143, 160]. One of the primary objectives of mobile-network planning is to efficiently use the allocated frequency band to ensure that some geographical area of interest can be satisfactorily reached with the

BSs of the network. To this end, radio-coverage prediction tools are of great importance, as they allow network engineers to test different network configurations before physically implementing the changes. Radio-coverage prediction is a complex task, mainly due to the several combinations of hardware and configuration parameters that have to be analyzed in the context of different environments. The complexity of the problem means that radio-coverage predictions are computationally-intensive and time-consuming, hence the importance of using fast and accurate tools (see Section 6.4.4 for a complexity analysis of the algorithm). Additionally, since the number of deployed transmitters keeps growing with the adoption of modern standards [131], there is a clear need for a radio-propagation tool that is able to cope with larger work loads in a feasible amount of time (see Section 6.4.4 for the running time of the serial version).

### 6.3.2 Radio-propagation model

PRATO uses a modified version of the well-known Okumura-Hata model for radio-propagation predictions [77]. Other accurate methods exist, like the ones based on ray tracing [37, 165]. However, these methods are more sensible to deviations in input data, like DEMs and buildings, and are still inefficient in terms of the computational effort required to achieve satisfying results. Empirical methods for radio-propagation predictions, like Okumura-Hata, give acceptable results within a feasible amount of time. For this reason, they have become the industry standard for non-deterministic signal-propagation calculations [15, 33, 77, 136, 145].

The Okumura-Hata model has been largely studied and shown to be suitable for predicting the radio propagation in LTE networks [4, 136]. In its primary form, the model distinguishes the distance from the receiver to the transmitter, the frequency used and the effective antenna height, i.e., the height of the antenna above the receiver's level. These variables are taken into account in order to calculate the path loss in open-area (OA) conditions. Additionally, for distinguishing non-line-of-sight (NLOS) conditions, the terrain profile and Earth shape are added to the original formula. In this context, a NLOS situation appears when the first Fresnel zone is obscured by at least one obstacle. A Fresnel zone is defined as one of the concentric ellipsoids which define volumes in the radiation pattern of a circular aperture [171].

As it was mentioned before, the path-loss evaluation employed here is based on the Okumura-Hata model. When in an OA environment, the path loss is described as in Equation (6.1):

$$\begin{aligned} L_{OA}(x, y, \vec{\beta}) = & \beta_0 + \beta_1 \log(d_{(x,y)}) + \beta_2 \log(H_A) + \beta_3 \log(d_{(x,y)}) \log(H_A) - \\ & 3.2 [\log(11.75 \cdot H_R)]^2 + 44.49 \log(F) - 4.78 [\log(F)]^2, \end{aligned} \quad (6.1)$$

where  $\vec{\beta} = (\beta_0, \beta_1, \beta_2, \beta_3)$  is the vector containing the control parameters of the model,  $d_{(x,y)}$  is the distance (in kilometers) from the transmitter to the topography point with coordinates  $(x, y)$ ,  $H_A$  is the effective antenna height (in meters) of the transmitter,  $H_R$  is the antenna height (in meters) of the receiver, and  $F$  is the frequency, expressed in MHz. In NLOS conditions, the additional path loss is calculated as in Equation (6.2):

$$L_{NLOS}(x, y) = \sqrt{[\alpha K(x, y)]^2 + E(x, y)^2}, \quad (6.2)$$

where  $\alpha$  is the knife-edge diffraction<sup>1</sup> control parameter, the value of which is calculated based on the level of obstruction of a Fresnel zone,  $K(x, y)$  is the knife-edge diffraction loss (in dB), and  $E(x, y)$  is the correction due to the Earth sphere (in dB). All three values depend on the characteristics of the topography point with coordinates  $(x, y)$ .

In this work, as well as in [53], the terrain profile is used for line-of-sight (LOS) determination, i.e., an obstacle obstruction in the first Fresnel zone of the transmitter. In order to adequately predict signal-loss effects due to foliage, buildings and other fabricated structures, additional loss factors based on the land usage (clutter data), are included. This technique is adopted by several propagation models for radio networks [3, 15, 113]. Consequently, an extra term is introduced for signal loss due to clutter, thus defining the model-predicted path loss as:

$$L(x, y, \vec{\beta}) = L_{\text{OA}}(x, y, \vec{\beta}) + L_{\text{NLOS}}(x, y) + L_{\text{CLUT}}(x, y), \quad (6.3)$$

where  $L_{\text{CLUT}}(x, y)$  represents a look-up function for the clutter loss at the topography point with coordinates  $(x, y)$ , expressed in dB, which are given as input data.

## 6.4 Design and implementation

### 6.4.1 Geographic resources analysis support system

GRASS [114], a free and open-source software project that implements a Geographical Information System (GIS), is used as the software environment for PRATO. This GIS software was originally developed at the US Army Construction Engineering Research Laboratories and is a full-featured system with a wide range of analytical, data-management, and visualization capabilities. Currently, the development of GRASS GIS is supported by a growing community of volunteer developers.

The use of GRASS GIS as an environment for PRATO presents many advantages. First, the current development of GRASS is primarily Linux-based. Since the field of HPC is dominated by Linux and UNIX systems, an environment with Linux support is critical for this work. Software licensing is another important consideration for choosing GRASS, since it is licensed under the GNU Public License [146] that imposes the availability of the source code. This allows to make potential modifications to the system, thus adapting it for the parallel computation environment. Moreover, GRASS provides a great deal of built-in functionality, capable of operating with raster and vector topological data that can be stored in an internal format or a DB.

### 6.4.2 Multi-paradigm parallel programming

The implementation methodology adopted for PRATO follows a multi-paradigm, parallel programming approach in order to fully exploit the resources of each node in a computing cluster. This approach combines a master-worker paradigm with an external DB. To efficiently use a shared memory multi-processor on the worker side, and to effectively overlap the calculation and communication, PRATO uses POSIX threads [26].

To use the computing resources of a distributed memory system, such as a cluster of processors, PRATO uses the MPI [69]. The MPI is a message-passing standard that defines the syntax and semantics designed to function on a wide variety of hardware. The MPI enables multiple processes, running on different processors of a computer cluster, to communicate

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<sup>1</sup>Knife-edge diffraction occurs when a portion of the incident radiation reaches a well-defined obstacle, such as a mountain or the edge of a building.

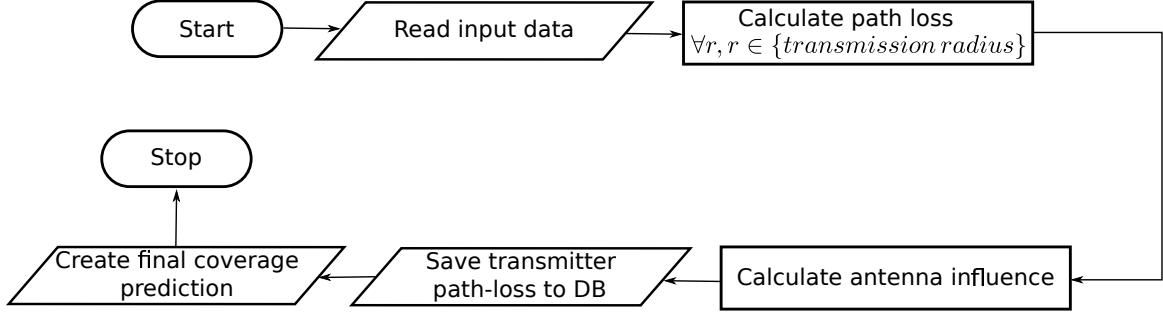


Figure 6.1: Flow diagram of the serial version.

with each other. It was designed for high performance on both massively parallel machines and on workstation clusters.

PRATO also supports the execution of the most computationally-intensive parts of the radio-propagation algorithm on a GPU. Moreover, the GPU hardware is used if it is available on the computing nodes that host the worker processes (see Section 6.4.5.3 later in this chapter for a discussion about the GPU implementation).

In order to make the text clearer and to differentiate between the programming paradigms used from here on, a POSIX thread will be referred to simply as a ‘thread’ and a MPI process as a ‘process’.

#### 6.4.3 Design of the serial version

This section describes the different functions contained in the serial version of PRATO, which is implemented as a GRASS module. Their connections and data flow are depicted in Figure 6.1, where the parallelograms of the flow diagram represent the input/output (I/O) operations.

The design follows a similar internal organization as the radio-planning tool presented in [82], but with some important differences. Specifically, the design presented here employs a direct connection to an external DB server for intermediate result saving, instead of the slow, built-in GRASS DB drivers. To explicitly avoid tight coupling with a specific DB vendor, the generated output is formatted in plain text, which is then forwarded to the DB. Any further processing is achieved by issuing a query over the DB tables that contain the path-loss results for each of the processed transmitters.

##### 6.4.3.1 Input parameters

All input data are read in the first step (see “Read input data” in Figure 6.1). Their formats differ based on the data they contain, i.e.:

- GRASS raster files are used for the DEM and clutter data, whereas
- a text file is used for the transmitter configurations and other service-dependent options.

Since the module accepts a considerable amount of input parameters, they are read from a text-based initialization (INI) file. This is far more practical than passing them as command-line parameters, which would make them error-prune and difficult to read. Besides, the INI file may contain configuration parameters for many transmitters. The user selects which one(s) to process at run-time by passing a command-line option.

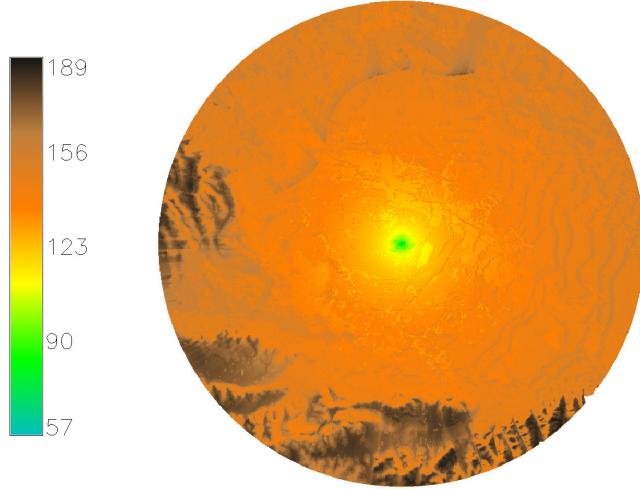


Figure 6.2: Example of a raster map, showing the result of a path-loss calculation from an isotropic antenna. The color scale is given in dB, indicating the path loss at a given point.

#### 6.4.3.2 Isotropic path-loss calculation

This step starts by calculating which receiver points,  $r$ , are within the specified transmission radius (see “*transmission radius*” in Figure 6.1). The transmission radius is defined around each transmitter in order to limit the radio-propagation calculation to a reasonable distance. For these points, the path loss for an isotropic source (or omni antenna) is calculated. This calculation is performed by applying the radio-propagation model, which was previously defined in Equation (6.3), to each of the points within the transmission radius around the transmitter (see “Calculate path loss” in Figure 6.1).

Figure 6.2 shows an example of the isotropic path-loss calculation, only including the map area within the transmission radius. The color scale is given in dB, indicating the signal loss from the isotropic source of the transmitter, located at the center. Notice the hilly terrain is clearly distinguished due to LOS and NLOS conditions from the signal source.

#### 6.4.3.3 Antenna diagram influence

This step considers the antenna radiation diagram of the current transmitter and its influence over the isotropic path-loss calculation (see “Calculate antenna influence” in Figure 6.1). Working on the in-memory results generated by the previous step, the radiation diagram of the antenna is taken into account, including the beam direction, the electrical and the mechanical tilt. Figure 6.3 shows the map area within the transmission radius, where this calculation step was applied to the results from Figure 6.2. Notice the distortion of the signal propagation that the antenna has introduced.

#### 6.4.3.4 Transmitter path-loss prediction

In this step, the path-loss prediction of the transmitter is saved in its own DB table (see “Save transmitter path-loss to DB” in Figure 6.1). This is accomplished by connecting the standard output of the GRASS module with the standard input of a DB client. Naturally, the generated plain text should be understood by the DB itself.

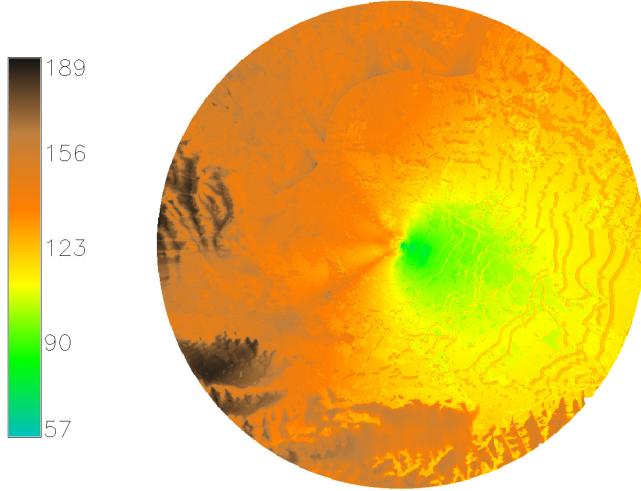


Figure 6.3: Example of a raster map, showing the influence of a directional antenna over the path-loss result depicted in Figure 6.2. The color scale is given in dB, indicating the path loss at a given point.

#### 6.4.3.5 Coverage prediction

The final radio-coverage prediction, containing the aggregation of the partial path-loss results of the involved transmitters, is created in this step (see “Create final coverage prediction” in Figure 6.1). The received signal strength from each of the transmitters is calculated as the difference between its transmit power and the path loss for the receiver’s corresponding position. This is done by executing a Structured Query Language (SQL) query over the tables containing the path-loss predictions of each of the processed transmitters. Finally, the output is generated using the GRASS built-in modules *v.in.ascii* and *v.to.rast*, which create a raster map using the query results as the input. The resulting raster map contains the maximum received signal strength for each individual point, as shown in Figure 6.4. In this case, the color scale is given in dBm, indicating the strongest received signal strength from the transmitters.

#### 6.4.4 Computational complexity

In this section, the time complexity of the radio-coverage prediction algorithm is presented, the pseudo-code of which is listed in Algorithm 6.1.

The algorithm starts by loading the input DEM and clutter data. Both RSGs should account for the same area and resolution, consequently containing the same number of pixels,  $M$ . The transmitter data is then loaded into set  $T$ , the cardinality of which is denoted as  $n = |T|$ . For each transmitter  $t \in T$ , a smaller subarea of the DEM and clutter data, denoted as  $DEM_t$  and  $Clut_t$ , respectively, is delimited around  $t$ , based on a given transmission radius. The number of pixels within this sub-area is denoted as  $m$ , and its value is the same for all  $t \in T$ . The visibility for an RSG pixel is computed using the *LineOfSight* function, by walking from the antenna of the transmitter to the given pixel element, along the elements intersected by a LOS, until either the visibility is blocked, or the target is reached [45]. Regarding the *PathLoss* function, whenever a receiver point is in NLOS, the walking path from the transmitter has to be inspected for obstacles, calculating the diffraction losses for each of them, i.e.,  $\alpha$  and  $K(d_{(x,y)})$  from Equation (6.2). Hence, its quadratic complexity,

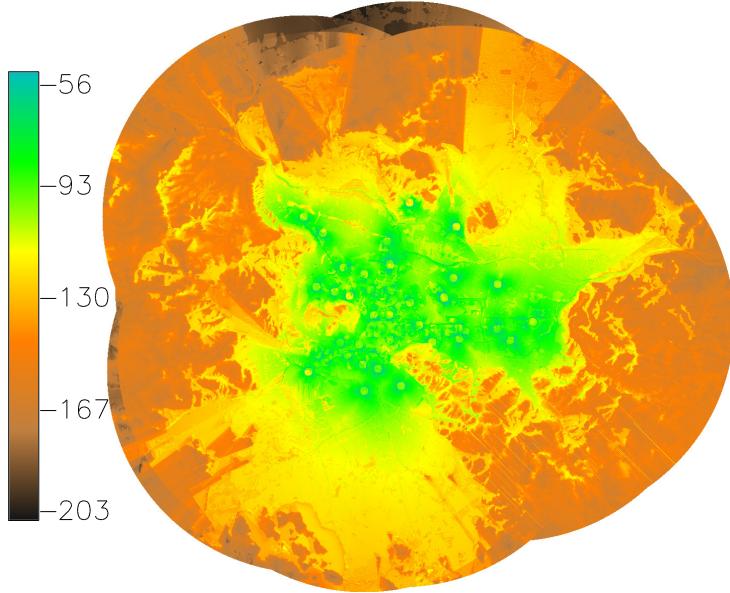


Figure 6.4: Example of a raster map, displaying the final coverage prediction of 136 transmitters over a geographical area. The color scale is given in dBm, indicating the received-signal strength. Darker colors denote areas with a reduced signal due to the fading effect of the hilly terrain and clutter.

