

Q4

(a)

n	Optimal sequence	M(n, 0)
1	[+ -]	2
2	[+ o -]	3
3	[+ o o -]	4
4	[+ + - -]	4
5	[+ + - o -]	5
6	[+ + o - -]	5
7	[+ + o - o -]	6
8	[+ + o o - -]	6
9	[+ + + - - -]	6
10	[+ + + - - o -]	7
11	[+ + + - o - -]	7
12	[+ + + o - - -]	7
13	[+ + + o - - o -]	8
14	[+ + + o - o - -]	8
15	[+ + + o o - - -]	8
16	[+ + + + - - - -]	8
17	[+ + + + - - - o -]	9
18	[+ + + + - - o - -]	9
19	[+ + + + - o - - -]	9
20	[+ + + + o - - - -]	9
21	[+ + + + o - - - o -]	10

(b)

Base case ($n = s^2$, $s \in \mathbb{Z}$)

Each item represent the moving distance during each time steps.

The number of items represent number of time steps, that is M.

$$1 + 2 + 3 + \dots + s + (s-1) + (s-2) + \dots + 1 + 0 = s^2$$

'+' * s

'-' * s

Let $s^2 = n$, so $s = \sqrt{n}$, and $M(n, 0) = 2\sqrt{n} = 2\sqrt{n}$

if $s(s+1) < n \leq (s+1)^2$, $s \in \mathbb{Z}$:

we change from the base case formula and avoid breaking the continuity.

formula	n
$1 + 2 + 3 + \dots + (s+1) + s + (s-1) + (s-2) + \dots + 1 + 0 = s^2$	$(s+1)^2$
-1	$(s+1)^2 - 1$
-1 -1	$(s+1)^2 - 2$
-1 -1 -1	$(s+1)^2 - 3$
...	...
-1 -1 -1 -1 -1 ... -1	$(s+1)^2 - s$

-1 * s

That is $s(s+1) + 1$

The number of items is still $2(s+1)$, means $M(n,0) = 2s+2$

if $s^2 < n \leq s(s+1)$, $s \in \mathbb{Z}$:

formula	n
$1 + 2 + 3 + \dots + (s+1) + s + (s-1) + (s-2) + \dots + 1 + 0 = s^2$	$(s+1)^2$
-(s+1)	$(s+1)^2 - (s+1)$
-(s+1) -1	$(s+1)^2 - (s+2)$
-(s+1) -1 -1	$(s+1)^2 - (s+3)$
...	...
-(s+1) -1 -1 -1 -1 ... -1	$(s+1)^2 - (s+s)$

That is $s^2 + 1$

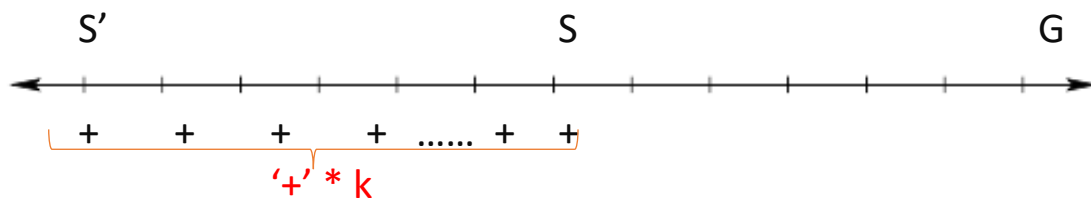
Remove this item !

The number of items is $2s+1$, means $M(n,0) = 2s+1$

We conclude that:

$$M(n, 0) = \left\lceil 2\sqrt{n} \right\rceil = \begin{cases} 2s + 1, & \text{if } s^2 < n \leq s(s+1) \\ 2s + 2, & \text{if } s(s+1) < n \leq (s+1)^2 \end{cases}$$

(c)



$$-\frac{1}{2}k(k+1)$$

0 k

n 0

We assume that the object start from S' with a velocity of 0, and keep speeding up to S, when it arrived S, it has the speed of k, the location of S' should be $(1 + 2 + 3 + \dots + k) = -\frac{1}{2}k(k+1)$ far away from S. Let $n' = S' \text{ to } G$, starts with 0 speed and ends with 0 speed.

$$M(n', 0) = M(n, k) + k$$

We use the formula from part(b):

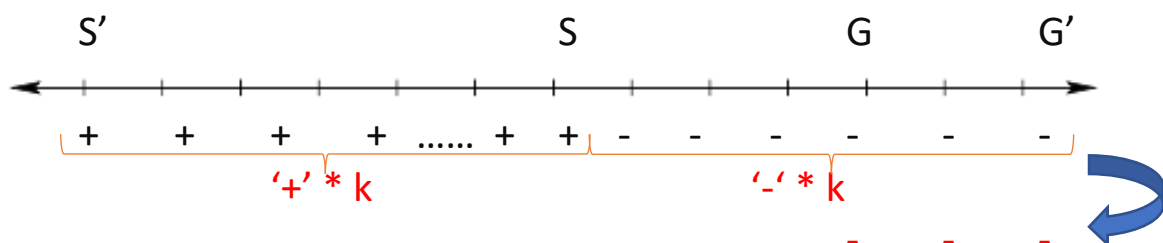
$$M(n', 0) = \left\lceil 2\sqrt{n'} \right\rceil$$

We already know that $n' = n + \frac{1}{2}k(k+1)$

$$\text{So } M(n, k) = M(n', 0) - k$$

$$= \left\lceil 2\sqrt{n + \frac{1}{2}k(k+1)} \right\rceil - k$$

(d)



$$-\frac{1}{2}k(k+1)$$

0 k

n 0

If the distance between S and G is smaller than S to G' (the smallest distance for decelerating from velocity of k) $= \frac{1}{2}k(k-1)$:

The object will first arrived at G' and turn back to G.

$$\text{Let } S'S = \frac{1}{2}k(k+1); SG = n; SG' = \frac{1}{2}k(k-1);$$

$$\text{So } GG' = SG' - SG = \frac{1}{2}k(k-1) - n$$

$$M(SG') = k$$

$$M(GG', 0) = \left\lceil 2\sqrt{\frac{1}{2}k(k-1) - n} \right\rceil$$

Because the object will first arrived at G' and turn back to G .

The real time steps is $M(SG', k) + M(GG', 0)$

$$\text{That is } M(n, k) = k + \left\lceil 2\sqrt{\frac{1}{2}k(k-1) - n} \right\rceil$$

(e)

$$h(r, c, u, v, r_G, c_G) = \max(M(r-r_G, u), M(c-c_G, v))$$