

CS240 Algorithm Design and Analysis  
Spring 2020  
Problem Set 4

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Due: 23:59, June 5, 2020

1. Submit your solutions to Gradescope ([www.gradescope.com](http://www.gradescope.com)).
2. In “Account Settings” of Gradescope, set your FULL NAME to your Chinese name and enter your STUDENT ID correctly.
3. If you want to submit a handwritten version, scan it clearly. CamScanner is recommended.
4. When submitting your homework, match each of your solution to the corresponding problem number.

## Problem 1:

Given an undirected graph  $G = (V, E)$ , consider the following randomized algorithm for finding a min-cut in  $G$ , which improves the probability of returning a min-cut compared to the basic min-cut algorithm presented in class. The improved algorithm works by performing two independent contractions of  $G$  until it has  $\lceil 1 + |V|/\sqrt{2} \rceil$  remaining vertices, producing graphs  $G_1$  and  $G_2$ , then recursively computing the min-cuts of  $G_1$  and  $G_2$  and returning the smaller of the two cuts.

Give a lower bound on the probability that *cut2* returns a correct min-cut.

*Hint:* Consider the probability that for a fixed min-cut  $C$  of  $G$ , there is no edge  $e \in C$  that is contracted by  $f(G, t)$ .

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1 Function  $f(G = (V, E), t)$ 
2   while  $|V| > t$  do
3     choose  $e \in E$  randomly;
4     contract edge  $e$ ;
5   end while
6   return  $G$  ;
7 end
8 Function  $cut2(G = (V, E))$ 
9   if  $|V| < 6$  then
10    return the min cut of  $G$ 
11  end if
12   $t \leftarrow \lceil 1 + |V|/\sqrt{2} \rceil$ ;
13   $G_1 \leftarrow f(G, t)$  ;
14   $G_2 \leftarrow f(G, t)$  ;
15  return smaller of  $cut2(G_1)$  and  $cut2(G_2)$  ;
16 end
```

## Problem 2:

Again consider an undirected graph  $G = (V, E)$

1. Give an upper bound on the maximum number of different min-cuts in  $G$ .

2. Give an efficient algorithm to find all min-cuts of  $G$  with probability, and argue that your algorithm is correct.

*Hint:* There is an elegant and simple solution which makes use of the basic min-cut algorithm presented in class.

### Problem 3:

Recall that the 3-SAT problem asks whether it is possible to satisfy all the clauses in a 3-CNF. Since 3-SAT is NP-complete, we will try to solve an easier problem, namely satisfying as many clauses in a 3-CNF as possible. Given a 3-CNF  $\phi$ , give an algorithm to find a truth assignment satisfying at least  $7/8$  of the maximum number of satisfiable clauses in  $\phi$ . Moreover, argue that there always exists an assignment satisfying at least  $7/8$  of the number of clauses in  $\phi$ .

### Problem 4:

Suppose we have a device that generates a sequence of independent fair random coin-flips, but what we want is a six-sided die that generates the values 1, 2, 3, 4, 5, 6 with *equal* probability. Give an algorithm that does so using  $\frac{11}{3}$  coin-flips on average.

### Problem 5:

In class, we saw that the expected time for randomized Quicksort to sort  $n$  numbers is  $O(n \log n)$ . Using Chernoff bounds, prove a high probability bound on this running time. That is, prove that for any  $c > 0$ , the probability randomized Quicksort finishes in  $O(n \log n)$  time is  $\geq 1 - \frac{1}{n^c}$ .

*Hint:* In each recursive call to Quicksort, consider the event that we get a "good split", i.e. that the smaller partition contains at least one quarter of all the elements.