---

**Algorithm 6.1** Pseudo-code of the radio-coverage prediction algorithm. The time complexity is given per line.

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```

 $DEM \leftarrow$  DEM of the whole area.                                 $\triangleright O(M)$ 
 $Clutter \leftarrow$  signal losses due to clutter of the whole area.     $\triangleright O(M)$ 
 $T \leftarrow$  transmitter configuration data.                           $\triangleright O(n)$ 
for all  $t \in T$  do                                          $\triangleright O(n \cdot m^2)$ 
     $DEM_t \leftarrow$  DEM area within transmission radius of  $t$            $\triangleright O(m)$ 
     $Clut_t \leftarrow$  Clutter area within transmission radius  $t$            $\triangleright O(m)$ 
     $LoS_t \leftarrow$  LineOfSight ( $DEM_t$ )                                      $\triangleright O(m^2)$ 
     $PL_t \leftarrow$  PathLoss ( $DEM_t, Clut_t, LoS_t$ )                       $\triangleright O(m^2)$ 
     $Diag_t \leftarrow$  Antenna diagram of  $t$                                   $\triangleright O(1)$ 
     $PL_t \leftarrow$  AntennaInfluence ( $Diag_t, PL_t$ )                       $\triangleright O(m)$ 
end for
for all  $t \in T$  do                                          $\triangleright O(n \cdot m)$ 
     $CoveragePrediction \leftarrow$  PathLossAggregation ( $t, PL_t$ )           $\triangleright O(m)$ 
end for
return  $CoveragePrediction$ 

```

---

which dominates the complexity of the algorithm, together with *LineOfSight*, resulting in an algorithmic complexity denoted by:

$$O(M + n \cdot m^2). \quad (6.4)$$

On the one hand,  $n$  will generally be many orders of magnitude smaller than  $m^2$ , although its computational-time complexity is relevant for practical use. For example, assuming the radio-coverage prediction for one transmitter completes in around 15 seconds using a serial implementation, the prediction for a mobile network comprising 10,240 transmitters would have an execution time of almost two days. On the other hand, when the input data correspond to a large geographical area, and only one transmitter is selected with a narrow calculation radius, i.e.,  $M$  is large,  $n = 1$  and  $m$  is small, then  $O(M)$  will dominate the time complexity of the algorithm.

#### 6.4.5 Design of the parallel version

The focus here is on the practical usability and performance of PRATO. The parallel implementation aims at overcoming the computational-time constraints that prevent a serial implementation from tackling bigger problem instances in a feasible amount of time.

A major drawback of the GRASS as a parallelization environment is that it is not thread-safe, meaning that concurrent changes to the same data set have an undefined behavior [24]. One technique to overcome this problem is to abstract the spatial data from the GRASS. For example, in [83], the authors achieved the GRASS abstraction by introducing a *Point* structure with four *double* attributes, where each pixel of the RSG is mapped to an instance of this structure. Another possibility is for one of the processes, e.g., the master, to read entire rows or columns of data before dispatching them for processing to the workers [5, 83]. In this case, an independence between row/column calculations is required, which is a problem-specific property. Here, abstraction from the GRASS is achieved by loading each spatial-data set into a separate 2D matrix of basic data-type elements, e.g., *float* or *double* depending on the desired accuracy. Each matrix is then assigned the geographical location of the closest corner to the origin of the map-projection system used, e.g., the lower-left corner for the transverse Mercator map projection over central Europe. It follows that the geographical location of any element within the matrix is calculated as the combination of the geographical location of the matrix and the offset of the target element (see Figure 6.5). The advantage of this technique is having the geographical location of a pixel readily available with a minimum memory footprint. Moreover, a convenient consequence of this abstraction schema is that worker processes are completely independent of the GRASS, thus significantly simplifying the deployment of the parallel implementation over multiple computing hosts.

In the area of geographical-information science, the master-worker paradigm has been successfully applied by several authors [5, 27, 72, 83, 153, 154]. However, this technique presents certain issues that prevent the full exploitation of the available computing resources when deployed over several networked computers. Specifically, the problem refers to network saturation and idle processes within the master-worker model. Generally speaking, a single communicating process, e.g., the master, is usually not able to saturate the network connection of a node. Using more than one MPI process per node might solve this problem, but possible rank-ordering problems may appear, thus restricting the full utilization of the network [126]. Additionally, the hardware-exploitation level is difficult to measure when the parallelization involves only one computing node [153, 154] (because no network communication is required), or only a few processes deployed over a handful of nodes [83].

Another issue appears when the master process executes the MPI code, in which case other processes sleep, making a serial use of the communication component of the system.

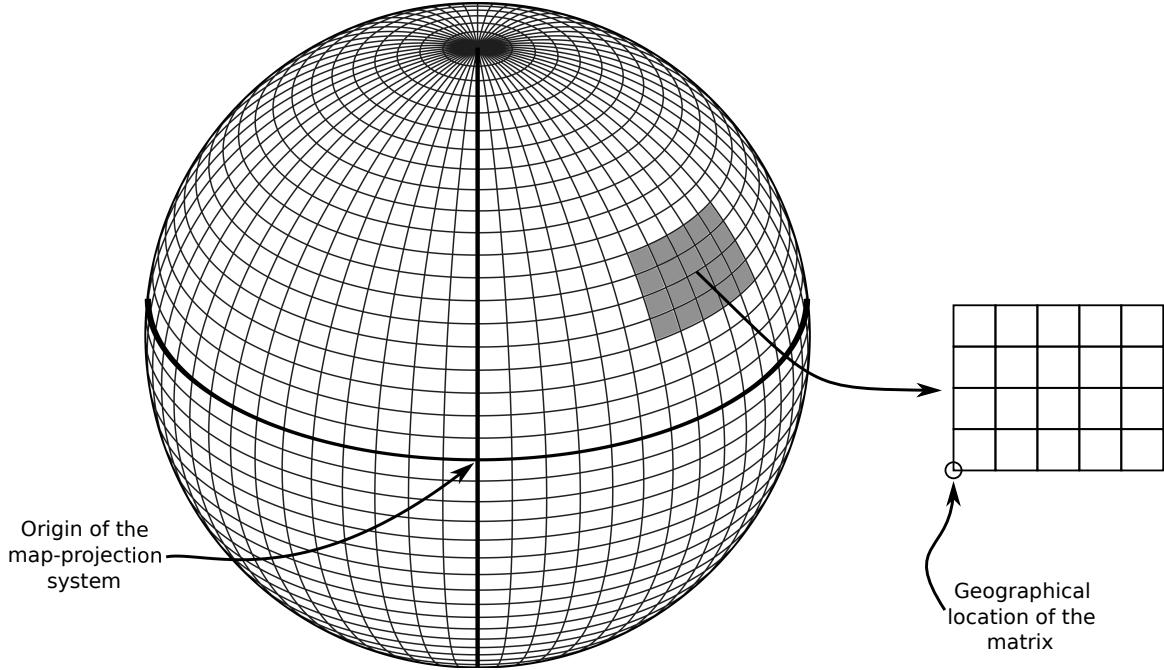


Figure 6.5: Example of a geographical-location mapping of the input-spatial data into a 2D matrix, the lower-left corner of which indicates the nearest point to the origin of the map-projection system.

Consequently, the master process becomes the bottleneck of the parallel implementation as the number of worker processes grows. This situation is also common when dealing with the metadata of a spatial region, which may relate to several elements of a RSG, making it a frequent cause of load imbalance [65, 78, 169]. In PRATO, the transmitter configuration and its antenna diagram represent metadata that are complementary to the sub-region that a transmitter covers.

Hybrid MPI-OpenMP implementations [153, 154], in which no MPI calls are issued inside the OpenMP-parallel regions, also fail to saturate the network [126]. A possible solution to this problem is to improve the communication overlap among the processes. To this end, PRATO features non-blocking point-to-point MPI operations, and an independent thread in the worker process that saves the intermediate results into a DB. One such DB system per computer cluster is used, which also serves the input data to the GRASS, in order to aggregate the partial results of the path-loss predictions or to visualize them. It is important to note that any kind of DB may be used, e.g., relational, distributed [121] or even those of the NoSQL type [148]. Nevertheless, a central, relational-DB system is used here, since they are the most popular and widely available ones. Additionally, the non-blocking message-passing technique used to distribute the work-load among the nodes provides support for heterogeneous environments. As a result, computing nodes featuring more capable hardware receive more work than those with weaker configurations, thus ensuring a better utilization of the available computing resources despite hardware diversity, and improved load balancing.

#### 6.4.5.1 Master process

The master process, the flow diagram of which is given in Figure 6.6, is the only component that runs within the GRASS environment. As soon as the master process starts, the input parameters are read. This step corresponds to “Read input data” in Figure 6.6, and it is carried out in a similar way as in the serial version. The next step delivers the metadata

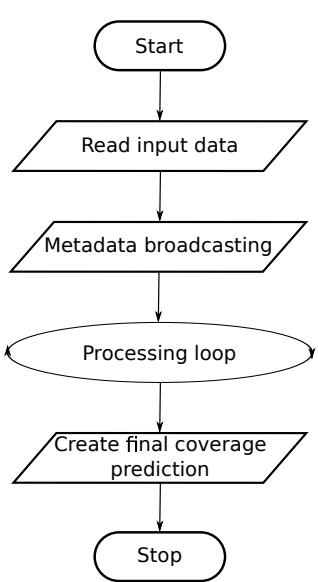


Figure 6.6: Flow diagram of the master process.

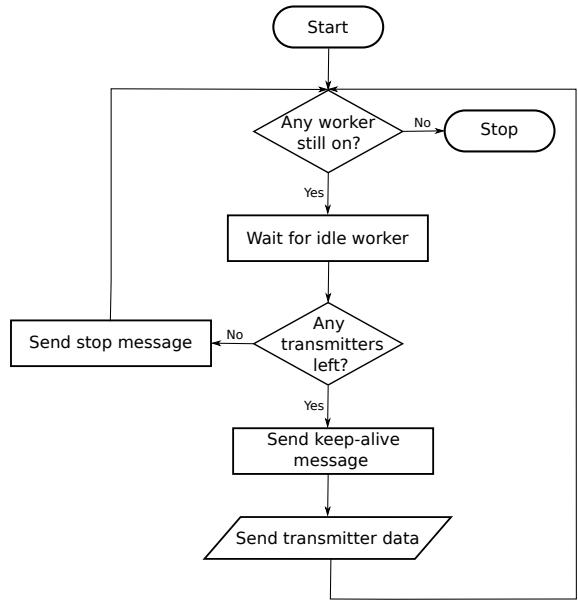


Figure 6.7: Flow diagram of the “Processing loop” step of the master process.

that is common to all the transmitters to all the processes (see “Metadata broadcasting” in Figure 6.6). Before distributing the work among the worker processes, the master process proceeds to decompose the loaded raster data into 2D matrices of basic-data-type elements, e.g., *float* or *double*, before dispatching them to the multiple worker processes. In this case, the decomposition applies to the DEM and the clutter data only, but it could be applied to any point-based data set. In the next step, the master process starts an asynchronous message-driven processing loop (see “Processing loop” in Figure 6.6), the main task of which is to assign and distribute the sub-region and configuration data of different transmitters among the idle worker processes.

The flow diagram shown in Figure 6.7 illustrates the “Processing loop” step of the master process. In the processing loop, the master process starts by checking the available worker processes that might calculate the radio-coverage prediction for the next transmitter. It is worth pointing out that this step also serves as a stopping condition for the processing loop itself (see “Any worker still on?” in Figure 6.7). The active worker processes inform the master process that they are ready to compute by sending an idle message (see “Wait for idle worker” in Figure 6.7). The master process then announces to the idle worker process that it is about to receive new data for the next calculation, and it dispatches the complete configuration of the transmitter to be processed (see “Send keep-alive message” and “Send transmitter data” steps, respectively, in Figure 6.7). This is only done in the case that there are transmitters for which the coverage prediction has yet to be calculated (see “Any transmitters left?” in Figure 6.7). The processing loop of the master process continues distributing the transmitter data among the worker processes, which asynchronously become idle as they finish the radio-prediction calculations they have been assigned by the master process. When there are no more transmitters left, all the worker processes announcing they are idle will receive a shutdown message from the master process, indicating to them that they should stop running (see “Send stop message” in Figure 6.7). The master process will keep doing this until all the worker processes have finished (see “Any worker still on?” in Figure 6.7), thus fulfilling the stopping condition of the processing loop.

Finally, the last step of the master process is devoted to creating the final output of the calculation, e.g., a raster map (see “Create final coverage prediction” in Figure 6.6). The final coverage prediction of all the transmitters is an aggregation from the individual path-loss results created by each of the worker processes during the “Processing loop” phase in Figure 6.6, which provides the source data for the final raster map. The aggregation of the individual path-loss results is accomplished by issuing an SQL query over the DB tables containing them, in a similar way as in the serial version.

#### 6.4.5.2 Worker processes on CPU

An essential characteristic of the worker processes is that they are completely independent of the GRASS, i.e., they do not have to run within the GRASS environment nor use any of the GRASS libraries to work. This aspect significantly simplifies the deployment phase to run PRATO on a computer cluster, since no GRASS installation is needed on the computing nodes hosting the worker processes.

One possibility to overcome the thread-safety limitation of the GRASS is to save the transmitter path-loss predictions through the master process, thus avoiding concurrent access. However, for the workers to send intermediate results back to the master process, e.g., as in [5, 83], is a major bottleneck for the scalability of a parallel implementation. In such case, the scalability is limited by the master process, because it must serially process the received results in order to avoid inconsistencies due to concurrent access. Instead, PRATO allows each of the worker processes to output its intermediate results into a DB, i.e., each path-loss prediction in its own table. Additionally, worker processes do this from an independent thread, which runs concurrently with the calculation of the next transmitter received from the master process. In this way, the overlap between the calculation and communication significantly hides the latency created by the result-dumping task, thus making better use of the available system resources.

The computations of the worker processes, the flow diagram of which is given in Figure 6.8, begin by receiving metadata about the transmitters and the geographical area from the master process during the initialization time (see “Receive broadcasted metadata” in Figure 6.8).

After the broadcasted metadata are received by all the worker processes, each one proceeds to inform the master process that it is ready (i.e., in an idle state) to receive the transmitter-configuration data that define which transmitter path-loss prediction to perform (see “Send idle message” in Figure 6.8). If the master process does not give the instruction to stop processing (see “Has stop message arrived?” in Figure 6.8), the worker process collects the sub-region spatial data and the transmitter configuration (see “Receive transmitter data” in Figure 6.8). In the event that a stop message is received, the worker process will wait for any result-dumping thread to finish (see “Wait for result-dump thread” in Figure 6.8) before shutting down. The coverage calculation itself follows a similar design as the serial version (see “Coverage calculation” in Figure 6.8).

As mentioned before, the worker process launches an independent thread to save the path-loss prediction of the target transmitter into a DB table (see “Threaded save path-loss to DB” in Figure 6.8). It is important to note that there is no possibility of data inconsistency due to the saving task being executed inside a thread, since path-loss data from different workers belong to different transmitters and are, at this point of the process, mutually exclusive.

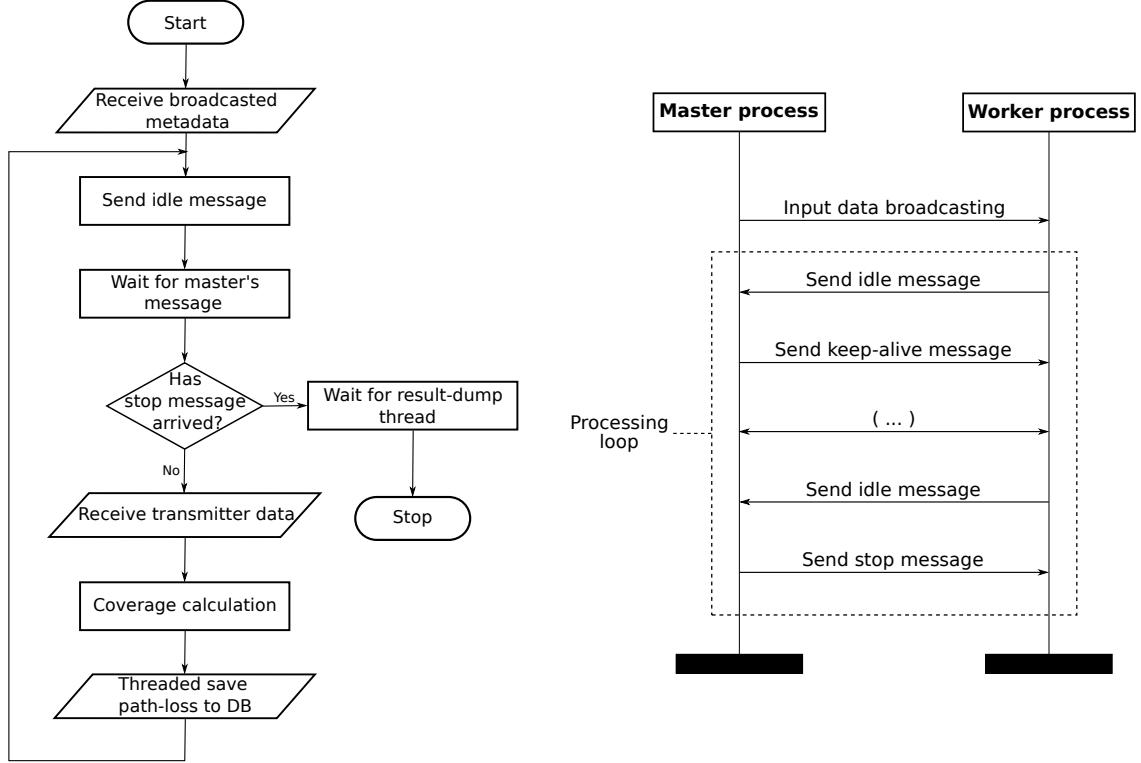


Figure 6.8: Flow diagram of a worker process.

Figure 6.9: Communication diagram, showing the message passing between the master and a worker process.

#### 6.4.5.3 Worker processes on GPU

PRATO provides multi-GPU support for improving the execution performance on the computing nodes hosting the worker processes. The algorithmic adaptation from a CPU to a GPU is not a trivial task. This section focuses on the main modifications made to the radio-propagation algorithm in order for it to work on GPU hardware.

It is well known that the bandwidth of the PCI Express bus can cause a throughput bottleneck when a significant amount of data is transferred between a CPU and a GPU in a heterogeneous system [68]. Some researchers acknowledged that unless a full working set of data can fit into the memory on a GPU, the PCI Express will become the bottleneck of the system [120, 134]. For this reason, it is imperative to have as much data as possible allocated on the GPU itself.

In order to minimize the CPU-to-GPU memory transfers, the spatial data used by the radio-propagation algorithm was organized as explained in Section 6.4.5, i.e., using geographically-located, offset-based 2D matrices. However, the internal representation of the matrix elements was changed to use less memory. To this end, the clutter-matrix elements were represented as *unsigned char*, since they express signal loss in dB. It follows that for the radio-propagation prediction of one transmitter, the following matrices should be allocated on the GPU:

- one 2D matrix containing DEM data for the target sub-region, the elements of which are *float* or *double*;
- one 2D matrix containing clutter data for the target sub region, the elements of which are *unsigned char*; and

- one 2D matrix containing the resulting path-loss prediction.

The dimension of all matrices is based on the transmission radius, within which the radio-propagation prediction should be calculated. The contents of the DEM and clutter matrices is constant throughout the calculation process. For this reason they were kept in texture memory to take advantage of the faster access time (see Section 5.1). Regarding the resulting path-loss matrix, each step of the radio-prediction algorithm is applied over the results of the previous step (see Figure 6.1 for a flow diagram of the steps involved), thus avoiding extra allocation on global memory or a data-transfer from/to the CPU.

#### 6.4.5.4 Master-worker communication

Similar to [153, 154], the message-passing technique used in this work enables a better use of the available computing resources, both in terms of scalability and load balancing, while introducing a negligible overhead. This last point is supported by the experimental results, introduced in Section 6.5.3.

The first reason to implement the message-passing technique is to support heterogeneous computing environments. In particular, our approach focuses on taking full advantage of the hardware of each computing node, thus explicitly avoiding the bottlenecks introduced by the slowest computing node in the cluster. This problem appears when evenly distributing the data among the worker processes on disparate hardware, e.g., as in [5, 83], being more noticeable with a larger number of computing nodes and processes. In other words, computing nodes that deliver better performance have more calculations assigned to them. Moreover, in real-world scenarios, it is often the case that a large number of dedicated computing nodes featuring exactly the same configuration is difficult to find, i.e., not every organization owns a computer cluster.

A second reason for selecting a message-passing technique is related to the flexibility it provides for load balancing, which is of greater importance when dealing with extra data or information besides spatial data [78]. This can be seen in Figure 6.7, where the master process, before delivering the spatial subset and transmitter-configuration data, sends a message to the worker process, indicating that it is about to receive more work. This *a priori* meaningless message plays a key role in correctly supporting the asynchronous process communication. Notice that the subset of spatial data that a worker process receives is directly related to the transmitter for which the prediction will be calculated. Similar to [153, 154], this problem-specific property enables the use of a data-decomposition technique based on a block partition of spatial data, e.g., the DEM and clutter data.

In general, there are many different ways a parallel program can be executed, because the steps from the different processes can be interleaved in various ways and a process can make non-deterministic choices [139], which may lead to situations such as race conditions [35] and deadlocks. A deadlock occurs whenever two or more running processes are waiting for each other to finish, and thus neither ever does. To prevent PRATO from deadlocking, message sending and receiving should be paired, i.e., an equal number of send and receive messages on the master and worker sides [139]. Figure 6.9 depicts the master-worker message passing, from which the transmitter-data transmission has been excluded for clarity. Notice how each idle message sent from the worker process is paired with an answer from the master process, whether it is a keep-alive or a stop message.

## 6.5 Simulations

Considering the large computational power needed for predicting the radio-coverage of a real radio network, the use of a computer cluster is recommended. A computer cluster

is a group of interconnected computers that work together as a single system. Computer clusters typically consist of several commodity PCs connected through a high-speed local-area network (LAN) with a distributed file system, like the network file system (NFS) [138]. One such system is the DEGIMA cluster [73] at the Nagasaki Advanced Computing Center (NACC) of the Nagasaki University in Japan. This system ranked in the TOP 500 list of supercomputers until June 2012<sup>1</sup>, and in June 2011 it held the third place in the Green 500 list<sup>2</sup> as one of the most energy-efficient supercomputers in the world.

This section presents the simulations, and an exhaustive analysis of the performance and scalability of the parallel implementation of PRATO. The most common usage case for PRATO is to perform a radio-coverage prediction for multiple transmitters. Therefore, a straight-forward parallel decomposition is to divide a given problem instance by transmitter, for which each coverage prediction is calculated by a separate worker process.

The following simulations were carried out on 34 computing nodes of the DEGIMA cluster. The computing nodes are connected by a LAN, over a Gigabit Ethernet interconnect. As mentioned before, the reason for using a high-end computer cluster such as DEGIMA is to explore, by experimentation, the advantages and drawbacks of the considered methods. However, this does not imply any loss of generality if applying these principles over a different group of networked computers that do not operate as a computer cluster. Moreover, PRATO also supports parallel calculation of radio-propagation predictions for multiple cells by distributing the processes among the individual cores of a single CPU.

Each computing node of DEGIMA features one of two possible configurations, namely:

- Intel Core i5-2500T quad-core processor CPU, clocked at 2.30 GHz, with 16 GB of RAM; and
- Intel Core i7-2600K quad-core processor CPU, clocked at 3.40 GHz, also with 16 GB of RAM.

During the simulation runs, the nodes equipped with the Intel i5 CPU hosted the worker processes, whereas the master process and the PostgreSQL DB server (version 9.1.4) each run on a different computing node, featuring an Intel i7 CPU. The DB server performed all its I/O operations on the local file system, which was mounted on an 8 GB RAM disk. During the simulations, the path-loss predictions of 5,120 transmitters occupied less than 4 GB of this partition. No GPU hardware was used for the following simulation sets.

A 64-bit Linux operating system (Fedora distribution) was the operating system used. The message-passing implementation OpenMPI, version 1.6.1, was manually compiled with the distribution-supplied *gcc* compiler, version 4.4.4.

### 6.5.1 Test networks

The parallel performance of PRATO was tested with real radio networks of different sizes. In order to create the synthetic test data sets with an arbitrary number of transmitters, a group of 2,000 transmitters of a real network was used. The configuration parameters of these transmitters resembled those of the LTE network deployed in Slovenia by Telekom Slovenije, d.d., which were, in turn, randomly replicated and distributed over the whole target area. The path-loss predictions were calculated using the radio-propagation model introduced in Section 6.3.2. The DEM area, as well as the clutter data, covered 20,270 km<sup>2</sup> with a pixel resolution of 25 m<sup>2</sup>. The clutter data contained different levels of signal loss due to land usage. For all the points within a radius of 20 km around each transmitter, a receiver positioned 1.5 m above the ground was assumed, and the frequency was set to 1,843 MHz.

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<sup>1</sup><http://www.top500.org>

<sup>2</sup><http://www.green500.org>

Table 6.1: Running-time gain (in percent) of the simulations for the weak-scalability of the MWD setup relative to the classic MW approach.

TX/core	Number of cores						
	1	2	4	8	16	32	64
5	-11.39	-10.42	-11.14	-0.95	11.75	26.15	32.53
10	-5.84	-7.78	-7.67	0.91	12.81	33.28	33.55
20	-8.59	-10.88	-1.04	1.95	14.29	35.23	35.27
40	-5.26	-6.90	-3.68	-0.67	17.27	36.23	36.65
80	-5.29	-7.11	-3.20	-0.31	17.94	36.32	36.57

### 6.5.2 Weak scalability

The weak-scalability experiments are meant to analyze the scalability of the parallel implementation in cases where the workload assigned to each process (one MPI process per processor core) remains constant as the number of processor cores is increased. It follows that the number of transmitters deployed over the target area is directly proportional to the number of processor cores and worker processes. This was accomplished by assigning a constant number of transmitters per core, while increasing the number of cores hosting the worker processes. Here, the following numbers of transmitters per worker/core were tested  $\{5, 10, 20, 40, 80\}$ , by progressively doubling the number of worker processes from 1 to 64.

Problems that are particularly well-suited for parallel computing exhibit computational costs that are linearly dependent on the size of the problem. This property, also referred to as algorithmic scalability, means that proportionally increasing both the problem size and the number of cores results in a roughly constant time to solution.

The master-worker (MW) configuration performs result aggregation continuously, i.e., while receiving the intermediate results from the worker processes. In contrast, the master-worker-DB (MWD) setup performs the result aggregation as the final step. This set of experiments is meant to investigate how the proposed MWD technique compares with the classic MW approach in terms of scalability when dealing with a constant computational load per core.

An important fact about the presented simulations when using multi-threaded implementations is to avoid oversubscribing a computing node. For example, if deploying four worker processes over a quad-core CPU, the extra threads will have a counter effect on the parallel efficiency, since the CPU resources would be exhausted, which slows the whole process down. For this reason, we have deployed three worker processes per computing node, leaving one core free for executing the extra threads.

## Results

The results represent the average running time out of a set of 20 independent simulation runs, for which the transmitters and the rank ordering of the worker processes were randomly selected. The collected running times for the weak-scalability experiments are shown in Figure 6.10. All the measurements express wall-clock times in seconds for each setup and problem instance, defined as the number of transmitters per process (Tx/process). The wall-clock time represents the real time that elapses from the start of the master process to its end, including the time that passes while waiting for the resources to become available. The running-time improvements of the MWD versus the MW setup are shown in Table 6.1.

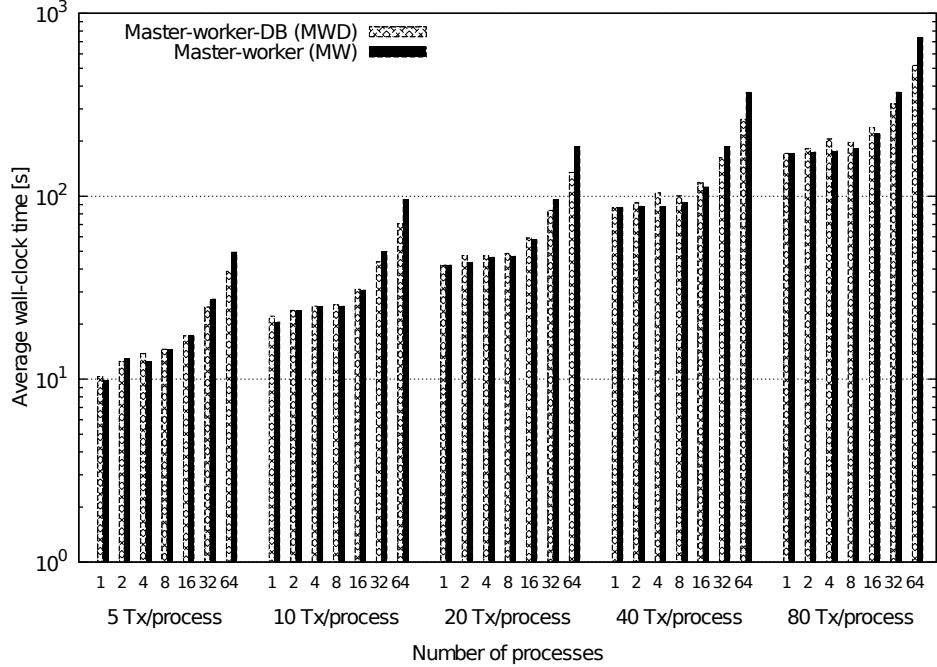


Figure 6.10: Measured wall-clock time for weak-scalability experiments, featuring MW and MWD setups. Experiments allocated one MPI worker process per core.

The time measurements observed from the weak-scalability results show that the classic MW approach performs well for up to four worker processes. When using eight worker processes, the MW setup is practically equivalent to the MWD approach, indicating that the master process is being fully exploited. When increasing the problem size and the number of worker processes to 16, the running-time gain is already clear, favoring the MWD configuration. This gain keeps growing, although slower, as we increase the number of worker processes to 32 and 64, confirming the hypothesis that in a classic MW approach, the parallel efficiency is bounded by the capacity of the master process to serve an increasing number of worker processes. Interestingly, the gain when using 32 and 64 worker processes is almost the same. After further investigation, the reason for this behavior was found: the new bottleneck was the LAN being completely saturated by the worker processes. Consequently, they have to wait for the network resources to become available before sending or receiving data, which is not the case when running the MW setup. Therefore, using the MWD approach a hardware constraint is hit, meaning that the bottleneck is no longer at the implementation level. Moreover, since the master process is far from overloaded when serving 64 worker processes, it can be expected that the MWD approach will keep scaling if a faster network infrastructure is used, e.g., 10-gigabit Ethernet or InfiniBand.

Certainly, the parallel version of PRATO scales better using the MWD approach, when challenged with a large number of transmitters (5,120 for the biggest instance) over 64 cores. This fact shows PRATO would be able to calculate the radio-coverage prediction for real networks in a feasible amount of time, since many operational radio networks have already deployed a comparable number of transmitters, e.g., the 3G network within the Greater London Authority area, in the UK [118]. For a more in-depth discussion and experimentation about real-world planning scenarios, see Chapter 10.

Not being able to achieve perfect weak scalability using the MWD setup is due to a number of factors. Specifically, the overhead time of the serial sections of the parallel process grows proportionally with the number of cores, e.g., aggregation of the intermediate

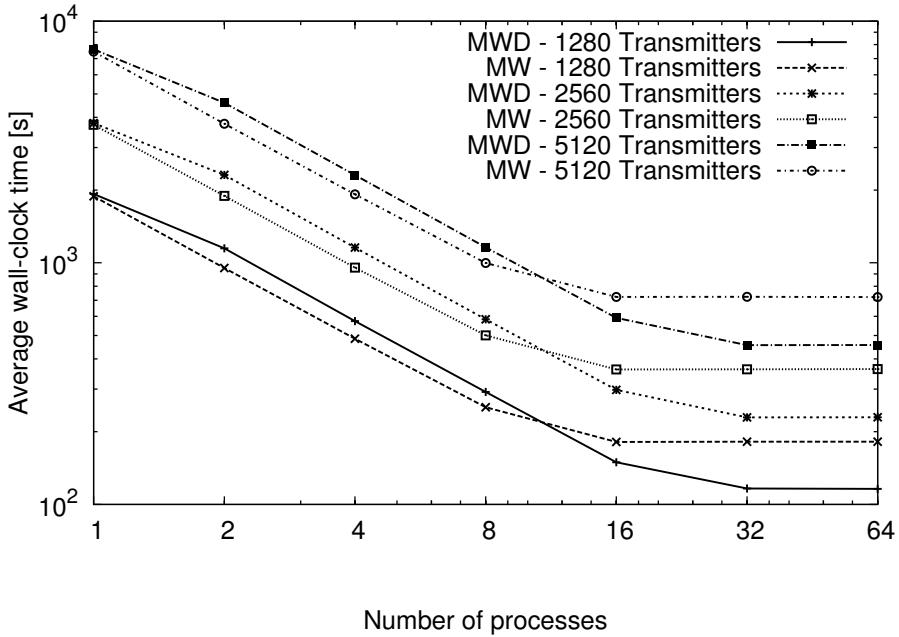


Figure 6.11: Measured wall-clock time for strong-scalability experiments, featuring MW and MWD setups. Experiments assigned one MPI worker process per core.

results, although the total contribution of this overhead remains low for large problem sizes. Moreover, the communication overhead grows linearly with the number of cores used. Consequently, the findings of Huang et al. [83] can be confirmed, who concluded that the data-set size should be large enough for the communication overhead to be hidden by the calculation time. This ensures profitable parallelization in terms of running-time reduction.

### 6.5.3 Strong scalability

This set of simulations is meant to analyze the impact of increasing the number of computing cores for a given problem size, i.e., the number of transmitters deployed over the target area does not change, but only the number of worker processes used is increased. Here, the following number of transmitters were tested  $\{1,280, 2,560, 5,120\}$ , by gradually doubling the number of workers from 1 to 64 for each problem size.

## Results

Similar to the weak-scalability experiments, the time measurements plotted in Figure 6.11 show that, when applying a classic MW approach, the running-time reduction starts flattening if more than eight worker processes were used. Moreover, the running times for 16, 32 and 64 worker processes are the same, i.e., they do not improve due to the master process being saturated. In contrast, when using the proposed MWD technique, the running-time reduction improves for up to 32 worker processes, after which there is no further improvement since the network was being fully exploited. These results clearly show that when applying parallelization using a larger number of worker processes, the master process becomes the bottleneck of the MW approach. When using the MWD configuration, a steady running-time reduction is observed, until a hardware constraint is hit, e.g., the network infrastructure.

The overhead of sending/receiving asynchronous messages in order to support heterogeneous systems was also measured. It was found that this overhead never exceeds 0.02 % of the total running time for the MW experiments, and 0.01 % for the MWD experimental set.

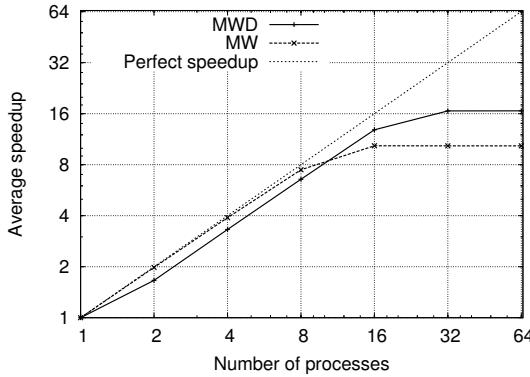


Figure 6.12: Average speedup for the strong-scalability experiments.

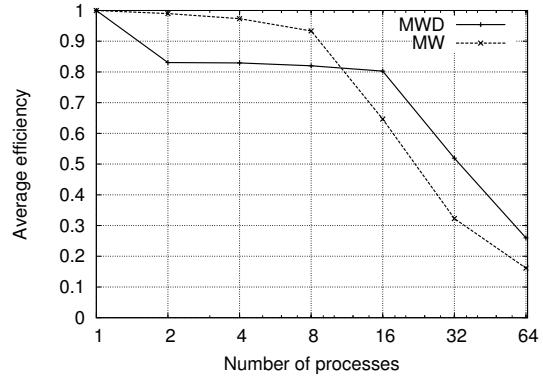


Figure 6.13: Average parallel efficiency for the strong-scalability experiments.

### 6.5.3.1 Speedup

In order to further analyze how well the PRATO scales using the MW and MWD approaches, the performance of the parallel implementation in terms of its speedup was measured, which is defined as:

$$S(NP) = \frac{\text{execution time for base case}}{\text{execution time for } NP \text{ cores}}, \quad (6.5)$$

where  $NP$  is the number of cores executing the worker processes. The parallel implementation running on only one core was the base case for comparisons. The serial implementation is not a good base comparison for the parallel results as it does not reuse the resources between each transmitter-coverage calculation and it does not overlap the I/O operations with the transmitter computations. In practice, this means that several concatenated runs of the serial version would be considerably slower than the single-worker configuration.

Using the speedup metric, linear scaling is achieved when the obtained speedup is equal to the total number of processors used. However, it should be noted that a perfect speedup is almost never achieved, due to the existence of serial stages within an algorithm and the communication overhead of the parallel implementation.

Figure 6.12 shows the average speedup of the parallel implementation for up to 64 worker processes, using the standard MW method and the proposed MWD approach. The average speedup was calculated for the three different problem instances, i.e., 1,280, 2,560, and 5,120 transmitters deployed over the target area. The number of transmitters used in these problem sizes is comparable to several real-world radio networks that were already deployed in England, e.g., Hampshire County with 227 BSs, West Midlands with 414 BSs, and Greater London Authority with 1,086 BSs [118]. Note that it is common for a single base station to host multiple transmitters.

The plotted average speedup clearly shows the minimal overhead of the MWD approach when using a small number of worker processes. This overhead accounts for the final aggregation of the intermediate results at the DB, which in the MW configuration is performed along worker processing. Like before, the DB component allows the parallel implementation to fully exploit the available computing resources when deploying a larger number of worker processes, until the network-capacity limit is met. Of course, these results are directly correlated with the wall-clock times shown in Figure 6.11.

### 6.5.3.2 Efficiency

Another measure to study how well PRATO utilizes the available computing resources considers the parallel efficiency of the implementation. The definition of parallel efficiency is as follows:

$$E(NP) = \frac{S(NP)}{NP}, \quad (6.6)$$

where  $S(NP)$  is the speedup as defined in Equation (6.5), and  $NP$  is the number of cores executing worker processes. Figure 6.13 shows the average parallel efficiency of the parallel implementation for different problem sizes, as the number of processing cores was increased. Like for the speedup measure, the average parallel efficiency from the same problem instances was calculated.

The ideal case for a parallel application would be to utilize all the available computing resources, in which case the parallel efficiency would always be equal to one as the core count increases. From the plot in Figure 6.13, it can be observed that the efficiency of the MWD approach is better than in the MW case for larger number of processes, and as long as there was available capacity at the LAN level. In accordance to the previous analysis, the under utilization of the computing resources is more significant when the master process is overloaded (in the MW case) than when the network infrastructure is saturated (in the MWD case). The lower efficiency is directly proportional to the number of idle worker processes that wait either for the master process (MW case) or for network access (MWD case).

Overall, the experimental results confirm that the objective of fully exploiting the available hardware resources is accomplished when applying the presented MWD approach, thus improving the scalability and efficiency of PRATO when compared to a traditional MW technique.

Nevertheless, some limitations still apply, e.g., DEM block partition and DB control. Regarding DEM division it is necessary for the partitioned blocks to fit into the memory of a worker node, or global memory of the GPU if used. If deploying several workers on one computing node, these restriction applies accordingly. When using the external DB in the MWD case, the SQL statements for the manipulation of the DB tables are included as part of the intermediate results. For this reason, the credentials used for accessing the DB should provide enough privileges for the successful completion of the SQL statements. Additionally, due to the multi-threaded implementation of the result-dumping task, it is very important to avoid oversubscribing a computing node, e.g., deploying two worker processes on a dual-core CPU. In such case, the result-dumping threads have to interrupt the execution of the worker processes in order to access the computing resources of the node, thus having a negative impact on system performance.

## 6.6 Summary

PRATO, a parallel radio-coverage prediction tool for radio networks, has been presented in this chapter. The tool is intended to be used for radio-network planning analysis and decision support. Its high-performance capabilities make it ideal for automatic-optimization tasks that require a large number of evaluations.

The parallel implementation of PRATO includes a novel parallel technique for master-worker configurations. The introduced MWD technique, which combines the use of a DB system with a work-pool approach, delivers improved performance when compared with a traditional MW setup. Moreover, the presented system provides parallel and asynchronous

computation that is completely independent of the GIS used, in this case the GRASS environment. Consequently, a GIS installation is only needed on the master node, thus simplifying the required system setup and greatly enhancing the applicability of this methodology in different environments.

The extensive simulations, performed on the DEGIMA cluster of the Nagasaki Advanced Computing Center, were analyzed to determine the level of scalability of the implementation, as well as the impact of the presented methods for parallel-algorithm design aimed at spatial-data processing. The conducted analyses show that when using the MWD approach, PRATO is able to calculate the radio-coverage prediction of real-world radio networks in a reduced amount of time. Moreover, the experimental results show that PRATO has a better scalability when using the MWD approach than the standard MW setup, since it is able to completely saturate the network infrastructure of the computer cluster. These promising results also show the great potential of the MWD approach for the parallelization of different time-consuming tasks dealing with spatial data, where DBs form an intrinsic part of almost all GIS. Furthermore, the automatic optimization of radio networks, where a large number of radio-propagation predictions take part in the evaluation step of the optimization process, can also benefit from the improved performance of PRATO. Indeed, this last point will be further discussed and validated in the following chapters.

The performance of the worker processes has been additionally improved by including the implementation of the radio-propagation algorithm on GPU. The use of GPU hardware is optional, i.e., it is exploited only if it is available on the computing nodes that host the worker processes.

To the best of the author's knowledge, neither the MWD parallel technique nor the parallel implementation of the radio-prediction algorithm as presented in this chapter, have yet been described in the related literature.



## 7 Service-Coverage Optimization

The high-performance of PRATO, the radio-coverage simulation framework presented in Chapter 6, allows dealing with big problem instances in a reduced amount of time. Additionally, it enables tackling optimization problems that, because of their size, are out-of-reach of traditional approaches, mainly due to the computational-time complexity of their objective-function evaluation.

In this chapter, the challenge is to exploit PRATO for solving one of the classic optimization problems of radio networks: the service-coverage problem. Considering the minimization of the total amount of pilot power subject to a full coverage constraint, a novel optimization approach is introduced. The presented method, based on parallel autonomous agents, gives very good solutions to the problem in an acceptable amount of time. The parallel implementation takes full advantage of GPU hardware in order to achieve considerable speedup. The analysis of the experimental results, considering six real-world, radio networks of different sizes, studies solution-quality and performance aspects.

The content of this chapter extends the research work published by the author in [19] and [21]. The rest of this chapter is organized as follows. Section 7.1 gives a description of the coverage problem and its motivation from the mobile operator's perspective. In Section 7.2, a short overview of related research works is given, before formally introducing the key elements of the service-coverage problem in Section 7.3. The parallel-agent approach, as well as the strategies used for result comparison, are presented in Section 7.5, followed by the simulations and their analyses in Section 7.6.

### 7.1 Motivation

Solving the service-coverage problem for radio networks has received a great deal of attention in the past years. Its complexity demands the confluence of different skills in areas such as propagation of radio signals, telecommunications and information systems, among others.

Even several decades after the launch of the first commercial GSM network, service-coverage planning remains a key problem that all mobile operators have to deal with. Its intricacy arises from the wide range of different combinations of configuration parameters and their evaluation-time complexity.

Regardless of the mobile technology used, e.g., GSM, UMTS or LTE, a lower transmit power generates less interference, which, in turn, translates into more capacity of the radio link (see Chapter 3). Additionally, reducing the transmit-power usage is also related to issues regarding human exposure to the electromagnetic fields generated by BS antennas [51]. During the past few years, public opinion has been extremely sensitive regarding this issue, and thus many countries have already imposed safety standards to limit the electromagnetic field levels produced by antennas in a given range.

From the UMTS perspective, minimizing pilot-power usage leaves more power available for increased network capacity. This is especially important if the traffic and other channels are configured relative to the pilot channel [79]. Moreover, as the demand for internet access

and data services increases [40], so does the pressure on existing network infrastructure, making parameter optimization the only viable solution in the short-term [112].

The idea of using autonomous agents for optimization is not new. It has proven to be a solid optimization approach for solving different types of problems, not only within the area of radio networks [32, 51], but also in other fields [160, 164]. The increased computational-time complexity, when dealing with big problem instances, is tackled using a parallel, agent-based algorithm on GPU. This minimizes the overhead when deploying a larger number of agents working in parallel over the service area, only limited by the amount of memory available on the GPU.

## 7.2 Related work

There are several approaches in the literature that address the service-coverage problem in radio networks [7, 112, 140]. Some of them even claim to achieve near-optimal solutions [143]. As a matter of fact, most formulations are only useful for small network instances and often fail when challenged with larger, real-world networks.

A genetic-algorithm approach for solving the service-coverage problem for GSM networks was presented in [100]. The proposed solution is based on the physical distribution of BSs in order to maximize coverage. The simulations were performed on a test network with 40 candidate sites for BS antennas.

In [143], Siomina and Yuan considered the problem of minimizing the total amount of pilot power for UMTS networks, subject to a full coverage constraint. They tackled the problem with an iterative linear-programming approach, reporting very good results for some test networks, containing from 15 to 65 base stations. The authors noted that bigger problem instances could not be solved because of hardware constraints on the target platform.

As for LTE networks, the service-coverage problem was addressed in [158]. The authors presented an algorithm, based on reinforcement learning, to tackle three aspects of the coverage problem, i.e., coverage holes, weak coverage and pilot pollution. The experimental simulations, performed on 3 BSs, used different antenna-tilt configurations as the proposed solutions.

The service-coverage problem, as presented in this chapter, corresponds to achieving full coverage of the target area, without coverage holes.

## 7.3 Radio-network model

Extending the representation of a radio-network model from [112], this section presents the definitions of all the elements included in the mathematical model used for the simulations.

The goal here is to analyze the state of the network in a given situation, i.e., a ‘snapshot’ at an arbitrary instance. Radio-network planning tools, including the commercial ones, typically rely on static analysis [116].

A snapshot consists of a set of UEs having individual properties, such as location, and equipment type that provides an estimate of the average network behavior. The static approach inherently ignores dynamic effects that influence the system, like fast power control, but the analysis relies on having multiple independent snapshots to produce an average behavior.

Another alternative is to consider using a dynamic tool for radio-network planning [74, 81]. A dynamic simulation considers all Radio-Resource Management (RRM) functionality, as well as user mobility. However, when compared to the static approach, a major disadvantage of the dynamic simulation is the large computational time it requires. This clearly

excludes the possibility of achieving fast performance evaluations of a network, which is one of the objectives of this thesis.

Considering that the performance of a static approach was demonstrated to provide sufficiently accurate results compared to a fully dynamic approach [94, 95], the former one is implemented in this work.

For additional information regarding mathematical models for radio networks and signal propagation, see [94, 112, 150].

### 7.3.1 Basic elements

Consider a radio network with a set of antenna installations (cells),  $C$ . A RSG of a given resolution represents a geographical area,  $A_{\text{total}}$ , within which a set of UEs,  $M$ , is spatially distributed over the pixels of  $A_{\text{total}}$ . Further,  $l_{cm}^{\downarrow}$  is defined as the downlink attenuation factor between cell  $c \in C$  and UE  $m \in M$ . Similarly,  $l_{mc}^{\uparrow}$  represents the uplink attenuation factor between UE  $m$  and cell  $c$ . The attenuation factor values are calculated by performing signal-propagation predictions for every pair  $(c, m)$ , using the radio-propagation model introduced in Section 6.3.2. These predictions already include losses and gains from cabling, hardware, and user equipment.

The amount of power allocated to the pilot signal of cell  $c$  is denoted as  $p_c$ , and it can adopt any value from the sorted set of available pilot power levels,  $P_c = \{p_c^1, p_c^2, \dots, p_c^k\}$ , where  $p_c^k$  is the maximum power.

Based on the introduced elements, the received pilot power from cell  $c$  to UE  $m$  is  $l_{cm}^{\downarrow} p_c$ .

### 7.3.2 Coverage

A UE  $m$  within the area  $A_{\text{covered}}$  is under service coverage if at least one cell  $c$  covers it. Cell coverage is provided to a UE  $m$  from a cell  $c$  if its signal-to-interference ratio,  $\text{SINR}(c, m)$ , at the RSG pixel where  $m$  is located, is not lower than a given threshold,  $\gamma^{\text{cov}}$ :

$$\text{SINR}(c, m) = \frac{l_{cm}^{\downarrow} p_c}{N_0 + \sum_{i \in C} l_{im}^{\downarrow} p_i} \geq \gamma^{\text{cov}}, \quad (7.1)$$

where  $N_0$  is the thermal noise [94]. For convenience, a binary function is defined to determine the coverage of a UE  $m$  by a cell  $c$ . So, for any pair  $(c, m)$ , the coverage of UE  $m$  by cell  $c$  is defined as:

$$\text{cov}(c, m) = \begin{cases} 1 & \text{if } \text{SINR}(c, m) \geq \gamma^{\text{cov}} \\ 0 & \text{otherwise} \end{cases}. \quad (7.2)$$

A set, denoted as  $C_m$ ,  $C_m \subset C$ , contains all the cells covering a UE  $m$ . From this set, the cell with the highest  $\text{SINR}(c, m)$  is referred to as the best server, and denoted as  $c_m^*$ .

Notice that the described radio-network model is easily adaptable for different mobile technologies, e.g., GSM, UMTS and LTE. For example, if solving the service-coverage problem for UMTS, it would be reasonable to assume that all cells in the network operate at maximum power, and adapt Equation (7.1) accordingly. This is, from the interference point of view, the worst-case scenario [28, 143]. This assumption guarantees that even under heavy user traffic full coverage of the service area is maintained, due to the cell-breathing principle [79].

## 7.4 Problem definition

In the problem of optimization of pilot powers for service coverage, the objective is to find a set of pilot-power settings for all cells in the network, such that the total pilot power used is minimized, and a given service coverage criteria is fulfilled. In other words, solving the service-coverage problem corresponds to finding the pilot power levels  $p_c$ , for all cells  $c \in C$ , such that coverage of at least  $b$  UEs is guaranteed, while the total amount of pilot power used is minimized. Here, full coverage of the service area is being considered, thus  $b = |M|$ . Consequently, the optimization objective is defined as follows:

$$P^* = \min \sum_{c \in C} p_c, \quad p_c \in P_c \quad (7.3)$$

subject to

$$\frac{\sum_{m \in M} \text{cov}(c_m^*, m)}{b} = 1. \quad (7.4)$$

It has been proved that the problem of pilot-power optimization for full coverage of the service area is  $NP$ -hard, since it can be reduced to the set-covering problem [163]. Consequently, as long as  $P \neq NP$ , it is unfeasible that a polynomial-time algorithm exists, which is able to find an exact solution to this problem.

## 7.5 Optimization approaches

Since some of the analyzed problem instances are part of a real mobile network deployed in Slovenia by Telekom Slovenije, d.d., there are no references in the literature of other optimization techniques dealing with exactly the same data set. For this reason, two different strategies for setting the pilot power are being presented. They should provide a basis for the comparison of the experimental results. The first strategy is the attenuation-based pilot power, which was used in [140, 143] for result-comparison criteria. In this strategy, a pixel of the service area is always covered by the cell with the maximum attenuation-factor value, i.e., the minimum path loss. The second strategy is the proposed parallel-agent approach, a detailed description of which is given in Section 7.5.2.

### 7.5.1 Attenuation-based approach

The first heuristic for setting the pilot power of all cells in the network is known as attenuation-based, since it relies on the downlink-attenuation factor,  $l_{cm}^k$ . A UE located on some pixel of the service area is always covered by the cell with the minimum path loss, i.e., the highest  $l_{cm}^k$  value. Whenever the maximum available power,  $p_c^k$ , is the same for all the cells in the network, this is equivalent to selecting the cell with the minimum required pilot power to cover a UE  $m$ . Hence, under this assumption, the cell  $c$  covering UE  $m$  is identified as:

$$p_{cm}^{\text{att}} = \min p_c \forall c \in C \iff \text{cov}(c, m) = 1 \quad (7.5)$$

Picking the cells conforming to Equation (7.5) and setting the pilot powers accordingly, full coverage of the service area is achieved. The solution exhibits a total pilot power defined as:

$$P^{\text{att}} = \sum_{c \in C} \max p_{cm}^{\text{att}}. \quad (7.6)$$

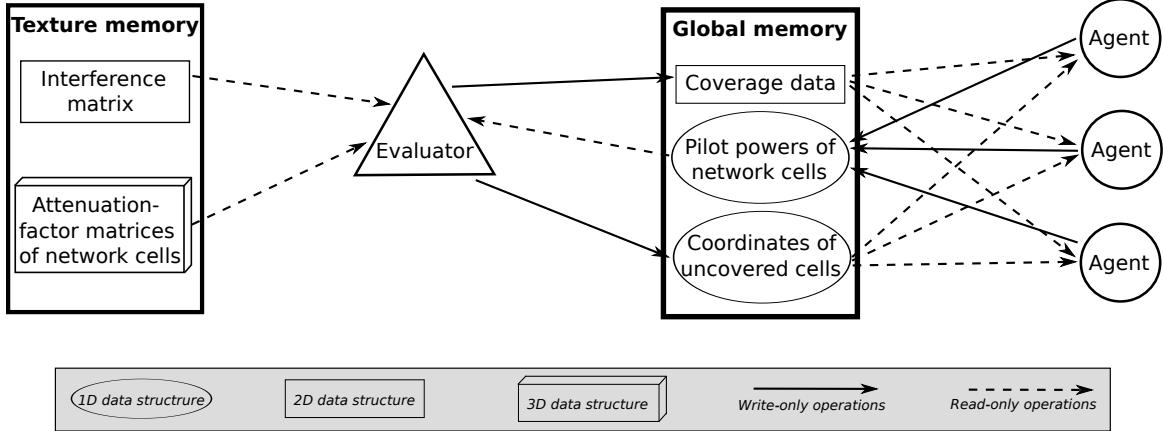


Figure 7.1: Architecture of the parallel, agent-based optimization system on GPU.

The procedure to find a cell  $c$  for every UE  $m$  in the service area consists in sorting, in descending order, all UEs by their attenuation-factor values,  $l_{cm}^{\downarrow}$ . The solution is thus established by the first  $b$  UEs of the sorted sequence, taking the maximum pilot-power setting for a cell into account, i.e.,  $p_{cm}^{\text{att}}$ .

### 7.5.2 Parallel-agent approach

In the parallel-agent approach, a set of autonomous worker agents explore the target geographical area,  $A_{\text{total}}$ , in order to optimize the pilot-power consumption. Each agent randomly moves over the  $A_{\text{total}}$  as it dictates different changes to the pilot power of the cells. PRATO, the radio-coverage framework presented in Chapter 6, performs the objective-function evaluations and radio-propagation predictions based on the proposed changes of the agents.

The moving process during the optimization is strictly random. However, several physical properties that are exclusive to the service-coverage problem are being exploited during the exploration of the search space. Additionally, whenever the current solution breaks any of the given constraints, the optimization process is guided back to the space of valid solutions, providing a mechanism for improving exploration and escaping from local optima.

Because the behavior of the agents is independent between each other, a parallel implementation is fairly straight-forward to achieve. Figure 7.1 shows the architecture of the agent-optimization system. In this GPU-only architecture, agents work in a parallel and autonomous manner, while the evaluator reacts to their changes.

#### 7.5.2.1 Objective-function evaluation

The evaluator represents a central component of the optimization system. It reacts to the pilot-power changes by recalculating the objective-function value. Recall that the objective-function evaluation involves the radio-coverage prediction of the service area and the calculation of the total pilot power used by the cells in the target network.

After a short initialization, during which the attenuation-factor matrices of all the cells and the interference matrix are calculated, the evaluator computes the coverage of the service area based on the pilot powers supplied as the initial solution. Initial solutions are randomly generated from valid pilot-power settings that conform to the full coverage constraint.

The evaluator also maintains a special part of the memory (see “Coordinates of uncovered cells” in Figure 7.1) that is intended for registering uncovered areas, i.e.,  $A_{\text{covered}}$ . If Equation

---

**Algorithm 7.1** Pseudo-code representing the behavior of an agent.

---

