



# Distributed Random Servers with Timed Labels for Synchronization over Named Data Network

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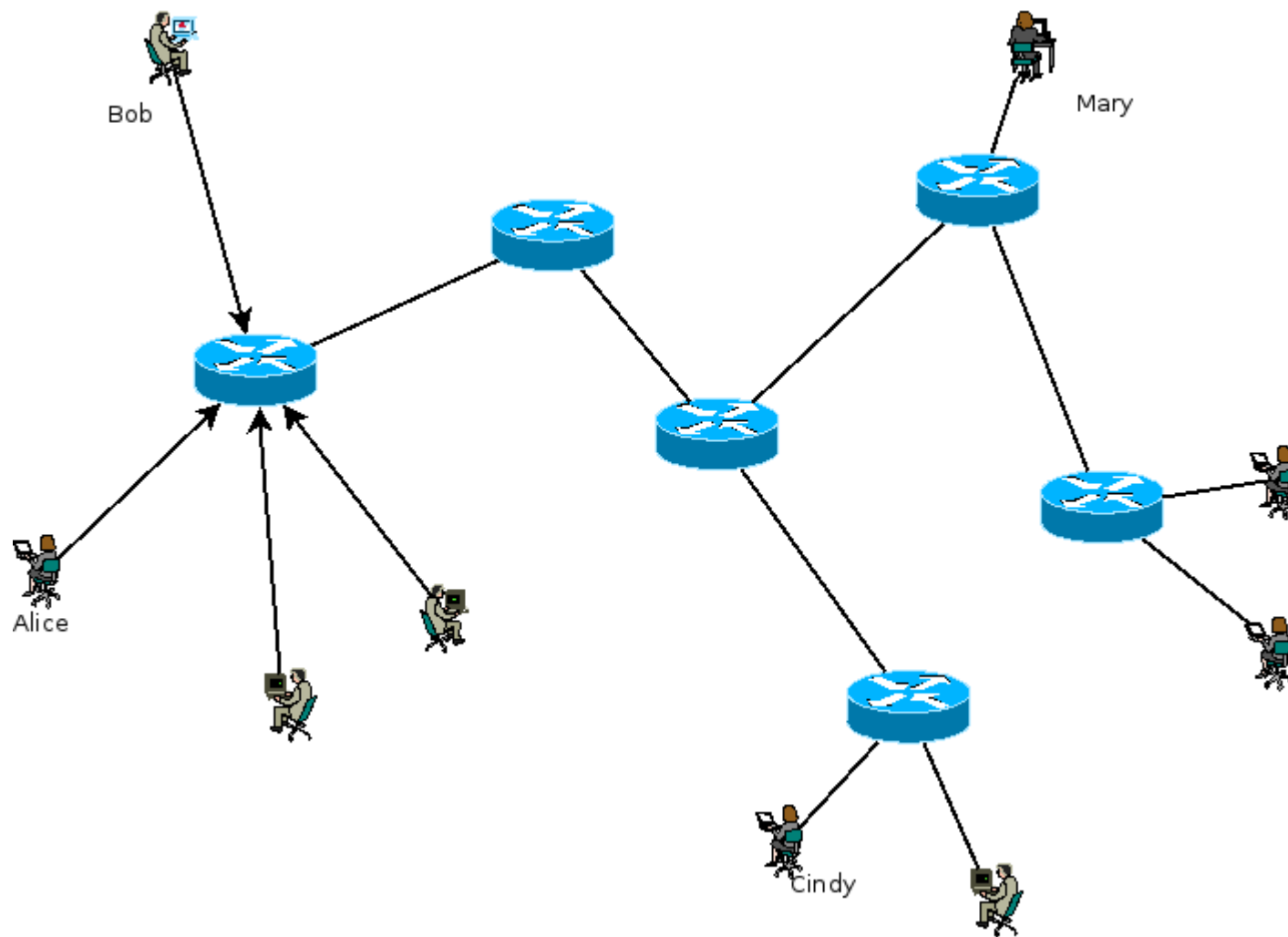


# 1. Problem

Let all participants share the same up-to-date data.

