

# Reflections on Prismatic Constructions

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The Calculus of Prismatic Constructions, upon which this platform is based, is an extension of the standard CoC with a mechanism for discriminating inductive constructors.

## The pure Calculus of Construction

It is already very well-described elsewhere, so I won't try to provide a full and correct history of the CoC. Suffice to say that it is a logically consistent programming language, that can prove properties within the framework of intuitionistic logic.

At its simplest, it provides five basic constructions :

- universes, of the form  $Set_n$ , are the “types of types”.  $Set_{n+1}$  is the type of  $Set_n$
- products, noted  $\forall(x : X), Y x$  – or  $X \rightarrow Y$  when  $Y$  doesn't depend on  $x$  – are the “types of functions”.  $\mathbb{N} \rightarrow \mathbb{R}$ , for instance is the type of functions from the natural numbers to the real numbers.
- functions or lambdas, noted  $\lambda(x : X), Y x$ , are the “proofs of products”. A valid lambda can be interpreted as the proof of a property, quantified over its variable.
- hypotheses, or variables, are the symbols introduced by surrounding quantifiers ( $\lambda$  and  $\forall$ ). In their context, they are valid proofs of their type.

For example, the identity function can be written  $\lambda(A : Set_0), \lambda(a : A), a$ , and it is a valid proof of  $\forall(A : Set_0), \forall(a : A), A$ , since  $a$  is a valid proof of  $A$  in its context.

- Applications, of the form  $f x$ , where  $f : \forall(x : X), Y x$  and  $x : X$ , signify the specialization of a quantified property over an object  $x$ .

For instance, given a proof  $f$  of  $\forall(x : \mathbb{N}), \exists(y : \mathbb{N}), y = x + 1$ , we can prove that  $\exists(y : \mathbb{N}), y = 10 + 1$ , by applying  $f$  to 10 (aka.  $f 10$ ).

Given these axioms, we can build many theorems and their proofs, in a verifiable manner (i.e. there exists an algorithm to automatically check whether a claim like  $x : X$  holds).

However, it's been known for a while that the CoC by itself is not capable of handling a large class of the proofs that modern mathematicians (and even ancient ones) take for granted.

## The Limits of Constructions

To illustrate the kind of reasoning that can't be carried out with raw intuitionistic logic, let's take an obvious statement : true is not false (and vice-versa).

We'd like to prove this statement using only the tools given by the CoC. For this, we have to define a few concepts, namely *true*, *false*, and what it means to “not be” something.

### Intuitionistic Booleans

In order for two things to be considered the same, they must at least belong to the same family. In this case, it means that *true* and *false* must have the same type.

By convention, we like to call the type of “true or false” the *Boolean* type, in honor of George Boole, who was one of the first logicians to start considering truth and falsity as concrete values rather than abstract concepts.

Given a Boolean  $b$ , we would like to be able to return different values from a function, depending on whether  $b$  is true or false. Otherwise, our Boolean wouldn't be much use in a computation.

With all that in mind, here is the definition I propose the *Boolean* type :

$$Boolean = \forall(P : Prop), \forall(ptrue : P), \forall(pfalse : P), P$$

That is, a Boolean is a way to produce any  $P$ , given two alternatives *ptrue* and *pfalse*, and nothing else.

There are, conveniently, two ways to construct a Boolean, given this definition :

- $true = \lambda(P : Prop). \lambda(ptrue : P). \lambda(pfalse : P). ptrue$
- $false = \lambda(P : Prop). \lambda(ptrue : P). \lambda(pfalse : P). pfalse$

## Sameness, and negation

Two values  $x$  and  $y$  can be said to be the same when everything that can be proven of  $x$  can also be proven of  $y$ . More formally, given a type  $A$  of things, and two values  $x$  and  $y$  of type  $A$  we have :

$$(x \text{ sameas } y) = \forall (P : A \rightarrow \text{Set}_n), Px \rightarrow Py$$

## Inductive Types

Inductive types can be described as enumerations of constructors. In Coq (and similarly in other proof assistants), an inductive type must be declared along with its constructors, using a syntax like :

```
Inductive T : forall A..., Type :=
| t0 : forall x0..., T (f0... x0...)
...
| tn : forall xn..., T (fn... xn...)
.
```

Here, we declare the inductive type  $T : \forall A..., Type$ , and its constructors called  $t_i$  ( $i \in \{0..n\}$ ).

As a more concrete example, here is how the type of Booleans can be defined inductively :

```
Inductive Boolean : Type := true : Boolean | false : Boolean.
```

The above definition is essentially a formal statement of the following description of Booleans : a Boolean can have one of two shapes, *true* or *false*, and cannot be any other thing.

This means that, if we want to prove a property  $Px$  for some unknown Boolean  $x$ , all we need is to prove  $Ptrue$  and  $Pfalse$ .

This exact information is summed up in what we call the *induction principle* for Booleans. In Coq, it will be given the name `Boolean_rect`, for instance, and have the type  $\forall (P : \text{Boolean} \rightarrow \text{Type}), Ptrue \rightarrow Pfalse \rightarrow \forall (b : \text{Boolean}), Pb$ .