```

repeat
    if is_special_agent() and  $\overline{A_{\text{covered}}} > 0$  then
         $l \leftarrow \text{pick\_random\_location}(\overline{A_{\text{covered}}})$ 
    else
         $l \leftarrow \text{pick\_random\_location}(A_{\text{total}})$ 
    end if
    move(l)
    if  $|C_m| = 0$  then
        apply(SS0)
    else
        if  $|C_m| \geq 1$  then
            apply(SS1)
        end if
    end if
until stopping_criterion()

```

---

(7.4) does not hold, the “special” agents randomly select a location from this portion of memory so that a valid solution may be reached again.

It is worth mentioning that the evaluator itself has no influence in the optimization process from a quality point-of-view. Its task is to provide feedback and updated information to the agents that move through the service area. From a performance point-of-view, the importance of the evaluator is significant, as it will be shown in the following sections.

The evaluation of the objective function was completely implemented on the GPU using OpenCL (see Section 5.2). The reason behind this decision is the impact objective-function evaluation has on the performance of the optimization system as a whole, as discussed in Section 2.4. The implementation of the agents is also based on the GPU, which drastically reduces the number of data transfers between CPU and GPU, since all problem elements are available on the GPU during the optimization process. Consequently, careful memory utilization and organization are critical to successfully accommodate all involved problem elements on the GPU, the memory of which is significantly smaller than the RAM available in desktop computers.

### 7.5.2.2 Autonomous agents

The agents apply the pilot-power changes only considering local information. Each of them encapsulates a set of steps that is consistently applied as it randomly moves through the service area of the network. Whenever an agent arrives at a new location, the set of covering cells is calculated, i.e.,  $C_m$ .

The step set an agent applies following this point is directly related to  $|C_m|$ , whereas its movement is determined by  $|\overline{A_{\text{covered}}}|$ , i.e., the area without service coverage.

The behavior of an agent is dictated by the pseudo-code shown in Algorithm 7.1. The first four steps are responsible for guiding its movements. The coordinates are randomly selected from two sets,  $A_{\text{total}}$  and  $\overline{A_{\text{covered}}}$ . Only “special” agents may move to a location without service coverage, and they apply the step set  $SS_0$  for as long as the solution is not valid. The portion of “special” agents used for correcting a solution is a parameter of the optimization process. During the following steps of Algorithm 7.1, the agent applies step sets  $SS_0$  and  $SS_1$  based on the number of cells in  $C_m$ .

---

**Algorithm 7.2** Pseudo-code representing the step set  $SS_0$ , which is applied by the agents in areas without service coverage.

---

```

repeat
     $c' \leftarrow \text{cell\_with\_min\_path\_loss}(m)$ 
     $p_{c'} \leftarrow \text{adjust\_power}(c', \text{inc\_rate})$ 
until  $p_{c'} \in P_{c'}$ 
```

---

**Algorithm 7.3** Pseudo-code representing the step set  $SS_1$ , which is applied by the agents in areas with service coverage.

---

