

# Design Document: ASGN-4 Adding aliases to the HTTP server

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## 1 Goals

The goal of this project is to implement aliases functionality to support different name pointing to one HTTP name. The aliases will adopt KVS method to store data and in-memory cache to support better performance.

Both client and server now support PATCH command which accompany the alias KVS file to implement. Compared to data KVS file, alias KVS file don't have data range that have be pointer by entries. It will be directly stored in KVS file.

Other than that, server now will delete the slash char '/' for all names.

## 2 Design

There are three parts in the design: (1) client: support PATCH command; (2) server: alias key-value store implementation; (3) multithreading and concurrency control.

### 2.1 Client: support PATCH command

Client now support PATCH command in the form of "a:existing\_name:new\_name". Where the new\_name is an alias to existing\_name.

The new\_name and existing\_name can be both alias names and existing\_name could be a real httpname. If either one is lacking in command line, client will print out error message.

```
1. Input : buf[4096] stored command arguments: a:existing_name:new_name
2. Obtain action, existing_name, new_name
3. if buf[0] == '/' then // delete slash '/'
4.     strncpy(httpname, httpname + 1, strlen(httpname))
5. end
6. if strlen(httpname) == 0 || strlen(filename) == 0 then
7.     fprintf(stderr, "HTTP/1.1 400 Bad Request\n")
8.     continue
9. end
10. if argv[i][0] == 'a' then
11.     send(sock, "PATCH existing new_name \r\n\r\n ALIAS existing new_name \r\n")
12. end
```

Algorithm 1: Client side additional code to send PATCH command

## 2.2 Server: alias key-value store implementation

Server now will support alias functionality. To achieve this goal, there are four functions need to be added into the program: (1) alias KVS file and alias cache initialization; (2) server side: handling PATCH request; (3) name\_handler to support mapping alias to existing name; (4) changes in PUT and GET execution.

### 2.2.1 alias KVS file and alias cache initialization

Server now support “-m filename” to open or create alias KVS file. If it open a KVS file, then the program will call `alias_init()` to reload entries into cache. If it creates a new KVS file, then do nothing.

```

Input          :char * optarg int32_t argc char * argv[]
Shared Variable: int32_t fd_alias
1. while (opt = getopt(argc,argv, "N:c:f:m:") != 1 do
2.     switch(opt)
3.         case 'm' :
4.             if (fd_alias = open(optarg, O_RDWR)) == -1 then
5.                 if (fd_alias = open(optarg, O_CREAT | O_RDWR | O_TRUNC, S_IRWXU |
S_IRWXG | S_IRWXO)) == -1 then
6.                     fprintf("SET UP FAILED: cannot open or creat the KVS file\n")
7.                     exit(EXIT_FAILURE);
8.                 end
9.             end
10.    alias_init(); // reload alias kvs entries to cache
11. end
12. done

```

Algorithm 2: open or create alias KVS file in main function

In order to keep data consistency, the program will reload all entries back into the cache after open the alias KVS file. The program adopts in-memory map `alias_map<string, uint32_t>` for cache usage.

The cache will store two elements `<entry, alias_end>`. The entry is `char[128]` contains `exisiting_name` and `new_name` with null terminator after each of two. The combined two names will two null terminators will not exceed 128 bytes. The `alias_end` is a shared variable pointing to the end of alias KVS file.

```

Input          :None
Shared Variables :int32_t fd_alias uint32_t alias_end
function void alias_init()
1. char entry[128]
2. while pread(fd_alias, entry, 128, alias_end) != 0 do
3.    alias_map.insert(make_pair(entry, alias_end))
4.    alias_end += 128;
5. done

