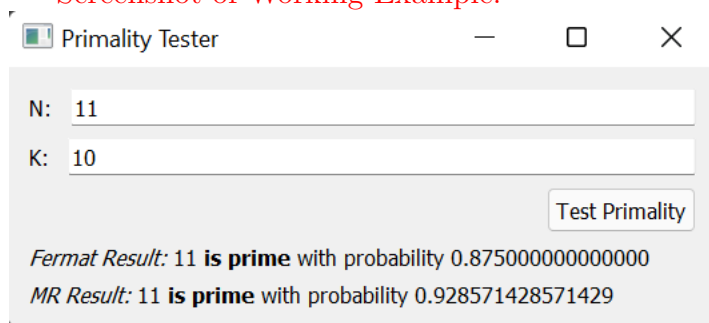


Project-1 Report

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Screenshot of Working Example:



Code that was written:

```
import random
```

```

import math

def prime_test(N, k):
    # This is main function, that is connected to the Test button. You don't need
    # to touch it.
    return ferat(N,k), miller_rabin(N,k)

def mod_exp(x, y, N):
    #Implement modular exponentiation from Figure 1.4, Pg. 19
    if y == 0:
        return 1
    """
    Every multiply and divide is  $O(n^2)$ 
    The amount of multiplies/divides (or depth) is determined by y
    Therefore,  $O(yn^2) == O(n^3)$ 
    """
    z = mod_exp(x, math.floor(y/2), N)
    if y % 2 == 0:
        return (z**2) % N
    else:
        return x * (z**2) % N

def fprobability(k):
    # Returns the probability that N is prime, given k trials, with Fermat
    # primality test
    fprob = 1 - (1/(2^k))
    return fprob

def mprobability(k):
    # Returns the probability that N is prime, given k trials, with
    # Miller-Rabin primality test
    mprob = 1 - (1/(4^k))
    return mprob

```

```

def fermat(N,k):
    #Implement Fermat primality test
    """
    In this algorithm, we use the modular exponentiation function that
    had a complexity of  $O(n^3)$ . We use this function 'k' times in the
    for loop. Everything else takes  $O(1)$ . Therefore, this primality test
    is  $O(kn^3) = O(n^3)$ .
    """
    #create for loop in range of k, or trial attempts
    for i in range(k):
        #generate random value for 'a' between 2 and N-1
        a = random.randint(2, N-1)
        #equivalent to:  $(a^{N-1}) \% N$ . if this equates to 1, there is a
        # chance the number is prime
        if mod_exp(a, N-1, N) == 1:
            return 'prime'
        else:
            break

    return 'composite'

def miller_rabin(N,k):
    # Implement Miller-Rabin primality test
    # create for loop in range of k, or trial attempts
    """
    In this algorithm, we use the modular exponentiation function that
    had a complexity of  $O(n^3)$ . We use this function 'k' times in the
    for loop. Everything else takes  $O(1)$ . Miller-Rabin differs from Fermat
    because there is another loop and further multiplies/divides. Therefore,
    this primality test is  $O(kn^3kn) = O(n^4)$ .
    """
    for i in range(k):
        # generate random value for 'a' between 2 and N-1
        a = random.randint(2, N - 1)
        # condition to check for primality
        #equivalent to:  $(a^{N-1}) \% N$ . if this equates to 1, there is

```

```

# a chance the number is prime
if mod_exp(a, N - 1, N) == 1:
    z = (a ** (N - 1))
    #taking sqrt of exponent and making sure it is an even value;
    # stop loop if odd
    while (z ** (1/2)) % 2 == 0:
        #further testing in order to eliminate Carmichael numbers
        if (z ** (1/2)) % N == 1 or -1:
            return 'prime'
else:
    break

return 'composite'

```

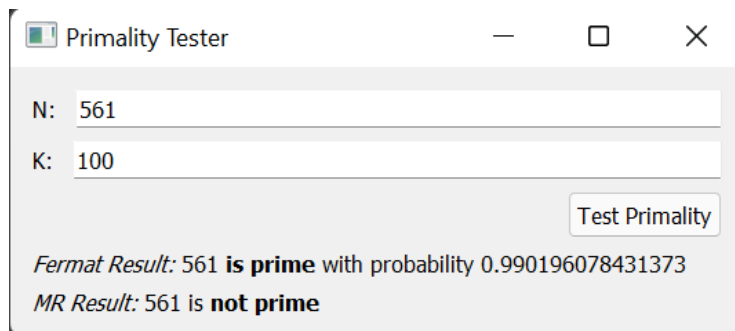
"""

I tried many different inputs to see what would yield different results for the two tests. I tested several large even numbers to make sure those came back as composite. I tried many different random numbers with varying trial size to see what the results would be.

For the most part, the tests did fairly similiar; however, the Miller-Rabin test did much better with Carmichael numbers. I tested 561 and 1105, both of which are Carmichael numbers. The Fermat test would give different answers, but the Miller-Rabin was much better about consistently removing these Carmichael numbers.

"""

Primality Test Disagreement:



Time and Space Complexity:

Modular Exponentiation:

Every multiply and divide is $O(n^2)$. The amount of multiplies/divides (or depth) is determined by y . Therefore, $O(yn^2) == O(n^3)$.

Fermat Primality Test:

In this algorithm, we use the modular exponentiation function that had a complexity of $O(n^3)$. We use this function 'k' times in the for loop.

Everything else takes $O(1)$. Therefore, this primality test is $O(kn^3) = O(n^3)$.

Miller-Rabin Primality Test:

In this algorithm, we use the modular exponentiation function that had a

complexity of $O(n^3)$. We use this function 'k' times in the for loop. Everything else takes $O(1)$. Miller-Rabin differs from Fermat because there is another loop and further multiplies/divides. Therefore, this primality test is $O(kn^3kn) = O(n^4)$.

Probabilities:

The probability of a number actually being prime with the Fermat test is: $1 - 1/2^k$.

The probability of a number actually being prime with the Miller-Rabin test is: $1 - 1/4^k$.

GitHub Link:

<https://github.com/liecchr2/Project-1.git>