```

repeat
     $c' \leftarrow \text{pick\_random\_cell}(C_m)$ 
     $p_{c'} \leftarrow \text{adjust\_power}(c', \text{dec\_rate})$ 
until  $p_{c'} \in P_{c'}$ 
```

---

If the current location of the agent is not covered by any cell, i.e.,  $|C_m| = 0$ , the step set  $SS_0$  is applied (see Algorithm 7.2). At the beginning, the cell with the lowest path loss,  $c'$ , that may cover a UE at this location, is selected. If several cells have the same  $l_{cm}^{\downarrow}$  value, one of them is randomly chosen. Once  $c'$  is uniquely identified, the agent changes its pilot power by  $\text{inc\_rate}$  dB.

The step set  $SS_1$ , the pseudo-code of which is listed in Algorithm 7.3, is applied if the location of the agent is under the coverage of one or more cells, i.e.,  $|C_m| \geq 1$ . The first step randomly selects a cell from the set  $C_m$ , followed by a decrease of the pilot power of cell  $c'$ . This practice keeps the coverage constraint valid over  $A_{\text{total}}$ , although it might potentially break it on other areas. Ideally, every pixel of the geographical area has to be covered by exactly one network cell, although this is just a representation of a perfect solution that is unreachable because of irregularities in the network topology and the terrain.

In both step sets,  $SS_0$  and  $SS_1$ , the agent makes sure that the new pilot power setting,  $p_{c'}$ , is an element of  $P_{c'}$ . If this is not the case, cell  $c'$  is discarded and another cell is repeatedly selected at the beginning of both step sets, until this condition is satisfied.

The values  $\text{inc\_rate}$  and  $\text{dec\_rate}$  are configurable parameters that should be set before starting the optimization process. They indicate the relative adjustment (expressed in dB) of the pilot power of cell  $c'$ . On the one hand, lowering the pilot power of a cell decreases the interference it creates within its coverage area and those of their neighbors. Since the  $\text{SINR}(c, m)$  value increases with lower interference, the coverage of  $m$  may be achieved by a neighbor cell with the same or lower pilot power. On the other hand, increasing the pilot power of the cell with the minimum path loss improves the coverage by evenly distributing the power among different network cells. This cell is, on average, the nearest one to the location of a UE  $m$ .

With the objective-function evaluation running on the GPU, a new performance bottleneck appeared. The limitation factor in this case was the CPU-to-GPU data transfers that occurred in each iteration of the optimization process (see Section 5.1).

The GPU kernel of the agents is launched as one thread block that contains one thread per deployed agent. The thread block is organized in a one-dimensional grid. The initial location of each agent is randomly generated using the current system time as a random seed. Since OpenCL provides no function for random-number generation, a simplified version of Marsaglia's generator [107] was implemented.

The analysis each agent does about the received signals at the current location is saved into the shared memory of the thread block. It contains the network cell and its pilot-power

setting. Since both numbers are of type *short*, each of which takes up two bytes, there is enough space in a 16 KB shared-memory block to allocate 4,096 agents. The last step involves saving the new pilot powers into global memory. This step is performed by only one of the threads within the thread block in order to avoid memory-access conflicts. Updated pilot powers are saved in negative form to indicate that coverage re-calculation is needed for these cells. In case there are several updated pilot powers for one network cell, the median is calculated and applied as the new pilot power.

Even though coalesced access is not achieved by the GPU kernel of the agents, its sole implementation provided enhanced performance. This performance gain appears because of the lower number of data transfers between the CPU and the GPU, since most data are available in global memory. Moreover, the GPU kernel also produces the truly parallel behavior of the agents, as they all apply the pilot-power changes at the same time.

## 7.6 Simulations

### 7.6.1 Test networks

The test networks,  $\text{Net}_1$ ,  $\text{Net}_2$  and  $\text{Net}_3$ , are subsets of the real radio network deployed by Telekom Slovenije, d.d. The path-loss predictions were calculated using the radio-propagation model presented in Section 6.3.2. A DEM with a  $25 \text{ m}^2$  resolution was used as the terrain-profile data. The requirements for the coverage threshold,  $\gamma^{\text{cov}}$ , were provided by experts of the Radio Network department of Telekom Slovenije, d.d.

$\text{Net}_1$  is deployed over a densely populated urban area. For this reason, the value of  $\gamma^{\text{cov}}$  is lower here, since network capacity is the dominating factor, whereas coverage is flexible because of a larger cell density, i.e., more BSs per surface unit.  $\text{Net}_2$  represents a network deployed over a rural area, meaning that the network capacity can be reduced at the cost of a better coverage, since the user density is lower. The last network,  $\text{Net}_3$ , represents a suburban area with a densely populated, but relatively small, downtown center, where a compromise between the network capacity and the coverage has to be achieved.

The second group of test networks, including  $\text{Net}_4$ ,  $\text{Net}_5$  and  $\text{Net}_6$ , is part of the publicly available MOMENTUM project [86]. Test network  $\text{Net}_4$  represents the city of Berlin (Germany),  $\text{Net}_5$  represents the city of The Hague (Netherlands), and  $\text{Net}_6$  is the largest network optimized in [143], representing a reduced version of  $\text{Net}_4$ . All networks include information about BS locations, path-loss predictions and realistic antennas, which are part of the scenarios provided by the MOMENTUM project.

Network configurations that represent what could be an initial-network setup by common-planning standards [79] were produced using the attenuation-based approach. Such configurations can be easily calculated by a network planner. Table 7.1 lists the number of BSs and cells per test network, as well as the size of the geographical area. Different network-parameter values used during the simulations are shown in Table 7.2.

### 7.6.2 Parameter settings of the parallel-agent approach

The parameter settings for the optimization algorithm were determined after some experimentation with the test networks. The parameter settings for each test networks are listed in Table 7.3.

Using a higher *inc\_rate* than *dec\_rate* reflects the behavior of the agents when full coverage of the service area is not guaranteed. In practice, areas without service coverage usually appear as irregular islands. The stopping criteria were set by limiting the total number of pilot-power changes an agent is allowed to make. The value was set to 10,000,

Table 7.1: Sizes of the test networks used for experimentation of the service-coverage problem, in terms of equipment and geographical area.

	Number of base stations	Number of cells	Surface [km <sup>2</sup> ]	Resolution [m <sup>2</sup> ]
Net <sub>1</sub>	26	77	100.00	25
Net <sub>2</sub>	8	23	306.25	25
Net <sub>3</sub>	45	129	405.00	25
Net <sub>4</sub>	65	193	56.25	50
Net <sub>5</sub>	12	36	16.00	50
Net <sub>6</sub>	50	148	56.25	50

Table 7.2: Network parameters of the test networks used for the service-coverage problem.

	$p_c^k$	$N_0$	$\gamma^{\text{cov}}$
Net <sub>1</sub>	15.00 W	$1.55 \cdot 10^{-14}$ W	0.010
Net <sub>2</sub>	19.95 W	$1.55 \cdot 10^{-14}$ W	0.020
Net <sub>3</sub>	15.00 W	$1.55 \cdot 10^{-14}$ W	0.015
Net <sub>4</sub>	19.95 W	$1.55 \cdot 10^{-14}$ W	0.010
Net <sub>5</sub>	19.95 W	$1.55 \cdot 10^{-14}$ W	0.010
Net <sub>6</sub>	19.95 W	$1.55 \cdot 10^{-14}$ W	0.010

Table 7.3: Parameter settings of the parallel-agent approach for each test network.

	Agents	<i>inc_rate</i> [dB]	<i>dec_rate</i> [dB]	Pilot-power changes
Net <sub>1</sub>	16	0.2	-0.1	10,000
Net <sub>2</sub>	16	0.2	-0.1	10,000
Net <sub>3</sub>	16	0.2	-0.1	10,000
Net <sub>4</sub>	6	1.0	-0.1	10,000
Net <sub>5</sub>	2	1.0	-0.1	10,000
Net <sub>6</sub>	6	1.0	-0.1	10,000

Table 7.4: Optimization results after applying two different approaches for solving the service-coverage problem. All values are expressed in Watts.

Attenuation-based			Parallel agents	
	Total power	Average pilot power	Total power	Average pilot power
Net <sub>1</sub>	419.292	5.445	137.064	1.780
Net <sub>2</sub>	78.297	3.404	33.344	1.450
Net <sub>3</sub>	1,014.113	7.861	582.954	4.519
Net <sub>4</sub>	179.876	0.932	145.715	0.755
Net <sub>5</sub>	73.872	2.052	34.884	0.969
Net <sub>6</sub>	147.014	0.993	112.332	0.759

even though for some of the test networks the best solutions were found in the first quarter of the experiment.

All experiments were performed on a multi-core Intel i7 2.67 GHz desktop computer with 6 GB of RAM running a 64-bit Linux operating system. The GPU hardware used was an ATI HD5570 with 1 GB of DDR3 RAM. The implementation language used was C, combined with OpenCL and OpenMPI extensions.

### 7.6.3 Results

The results achieved by the parallel-agent approach, listed in Table 7.4, improved the optimization objective significantly. They show that the pilot-power usage was reduced in all networks while the service area was kept under full coverage. Moreover, the parallel-agent solution for Net<sub>1</sub> improved the attenuation-based setting by more than 300 %. As for Net<sub>2</sub>, the observed improvement is around 232 %, while the improvement for Net<sub>3</sub> is more than 170 %.

The last test network, Net<sub>6</sub>, is the same as N6 in [143]. When comparing these results to those of [143], an improvement of almost 3 % can be observed in the solution provided by the parallel-agent approach, i.e., 0.778 against 0.759 for the average pilot power of Net<sub>6</sub>.

### 7.6.4 Performance analysis

The graphs shown in Figures 7.2, 7.3 and 7.4 depict the convergence of the parallel-agent approach after ten independent runs for test networks Net<sub>1</sub>, Net<sub>2</sub> and Net<sub>3</sub>, respectively. Only feasible solutions were plotted, i.e., the solutions that meet the full-coverage constraint. Unfeasible solutions were marked with a value of inferior quality than the worst solution found: 428 for Net<sub>1</sub>, 129 for Net<sub>2</sub>, and 1,435 for Net<sub>3</sub>.

From the graphs of Net<sub>1</sub> (see Figure 7.2) and Net<sub>2</sub> (see Figure 7.3), a good initial convergence can be observed. This is followed by a steady improvement of the intermediate solutions. In Net<sub>1</sub>, no additional solution improvement is noticed towards the end of the optimization process. This fact suggests that the stopping criteria is suitable for this problem instance. A similar situation is observed for Net<sub>2</sub> that also shows a flat profile towards the end. From the graph of Net<sub>3</sub> (see Figure 7.4), a slower initial convergence, followed by a steady improvement of intermediate solutions and no significant solution enhancement towards the end, can be observed. This convergence profile suggests that this problem instance presents a more difficult optimization case than for Net<sub>1</sub> and Net<sub>2</sub>. Indeed, this is the largest test network in terms of surface area. However, further investigation is needed to

confirm this hypothesis. Nevertheless, the parallel-agent approach improved the pilot-power usage of this test network by almost 75 %.

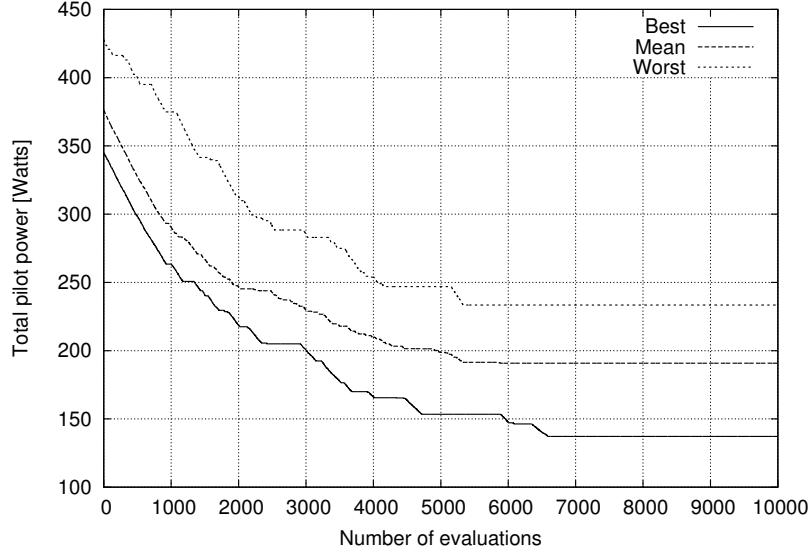


Figure 7.2: Convergence profile of the parallel-agent approach for the test network  $\text{Net}_1$ , deployed over an urban area.

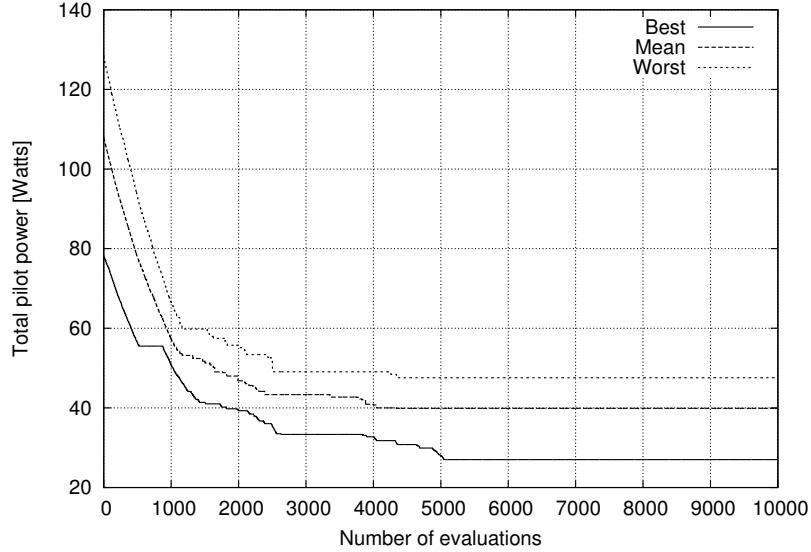


Figure 7.3: Convergence profile of the parallel-agent approach for the test network  $\text{Net}_2$ , deployed over a rural area.

In the following, the speed-performance analysis of the experimental simulations is presented. This analysis covers the running times during the optimization of the first three test networks, i.e.,  $\text{Net}_1$ ,  $\text{Net}_2$  and  $\text{Net}_3$ . The running times were measured for each implementation, and the average times, calculated after ten independent runs, are given. The number of pilot-power changes per agent was limited to 1,000, while all other algorithm parameters were kept at the same values as in Section 7.6.2.

Table 7.5 lists the average wall-clock times in seconds for the different implementations and test networks. The implementations include: the CPU-MPI implementation that con-

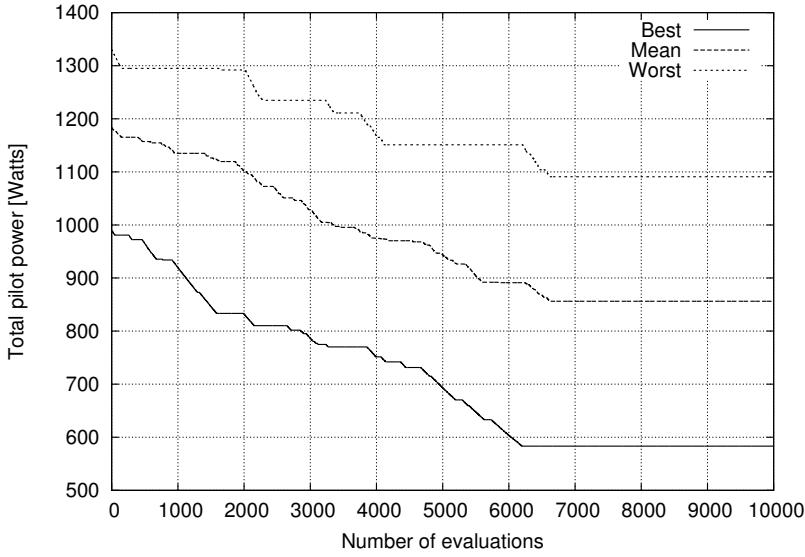


Figure 7.4: Convergence profile of the parallel-agent approach for the test network Net<sub>3</sub>, deployed over a suburban area.

Table 7.5: Wall-clock times (in seconds) and speedup factors for the different implementations of the objective-function evaluation and the parallel agents, as measured during the experimentation of the service-coverage problem.

CPU-MPI		GPU-MPI		GPU-GPU	
	Avg. time [s]	Avg. time [s]	Speedup	Avg. time [s]	Speedup
Net <sub>1</sub>	4,105	872	4.71	206	19.93
Net <sub>2</sub>	774	304	2.55	124	6.24
Net <sub>3</sub>	5,981	1,283	4.66	249	24.02

sists of objective-function evaluation on CPU and parallel agents over MPI, the GPU-MPI implementation that consists of objective-function evaluation on GPU and parallel agents over MPI, and the GPU-GPU implementation that consists of objective-function evaluation and parallel agents on the same GPU. The CPU-MPI implementation is the basis for the speedup calculation of the other two implementations.

The function evaluation on the GPU that communicates with the agents over MPI provides the second measured setup. The evaluator implementation takes advantage of shared memory for thread collaboration within a thread block and texture memory for constant elements, as is it was explained in Section 6.4.5.3. Still, the speedup is considerable but improvable, since numerous data transfers between CPU and GPU are needed for the agents to access optimization-related information. The last result set presents measurements for the complete GPU implementation, including objective-function evaluation and agents on the same device. The improved speedup delivered by this combination highlights the impact that CPU-to-GPU memory transfers have on the overall system performance. This fact is supported by the second and third measured setups, the speedups of which exhibit, on average, a four-fold improvement. It is also interesting to note how the speedup gain increases with the problem-instance size. This fact confirms that larger problem instances have a greater benefit from a parallel implementation in terms of computational time.

## 7.7 Summary

This chapter presented a novel optimization approach for solving the well-known service-coverage problem in radio networks. The problem addressed the full coverage of a geographical area using a minimum amount of pilot power. The newly introduced parallel-agent approach was successfully tested in six networks that represent real-world scenarios. The experimental results show that the parallel-agent approach is able to find better solutions than some heuristics, like the presented attenuation-based approach. Moreover, the algorithm successfully tackled larger networks, thus overcoming the obstacles of other state-of-the-art optimization methods regarding problem-instance size [140, 143].

Compared to a different optimization approach in the literature [143], the solution-quality of the parallel-agent approach showed a quality improvement. The proposed solutions, calculated for the same problem instance as in [143], were improved at the cost of a longer running time. It is worth mentioning that it is feasible for the optimization algorithm to take a longer time to reach the solution, since design problems, as the service-coverage one, are usually solved offline. A comparison and analysis of the performance of the radio-coverage prediction for real-world, radio-network planning is later provided in Chapter 10.

Different implementations of the parallel-agent approach, combining a serial version on CPU, parallel processes over MPI and GPU kernels, were presented. In particular, GPU architectures enable the implementation of parallel heuristics in a natural way while substantially improving the computational-time performance. To the best of the author's knowledge, the parallel-agent approach as presented in this chapter, has not yet been described in the related literature.



## 8 Soft-Handover Optimization

In Chapter 7, an application exploiting the advantages of faster evaluation methods has been presented. Solving the service-coverage problem for real-world networks capitalizes on the ability to tackle bigger problem instances. Because of their size, such problems were previously unsolvable in a feasible amount of time. This improved performance also allows solving optimization problems with a higher degree of complexity, usually represented by the evaluation of multi-dimensional, non-convex objective functions.

This chapter focuses on solving a new optimization problem for 3G networks that deals with downlink and uplink SHO areas (see Section 3.1). By introducing a penalty-based objective function and some hard constraints, the formal definition of the SHO-balancing problem in UMTS networks is given. The state-of-the-art mathematical model used and the penalty scores of the objective function are set according to the configuration and layout of a real mobile network, deployed in Slovenia by Telekom Slovenije, d.d. The balancing problem is then tackled by three optimization algorithms, each of them belonging to a different class of metaheuristic approaches.

To the best of the author's knowledge, there is no reference in the literature of a simulation-based approach to find active downlink and uplink SHO areas. Additionally, there are no formal optimization methods known to the author that tackle the SHO-balancing problem as described here, which extends the research work published by the author in [20].

The remainder of this chapter is organized as follows. Section 8.1 describes the motivation behind the SHO-balancing problem, whereas Section 8.2 gives an overview of other works related to pilot-power and SHO optimization in UMTS. The static network model is presented in Section 8.3, where all the elements of the mathematical model and the objective function are defined. In Section 8.4, the problem is formally defined, followed by a short description of the optimization algorithms used in Section 8.5. The simulations, including their environment, parameter setup and result analysis, are introduced in Section 8.6.

### 8.1 Motivation

Despite several built-in mechanisms that allow a radio network to overcome different problems due to the lack of SHO during a HSDPA connection, some abnormal cases do arise, especially in areas where there is SHO capability in the uplink, but none in the downlink. An example of such a case is depicted in Figure 8.1, which shows the interference behavior (denoted as UL RSSI) during a HSPA connection in: (a) normal SHO conditions, and (b) unbalanced SHO conditions. The plotted data are actual radio-network statistics from the mobile network deployed in Slovenia by Telekom Slovenije, d.d. The graph (a) shows a normal HSUPA-enabled service situation, in which the measured interference is proportional to the traffic being served. Note how the noise rises with the increased traffic on cell 1, while its neighbor (cell 2) has almost no interference nor traffic. Moreover, the graph profile for both traffic and noise of cell 1 are almost identical. The graph (b) depicts a problematic situation, where the noise level does not only rise on the cell serving the HSUPA services (cell 1), but also on the neighboring one. Notice how the interference level rises on the cell

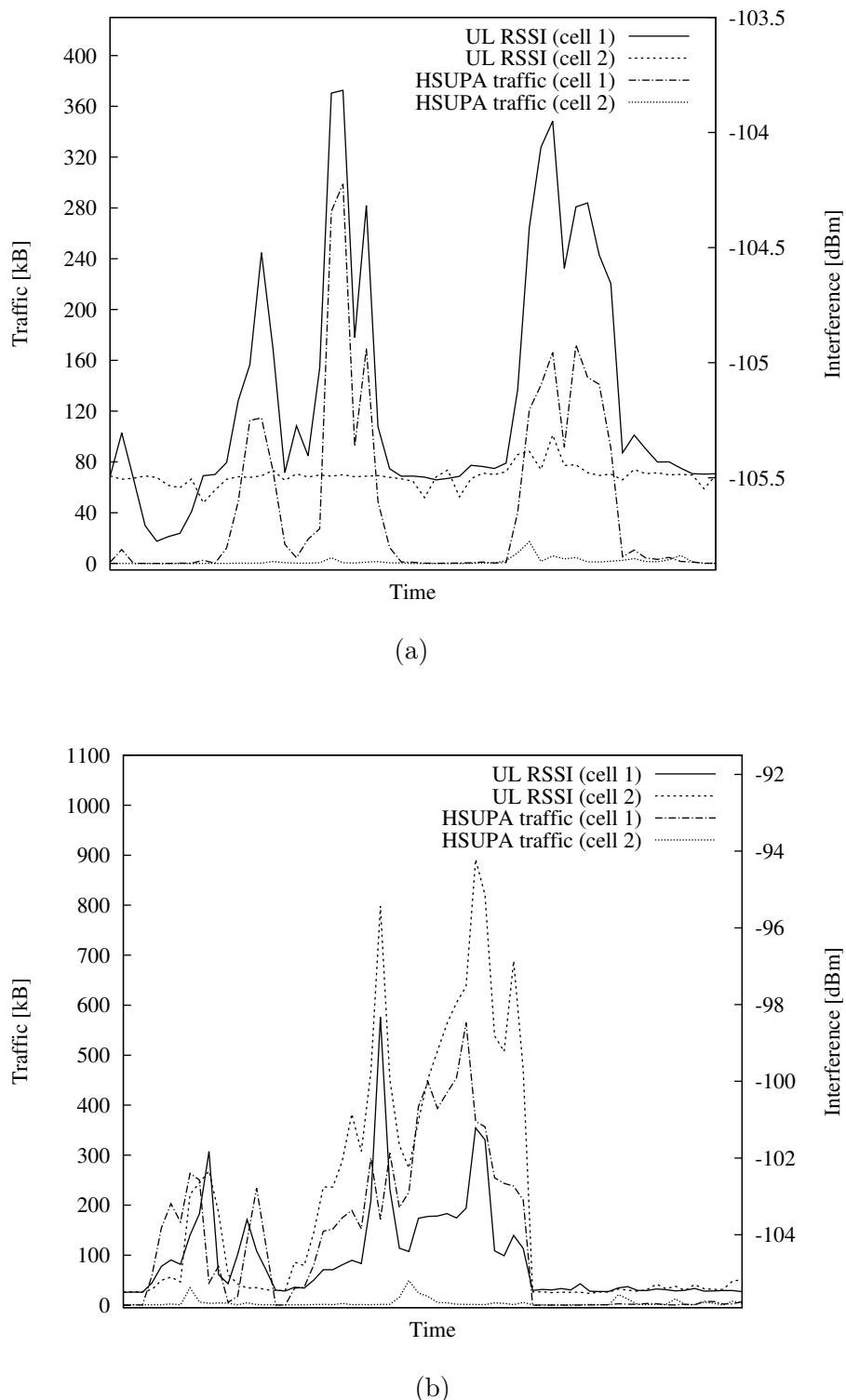


Figure 8.1: HSUPA traffic and uplink interference with: (a) balanced downlink and uplink SHO conditions, and (b) unbalanced downlink and uplink SHO conditions.

that has almost no traffic (cell 2). It is clear that the source of this noise rise is generated by the active connection on cell 1, which shows an increase in HSUPA traffic. However, the noise-level profile on cell 2 does not follow its traffic, as it did in the normal situation (a). This is due to cell 2 not being part of the active set. Such situations appear when the uplink coverage is larger than the downlink coverage. Interestingly enough, this seems to be an exceptional case, as Holma and Toskala write in [80], when describing the SHO in chapter 5: “... There is no obvious reason why the serving E-DCH cell would not be the same as the serving HSDPA cell, and this is also required to be the case in the specifications.”

Given the described context, the challenge is to achieve the correct balance or distribution of downlink and uplink SHO areas within a working UMTS network. Therefore, the network has to be fine-tuned to improve the SHO-area balancing, thus to avoid the exceptional appearance of problematic situations, as shown in Figure 8.1. This clearly implies that the mobile network configuration should not be excessively altered, since other aspects of the network are working well before starting the optimization process. Hence, the objective of the optimization problem is to find a pilot-power configuration for all the cells in the target network, such that the balance of downlink and uplink SHO areas is improved and other network aspects are preserved. The optimization process takes into account different kinds of hardware, e.g., amplifiers, cables, and antennas, adjusting the pilot powers of the cells.

PRATO, as defined in Chapter 6, is used as the evaluation framework of the SHO-balancing problem. A state-of-the-art mathematical model [112] describes the downlink and uplink SHO areas. By introducing a penalty-based objective function and some hard constraints, a formal definition of the SHO-balancing problem in UMTS networks is given. The mathematical model and the penalty scores of the objective function are set according to the configuration and layout of a real mobile network, deployed in Slovenia by Telekom Slovenije, d.d. The SHO settings are also taken from the actual network configuration, still they were adapted to closely model interference and other dynamic aspects of the network.

## 8.2 Related work

The SHO optimization has received quite some attention from the scientific community during the past years. This mainly relates to the importance it has within the deployed networks that provide high-speed services, such as video telephony [103] and data services by means of HSPA [29, 31].

Some authors tackled optimization problems at the planning stage of the network [50, 62], considering, among other variables, BS locations and hardware. However, most mobile operators are unable to apply such methods to a live network, since the planning phase of the new installation has long been concluded. Moreover, the great majority of the BSs has already been deployed and their hardware also installed. Therefore, from the mobile operator’s point of view, mainly parameter and software optimization are the only tools available, when it comes to QoS improvement (see Chapter 4) and network troubleshooting in the short term.

Optimizing SHO by means of pilot-power adjustment is an established way of enhancing network capacity, when high-speed services like HSDPA and HSUPA coexist with legacy technologies [29]. In the UMTS, the pilot power is typically between 5 % to 10 % of the total downlink transmit power of the BS [94], but there is no standardized method to find a pilot-power setting. A number of existing approaches to resolve this issue exist in the related literature [79, 141, 152], those based on optimization methods being the most effective ones [50, 56, 94, 98, 143]. Such a wide spectrum of available procedures is directly related

to the diverse criteria taken into account when assigning the pilot power of a cell. The fundamental reason behind this fact is that the pilot power is a common adjustment variable of various optimization problems in radio networks. This is especially true for UMTS networks, due to their frequency-reuse factor of one [79].

### 8.3 Radio-network model

This section extends the representation of the radio-network model, previously introduced in Section 7.3, in order to include the SHO functionality. In this context, the mathematical model links the SHO settings with the pilot-power level of each cell, the best-server pattern, and the network coverage.

By introducing a change step of 0.01 dB and bounding the pilot power of a cell  $c$ ,  $c \in C$ , to  $\pm 2$  dB (relative to the pilot-power setting the cell had before optimization), the number of elements in the set  $P_c$  is reduced. A smaller value for the change step would have practically no influence over the transmit-power adaptation of the cells, whereas bigger changes would suffer from excessive granularity. These reasons were discussed and confirmed by the radio engineers at Telekom Slovenije, d.d.

From the optimization point of view, the purpose of the change step is twofold. First, since the optimization targets a live network, there is no need for the algorithms to create complete new configurations, but just to fine-tune the existing ones. Second, the problem complexity is lowered, because the size of the search space is smaller and discrete.

#### 8.3.1 Soft-handover areas

To obtain a realistic outline of the areas where a UE may potentially maintain connections to more than one cell, a static version of the active set, as defined in [112], is used. To this end, a SHO window,  $\gamma^{\text{sho}}$ , and a maximum active-set size,  $as^{\max}$ , are introduced. Both parameters are taken from the working configuration of the real network. It follows that the cells to which a UE  $m$ ,  $m \in M$ , may maintain concurrent downlink connections are part of the set:

$$\begin{aligned} SHO_m^\downarrow &= \left\{ c \mid l_{c^*m}^\downarrow p_{c^*} - l_{cm}^\downarrow p_c \leq \gamma^{\text{sho}} \right\}, \\ |SHO_m^\downarrow| &\leq as^{\max}, \end{aligned} \quad (8.1)$$

where  $c$  is a network cell,  $l_{c^*m}^\downarrow$  is the downlink attenuation factor of the best-serving cell, and  $p_{c^*}$  is its pilot power. Since the number of elements in  $SHO_m^\downarrow$  is at most  $as^{\max}$ , the weakest links are removed if there are several present. This method is well suited for configurations with no hysteresis, since dynamic effects are ignored in static models [112].

Additionally, in the uplink, the set of cells to which a UE can potentially be in SHO is defined as:

$$SHO_m^\uparrow = \left\{ c \mid l_{mc}^\uparrow p_m^\uparrow \geq 3.16227766 \cdot 10^{-12} \text{ mW} \right\}, \quad (8.2)$$

where  $l_{mc}^\uparrow$  is the uplink attenuation factor from a UE  $m$  to a cell  $c$ , and  $p_m^\uparrow$  is the uplink transmit power of  $m$ .

The static nature of the model intentionally neglects mobility and dynamic interference by narrowing  $\gamma^{\text{sho}}$  down to 2 dB [112].

## 8.4 Problem definition

Using the elements defined in Section 8.3, an objective function was formulated in cooperation with a team of radio engineers of the Radio Network Department at Telekom Slovenije, d.d. The objective function is constructed as a weighted sum, containing different costs that penalize the occurrence of specific SHO conditions in downlink and uplink, which may potentially cause the aforementioned malfunctioning, introduced in Section 8.1.

A cost-based objective function is the most natural and straight-forward way of defining the optimization objective. Besides it is easily extendable to include other future situations, also defining the mutual importance of the different phenomena taken into account at the optimization phase.

Hence, the definition of the objective function for the SHO-balancing problem is the minimization of the sum of penalty scores given as:

$$f_{\text{sho}}^* = \min \sum_{c \in C} \sum_{m \in M} p f_{\text{cov}} (1 - \text{cov}(c, m)) + p f_{\text{sho}}^\uparrow sho_{cm}^\uparrow (1 - sho_{cm}^\downarrow) + p f_{\text{sho}}^\downarrow sho_{cm}^\downarrow (1 - sho_{cm}^\uparrow), \quad (8.3)$$

where

$$sho_{cm}^\downarrow = \begin{cases} 1 & c \in SHO_m^\downarrow \\ 0 & \text{otherwise} \end{cases}, \quad (8.4)$$

$$sho_{cm}^\uparrow = \begin{cases} 1 & c \in SHO_m^\uparrow \\ 0 & \text{otherwise} \end{cases}, \quad (8.5)$$

and

- $\text{cov}(c, m)$  is the binary function used to determine the coverage of a UE  $m$  by a cell  $c$ ,
- $p f_{\text{cov}}$  represents the penalty factor for uncovered areas,
- $p f_{\text{sho}}^\uparrow$  represents the penalty factor for uplink SHO areas where SHO is not possible in the downlink, and
- $p f_{\text{sho}}^\downarrow$  represents the penalty factor for downlink SHO areas where SHO is not possible in the uplink.

After extensive experimentation, and working in cooperation with the radio engineers from the Radio Network Department at Telekom Slovenije, d.d., the penalty factors from Equation (8.3) are set to the following values:

- $p f_{\text{cov}} = 15$ ,
- $p f_{\text{sho}}^\uparrow = 13$ , and
- $p f_{\text{sho}}^\downarrow = 3$ .

It is clear that the coverage is the most important quality aspect from the network point of view (penalty factor  $p f_{\text{cov}}$ ). Moreover, it imposes the biggest constraint to the optimization process, since the balance between SHO areas should not sacrifice network coverage. Another important characteristic that emerges from these values is the preference for minimizing areas where SHO capability is available in the uplink, but not in the downlink (penalty factor

$pf_{sho}^{\uparrow}$ ). As it has been described in Section 8.1, the consequences of such SHO arrangement produce severe interference in neighboring cells (Figure 8.1), which may also result in service inaccessibility. The last factor  $pf_{sho}^{\downarrow}$  imposes a penalty value over areas where the SHO capability is available in the downlink, but not in the uplink. Recall that when accessing HSPA services, SHO is available only in the uplink. For this reason, the link throughput may benefit from the SHO in the uplink if it is available. The relative lower importance of the last penalty factor, when compared to the other ones, is directly related to the consequences of the unbalancing that such SHO areas may have on the network. In this case, only the HSPA throughput is affected, while the service accessibility should not be an issue, given there is enough uplink coverage [80].

## 8.5 Optimization approaches

The SHO-balancing problem has been tackled using three fundamentally different optimization algorithms, namely:

- Differential evolution (see Section 2.3.1), from the family of evolutionary algorithms;
- Differential ant-stigmergy algorithm (see Section 2.3.2), from the family of swarm-intelligence algorithms; and
- Simulated annealing (see Section 2.3.3), from the group of classic metaheuristic algorithms.

Each of these algorithms shall minimize the objective function value by adopting essentially disparate approaches, hence the diversity of applying algorithms belonging to different families to solve the same optimization problem. Therefore, the result analysis shall establish which of the presented approaches is better suited for solving the SHO-balancing problem.

The following sections describe how the SHO-balancing problem is represented by the internal structure of each of the selected algorithms and their controlling parameters.

### 8.5.1 Differential evolution

The DE algorithm features a parallel direct search method, which utilizes a population of  $D$ -dimensional parameter vectors. The SHO-balancing problem is expressed in each component of a vector  $X$  of the population, which represents the pilot power of a target cell, i.e.:

$$X_{ag} = \{x_1, x_2, \dots, x_c, \dots, x_D\}, \quad (8.6)$$

where  $x_c \in P_c$  represents a candidate pilot-power setting of cell  $c$ , and  $g$  indicates the generation of an individual  $a$  in the population. Since there are  $|C|$  cells in a mobile network, it follows that the population size,  $D = |C|$ .

From the different variants of DE, the most popular one is used here, called *DE/rand/1/bin*. The nomenclature used to name this variant indicates the way the algorithm works:

- *DE* denotes the differential evolution algorithm,
- *rand* indicates that the individuals selected to compute the mutation values are randomly chosen,
- 1 specifies the number of pairs of selected solutions used to calculate the crossover vector, and
- *bin* means that a binomial recombination operator is used.

---

**Algorithm 8.1** A move in the search space of SA for solving the SHO-balancing problem.

---