```

Algorithm 3: `alias_init()` reloads all entries in KVS file back to cache

### 2.2.2 server side: handling PATCH request

When the request header contains a PATCH, the program will go into PATCH processing procedure. This procedure will first parse the request header and delete the / char if it exists in front of the name.

```
Relevant Variables : char * existing_name char * new_name
1. recv(cl, buf, sizeof(buf), 0)
2. ptr = strstr(buf, " ")
3. strncpy(action, buf, ptr - buf) // obtain action code
4. if strcmp(action, "PATCH") == 0 then // PATCH command
5.     ptr_tem = strstr(ptr + 1, " ")
6.     if buf[ptr + 1 - buf] == '/' then // delete '/' char
7.         ptr += 1;
8.     end
9.     memcpy(existing_name, ptr + 1, ptr_tem - (ptr + 1)) // obtain existing_name
10.    ptr = strstr(ptr_tem + 1, "\r\n\r\n")
11.    if buf[ptr_tem + 1 - buf] == '/' then // delete '/' char
12.        ptr_tem += 1
13.    end
14.    memcpy(new_name, ptr_tem + 1, ptr - (ptr_tem + 1)) // obtain new_name
15. end
```

Algorithm 4: parse the PATCH request header and delete '/' char

The procedure will first check if the new\_name which is alias exists in cache. If alias not found, the entry including "new\_name\0existing\_name\0" will be written into alias KVS file and will be updated in the cache. If alias found, the new entry will be written into the alias KVS file and leave cache unchanged.

To avoid data race, when the program writes entry into the KVS file, the mutex lock will control one thread can access alias\_end and write into the KVS file.

```
Relevant Variables : map<string, uint32_t> alias_map
Shared Variables : int32_t fd_alias uint32_t alias_end mutex alias_end_lock
                    mutex alias_busy
1. map<string, uint32_t>::iterator iter;
2. memcpy(file_buf, new_name, strlen(new_name));
3. memcpy(file_buf+strlen(new_name)+1, existing_name, strlen(existing_name));
4. if (iter = alias_map.find(new_name)) == alias_map.end() then // not found
5.     // store alias in KVS file and update cache
6.     pthread_mutex_lock(&alias_end_lock);
7.     pwrite(fd_alias, file_buf, 128, alias_end);
8.     alias_map.insert(make_pair(new_name, alias_end));
9.     alias_end += 128;
10.    pthread_mutex_unlock(&alias_end_lock);
```

```

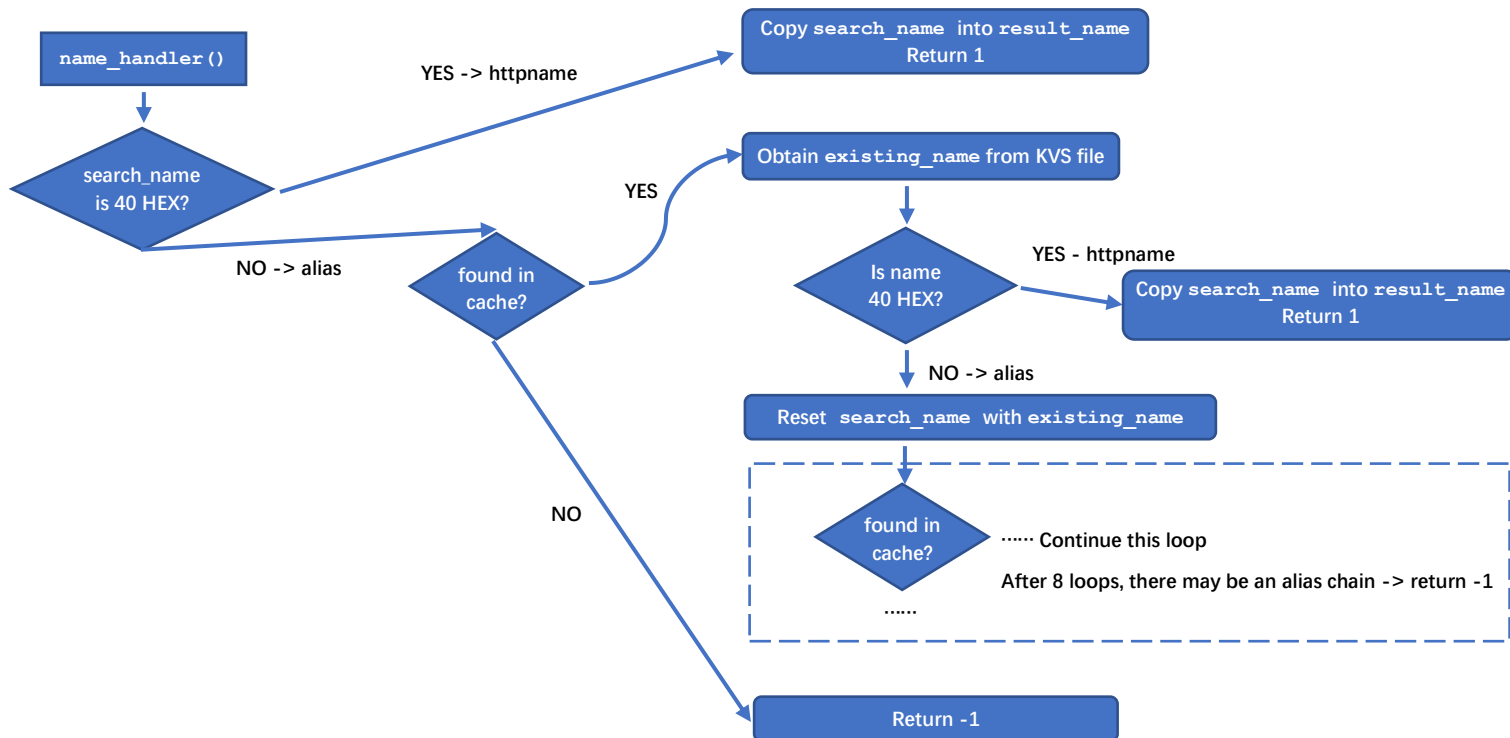
11. else // found -> update alias kvs file
12.   pthread_mutex_lock(&alias_busy);
13.   pwrite(fd_alias, file_buf, 128, iter->second);
14.   pthread_mutex_unlock(&alias_busy);
15. end