```

 $c' \leftarrow \text{pick\_random\_cell}(C)$ 
repeat
    if  $\text{uniform}[0, 1] < 0.5$  then
         $p_{c'}^{\text{new}} \leftarrow p_{c'} + 0.01$ 
    else
         $p_{c'}^{\text{new}} \leftarrow p_{c'} - 0.01$ 
    end if
until  $p_{c'}^{\text{new}} \in P_{c'}$ 
 $p_{c'} \leftarrow p_{c'}^{\text{new}}$ 

```

---

### 8.5.2 Differential ant-stigmergy algorithm

The mapping between the balancing problem and DASA is similar to the one for DE:

$$X_a = \{x_1, x_2, \dots, x_i, \dots, x_D\}. \quad (8.7)$$

In this case, each ant,  $a$ , creates its own solution vector,  $X_a$ , during the minimization process. At the end of every iteration, and after all the ants have created solutions, they are evaluated to establish if any of them is better than the best solution found so far.

### 8.5.3 Simulated annealing

From the SA perspective, the system under optimization is in a given *state* at each time step during the process. The objective function maps a system state to a value known as the *energy* of the system in that state. A *move* in the search space represents a change in the state of the system. Whether a move is performed or not depends on the values of the objective function and the temperature parameter (see Section 2.3.3). After making a move, the system may exhibit lower or higher energy, depending on the current objective-function value.

Algorithm 8.1 shows the pseudo-code of a move in the search space of possible pilot-power settings, resulting in a new state of the system.

At the beginning, a cell,  $c'$ , is randomly selected from the set of all cells in the network,  $C$ . Next, a change of  $+0.01$  dB or  $-0.01$  dB is applied with 50 % probability to  $p_{c'}$ . The pilot power of cell  $c'$  is expressed in dBm. The randomly generated pilot-power setting,  $p_{c'}^{\text{new}}$ , is then validated, i.e., it must be an element of the set  $P_{c'}$ . If  $p_{c'}^{\text{new}}$  is not a valid pilot power, the loop is executed again, generating another random pilot power. Finally, the pilot power of cell  $c$  is replaced by  $p_{c'}^{\text{new}}$ .

It is important to note that, as long as  $|P_{c'}| > 1$ , the pseudo-code shown in Algorithm 8.1 shall never be trapped in an endless loop. On the other hand, if  $|P_{c'}| < 2$ , there are no candidate pilot powers for cell  $c'$ , and thus there is no possibility of optimization. Notice also that the acceptance of a move in the search space is left to SA and its stochastic components.

## 8.6 Simulations

The simulations were performed over the target geographical area, for which DEM and clutter data were available. The mobile users were assumed to be uniformly distributed. The SHO conditions were determined by the relative received-signal quality from different cells, and the SHO window, which triggers the addition of a cell to a user's active set [79].

Table 8.1: Technical characteristics of Net<sub>7</sub>, the test network used for the SHO-balancing problem.

Parameter	Value
Number of cells	25
Coverage threshold	-115 dBm
SHO window ( $\gamma^{\text{sho}}$ )	2 dB
User equipment ( $p_m^\uparrow$ )	21 dBm, power class 4
Pixel resolution	25 m <sup>2</sup>
Population density	398/km <sup>2</sup>

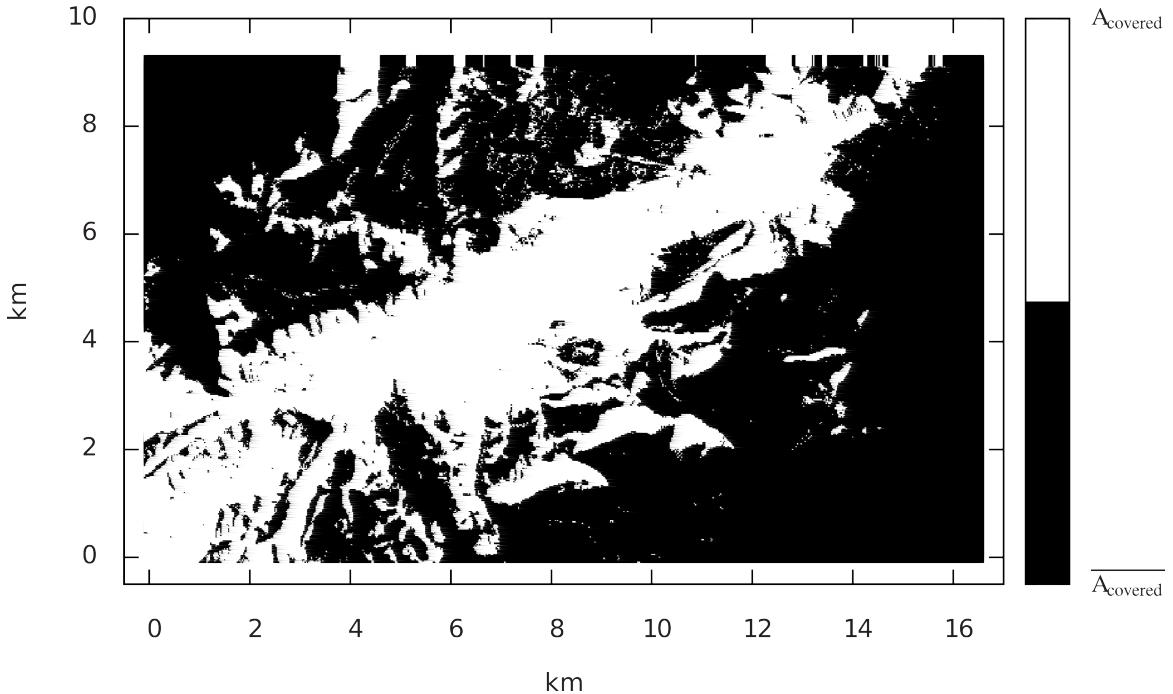


Figure 8.2: Area under radio coverage,  $A_{\text{covered}}$ , and without radio coverage,  $\overline{A_{\text{covered}}}$ , within the complete geographical area,  $A_{\text{total}}$ , of test network Net<sub>7</sub>.

### 8.6.1 Test network

The test network used for the simulations, Net<sub>7</sub>, is a subset of the real UMTS network deployed in Slovenia by Telekom Slovenije, d.d. It represents a network extending over a hilly terrain, combining both rural and middle-dense suburban areas, which contains 25 cells within an area of more than 150 km<sup>2</sup>. Table 8.1 shows some characteristics of the test network used, and Figure 8.2 shows the area under radio coverage,  $A_{\text{covered}}$ , within  $A_{\text{total}}$ .

### 8.6.2 Algorithm parameters

In this section, the algorithm-parameter setup used during the simulations is given. In all three cases, the parameter names are given with their respective values and descriptions.

The parameters controlling the behavior of the DE algorithm were set as follows:

- $NP = 100$ , the population size;
- $g_{\max} = 1000$ , the maximum number of generations for the algorithm to run;

Table 8.2: Solution-quality performance of the three algorithms, after 30 independent runs.

	Best	Worst	Mean	Std. deviation
DE	2,286,292.00	2,286,541.00	2,286,517.09	62.06
DASA	2,286,446.00	2,286,633.00	2,286,592.00	26.19
SA	2,293,350.00	2,295,570.00	2,294,626.50	159.61

- $CR = 0.8$ , the crossover constant; and
- $F = 0.5$ , the mutation-scaling factor.

As for DASA, the parameters were set to the following values:

- $m = 10$ , the number of ants;
- $b = 10$ , the discrete base;
- $\rho = 0.2$ , the pheromone dispersion factor;
- $s_+ = 0.01$ , the global scale-increasing factor;
- $s_- = 0.01$ , the global scale-decreasing factor; and
- $\epsilon = 10^{-2}$ , the maximum parameter precision.

The two parameters that control the behavior of SA are:

- $t_{\text{initial}} = 125$ , the initial temperature; and
- $it = 100,000$ , the total number of iterations.

In this case, the exponential-lowering schema was chosen as the way the temperature was lowered during the SA searches. Note that the probability of an increase/decrease of the pilot power and the change step, the value of which were set to 50 % and 0.01 dB, respectively, can also be considered algorithm parameters. For this reason, the presented approach is a special case where the values of both parameters were set in advance.

All experiments were carried out on a 4-core Intel i7 2.67 GHz desktop computer with 6 GB of RAM running a 64-bit Linux operating system. The implementation languages used were C and Python, with the latter mostly used as ‘glue’ to hold the different implementation parts together, as well as for I/O operations. To lower the time needed to run one optimization round, the entire objective-function evaluation was implemented using OpenCL and executed on a nVidia GeForce GTX 260. This individual change exhibited more than 15-fold, execution-time speedup, when compared to the original CPU-only version.

### 8.6.3 Results

In this section, the performance of the selected algorithms is presented. The analysis includes aspects related to solution quality and convergence speed. All experimental results were obtained after 30 independent runs, each of them limited to a maximum of 100,000 evaluations. The gathered results are shown in Table 8.2.

It may be observed that DE reached the lowest objective-function value, closely followed by DASA. Likewise, both algorithms arrived at very similar results for the worst, mean and

standard-deviation values. SA, on the other hand, did not achieve comparable values, since its results are behind those of DE and DASA. Notice that even the best SA solution is no better than the worst solution of DASA. Moreover, the standard deviation exhibited by SA is one order of magnitude bigger than those of DASA and DE.

The convergences of the best, mean and worst objective-function values are shown in Figures 8.3, 8.4, and 8.5, respectively. It is worth mentioning that every optimization run starts from a different solution, randomly constructed by picking a pilot-power setting,  $p_c^k$ , from every  $P_c$ ,  $1 \leq k \leq |P_c|$ ,  $\forall c \in C$ . Notice how fast DE converged to a good solution. DASA also converged considerably fast, although not as fast as DE did. After a number of evaluations without improvement, DASA resets itself and continues searching from a new random point within the search space [92]. Despite this, and based on the flat profile the graph exhibits towards the end of the optimization run, it is clear that 100,000 evaluations is an adequate stopping criterion for all algorithms. The third algorithm, SA, slowly converged towards the best solution found, even though it was not as good as the solutions found by DE and DASA.

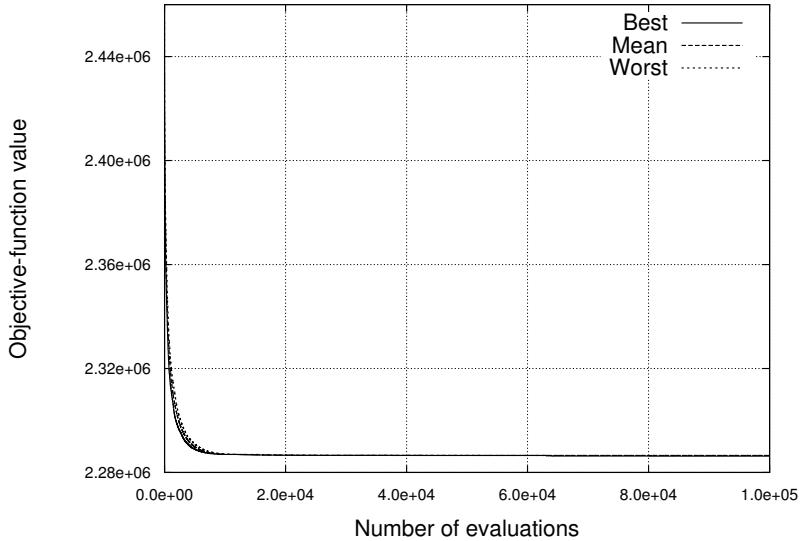


Figure 8.3: Convergence analysis for DE, showing the best, mean and worst objective-function values in each evaluation of the SHO-balancing problem.

The three convergence profiles shown in Figures 8.3, 8.4, and 8.5 give a clearer notion about the way these algorithms explore the search space of the SHO-balancing problem, and the variations occurred during the optimization process.

The simulation-running times have been intentionally omitted, since the algorithms are implemented in different programming languages and are therefore not suitable for comparison with each other.

#### 8.6.4 Performance analysis

Table 8.3 presents the analysis of the obtained results from the network point of view. After 30 independent runs of each of the three algorithms, the best results obtained were evaluated for the improvement and the decline of each of the measured network-performance aspects. The results are shown in Table 8.3, where '+' indicates improvement and '-' indicates a decline of a given criteria. Overall, it may be observed that the measured criteria have been significantly improved. The only exception is the measure for downlink SHO, without SHO

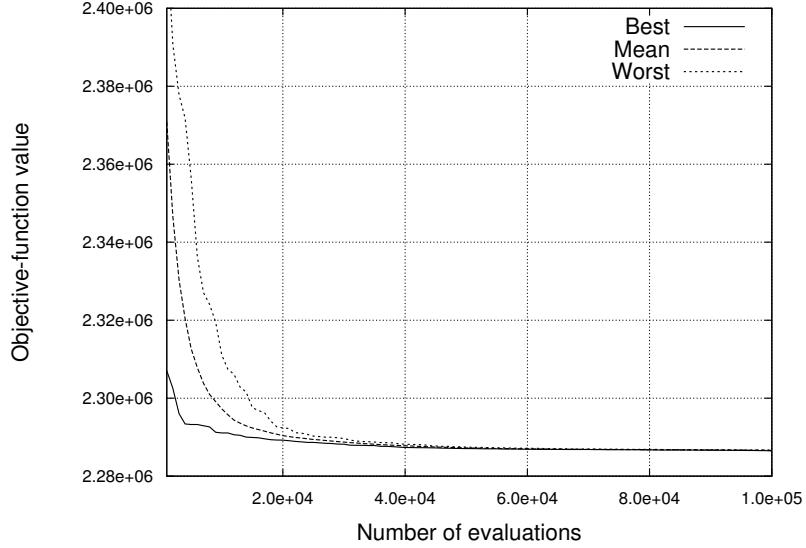


Figure 8.4: Convergence analysis for DASA, showing the best, mean and worst objective-function values in each evaluation of the SHO-balancing problem.

Table 8.3: Improvement analysis of the best solution that each algorithm achieved for the SHO-balancing problem. All values are expressed in percent.

	Uncovered	Covered, no SHO	SHO	$\text{no SHO}^\downarrow, \text{SHO}^\uparrow$	$\text{SHO}^\downarrow, \text{no SHO}^\uparrow$	Total
Before opt.	63.00	15.11	15.73	1.80	4.36	100.00
DE sol.	60.23	16.13	16.09	1.47	6.08	100.00
DASA sol.	60.24	16.16	16.90	1.46	5.24	100.00
SA sol.	60.42	16.55	15.97	1.56	5.50	100.00
DE impr.	+4.40	+6.75	+2.29	+18.33	-39.45	—
DASA impr.	+4.38	+6.95	+7.44	+18.88	-20.18	—
SA impr.	+4.09	+9.53	+1.52	+13.33	-26.15	—
Avg. impr.	+4.29	+7.74	+3.75	+16.85	-28.59	—

in the uplink (labeled as ‘ $\text{SHO}^\downarrow, \text{no SHO}^\uparrow$ ’), which shows an expected decline, since it is the optimization aspect with the lowest penalty-factor value.

The coverage has been improved with an average of 4.29 %, whereas the coverage area where there is no SHO capability, has been increased 7.74 % on average. Areas where SHO is available in both downlink and uplink have also been improved, i.e., 3.75 % on average. This particular improvement is interesting from the optimization point of view, because it had no explicit penalty factor set. Therefore, this may be understood as a consequence of the correct representation of the different network aspects in the objective function.

The second most important optimized aspect in the SHO-balancing problem is the proportion of areas with uplink SHO and no SHO in the downlink (labeled as ‘ $\text{no SHO}^\downarrow, \text{SHO}^\uparrow$ ’ in Table 8.3). This particular condition has been improved by almost 17 % on average, greatly reducing the possibility of interference in neighboring cells when serving HSPA traffic. The last measured aspect takes into account areas with downlink SHO and no SHO in the uplink (labeled as ‘ $\text{SHO}^\downarrow, \text{no SHO}^\uparrow$ ’ in Table 8.3). This condition, although it has not improved, does not expose the mobile network to malfunctioning, only to reduced through-

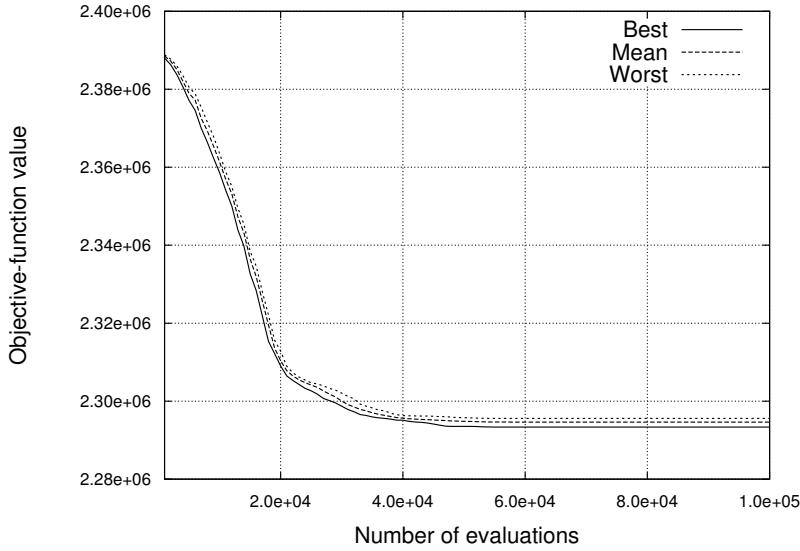


Figure 8.5: Convergence analysis for SA, showing the best, mean and worst objective-function values in each evaluation of the SHO-balancing problem.

put within these specific areas. However, the reduced throughput is relative, since there are many cells capable of serving HSDPA data access, as the downlink SHO condition confirms. For this reason, the serving cell should not only deliver HSDPA, but also take care of the user signaling and power control, received in the uplink. Clearly, this is only feasible in areas where uplink coverage is guaranteed.

It is worth mentioning that the simulation results were obtained for a real radio network with actual configuration data. Moreover, the hard constraints imposed to the optimization process (the pilot power limited to the  $\pm 2$  dB interval) ensure that the resulting configuration may be immediately applied to a mobile network. This fact can be contrasted with the spatial distribution of each of the optimized aspects, before and after applying the optimization results, as it is shown in Figures 8.6 and 8.7. The lack of any prominent visual change in Figures 8.6 and 8.7 is a desired consequence of the fine-tuning procedure the network has been exposed to. Still, the improvements are present precisely over the areas that are most exposed to malfunctioning due to unbalanced SHO, e.g., the cell-coverage borders.

## 8.7 Summary

This chapter formally introduced a new optimization problem for 3G networks: the SHO-balancing problem. A characterization of the consequences that unbalanced SHO areas have on the quality of HSPA services was also given. Particularly, tackling the SHO-balancing problem was possible due to the improved performance delivered by the evaluation framework PRATO (see Chapter 6).

Using an extension of the radio-network model presented in Section 7.3, the penalty scores of the objective function were set according to the configuration and layout of a real mobile network, deployed in Slovenia by Telekom Slovenije, d.d.

The balancing problem has been tackled by three optimization algorithms, namely DE, DASA and SA. All three algorithms were able to improve the given network configuration, DE being the most successful one. The presented results confirm that a great proportion of the SHO areas that were not balanced before the optimization were corrected, therefore

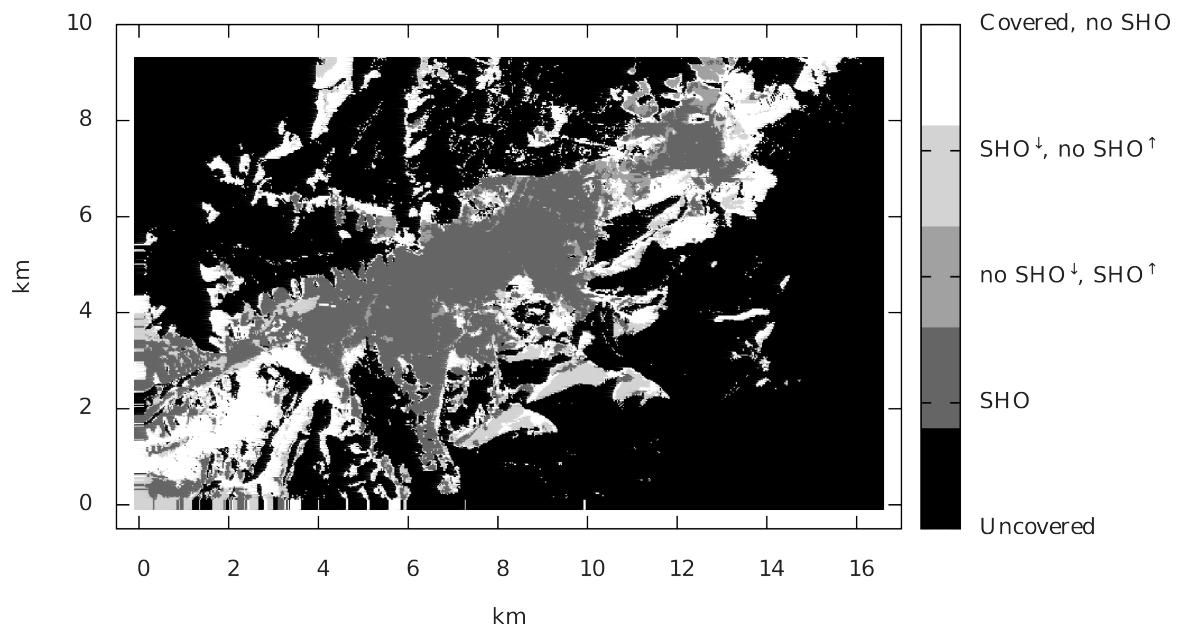


Figure 8.6: Spatial distribution of the SHO areas before the optimization.

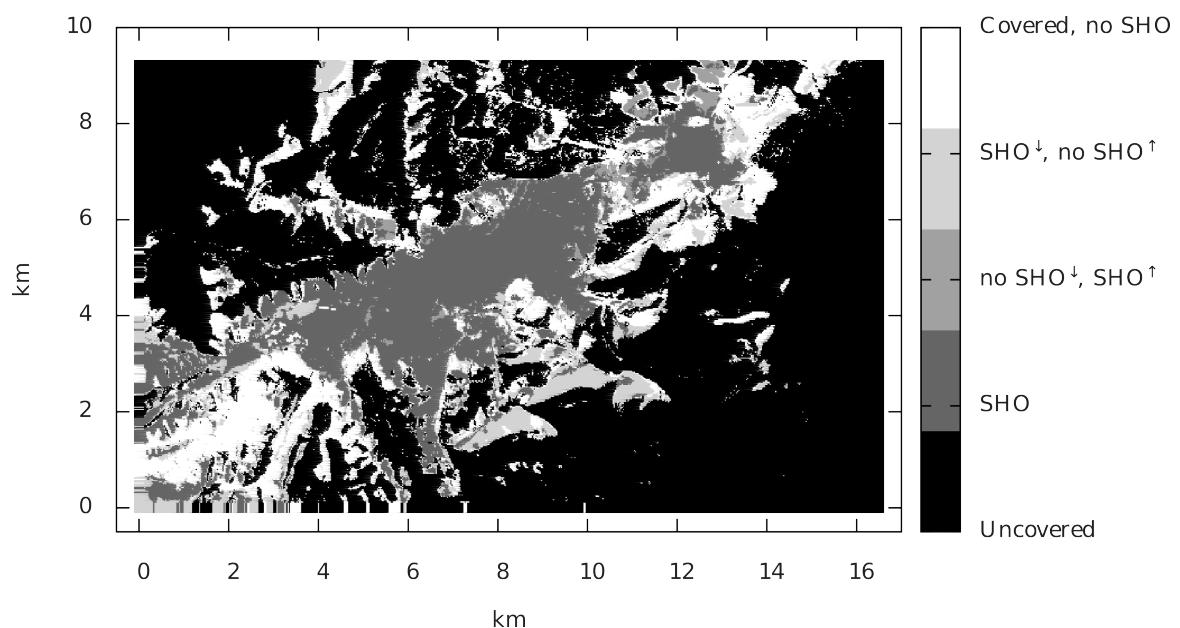


Figure 8.7: Spatial distribution of the SHO areas after the optimization.

significantly reducing the possibility of HSPA-service failures. Additionally, radio coverage was improved, while all other essential network services were not altered.

One of the key advantages of the presented method is that it targets the optimization of a deployed network, for which the focus is to fine-tune the existing configuration instead of creating complete new solutions. Furthermore, a deployed network has a great number of hard-constraints that should be taken into account at the optimization stage. Yet, the presented approach is simple and versatile enough to be used in practically any working UMTS network. Moreover, the introduced model is applicable for mobile networks in heterogeneous environments, because it imposes no restrictions regarding cell layout or radio-propagation characteristics, which are completely adaptable through PRATO.

It is important to note that, to the best of the author's knowledge, the formal definition of the SHO-balancing problem as presented here has not yet appeared in the related literature. Also note that although some methods proposed in this chapter have been particularly designed for problems that emerge during the planning of 3G radio networks, they may be adapted to other standards, e.g., 2G and 4G, without loss of generality.

## 9 Framework Automated Tuning

The models and algorithms from Chapter 6 constitute the basis for the radio-coverage framework used throughout this thesis. Until here, the framework has been studied in the context of optimization-problem solving for radio networks. In this chapter, the focus is shifted towards radio-network planning activities, and how the framework can aid the network-planning engineer in his or her everyday tasks. The objective is to facilitate refined network planning, the complexity of which is generally beyond the scope of any manual approach.

A central part of a radio-planning tool is its radio-propagation model. Generally speaking, the signal-propagation predictions will be as accurate as the input data used for their estimation. Moreover, acquiring and constantly upgrading the necessary data to support the decision making in this context is an expensive and challenging task. In practical situations, an emitted signal propagates by interacting with the surrounding environment. Consequently, the ability of a propagation model to adapt to the environment where it is used improves the accuracy of the calculated signal-propagation predictions.

In this context, this chapter presents two automated-tuning capabilities for PRATO. The first one involves the parameter tuning of the empirical radio-propagation model using a snapshot of field measurements. The second one involves the optimization of clutter losses over different regions of the country, therefore adapting the loss factors to the local conditions of each region. The results of the experimental simulations, performed over three regions of a real LTE network deployed by Telekom Slovenije, d.d., show the suitability of the presented methods to improve the accuracy of the calculated radio-propagation predictions.

To the best of the author's knowledge, there is no reference in the literature of an optimization-based approach to automatically adapt the signal losses due to clutter.

The rest of this chapter is organized as follows. Section 9.1 describes the benefits of the presented approach from the radio-planning point of view. After giving an overview of relevant publications in Section 9.2, the parameter-tuning problem and two solution approaches are presented in Section 9.3, including the experimental simulations and their results. Section 9.4 concentrates on the description of the optimization problem involving the regional adaptation of signal losses due to clutter, including an extensive analysis of the performed simulations on real LTE networks.

### 9.1 Motivation

With the advent of LTE as part of the 4G in cellular technology, mobile operators are facing the challenges of deploying a new network. LTE follows the well established UMTS/HSPA combo, targeting higher peak data rates, higher spectral efficiency and lower latency [145].

The deployment of a new radio network is always a challenge for mobile operators, who constantly struggle to find the optimal investment in order to provide a competitive network in terms of coverage and QoS. Indeed, coverage planning remains a key problem that all operators have to deal with.

Although different mathematical models have been proposed for radio-propagation modeling, none of them excels in a network-wide scenario [136]. Empirical propagation models

usually give good results with a limited computational effort. However, for improved accuracy, the model parameters have to conform to a specific environment within the network, mainly because of the inaccuracies in input data and the environmental changes in the region, e.g., foliage of trees or snow. Consequently, a combination of different parameters is generally needed in order to reliably calculate radio-propagation predictions of large networks that cover different environments.

To address the aforementioned issues, the parameters of an empirical propagation model are adapted based on a set of field measurements. A subset of parameters is analytically tuned by computing their values per cell, in order to increase the accuracy of the calculated predictions. Also, the full set of parameters is optimized by applying a metaheuristic approach. Moreover, by applying the same optimization algorithm, the signal losses due to clutter are automatically adjusted in a regional basis.

As a simulation framework to evaluate the presented problems, PRATO, the parallel radio-prediction tool presented in Chapter 6, is used. Therefore, the suitability of the framework for planning and optimization of LTE networks is also validated. Specifically, PRATO should be capable of handling a large number of radio-propagation predictions using a metaheuristic algorithm, and a distributed objective-function evaluation.

## 9.2 Related work

Following an optimization-oriented approach, the authors of [3] studied the effects on location accuracy while performing semi-automated optimization of the parameters of a radio-propagation model. While their optimization component improves the accuracy of the radio predictions, it does so requiring human intervention, hence the term semi-automated optimization. In terms of the effects of location accuracy, they concluded that locations with a mean accuracy of 60 m may be used for parameter tuning. Also, they notice that although the model accuracy improved after the parameter tuning, it gives inadequate results when used for predicting radio propagation over distant areas.

## 9.3 Parameter tuning of the radio-propagation model

The effectiveness of the decision-making process during radio-network planning is tightly coupled with the precision achieved by the propagation model used. In order to obtain a radio-propagation model that most accurately reflects the propagation characteristics of the area covered by each radio cell in the network, the parameters of the mathematical model are adapted to the target environment. Current state-of-the-art methods for such parameter tuning depend on existing field-measurement data [3, 173], which are collected in advance for the area covered by the target network. Starting from an a-priori best-known set of parameters, manually calculated by the radio engineers, this approach adapts the model parameters so that the deviation of the radio-propagation prediction to a given set of field measurements is minimized.

To calculate the radio-propagation predictions, the OA part of the empirical model, previously introduced in Section 6.3.2, is used. Recall that the model contains a vector of adaptable parameters,  $\vec{\beta}$ . For this reason, this mathematical model is especially appropriate for tuning, since it can be adapted to a given scenario and its local conditions by adjusting the values of the vector  $\vec{\beta} = (\beta_0, \beta_1, \beta_2, \beta_3)$ , the elements of which represent:

$\beta_0$  the reference loss or offset,

$\beta_1$  the loss slope due to distance of the receiver from the transmitter,

- $\beta_2$  the loss slope due to height of the transmitter antenna, and
- $\beta_3$  the loss slope due to the combined effect of the distance and height of the antenna.

The parameter tuning is performed per cell to improve the local fitting of the radio predictions, being its resulting solution a vector  $\vec{\beta}_c^*$  for a target cell  $c$ ,  $c \in C$ .

### 9.3.1 Field measurements

In radio networks, a UE constantly performs cell selection/reselection procedures and handover (see Section 3.1), in order to keep the best possible connection to the network. Within this context, the best connection is selected by probing a QoS measure of the neighboring cells. In LTE networks, the UE measures two parameters from the reference signal of the network, namely the Reference Signal Received Power (RSRP) and the Reference Signal Received Quality (RSRQ).

For a certain frequency bandwidth, RSRP measures the average received power over the resource elements that carry cell-specific reference signals. RSRP is applicable in both idle mode (e.g., waiting for a call) and connected mode (e.g., during a call). During the procedure of cell selection/reselection in idle mode, RSRP is used, whereas RSRQ is only applicable when the UE is in connected mode.

The radio-propagation prediction involves calculating the network coverage over a certain region. Hence, in the first place, the focus is on accurately predicting the best connection a UE would select in idle mode and the RSRP measurements it uses.

Here, the field measurements representing the RSRP at a given location were collected using a small truck that was equipped with a spectrum analyzer, the functionality of which supports LTE signal analysis. The spectrum analyzer was connected to an external omni antenna mounted on the roof of the truck, at roughly 2 m above the ground, taking measurements at a rate of 2 Hz. To accurately establish the measurement-location points, a Global-Positioning-System (GPS) unit was used. These GPS-informed locations have been tested to be compliant with the 60 m limit mentioned in [3]. The measurements cover most of the streets within the target area, with over 300,000 individual points collected for more than 190 network cells.

To minimize the error in measured RSRP values, and the impact that small-scale fading has in larger-scale path loss [47], all field measurements were post-processed so that a single value, the median, was calculated for each of the measured locations. This processing step improves measured-data quality in terms of possible deviations due to external factors, e.g., vehicle driving speed. The resulting RSRP was then used to estimate the path-loss prediction at the corresponding location, the resolution of which matches that of the DEM and clutter data used.

### 9.3.2 Linear least squares

The analytical approach for tailoring the radio-prediction model is to correlate the RSRP field measurements with the predicted radio-propagation. The new parameter set originates from the minimization of an error criterion. Similar to [3, 84, 173], the minimization criterion is the squared-sum difference between the predicted and the observed RSRP levels, i.e.:

$$E(\vec{\beta}_c) = \sum_{fm \in FM_c} (p_c - L_c(\text{coord}(fm), \vec{\beta}_c) - fm)^2 \quad \forall c \in C, \quad (9.1)$$

where  $E(\vec{\beta}_c)$  is the observed error (in dB) for cell  $c$  given the parameters  $\vec{\beta}_c$ ,  $C$  is the cell set of the target network,  $p_c$  is the transmit power of cell  $c$ , and  $L_c(\text{coord}(fm), \vec{\beta}_c)$  represents

the path loss of cell  $c$  at the same geographical point of the field measurement  $fm$ , the set of which is  $FM_c$ . Note that  $E(\vec{\beta}_c)$  is independently calculated for each  $c \in C$ .

As a general rule when applying this approach, only the first two components of the vector  $\vec{\beta}$  are adapted, i.e.,  $\beta_1$  and  $\beta_2$ , whereas the values of  $\beta_3$  and  $\beta_4$  are kept constant [84, 173]. Therefore, the analytical method consists in fitting the following expression, which is reduced from the path-loss definition presented in Equation (6.1):

$$\Delta L(x, y, \vec{\beta}) = \beta_1 + \beta_2 \log(d_{(x,y)}). \quad (9.2)$$

The expression in Equation (9.2) does not take the terrain height into account, which is feasible when the field measurements are taken at a roughly constant height relative to the base station [3, 42, 84, 173]. However, when these heights fluctuate within the coverage area of a cell, the other two parameters, i.e.,  $\beta_3$  and  $\beta_4$ , have an effect on the adaptation of the signal-propagation model, as it will be shown in the following sections.

### 9.3.3 Optimization objective

The objective of the parameter-optimization problem consists in adjusting the values of all four components of the vector  $\vec{\beta}$ , instead of only two as in the previous section, according to a set of field measurements of a given cell. Similar to the linear-least-squares approach, each network cell is independently optimized, so that its radio-propagation prediction minimizes the mean-squared error against the field measurements, as defined in Equation (9.3).

$$f_{\text{param}}^*(c) = \min \sum_{fm \in FM_c} \frac{(p_c - L_c(\text{coord}(fm), \vec{\beta}_c) - fm)^2}{|FM_c|} \quad \forall c \in C, \quad (9.3)$$

where  $f_{\text{param}}^*(c)$  is the optimization objective to be minimized for cell  $c$ ,  $FM_c$  is the set of all field measurements of cell  $c$ , the cardinality of which is  $|FM_c|$ , and the other parameters are as denoted previously for Equation (9.2). Note that  $f_{\text{param}}^*(c)$  is independently calculated for each  $c \in C$ .

### 9.3.4 Differential ant-stigmergy algorithm

The chosen optimization algorithm for the parameter-optimization problem is the DASA (see Section 2.3.2). The mapping between this problem and the DASA is as defined in Equation (9.4).

$$X_a = \{x_1, x_2, x_3, x_4\}, \quad (9.4)$$

where  $X_a$  is the solution vector of ant  $a$  during the minimization process, and  $x_j$ ,  $1 \leq j \leq 4$ , represents the  $j$ -th component of vector  $\vec{\beta}$  for the signal-propagation model of a given cell. At the end of every iteration, and after all the ants have created solutions, they are evaluated to establish if any of them is better than the best solution found so far.

There are several reasons for choosing the DASA as the optimization algorithm in the context of this problem. First, the benefits of metaheuristic algorithms for solving optimization problems, particularly in the context of radio networks, was demonstrated by several authors in general [20, 55, 84, 105], and in this thesis in particular (see Chapters 4, 7 and 8). Second, in [92], the authors validated the suitability of the algorithm for solving numerical-optimization problems.

After some trial-optimization runs, the six parameters that control the way the DASA explores the search space were set to the following values:

- $m = 80$ , the number of ants;

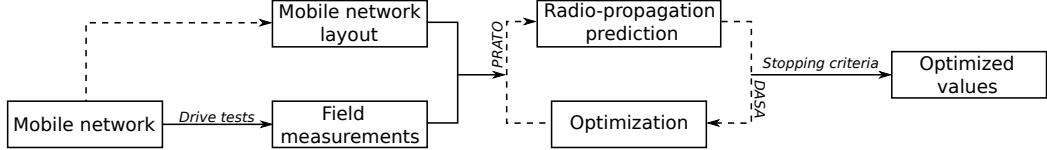


Figure 9.1: The automatic-tuning system, including its elements and their relationships.

- $b = 10$ , the discrete base;
- $\rho = 0.2$ , the pheromone dispersion factor;
- $s_+ = 0.01$ , the global scale-increasing factor;
- $s_- = 0.01$ , the global scale-decreasing factor; and
- $\epsilon = 10^{-5}$ , the maximum parameter precision.

The trial runs consisted in doubling  $m$  from 5 to 640, and verifying the convergence profile and best solution found. The values of the other parameters were left unchanged.

### 9.3.5 Simulations

The first simulation round consisted in building the matrices of the observed-error values,  $E(\vec{\beta}_c)$ , for each cell  $c$  in the target network. The linear systems of equations were then individually solved by applying the linear least-squares method, which involves the evaluation of one radio-coverage prediction per cell. Each solved system holds a unique solution for a cell  $c$ ,  $c \in C$ , denoted by the vector  $\vec{\beta}_c^*$ .

The second simulation round involved multiple iterations of several steps. The process began with a cell set of a mobile network and a set of field measurements, which were gathered before-hand through drive tests. Together with the mobile-network configuration and layout, they provide the input data for the optimization process itself. An iteration begins when the DASA generates a solution vector for each of the ants in the colony. The following step involves the parallel evaluation of the solution vector carried by an ant, i.e., one radio-propagation prediction per worker process of PRATO. The objective-function value is calculated as defined in Equation (9.3), and sent back to the DASA for it to generate the next set of solutions. The optimization process involves multiple iterations, which are repeated until some stopping criteria are met. Then, the best solution found represents the optimized values of the parameters of the radio-propagation model of the target cell. Figure 9.1 depicts the described system elements and their relationships for the parameter-optimization problem.

Compared with the analytical approach, that requires the solution of a linear system of equations, a large number of evaluations is needed for the DASA to converge to a solution. Therefore, it is essential to exploit the parallel nature of PRATO in order to evaluate the radio-coverage prediction of multiple cells within a network. Otherwise, such approach would not be feasible, since the time required to reach a reasonable solution would be excessive.

The stopping criteria for the optimization runs limited the maximum number of objective-function evaluations, the value of which was set to 50 for all the test networks.

### Test networks

The test networks, Net<sub>8</sub>, Net<sub>9</sub>, and Net<sub>10</sub>, are subsets of a real LTE network deployed in Slovenia by Telekom Slovenije, d.d. The path-loss predictions are calculated using PRATO,

Table 9.1: Parameters of the test networks Net<sub>8</sub>, Net<sub>9</sub> and Net<sub>10</sub> that were used for the experimental simulations during the automated tuning of the radio-propagation model.

	Number of cells	Transmission radius [km]	Area type	Area [km <sup>2</sup> ]	Drive-tests area [%]
Net <sub>8</sub>	12	16.00	Rural	82.90	4.75
Net <sub>9</sub>	156	8.00	Urban	468.30	22.31
Net <sub>10</sub>	26	16.00	Hilly	133.47	10.92

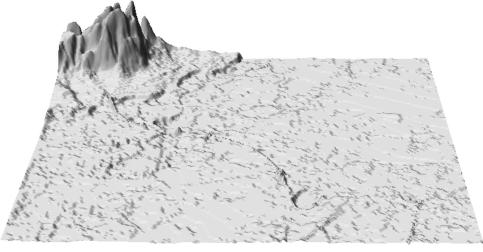


Figure 9.2: Terrain profile of the test network Net<sub>8</sub>, dominated by a flat agricultural area.

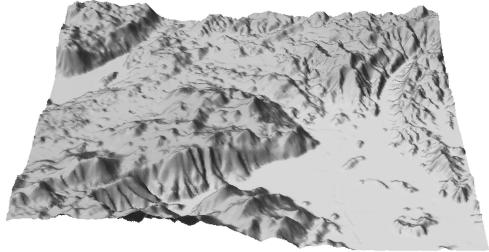


Figure 9.3: Terrain profile of the test network Net<sub>10</sub>, dominated by forested hills.

with a DEM and clutter map of 25 m<sup>2</sup> resolution, and a receiver height of 2 m above ground level. A transmission radius defines the coverage-prediction area around each network transmitter, thus limiting the path-loss prediction to a distance where it is feasible for a UE to connect to a cell, with a RSRP greater or equal to -124 dBm [115]. At the same time, the selected transmission radius provides enough overlap among neighboring cells to calculate the network coverage over the whole region. Table 9.1 provides detailed information about the test networks used, including the number of network cells, the area surface, and the covering proportion of the collected field measurements in terms of the total area of each test network.

Net<sub>8</sub> represents a network deployed over a dominant agricultural area with almost flat terrain, some forests and water streams. Net<sub>9</sub> is deployed over a densely populated urban area, containing high buildings, parks and avenues. The last one, Net<sub>10</sub>, represents a network deployed over hilly terrain, including some smaller villages and vast forests.

Since the terrain profiles are most relevant for Net<sub>8</sub> and Net<sub>10</sub>, they are shown in Figures 9.2 and 9.3, respectively. Note that the terrain shown in Figure 9.2 is mostly flat, since the agricultural area prevails in Net<sub>8</sub>. In contrast, the terrain for Net<sub>10</sub> is dominated by hills, which are mostly covered by forests, including some small villages in the valleys (see Figure 9.3).

In the following, the manually-calculated, default parameters of the radio-propagation model correspond to the values provided by the engineers of the Radio Network department at Telekom Slovenije, d.d.

The simulations were carried out on several computing nodes of the previously presented DEGIMA cluster [73] at the NACC of the Nagasaki University in Japan (see Section 6.5). Groups of 4, 12, and 40 nodes were used for executing the simulations of the different problem instances, i.e., Net<sub>8</sub>, Net<sub>9</sub> and Net<sub>10</sub>, respectively.

### 9.3.6 Results

The results of calculating the radio-propagation predictions with different parameter sets are presented in this section. The analysis to fit the parameters of the radio-propagation model to a set of field measurements is given with the default parameters, after applying the linear least-squares method, and after applying the metaheuristic-optimization method. The mean and standard-deviation values were calculated correlating the field measurements with the radio-propagation prediction of each test network.

Bar charts were prepared to show the probability distribution functions (PDFs) of the difference between the radio-propagation predictions and the field measurements (see Figures 9.4, 9.5, and 9.6). Each bar represents an open interval denoting the proportion of measurement points that deviate from the prediction in the given number of dB.

Figure 9.4 (a) depicts the PDF of the coverage prediction for test network Net<sub>8</sub> using the empirically-calculated parameters, Figure 9.4 (b) shows the difference distribution for the same test network, but using the analytically-calculated parameters, and Figure 9.4 (c) shows the difference distribution using the DASA-optimized parameters. Notice how the difference distributions show an improvement when the analytically-calculated parameters are used, lowering the largest (outer) deviations, and raising the lowest (inner) ones. Additionally, the difference is negligible when using the optimized parameters, thus confirming that it is sufficient to only adapt  $\beta_0$  and  $\beta_1$  in environments where the height difference between the base station and the receiver (i.e.,  $\log(H_A)$  in Equation (6.1)) is roughly constant, e.g., a rural area.

The PDFs test network Net<sub>9</sub> using the default, fitted and optimized parameters are shown in Figures 9.5 (a), 9.5 (b) and 9.5 (c), respectively. The fitted parameters represent a significant accuracy improvement when compared with the default ones. Similar to Net<sub>8</sub>, the optimized parameters show a negligible improvement when compared to the fitted ones. This is due to the height difference between the base-station antenna and the receiver having almost no variations throughout the area covered by the drive tests.

The difference distributions of the radio-propagation predictions for test network Net<sub>10</sub> using the manually-calculated parameters, the analytically-calculated, and the optimized ones are shown in Figures 9.6 (a), 9.6 (b), and 9.6 (c), respectively. Similar to Net<sub>8</sub>, the improvement appears in the largest deviations, since their values are lower than when using the manually-calculated parameters. In this case, we may also observe the improvement achieved by the DASA-optimized parameters, the optimization of which included all four components of the vector  $\vec{\beta}$ , thus better reflecting the signal propagation over the hilly terrain.

The presented results confirm that the parameter optimization of the signal-propagation model with respect to field measurements improves the quality of the calculated radio-propagation predictions. Considering the default parameter values were manually calculated by the radio engineers for the whole network, the convenience of the automated optimization procedure is clear. Indeed, these advantages are a consequence of a simpler method that delivers radio-predictions of superior quality, thus accurately representing the physical properties of a given environment.

### 9.3.7 Performance analysis

Table 9.2 shows the solutions reached by each of the compared approaches for the three test networks. The calculated mean and standard-deviation values are depicted in terms of the difference between the calculated radio-propagation predictions and the field measurements.

We may observe that the standard deviation improved for all test networks when using DASA as the optimization approach. Although the value of the mean measure is higher in

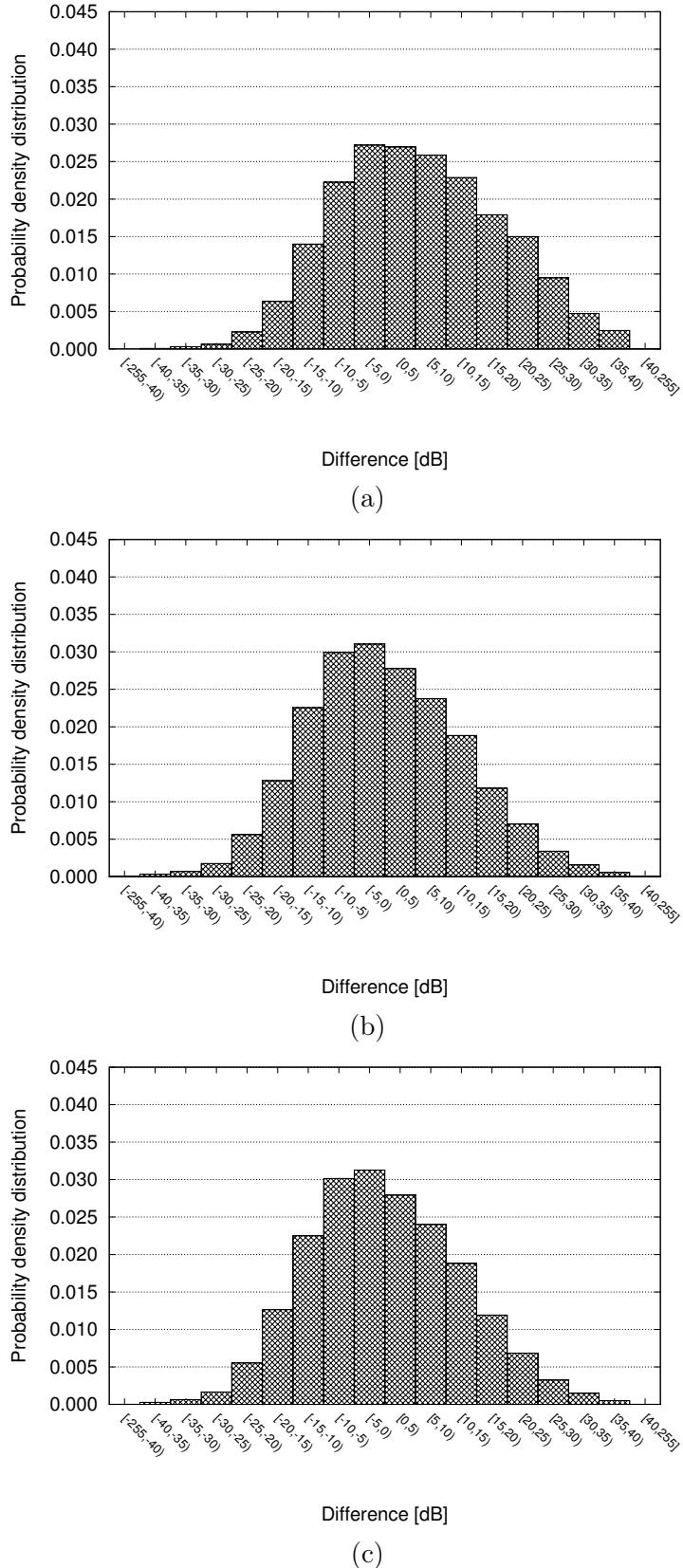


Figure 9.4: Probability-density distribution of radio predictions against the field measurements of network Net<sub>8</sub> over a rural area using the: (a) empirically-calculated parameters, (b) analytically-calculated parameters, and (c) optimized parameters.

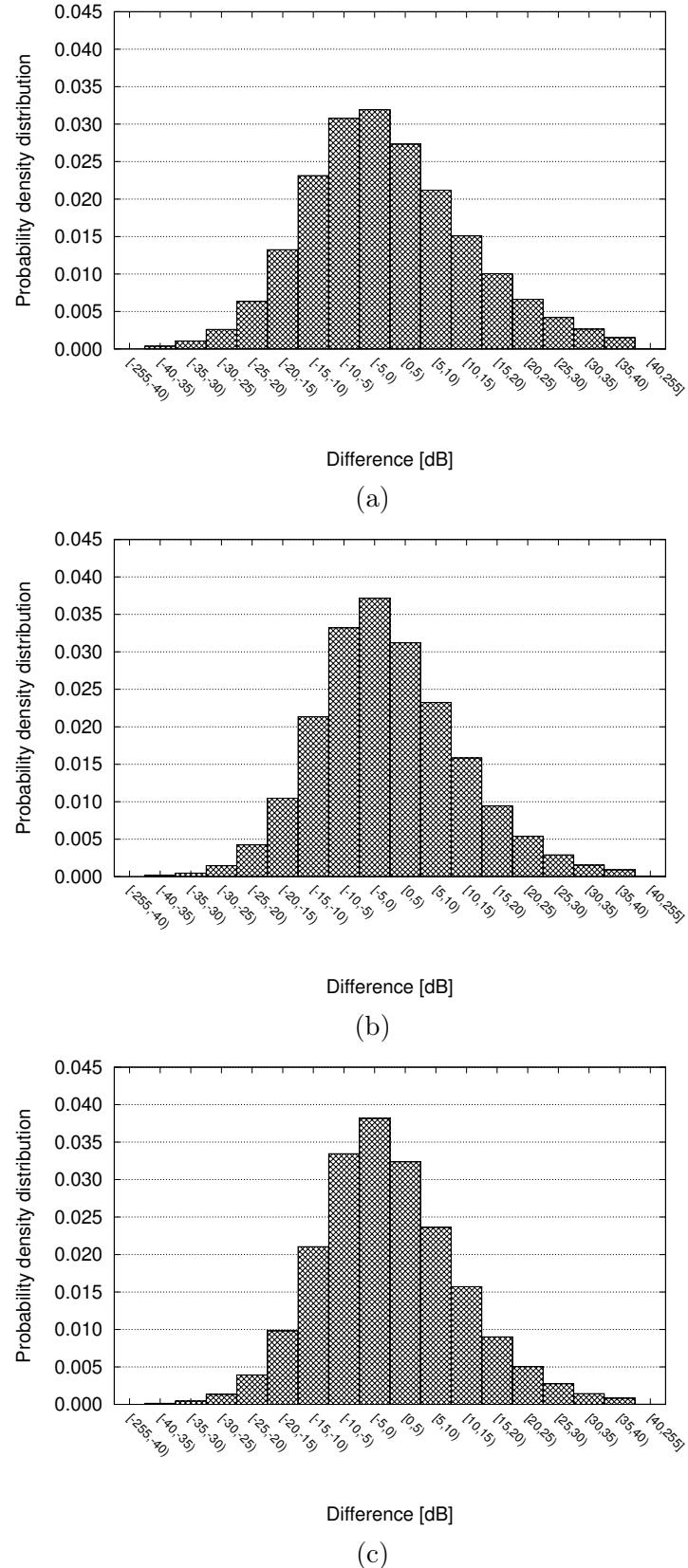


Figure 9.5: Probability-density distribution of radio predictions against the field measurements of network Net<sub>9</sub> over an urban area using the: (a) empirically-calculated parameters, (b) analytically-calculated parameters, and (c) optimized parameters.

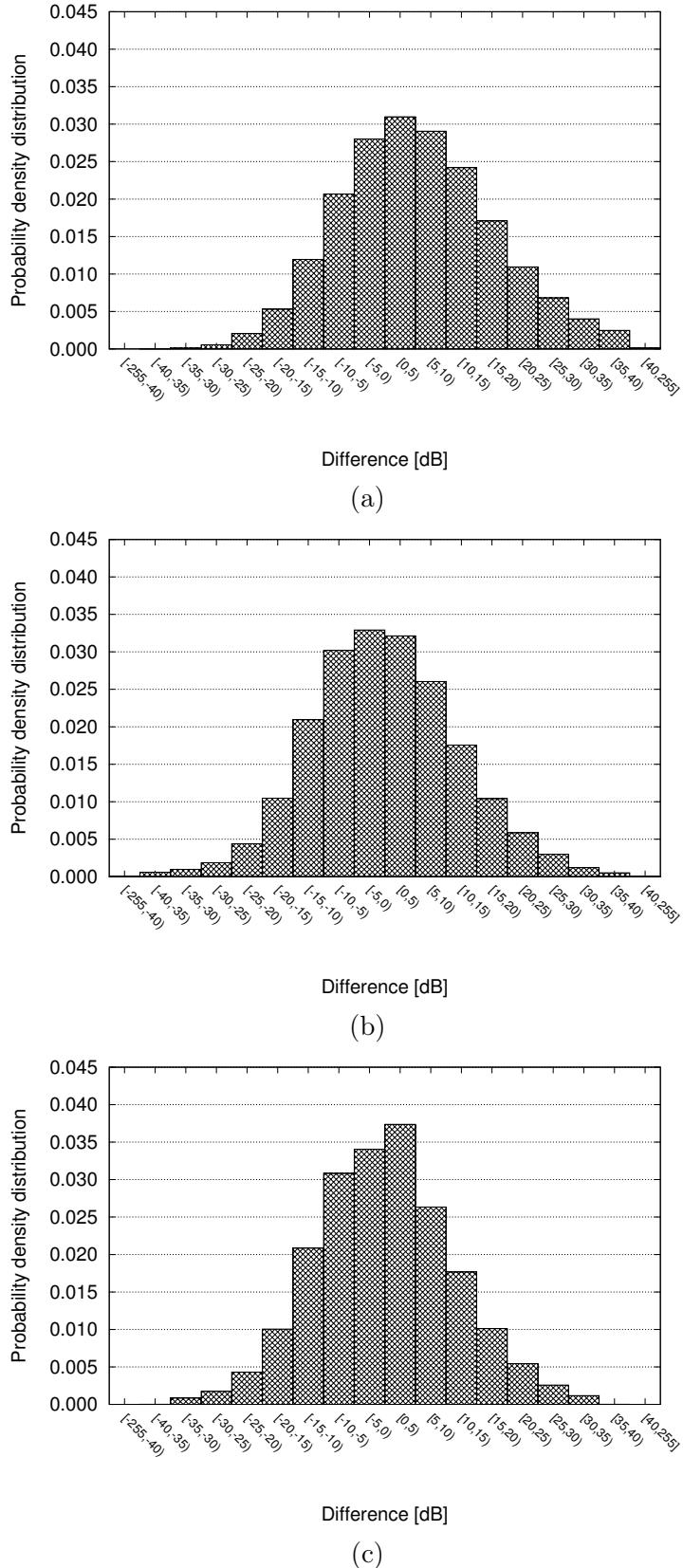


Figure 9.6: Probability-density distribution of radio predictions against the field measurements of network Net<sub>10</sub> over a hilly area using the: (a) empirically-calculated parameters, (b) analytically-calculated parameters, and (c) optimized parameters.

Table 9.2: Mean and standard-deviation values of the radio-propagation prediction against the field measurements. The values, expressed in dB, are given when using the manual, analytical and DASA approaches for each test network.

Test network	Manual		Analytical		DASA	
	Mean	Std. dev.	Mean	Std. dev.	Mean	Std. dev.
Net <sub>8</sub>	5.88496	13.69240	0.00001	12.64708	0.01974	12.13723
Net <sub>9</sub>	0.07172	13.94477	0.00001	12.13115	0.01504	11.82010
Net <sub>10</sub>	6.52300	14.35561	0.00007	12.30833	0.01372	10.59590

this case than when applying the analytical method, it is consistently lower than 0.05 dB, which is a negligible difference in the context of radio-propagation predictions. Regarding the lower standard-deviation measure for test network Net<sub>10</sub>, the value of which improved by almost 2 dB, it shows a significant gain in the accuracy of the radio-propagation predictions. This is especially important on the borders of cell coverage, where a 2 dB difference in the received-signal strength could mean predicting sufficient network coverage where there will be none.

## 9.4 Clutter optimization

In order to further improve the accuracy of the radio-prediction calculation over a given regional environment, the signal losses due to clutter are optimized in this section.

As it was mentioned before, there are several reasons for the predicted signal-loss values to be inaccurate. The seasonal changes are among them, like tree foliage and snow. Also, changes related to urban development, like demolition or construction of buildings and parks, and different kinds of forests or agricultural areas, etc. These changes are only noticeable through regular updates of accurate land-usage data. However, in the short term, updates are only available in the form of feedback by means of drive-tests campaigns for field-measurement gathering.

In the following, a metaheuristic algorithm is used to optimize the clutter losses of a target radio network. This is done over groups of network cells within different regions of the target network, e.g., agricultural, urban or hilly. In terms of coverage planning, a regional classification of the signal losses due to clutter improves the accuracy of the radio-coverage prediction.

Similar to the parameter tuning of the mathematical model presented in Section 9.3, the DASA metaheuristic algorithm is once again the tool of choice for optimizing the clutter losses.

Recall that in Equation (6.3), the model includes an extra term in order to adequately predict signal-loss effects due to foliage, buildings and other fabricated structures. These loss factors are based on the land usage, known as clutter data. Here, eleven different clutter categories are recognized. Table 9.3 lists these categories, including their label numbers and descriptions.

### 9.4.1 Optimization objective

The optimization objective consists of adjusting the loss values of the different clutter categories, i.e., LCLUT( $d_{(x,y)}$ ) as used in Equation (6.3), according to a set of field measurements of a given geographical region. The same three data sets used in Section 9.3.5 were used for the clutter-optimization problem: the first for Net<sub>8</sub>, the second for Net<sub>9</sub>, and the third for Net<sub>10</sub>. Each region was independently optimized, so that the radio-propagation predictions

Table 9.3: Clutter-category label numbers and descriptions for the signal loss due to clutter, as reflected by a radio-propagation model.

Clutter category	Description
1	Suburban area
2	Urban area
3	Dense urban area
4	Agricultural area
5	Forest area
6	Swamp area
7	Dry open land area with special vegetation
8	Dry open land area without special vegetation
9	Water area
10	Industrial area
11	Park area

of each set of cells per test network minimized the total mean-squared error against the field measurements, based on the previously defined objective function in Equation (9.3), i.e.:

$$f_{\text{clut}}^* = \min \sum_{c \in C} f_{\text{param}}^*(c, \vec{\beta}_c^*), \quad (9.5)$$

where  $\vec{\beta}_c^*$  is the optimized parameter vector for cell  $c$  as calculated previously in Section 9.3.6.

In this case, the mapping between the clutter-optimization problem and the DASA is as follows:

$$X_a = \{x_1, \dots, x_i, \dots, x_{11}\}, \quad (9.6)$$

where  $X_a$  is the solution vector of ant  $a$  during the minimization process, and  $x_i$  represents the  $i$ -th clutter category within a given region or network.

#### 9.4.2 Simulations

In this case, only one simulation round was performed for the DASA. The work flow of the optimization process, and the test networks used, are the same as for the parameter-optimization problem described previously in Section 9.3.5 and Figure 9.1.

For a clearer characterization of the test networks in terms of the available field measurements and terrain types, the proportions of each clutter category with respect to the total area covered by the drive tests are shown in Table 9.4.

In agreement with previous related research about clutter losses [129, 137], and since no clutter generates greater attenuation than 40 dB, we limited the optimized clutter losses to the interval [0,40].

The stopping criteria for the optimization runs was set by limiting the maximum number of objective-function evaluations. In this sense, the limits for Net<sub>8</sub> and Net<sub>10</sub> were set the 200, whereas for Net<sub>9</sub> the value was 500, since this network contains the largest number of cells. Overall, the framework completed 48,000 objective-function evaluations, i.e., 576,000 radio-coverage predictions for Net<sub>8</sub> and 288,000 for Net<sub>10</sub>, whereas for Net<sub>9</sub>, the number of objective-function evaluations was 120,000, for a total of 15,600,000 radio-coverage predictions.

Regarding the parameters that control the behavior of the DASA, they were set to the following values after some trial-optimization runs:

Table 9.4: Field-measurement proportions with respect to each clutter category, the percentages of which are given in terms of the area covered by the drive tests for each of the test networks. The clutter category legend is given in Table 9.3.

Clutter category	Net <sub>1</sub> [%]	Net <sub>2</sub> [%]	Net <sub>3</sub> [%]
1	12.77	9.23	8.34
2	27.57	35.99	31.17
3	3.10	7.34	4.15
4	48.24	33.05	41.13
5	5.19	8.29	8.59
6	0.00	0.00	0.10
7	0.00	0.05	0.00
8	0.39	0.02	0.04
9	0.21	0.58	3.00
10	2.53	4.74	3.37
11	0.00	0.71	0.11
Drive-test area	100.00	100.00	100.00

Table 9.5: Clutter-category losses after the optimization. The default losses for each clutter category are given along the solutions for each of the test networks. All values are expressed in dB.

Clutter category	Default	Net <sub>8</sub>	Net <sub>9</sub>	Net <sub>10</sub>
1	13.0	18.43	15.10	14.79
2	15.0	19.66	16.34	15.94
3	28.0	16.92	19.83	22.21
4	12.0	9.50	11.58	11.85
5	20.0	14.63	13.22	11.79
6	15.0	-	-	14.57
7	8.0	-	12.79	-
8	5.0	3.85	9.78	0.68
9	1.0	3.99	9.67	7.10
10	20.0	12.85	16.49	15.62
11	8.0	-	19.16	15.01

- $m = 240$ , the number of ants;
- $b = 10$ , the discrete base;
- $\rho = 0.2$ , the pheromone dispersion factor;
- $s_+ = 0.01$ , the global scale-increasing factor;
- $s_- = 0.01$ , the global scale-decreasing factor; and
- $\epsilon = 10^{-2}$ , the maximum parameter precision.

The trial runs consisted in doubling  $m$  from 15 to 480, and verifying the convergence profile and best solution found. The values of the other parameters were left unchanged.