```

Algorithm 5: handling PATCH command in the processing function

### 2.2.3 name\_handler to support mapping alias to existing name

Basically, the name\_handler will be given a search\_name and return success code and a result\_name which is the 40 HEX httpname. If no httpname indicated or there is an alias chain loop, the name\_handler will return a fail code.



In terms of the return value of `name_handler()`, 1 indicates successful obtaining httpname, and -1 indicates no valid httpname is found. Based on the return value, the caller is able to know the status of the `result_name`, valid or not valid.

```

1. Input :char * search_name char * result_name
2. Shared Variables:map<string, uint32_t> alias_map int32_t fd_alias
3. if strlen(search_name) == 40 then
4.     memcpy(result_name, search_name, 40)
5.     return 1
6. end
7. map<string, uint32_t>::iterator iter
8. char temp[128]
9. for i = 1 -> 8 then
10.    if (iter = alias_map.find(search_name)) == alias_map.end() // not found
11.        return -1
12.    else // found
13.        pread(fd_alias, temp, 128, iter->second)
14.        char * ptr = temp + strlen(temp) + 1
15.        memset(search_name, 0, 128)
16.        strcpy(search_name, ptr) // obtain existing_name
17.        if strlen(search_name) == 40 then
18.            strcpy(result_name, search_name)
19.            return 1
20.        end
21.    end
22. end

```

Algorithm 6: name\_handler() obtain httpname by given search\_name

### 2.2.4 changes in PUT and GET execution

Call function name\_handler() before every execution start. If the function returns -1, then send 404 Not Found response back to client; otherwise, the program has 40 HEX httpname to use.

```

Relevant Variables :char * new_name char * httpname
1. if name_handler(new_name, httpname) < 0 then
2.     send(cl, "HTTP/1.1 404 Not Found\r\n", 24, 0)
3. else
4.     ..... other execution in GET and PUT

```

Algorithm 7: call name\_handler() before regular execution in GET and PUT

## 2.3 Multithreading and concurrency control

Because this design include writing into a file, two new shared variables which are mutex alias\_end\_lock and alias\_busy are introduced to control only one thread can write into the file and only one thread can access the end of alias KVS file to insert a new entry simultaneously.

There is no risk for multiple threads read an entry in alias KVS file since we don't need to control the consistency.