### 9.4.3 Results

The results achieved by the optimization process are shown in Table 9.5. The solutions are given for each of the test networks, along with the manually-calculated (default) loss values. These default values were provided by the radio experts of the Radio Network department at Telekom Slovenije, d.d. Hyphens represent clutter categories for which there were no field measurements available. Consequently, it was not possible to evaluate the objective-function for them.

The optimized loss for the first clutter category, 1, representing suburban area, is above the default one for all test networks, indicating a building density above the average. For category 2, representing urban area, the optimized values are again above the default ones, clearly showing underestimation of the manual approach. Despite this, the relation with the suburban clutter losses is correctly kept, i.e., the signal losses are higher in each test network. In contrast with category 1, the optimized values of category 3, representing dense urban area, are lower than the default one. This indicates that dense urban areas in these regions have a lower density than the average case. Representing agricultural area, category 4 gets a value very close to the default one for Net<sub>9</sub> and Net<sub>10</sub>, whereas for Net<sub>8</sub> the value is lower, indicating that this type of land is mostly open there. As for category 5, representing forests, the results correspond with the type of forest dominating each of the test networks. Namely, Net<sub>8</sub> and Net<sub>9</sub> are dominated by dense forests presenting leave foliage, whereas in Net<sub>10</sub> the forests are mostly coniferous and more sparse. Keeping the default loss values for categories 6 and 7, category 8, representing open land without special vegetation, the optimized values of which vary significantly, mostly influenced by the low proportion of field measurements over this type of land in each region. As for category 9, representing water, the results of which indicate creeks and rivers in all regions are surrounded by forests (Net<sub>8</sub>), buildings (Net<sub>9</sub>), and hills and forests (Net<sub>10</sub>), since none of the regions lays by the sea. As for the industrial area, denoted by clutter category 10, lower loss values than the manually-calculated default appear. This indicates the presence of sparse industrial buildings in Net<sub>8</sub>, and a higher density of mostly commercial buildings for Net<sub>9</sub> and Net<sub>10</sub>. Finally, the last clutter category 11 could not be calculated for Net<sub>8</sub> since there are not parks there, whereas the values for Net<sub>9</sub> and Net<sub>10</sub> agree with the region type, i.e., urban and hilly, respectively.

It is important to note that the relation among different clutter categories within a region is correctly kept. For example, we may observe that the clutter loss for urban area (category 2) is higher than those for the suburban area (category 1), as well as for category 4 (agricultural area). Hence, the results reflect physically feasible losses, despite the higher deviation from the default losses shown by category 3, and the lower deviation for category 4, again with respect to the default losses. Such relations hold for different categories in all test networks, strongly suggesting the correctness of the applied optimization approach and the evaluation methodology used.

Based on these observations and after summarizing the extensive experimentation, the results obtained indicate that the use of a metaheuristic algorithm to perform automatic optimization of clutter losses is clearly applicable for radio-propagation predictions in LTE networks. Most notably, the proposed approach is capable of reflecting the physical phenomena appearing in real-world conditions within geographically-different network instances that contain a large number of cells.

Similar to Section 9.3.6, charts were prepared to show the PDFs of the difference between the signal-propagation predictions and the field measurements, the values of which are expressed in dBm. Figure 9.7 depicts the distribution of the difference between the coverage prediction for test network Net<sub>8</sub> using the default clutter losses, and using the optimized clutter losses (see Table 9.5, column Net<sub>8</sub>). Notice how the distribution shows an improve-

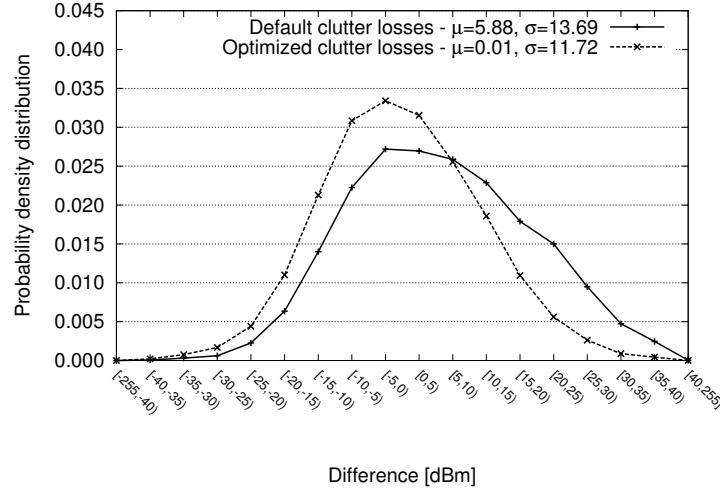


Figure 9.7: Probability density function of the difference between the radio-propagation prediction and the field measurements for network Net<sub>8</sub> over a rural area.

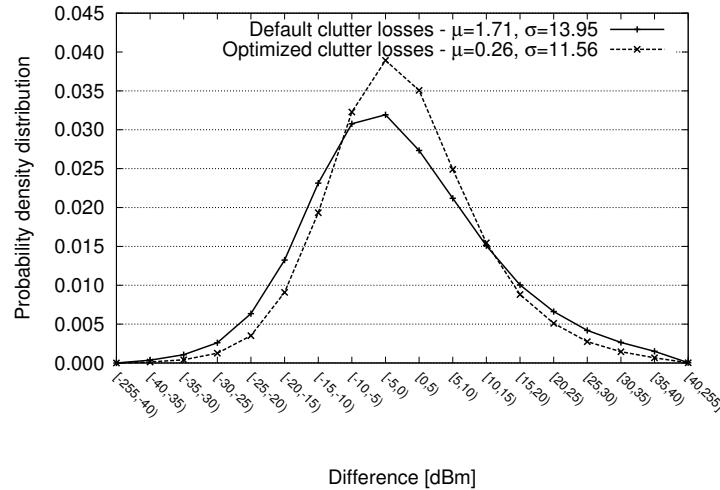


Figure 9.8: Probability density function of the difference between the radio-propagation prediction and the field measurements for network Net<sub>9</sub> over an urban area.

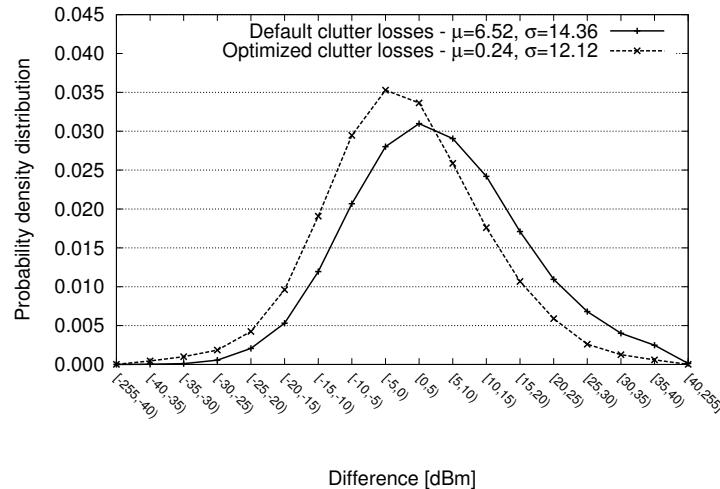


Figure 9.9: Probability density function of the difference between the radio-propagation prediction and the field measurements for network Net<sub>10</sub> over a hilly area.

Table 9.6: Statistical analysis of the solutions for the clutter-optimization problem. All values are expressed in dB. The corresponding box plots are depicted in Figure 9.10.

Cat.	Nets <sub>8</sub>				Net <sub>9</sub>				Net <sub>10</sub>			
	Min	Max	Avg	St.dev.	Min	Max	Avg	St.dev.	Min	Max	Avg	St.dev.
1	18.37	18.49	18.43	0.04	15.83	16.27	16.11	0.14	14.42	15.26	14.79	0.24
2	19.61	19.71	19.66	0.02	17.05	17.15	17.10	0.02	15.66	16.09	15.94	0.13
3	16.89	16.97	16.92	0.12	19.80	21.22	20.28	0.28	21.78	22.57	22.21	0.26
4	9.47	9.52	9.50	0.01	11.55	11.72	11.64	0.06	11.64	12.05	11.85	0.11
5	14.57	14.69	14.63	0.04	10.20	10.85	10.49	0.21	11.50	12.17	11.79	0.22
6	-	-	-	-	-	-	-	-	11.68	16.99	14.57	1.67
7	-	-	-	-	9.51	15.06	12.76	1.81	-	-	-	-
8	3.37	4.79	3.85	0.27	2.15	8.55	4.75	1.98	0.00	2.28	0.68	1.90
9	3.81	4.18	3.99	0.09	9.02	14.17	11.77	1.25	6.45	7.43	7.10	0.34
10	12.76	12.94	12.85	0.07	16.20	17.36	16.69	0.50	15.33	16.02	15.62	0.31
11	-	-	-	-	17.89	20.26	19.05	0.81	11.97	17.51	15.01	1.43

ment when the optimized clutter losses are used, lowering the largest (outer) differences, and raising the lowest (inner) ones. The mean of the PDF was reduced from 5.88 when using the default clutter losses, to 0.01 with the optimized losses; whereas the standard deviation was reduced from 13.69 to 11.72, respectively.

The difference distributions of the radio-propagation predictions for test network Net<sub>9</sub> using default and optimized clutter losses are shown in Figure 9.8. Again, an improvement can be noticed in its shape, clearly showing the favorable effects of the optimization process. In this case, the mean value dropped from 1.71 to 0.26, whereas the standard deviation was reduced from 13.95 to 11.56.

For the last test network, Net<sub>10</sub>, the difference distributions are depicted in Figure 9.9. Similar to Net<sub>8</sub>, the variations appear in the largest difference values, since they are lower than when using the default clutter losses. Regarding the improvement of the mean of the PDFs with the default and optimized clutter losses, they were 6.52, and 0.23, respectively; whereas the standard deviation was reduced from 14.36 to 12.12, respectively.

The presented results confirm that the optimization of clutter losses with respect to field measurements improves the accuracy of the calculated radio-propagation predictions. Considering the default clutter losses were manually calculated by the radio engineers for the whole network, there is a clear advantage for the presented automatic optimization procedure. It follows that the accurate representation of the physical properties in a given environment produces more precise radio-propagation predictions.

#### 9.4.4 Performance analysis

Because of the stochastic nature of the DASA, the results of 30 independent runs were collected in order to have enough data for them to be statistically relevant. In other words, the robustness of the solutions that were presented in the previous section is analyzed here.

To this end, Table 9.6 shows the solutions reached by the DASA for each of three test networks. The calculated signal losses are depicted with the minimum, maximum and average values for every clutter category, along with their standard deviations. Again, hyphens represent clutter categories for which there were no field measurements available, and thus they could not be optimized. In order to easily visualize the data shown in Table 9.6, box plots based on the same values are depicted in Figure 9.10.

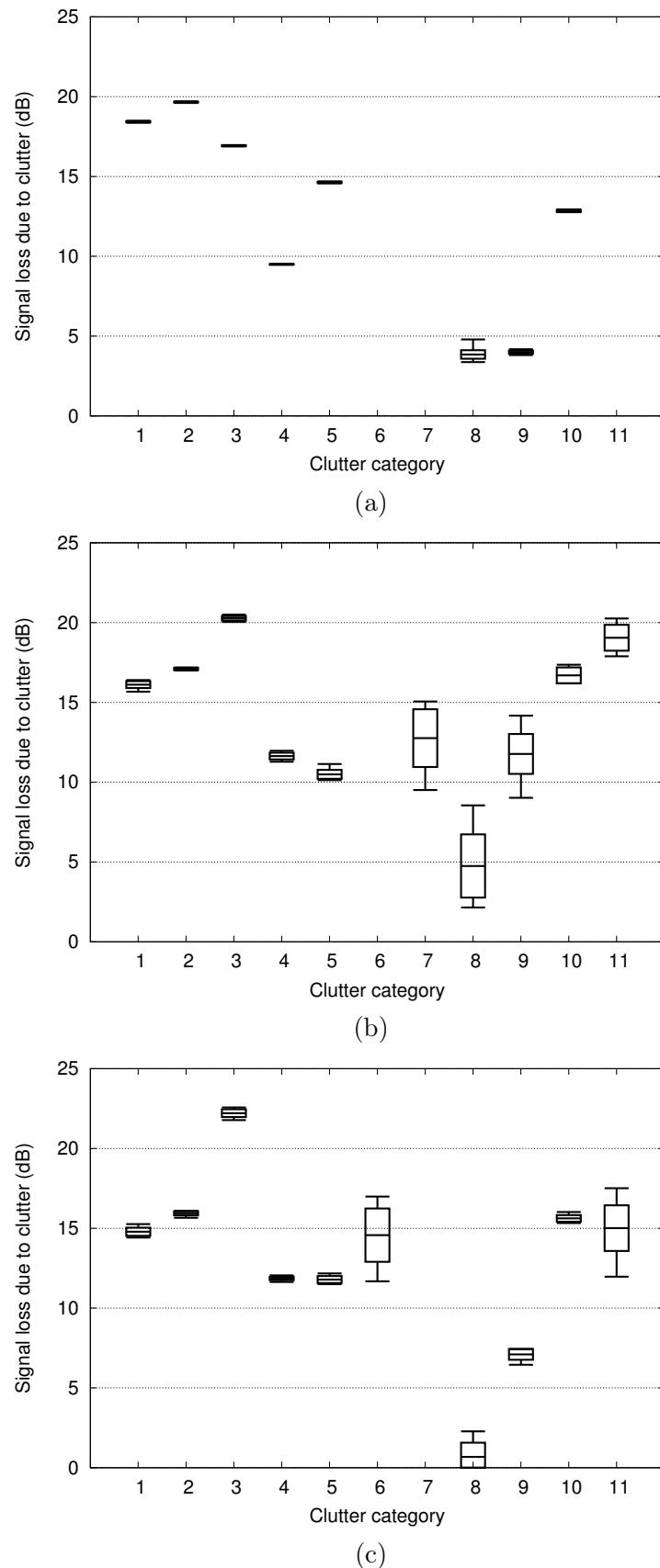


Figure 9.10: Box plots representing the statistical-analysis values of Table 9.6, for the solutions of the clutter-optimization process of each test network: (a) Net<sub>8</sub>, (b) Net<sub>9</sub>, and (c) Net<sub>10</sub>.

It can be observed that the standard deviation is low for almost all optimized categories, indicating a consistent convergence of the optimization algorithm. However, we notice some exceptions, e.g., categories 7 and 8 in Net<sub>9</sub>, and 6, 8 and 11 in Net<sub>10</sub>. There is a particularly low proportion of field measurements for the enumerated clutter categories (see Table 9.4). Indeed, higher standard deviation values consistently appear for clutter categories with fewer field measurements. Therefore, and based on the obtained results, the standard deviation of the optimized clutter losses provides a quantity indicator of the field measurements required by the optimization process, i.e., a higher standard deviation denotes more field measurements are needed over these areas in order to improve the optimization results the target clutter category.

## 9.5 Summary

This chapter illustrated the suitability of PRATO as a network-planning tool by tackling two coverage-planning problems, which were tested over the newly deployed LTE network in Slovenia. The first one involved the parameter tuning of the empirical radio-propagation model using a snapshot of field measurements. The second one considered the optimization of the clutter losses over different sub-networks of the country, therefore automatically adapting them to the local conditions of the geographical region of each network.

The combination of the aforementioned techniques with PRATO provides an environment-adaptable framework for radio-network planning. This delivers a tool with considerably improved solution accuracy of the analyzed instances, especially if compared to traditional, i.e., manual or semi-automated, coverage-planning methods.

Additionally, the simulation results indicate that PRATO is applicable for planning and optimization of real-world radio networks, since it is capable of simulating a large number of coverage configurations in a feasible amount of time. In particular, several of the presented instances show a large computational-time complexity, which is beyond reach for a serial implementation of an automated approach. Moreover, the parallelization capabilities provided by the framework create new problem-solving possibilities, together with the automation of tasks that have been traditionally performed manually by the network engineers. Furthermore, even computational-intensive tasks such as objective-function evaluations for stochastic optimization are feasible if PRATO is used.

To the best of the author's knowledge, an automatic-optimization method for the clutter losses of a geographical area has not yet been presented in the related literature.

## 10 Framework Verification

The objective of Chapter 9 was to ease the execution of radio-network planning activities, the complexity of which is generally beyond the scope of a manual approach. In this context, PRATO was presented as a tool that can help an engineer in realizing his or her everyday network-planning tasks.

Another important factor to further validate the adoption of the presented framework is to verify its accuracy. This is especially important in real-world scenarios, where it is not feasible to improve the performance of radio-coverage predictions at the cost of precision loss. The scope of this chapter is therefore to establish that the accuracy of the radio-propagation predictions of PRATO, as described previously in Chapter 6, is adequate for real-world, radio-network planning purposes. To this end, and with the help of the radio engineers at Telekom Slovenije, d.d., some real-world radio-propagation scenarios are calculated using PRATO and an enterprise, industrial software. During the experimentation phase, as well as when comparing the outcome, the engineers provided guidelines to assess the results from a practical radio-network planning perspective. The objective is to compare both tools in terms of solution quality and computational-time performance.

The rest of this chapter is organized as follows. Section 10.1 gives an overview of the reasons behind the verification of the framework. A description of the testing environment, including three networks, is given in Section 10.2, followed by an extensive performance analysis of the simulation results in Section 10.3.

### 10.1 Motivation

For a mobile operator, the utilization of a radio-network planning tool has clear economical and technical benefits. As it has been pointed out throughout the previous chapters, the usage of accurate planning tools minimizes the operator's costs and effort, and also automates manual processes. In this sense, the important role that a radio-planning tool has during the optimization process of a network was also presented.

However, up to this stage, little has been said regarding the reliability of PRATO as a tool for everyday coverage planning. In other words, a question has not yet been answered: is PRATO able to provide sufficiently accurate estimates of a network-coverage performance?

As it was mentioned in Section 6.1, the accuracy of the coverage predictions has a fundamental impact on the performance accuracy of the framework. For this reason, one of the objectives of this chapter is to assess the precision of several coverage-propagation predictions of PRATO and an enterprise radio-planning tool, the design of which is tailored to be used in industrial environments, using field measurements as a reference.

The second objective of this chapter is related to the the high-performance characteristics of PRATO. To this end, an insight to the computational-time performance of PRATO, compared to the commercial radio-planning tool<sup>1</sup>, is also given.

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<sup>1</sup>Due to the currently applicable business-secrecy policy of Telekom Slovenije, d.d., the author is not able to reveal any details about the commercial tool that was used during experimentation.

Table 10.1: Parameter values used during the coverage-prediction calculations for each test network. The same values were selected on PRATO and the commercial tool.

Parameter	Net <sub>11</sub>	Net <sub>12</sub>	Net <sub>13</sub>
frequency	900 MHz	2140 MHz	1800 MHz
receiver height	1.5 m	1.5 m	1.5 m
calculation radius	35 km	25 km	25 km

## 10.2 Radio-environment setup

Field measurements are used as the reference for analyzing the accuracy of a set of coverage-prediction calculations. The reference measurements were conducted in three commercial networks of different sizes and technologies. Namely, Net<sub>11</sub> denotes a GSM network that contains 830 BSs with 1240 cells, Net<sub>12</sub> is a UMTS network with 700 BSs and 2000 cells, and Net<sub>13</sub> is a LTE network featuring 120 BSs with 350 cells. The selected networks extend throughout diverse geographical regions, thus covering different terrains and representing various environmental characteristics, including urban, suburban and rural.

Using a similar setup as presented in Section 9.3.1, the average-received power was measured on the field for each network. The field measurements were captured with the air-interface measurement tool for the corresponding network technology. The captured measurements are lower-bounded by the receiver sensitivity of a given technology.

In order to minimize the error impact in the measured signals, all field measurements were processed so that a single value, the median, was calculated for each measured location. Similar to Section 9.3.1, this step improves the data quality in terms of possible deviations due to external factors during the measurement gathering on the field.

Regarding the radio-propagation models, the coverage predictions were calculated using the proprietary model of the commercial tool, and the previously introduced model for PRATO (see Section 6.3.2). Additionally, the coverage-prediction parameters were equally set in both tools (see Table 10.1).

Finally, in order to have comparable results, the same DEM and clutter data were used, with a resolution of 25 m.

## 10.3 Performance analysis

In this section, the performance of both tools is presented, the analysis of which is focused on coverage examination. The accuracy of PRATO as a radio-coverage prediction tool was investigated by comparing the simulation results and the field measurements. Its performance was investigated for different network types (GSM, UMTS and LTE) and terrains (hilly, almost flat rural, urban, and suburban).

The accuracy of PRATO can be verified from Figures 10.1, 10.2 and 10.3, the graphs of which present the simulation results of both tools for Net<sub>11</sub>, Net<sub>12</sub> and Net<sub>13</sub>, respectively. The graphs labeled as (a) show the analysis comparison between PRATO and the field measurements, whereas the ones labeled as (b) show the same analysis for the commercial tool. The field measurements used for this analysis are selected according to the spatial arrangement of the given coverage prediction.

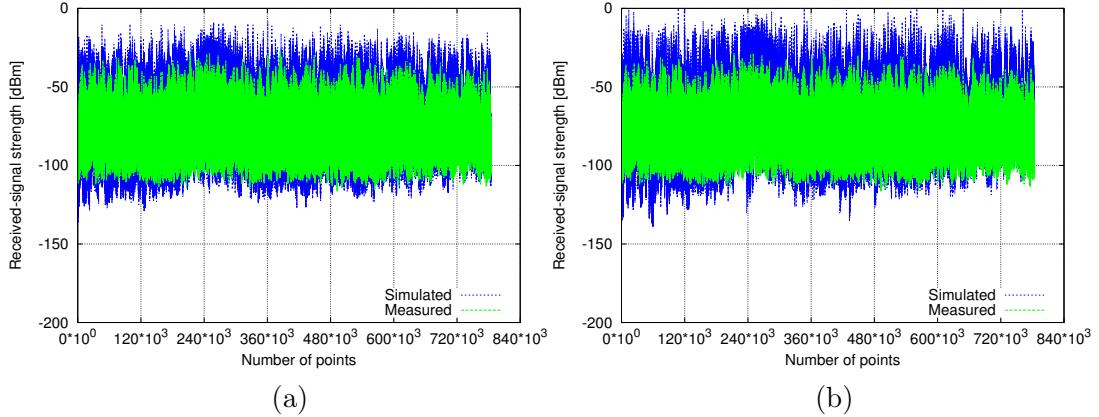


Figure 10.1: Net<sub>11</sub> distribution of the predicted received-signal powers (GSM) compared to field measurements for: (a) PRATO, and (b) the commercial radio-planning tool.

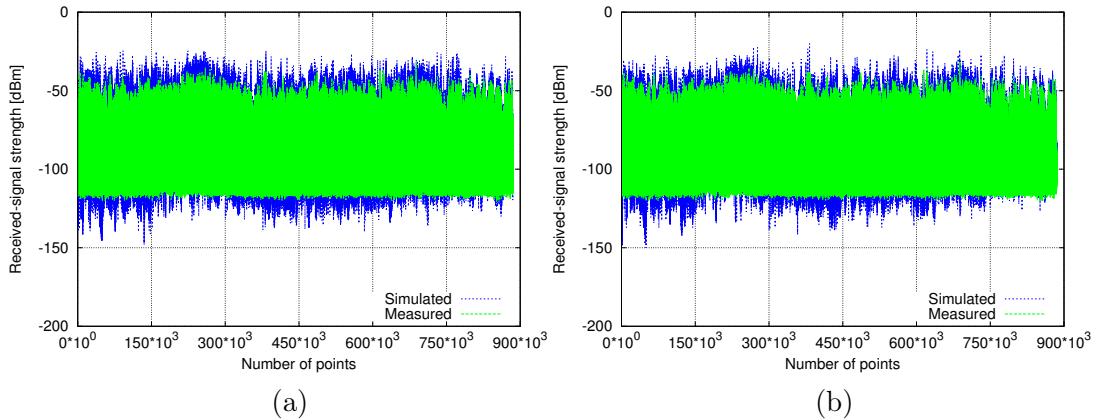


Figure 10.2: Net<sub>12</sub> distribution of the predicted received-signal powers (UMTS) compared to field measurements for: (a) PRATO, and (b) the commercial radio-planning tool.

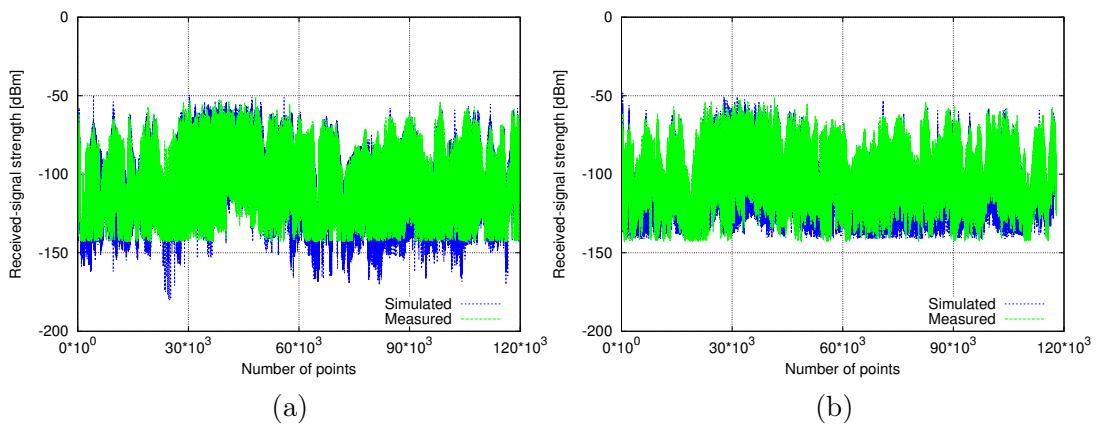


Figure 10.3: Net<sub>13</sub> distribution of the predicted received-signal powers (LTE) compared to field measurements for: (a) PRATO, and (b) the commercial radio-planning tool.

Figure 10.1 (a), that depicts the received-power levels for Net<sub>11</sub> (GSM), shows that the prediction results match the field measurements rather well. A similar result arrangement can be observed in Figure 10.2 (a) for Net<sub>12</sub> (UMTS), whereas the prediction results for

Net<sub>13</sub> (LTE) show a slightly increased deviation from the field measurements, as presented in Figure 10.3 (a).

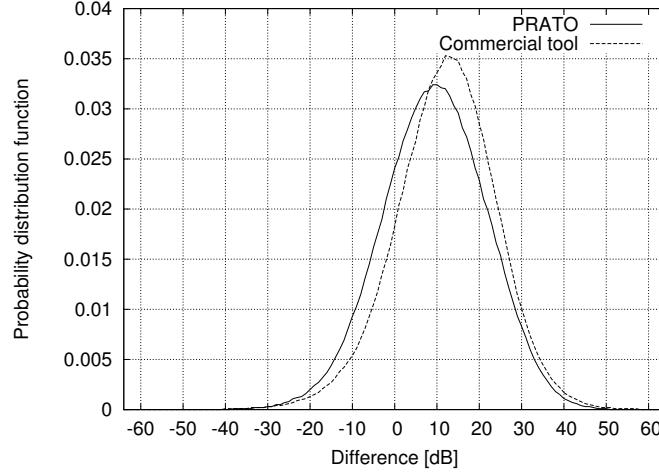


Figure 10.4: Probability distribution function for Net<sub>11</sub> (GSM) of the difference between the field measurements and the simulation results of PRATO, and the commercial tool.

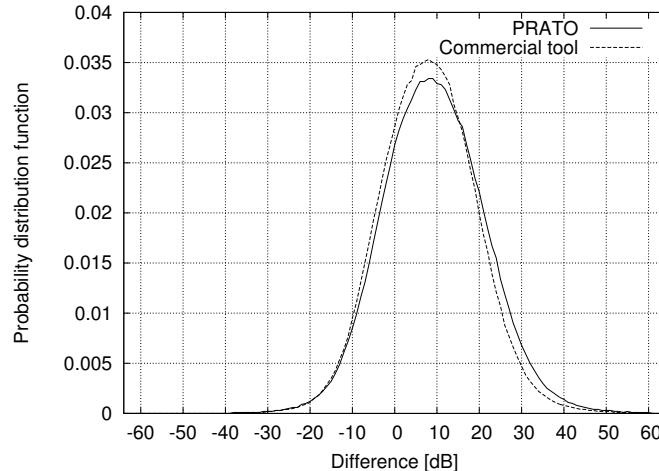


Figure 10.5: Probability distribution function for Net<sub>12</sub> (UMTS) of the difference between the field measurements and the simulation results of PRATO, and the commercial tool.

Plots showing the probability distribution function (PDF) of the difference between the simulation results and the field measurements were also produced. In this case, Figures 10.4, 10.5, and 10.6 show graphs representing the difference, expressed in dB, between the predictions and measurements of Net<sub>11</sub>, Net<sub>12</sub> and Net<sub>13</sub>, respectively. Again, the graphs labeled as (a) show the analysis for PRATO, whereas the ones labeled as (b) show the same analysis for the commercial tool. Table 10.2 lists the mean and standard deviation of the PDFs for each network and tool tested, the values of which are expressed in dB. It is important to note that the parameter set used for both tools intentionally generate pessimistic results in terms of coverage prediction. This is clearly observed from the mean-difference values of all three test cases.

Comparing the plots for each of the test networks (see Figures 10.4, 10.5, and 10.6), it is clear that the difference between each pair of (a) and (b) diagrams is minor. A small difference is present on the commercial tool, the predictions of which show a slightly greater

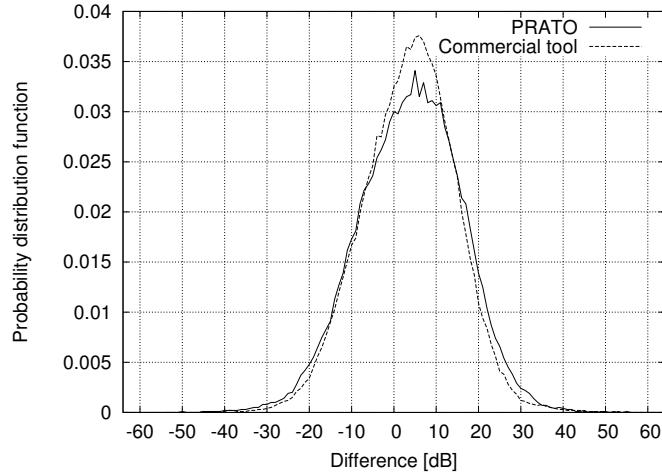


Figure 10.6: Probability distribution function for Net<sub>13</sub> (LTE) of the difference between the field measurements and the simulation results of PRATO, and the commercial tool.

Table 10.2: Mean and standard-deviation values for the PDFs of the difference between simulation and measurement results.

	Mean [dB]		Std. deviation [dB]	
	PRATO	Commercial tool	PRATO	Commercial tool
Net <sub>11</sub>	9.37	12.04	12.29	11.98
Net <sub>12</sub>	9.31	8.71	11.94	11.23
Net <sub>13</sub>	3.84	3.63	12.20	11.06

deviation with respect to the measurements. Therefore, it can be concluded that the prediction results of PRATO are comparable with the results of the commercial tool for the three test networks. Moreover, the presented analysis also confirms that the calculated results are independent of the frequency band used, since each test network operates in a different frequency.

Additionally, each of the test networks extended over different terrain types and environments, e.g., urban and suburban areas. Since the curves of the presented charts show similar profiles for the difference between the measurements and simulations for both tools, the applicability of PRATO for arbitrary terrain types can also be expected.

Notice also that PRATO generated similar results as the commercial tool in all three cases, irrespective of the operational frequency or chosen terrain type. The predicted values are also comparable for different distances between the BSs and UEs. Moreover, a slight improvement can be observed in some of the prediction values of PRATO, because they better resemble the profile shown by the field measurements.

### Computational-time performance

In the following, the computational-time performance of both tools is analyzed. The analysis focused on the required processing times of both tools on the same system, the hardware of which consisted of a 4-core Intel i7 2.67 GHz CPU, 24 GB of RAM and a dual nVidia GeForce GTX 590 GPU. The simulations for PRATO were performed on a Linux operating system, using multiple processes on the CPU. The commercial tool required a Windows

Server operating system, and provided single-process, multi-threading support to use all the cores of the CPU.

The simulations included the test networks presented above, all of which extend over a geographical region of  $285 \times 185 \text{ km}^2$ , with a resolution of 25 m. The processing times for a set of simulations are given in Figures 10.7, 10.8, and 10.9, for test networks  $\text{Net}_{11}$ ,  $\text{Net}_{12}$  and  $\text{Net}_{13}$ , respectively.

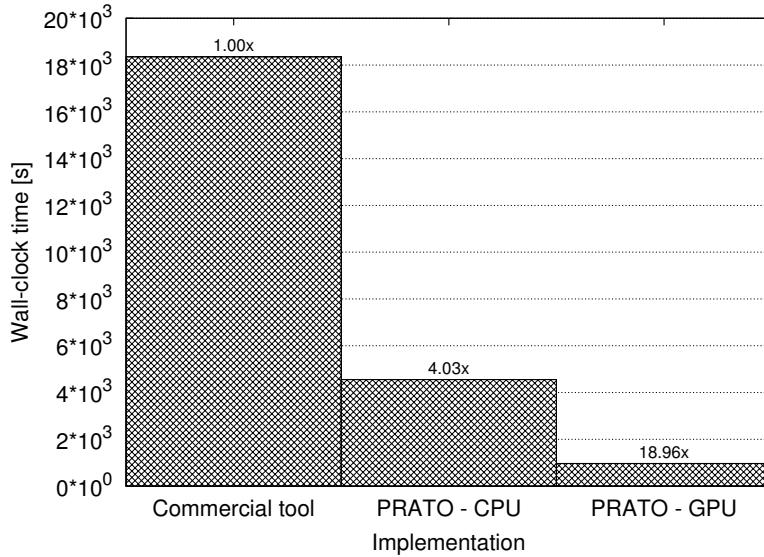


Figure 10.7: Simulation-processing times and speedup factors of  $\text{Net}_{11}$  (GSM) for the commercial tool and two implementations of PRATO.

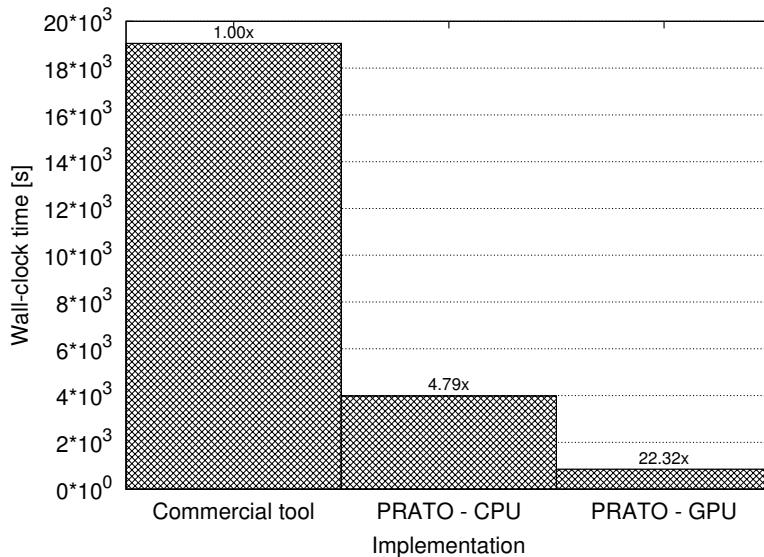


Figure 10.8: Simulation-processing times and speedup factors of  $\text{Net}_{12}$  (UMTS) for the commercial tool and two implementations of PRATO.

The plotted values represent the average simulation-processing time after performing 10 independent measurements. The time-measurement gathering was performed during the simulations presented in the previous section. The PRATO-CPU deployed used six workers and one master process. For the PRATO-GPU configuration, the same process deployment

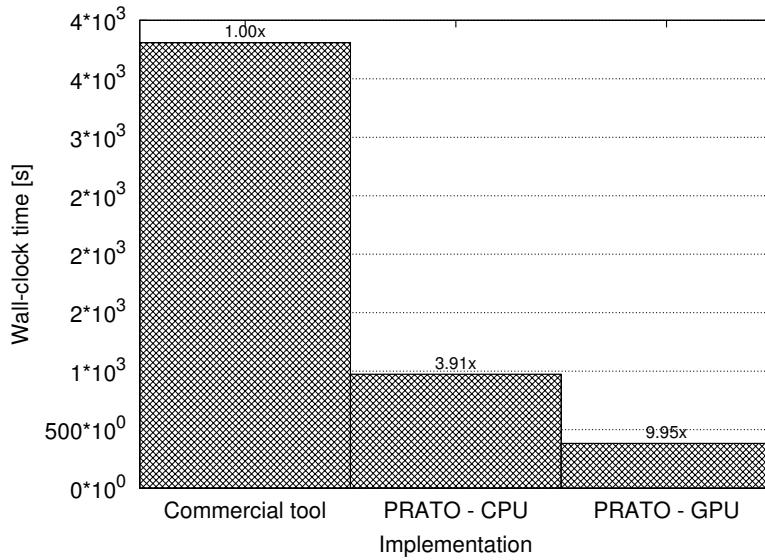


Figure 10.9: Simulation-processing times and speedup factors of Net<sub>13</sub> (LTE) for the commercial tool and two implementations of PRATO.

was used, and the worker processes operated on the dual GPU, i.e., two GPUs on one board, each of which featured 1.5 GB of DRAM.

The benefits of the parallel implementation of PRATO is clear in all three cases. Moreover, the multi-GPU support increases the running-time performance even further, achieving speedup factors of 18.96, 22.32, and 9.95, for Net<sub>11</sub>, Net<sub>12</sub>, and Net<sub>13</sub>, respectively. These results confirm that the use PRATO as a radio-planning tool in a real-world environment is feasible and it outperforms the compared commercial tool in terms of computational-time performance.

## 10.4 Summary

The radio planning of modern cellular networks requires efficient and exact radio-signal coverage calculations. In this context, the unified framework PRATO was evaluated from a radio-network planning point-of-view by comparing its simulation results with field measurements. The same analysis was applied to a professional software application in order to assess the results of PRATO from a quality and performance perspectives.

The extensive analyses presented in this chapter showed satisfactory results. Compared to a professional network-planning tool, the result accuracy achieved is completely comparable irrespective of the terrain type or operational frequency, while the computational speed is many times higher. These results confirm that PRATO does not reduce the solution quality due to the increased performance. Indeed, its suitability for use in a real-world environment, addressing different radio-planning activities by simulation, was also confirmed. This fact makes the framework interesting for researchers as well as for radio-network engineers.



## 11 Conclusion and Further Work

Fast and accurate performance evaluation is of essential importance for the planning and optimization of mobile radio networks. In this thesis, several high-performance and novel methods for the analysis and optimization of radio networks were developed and discussed. Some of them even outperform state-of-the-art methods to which they were compared in terms of result accuracy and instance-size complexity.

The snapshot analysis is best suited for the detailed performance evaluation of radio networks. The performance of a snapshot method is mainly influenced by two factors. First, the performance of the method for modeling a single snapshot and, second, the quality level of the estimate applied for obtaining the interest figures from the individual snapshot results. Both topics were comprehensively addressed in this thesis.

While providing detailed and accurate results for various applications in network planning, the snapshot analysis is too time-consuming to be applied in radio-network optimization, where typically a large number of different configurations need to be compared in a shorter time-frame. To this end, a parallel framework for radio-network planning was presented in this thesis. The framework is very flexible in terms of air-interface modeling, e.g., different QoS schemes, and is of significantly higher computational-time performance compared with any currently available solution known to the author. Moreover, it incorporates multi-GPU support and the corresponding parallel-programming techniques that are required to exploit the computational power of such hardware. The performance gain results mainly due to the combination of a novel parallel-programming approach and some state-of-the-art methods. The achieved performance improvement excels in radio-network optimization environments. Further research in this field will include abstracting the introduced master-worker-database principle into a multi-purpose parallel framework such as Charm++ [87], which provides a functionality for overlapping execution and communication, as well as fault tolerance.

When applying automatic network optimization, a fast and accurate performance analysis is of crucial importance. However, an overview of the related literature shown that this fact seems to have been only partially taken into consideration in the design of several optimization methods for radio networks. A common approach is to apply a detailed and too time-consuming method for evaluating different candidate solutions, resulting in an unacceptable computational-time complexity for medium to large-sized problem instances. The opposite approach, which is even less satisfactory, is to build fast but inaccurate models, the results of which do not sufficiently correlate with reality. To this end, the quality and speed performance of the presented framework was empirically verified with an industrial software tool for radio-network planning. The results clearly show a very good agreement in terms of accuracy of the radio-propagation predictions, compared to those obtained with the commercial tool. Additionally, the results demonstrated the performance advantage of the framework compared to the computational time of the enterprise software. It is important to note that these results also apply for an application in everyday network planning. However, to validate the complete array of radio-planning activities the framework can handle further research is required.

Increasing the performance of the simulations involved during the objective-function evaluation is only the first step towards a practical running-time reduction for radio-network optimization. In this sense, the performance of the novel agent-based algorithm presented in this thesis was tested while solving the service-coverage problem in radio networks. The results show significant gains with respect to the size of problem instances, as well as regarding its speed performance and solution quality, even outperforming a state-of-the-art method, to which it was compared. Further research will include experimentation with different parameters and optimization problems, in order to gain better understanding of the dynamics that guide the algorithm through the search space of the problem.

The new optimization problem for 3G radio networks identified in this thesis deals with SHO balancing of downlink and uplink areas. The problem was tackled by three different metaheuristic algorithms, the solutions of which show a substantial improvement of downlink and uplink balance. A challenge for future work is to evaluate this optimization problem in a dynamic context. This requires using a full-stack simulation tool that includes dynamic effects, such as fast power control.

The clutter-loss optimization method developed in this thesis makes use of the tools discussed above, providing a faster, more accurate and simpler method that replaces a manual approach. It applies a metaheuristic algorithm which has been tailored for the special requirements of the automatic clutter-loss optimization. The presented results make the benefits of an automated-optimization approach evident. Furthermore, in the context of other radio-coverage planning activities carried out at the Radio Network department of Telekom Slovenije, d.d., supplementary testing of the framework as a coverage-planning tool is currently being conducted.

PRATO, the radio-planning framework presented in this thesis, is a free and open-source software. For this reason, it can be readily modified and extended to support, for example, other propagation models and post-processing algorithms. This characteristic provides it with a clear advantage when compared to commercial and closed-source tools. The source code is available for download from the author's home page, <http://cs.ijs.si/benedicic/>.

The constantly improving performance of computer hardware might allow an even more detailed analysis of networks in the near future. In this sense, a major challenge that will require further research is to extend the presented models and algorithms, e.g., by incorporating user mobility and traffic dynamics. This is especially important due to different existing and emerging mobile technologies. Indeed, the co-existence of multiple technologies will motivate an increasing interest about the combined analysis and optimization of different technologies and their mutual dependency.

## 11.1 Scientific contributions

The work in this thesis has led to the following original contributions to science:

1. Design and development of a unified framework for radio-network planning and optimization, incorporating the second and third contribution items, as well as experimentation on real-world radio networks. The experimentation includes a comparison with a commercial tool that is currently being used for real radio-network planning. Moreover, the parallel implementation of the framework exploits the computing resources of computer clusters and graphic-processing units.
2. Proposal of a new approach for parallel programming that combines the well-known master-worker scheme with an external database. The proposed approach relieves the problem of congestion at the master node, which is responsible for result collection and selection, thus improving the scalability of a parallel system.

3. Quality improvement of radio-propagation predictions by applying metaheuristic optimization to the parameters of radio-propagation models. This technique enables the adaptation of radio-propagation models to the local environment over which a radio network is deployed, as well as the automatic optimization of signal losses due to clutter.
4. A new algorithm, based on autonomous agents, to tackle the service-coverage problem in radio networks. The algorithm deals with problem instances that are out-of-reach of other compared state-of-the-art techniques used as reference. It also reaches good quality of solutions if compared to classic network-planning techniques. The proposed approach is especially suitable for optimizing large radio networks.
5. Identification and formalization of a new optimization problem in third-generation radio networks that deals with soft-handover alignment of downlink and uplink areas. By solving this problem, network malfunctioning is avoided in areas where there is soft-handover capability in the uplink, but none in the downlink. So far, this problem has not been formalized nor tackled by means of automatic optimization.



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## Index of Figures

127

Figure 1.1 A diagram of the classic decision-making process. Adapted from [156].	1
Figure 2.1 Gradient descent with a negative slope, i.e. $x$ is increasing.	8
Figure 2.2 A saddle point or point of inflection, where the derivative is zero.	9
Figure 2.3 Graphical representation of a linear-programming example with two constraints, $c_1$ and $c_2$ . The grayed area is the polytope representing the region of feasible solutions.	10
Figure 2.4 A metaheuristic algorithm process using a black box for the objective-function evaluation, $f(x)$ , of a solution, $x$ .	15
Figure 3.1 An example of the soft-handover region, showed as a grayed area. The area boundaries are configurable relative to the pilot-signal level of each cell. Adapted from [150].	19
Figure 4.1 A typical optimization cycle for radio networks. This sequence is repeated until the achieved results are acceptable.	21
Figure 4.2 A graphical representation of the set-covering problem: (a) the problem input and (b) the solution.	22
Figure 4.3 An example of an antenna-azimuth pattern, showing the gain in each direction.	24
Figure 4.4 The antenna tilt, showing its angle with the horizontal plane. Adapted from [79].	24
Figure 5.1 A simplified diagram of the architecture of a modern GPU. Adapted from [130].	28
Figure 6.1 Flow diagram of the serial version.	37
Figure 6.2 Example of a raster map, showing the result of a path-loss calculation from an isotropic antenna. The color scale is given in dB, indicating the path loss at a given point.	38
Figure 6.3 Example of a raster map, showing the influence of a directional antenna over the path-loss result depicted in Figure 6.2. The color scale is given in dB, indicating the path loss at a given point.	39
Figure 6.4 Example of a raster map, displaying the final coverage prediction of 136 transmitters over a geographical area. The color scale is given in dBm, indicating the received-signal strength. Darker colors denote areas with a reduced signal due to the fading effect of the hilly terrain and clutter.	40
Figure 6.5 Example of a geographical-location mapping of the input-spatial data into a 2D matrix, the lower-left corner of which indicates the nearest point to the origin of the map-projection system.	42
Figure 6.6 Flow diagram of the master process.	43
Figure 6.7 Flow diagram of the “Processing loop” step of the master process.	43
Figure 6.8 Flow diagram of a worker process.	45

Figure 6.9 Communication diagram, showing the message passing between the master and a worker process. . . . .	45
Figure 6.10 Measured wall-clock time for weak-scalability experiments, featuring MW and MWD setups. Experiments allocated one MPI worker process per core. . . . .	49
Figure 6.11 Measured wall-clock time for strong-scalability experiments, featuring MW and MWD setups. Experiments assigned one MPI worker process per core. . . . .	50
Figure 6.12 Average speedup for the strong-scalability experiments. . . . .	51
Figure 6.13 Average parallel efficiency for the strong-scalability experiments. . . . .	51
Figure 7.1 Architecture of the parallel, agent-based optimization system on GPU. .	59
Figure 7.2 Convergence profile of the parallel-agent approach for the test network Net <sub>1</sub> , deployed over an urban area. . . . .	65
Figure 7.3 Convergence profile of the parallel-agent approach for the test network Net <sub>2</sub> , deployed over a rural area. . . . .	65
Figure 7.4 Convergence profile of the parallel-agent approach for the test network Net <sub>3</sub> , deployed over a suburban area. . . . .	66
Figure 8.1 HSUPA traffic and uplink interference with: (a) balanced downlink and uplink SHO conditions, and (b) unbalanced downlink and uplink SHO conditions. . . . .	70
Figure 8.2 Area under radio coverage, $A_{\text{covered}}$ , and without radio coverage, $\overline{A}_{\text{covered}}$ , within the complete geographical area, $A_{\text{total}}$ , of test network Net <sub>7</sub> . . . . .	76
Figure 8.3 Convergence analysis for DE, showing the best, mean and worst objective-function values in each evaluation of the SHO-balancing problem. . . . .	78
Figure 8.4 Convergence analysis for DASA, showing the best, mean and worst objective-function values in each evaluation of the SHO-balancing problem. . . . .	79
Figure 8.5 Convergence analysis for SA, showing the best, mean and worst objective-function values in each evaluation of the SHO-balancing problem. . . . .	80
Figure 8.6 Spatial distribution of the SHO areas before the optimization. . . . .	81
Figure 8.7 Spatial distribution of the SHO areas after the optimization. . . . .	81
Figure 9.1 The automatic-tuning system, including its elements and their relationships. . . . .	87
Figure 9.2 Terrain profile of the test network Net <sub>8</sub> , dominated by a flat agricultural area. . . . .	88
Figure 9.3 Terrain profile of the test network Net <sub>10</sub> , dominated by forested hills. .	88
Figure 9.4 Probability-density distribution of radio predictions against the field measurements of network Net <sub>8</sub> over a rural area using the: (a) empirically-calculated parameters, (b) analytically-calculated parameters, and (c) optimized parameters. . . . .	90
Figure 9.5 Probability-density distribution of radio predictions against the field measurements of network Net <sub>9</sub> over an urban area using the: (a) empirically-calculated parameters, (b) analytically-calculated parameters, and (c) optimized parameters. . . . .	91
Figure 9.6 Probability-density distribution of radio predictions against the field measurements of network Net <sub>10</sub> over a hilly area using the: (a) empirically-calculated parameters, (b) analytically-calculated parameters, and (c) optimized parameters. . . . .	92





Table 5.1 Properties of different memory types on a GeForce 8800 GPU. Adapted from [130]. . . . .	29
Table 5.2 Naming-convention translation between OpenCL and CUDA. Adapted from [91]. . . . .	29
Table 6.1 Running-time gain (in percent) of the simulations for the weak-scalability of the MWD setup relative to the classic MW approach. . . . .	48
Table 7.1 Sizes of the test networks used for experimentation of the service-coverage problem, in terms of equipment and geographical area. . . . .	63
Table 7.2 Network parameters of the test networks used for the service-coverage problem. . . . .	63
Table 7.3 Parameter settings of the parallel-agent approach for each test network. . . . .	63
Table 7.4 Optimization results after applying two different approaches for solving the service-coverage problem. All values are expressed in Watts. . . . .	64
Table 7.5 Wall-clock times (in seconds) and speedup factors for the different implementations of the objective-function evaluation and the parallel agents, as measured during the experimentation of the service-coverage problem. . . . .	66
Table 8.1 Technical characteristics of Net <sub>7</sub> , the test network used for the SHO-balancing problem. . . . .	76
Table 8.2 Solution-quality performance of the three algorithms, after 30 independent runs. . . . .	77
Table 8.3 Improvement analysis of the best solution that each algorithm achieved for the SHO-balancing problem. All values are expressed in percent. . . . .	79
Table 9.1 Parameters of the test networks Net <sub>8</sub> , Net <sub>9</sub> and Net <sub>10</sub> that were used for the experimental simulations during the automated tuning of the radio-propagation model. . . . .	88
Table 9.2 Mean and standard-deviation values of the radio-propagation prediction against the field measurements. The values, expressed in dB, are given when using the manual, analytical and DASA approaches for each test network. . . . .	93
Table 9.3 Clutter-category label numbers and descriptions for the signal loss due to clutter, as reflected by a radio-propagation model. . . . .	94
Table 9.4 Field-measurement proportions with respect to each clutter category, the percentages of which are given in terms of the area covered by the drive tests for each of the test networks. The clutter category legend is given in Table 9.3. . . . .	95
Table 9.5 Clutter-category losses after the optimization. The default losses for each clutter category are given along the solutions for each of the test networks. All values are expressed in dB. . . . .	95

Table 9.6 Statistical analysis of the solutions for the clutter-optimization problem. All values are expressed in dB. The corresponding box plots are depicted in Figure 9.10. . . . .	98
Table 10.1 Parameter values used during the coverage-prediction calculations for each test network. The same values were selected on PRATO and the com- mercial tool. . . . .	102
Table 10.2 Mean and standard-deviation values for the PDFs of the difference between simulation and measurement results. . . . .	105

Algorithm 2.1 Pseudo-code of the simulated-annealing metaheuristic. Adapted from [23]. . . . .	14
Algorithm 6.1 Pseudo-code of the radio-coverage prediction algorithm. The time complexity is given per line. . . . .	40
Algorithm 7.1 Pseudo-code representing the behavior of an agent. . . . .	60
Algorithm 7.2 Pseudo-code representing the step set $SS_0$ , which is applied by the agents in areas without service coverage. . . . .	61
Algorithm 7.3 Pseudo-code representing the step set $SS_1$ , which is applied by the agents in areas with service coverage. . . . .	61
Algorithm 8.1 A move in the search space of SA for solving the SHO-balancing problem. . . . .	75



## Appendix A: Bibliography

### A.1 Publications related to this thesis

#### A.1.1 Journal articles

- Lucas Benedičič, Felipe A. Cruz, Tsuyoshi Hamada, and Peter Korošec. A GRASS GIS parallel module for radio-propagation predictions. *International Journal of Geographical Information Science*, (in press), 2014.
- Lucas Benedičič, Mitja Štular, and Peter Korošec. A GPU-based parallel-agent optimization approach for the service-coverage problem in UMTS networks. *Computing and Informatics*, (in press), 2013.

#### A.1.2 Conference papers

- Lucas Benedičič, Mitja Štular, and Peter Korošec. Balancing downlink and uplink soft-handover areas in UMTS networks. In *Evolutionary Computation (CEC), 2012 IEEE Congress on*, pages 1–8. IEEE, 2012.
- Lucas Benedičič, Mitja Štular, and Peter Korošec. Pilot power optimization in UMTS: a multi-agent approach. In *Proceedings of the 13th International Multiconference Information Society - IS 2010*, volume A. Jožef Stefan Institute, October 2010.

#### A.1.2 Workshop papers and posters

- Lucas Benedičič. Parallelization and Optimization Methods for Radio-Network Planning. In *Zbornik povzetkov delavnic "Algoritmi po vzorih iz narave" v študijskem letu 2012/2013*. Jožef Stefan Institute, 2013.
- Lucas Benedičič. Pilot power optimization in UMTS: a multi-agent approach. In *Zbornik povzetkov delavnic "Algoritmi po vzorih iz narave" v študijskem letu 2010/2011*, page 7. Jožef Stefan Institute, 2011.

### B.1 Other publications

#### A.1.2 Conference papers

- Lucas Benedičič and Peter Korošec. Reducing costs with computer power management. In *Proceedings of the 4th Jožef Stefan International Postgraduate School Students Conference*. Jožef Stefan International Postgraduate School, May 2012.



## Appendix B: Biography

Lucas Benedičić was born on December 12, 1977 in Buenos Aires, Argentina. He was a student at the University of Buenos Aires, where he completed the first 4 years of the undergraduate programme in computer sciences, at the Faculty of Exact and Natural Sciences.

After moving to Slovenia in 2002, he enrolled into the University of Ljubljana, Faculty of Computer and Information Science, where he graduated in 2007 with the work titled “Application development using open-source technologies”.

In September 2007, he started his postgraduate studies at the University of Primorska, Faculty of Mathematics, Natural Sciences and Information Technologies, after enrolling into the Masters study programme of computer sciences. He graduated in 2009 with the work titled “Optimization of Common Pilot Channels in UMTS Networks”.

Towards the end of 2007, he was awarded a scholarship for continuing his postgraduate studies in Slovenia. The award was received after winning a problem-solving contest, organized by Halcom, d.d. In the fall of 2009, he enrolled in the PhD program entitled “Information and Communication Technologies” under the supervision of Assoc. Prof. Peter Korošec and the co-supervision of Assist. Prof. Tomaž Javornik, at the Jožef Stefan International Postgraduate School in Ljubljana, Slovenia.

In November 2009, he became a junior researcher, partially funded by the European Union through the European Social Fund. During this period he held working positions at Telekom Slovenije, d.d., in research and development within the Radio Network Department, and at the Jožef Stefan Institute, in the Computer Systems Department.

In 2011, he attended and successfully completed the summer institute “Scientific Computing in the Americas: the challenge of massive parallelism”, organized by the Pan-American Advanced Studies Institute, Boston University. The focus of the institute was on advanced theoretical topics and programming of massively parallel processors.

In early 2012, he received an invitation to work as a visiting researcher at the Nagasaki Advanced Computer Center (NACC) at the University of Nagasaki, Japan. His joint research work at NACC was completed by the end of 2012.

Over the past few years, he has presented his research work at several international conferences and workshops in the areas of evolutionary computation, optimization and parallel computing.

His research interests include combinatorial and numerical optimization, parallel and GPU computing, and metaheuristic algorithms.



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