For example:

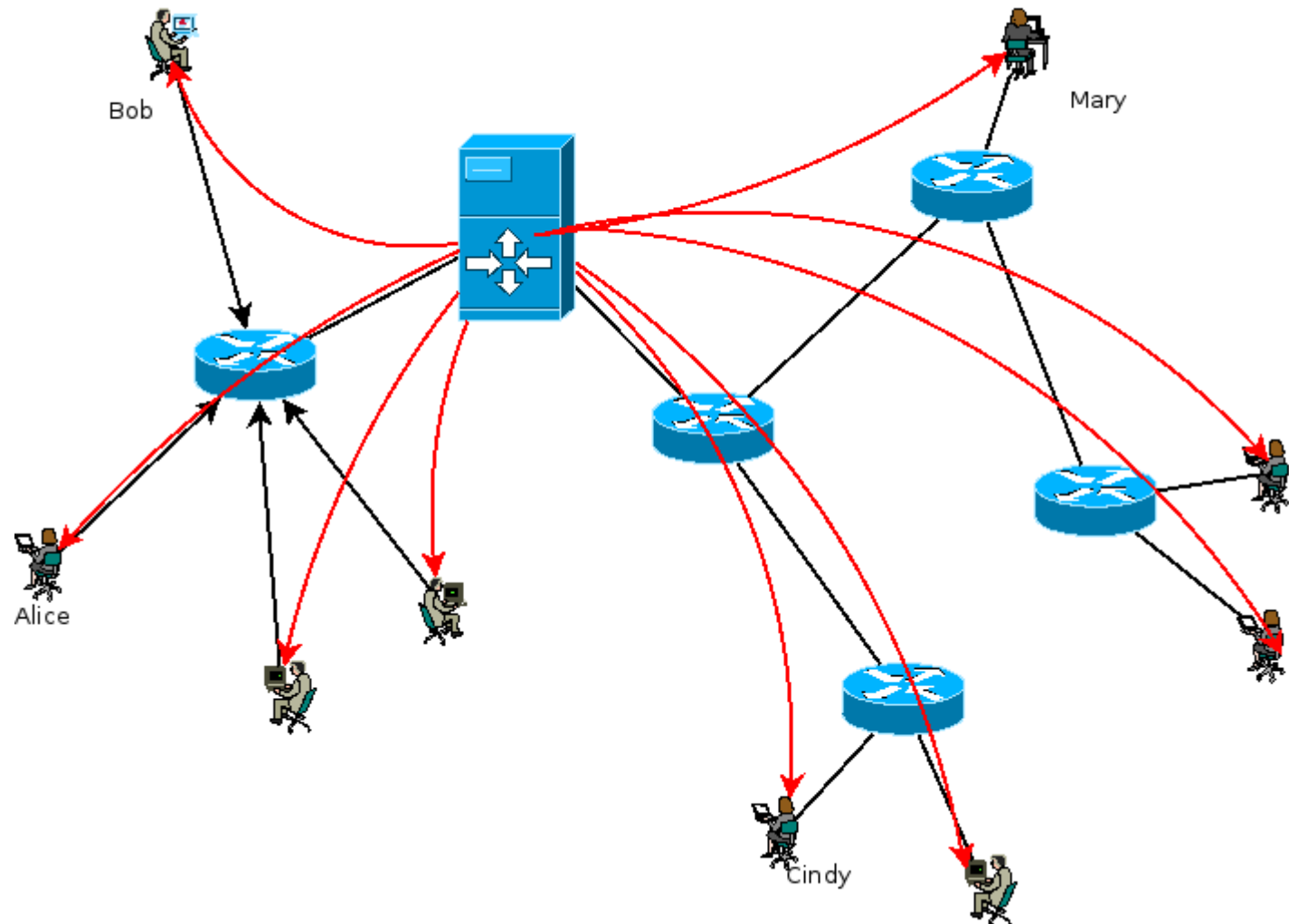
1. group chat
2. video conference



## 2. Background

- Why this matters?
  - Traditional IP based Network

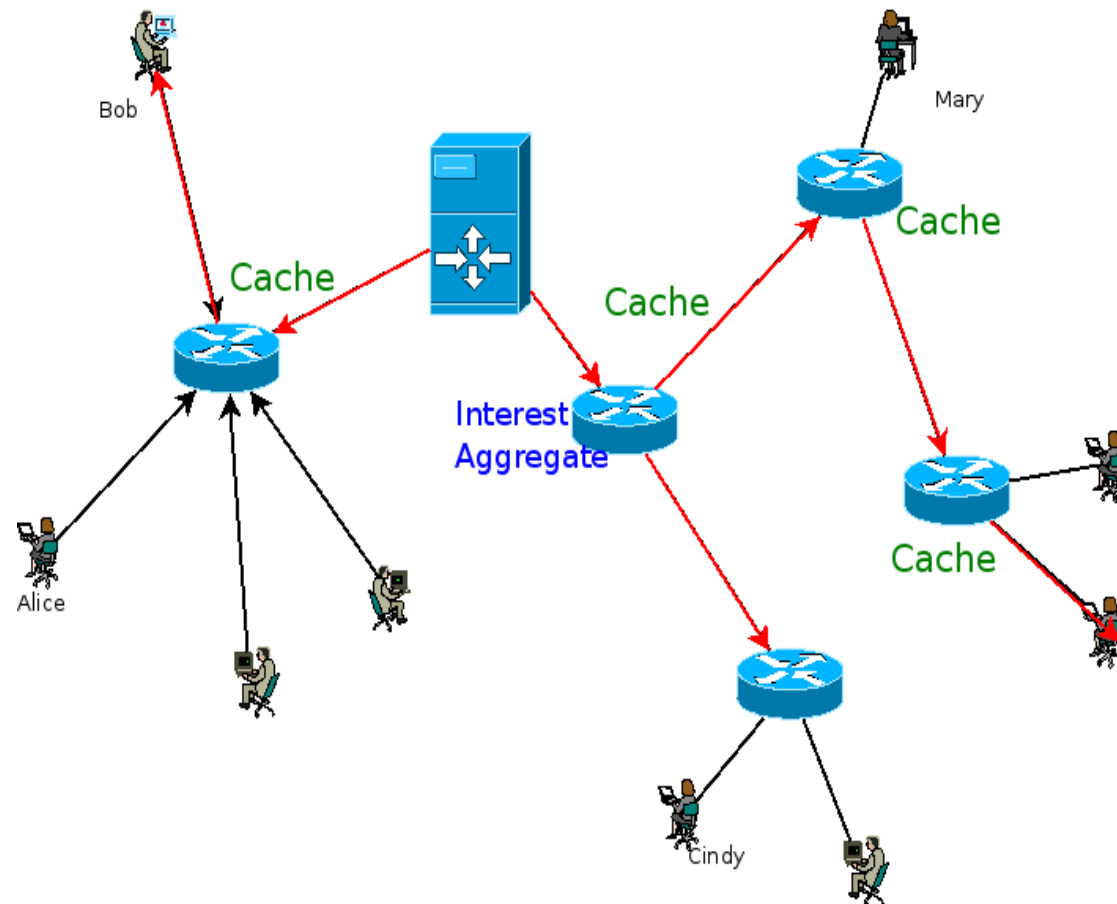
- Downsides:
  - Link Burden
  - overhead
  - Robust
  - Non Local





## 2. Background

- What NDN brings to us?
  - Traditional IP based Network
  - **Named Data Network**
    - **Cache**
    - **Interest Aggregate**



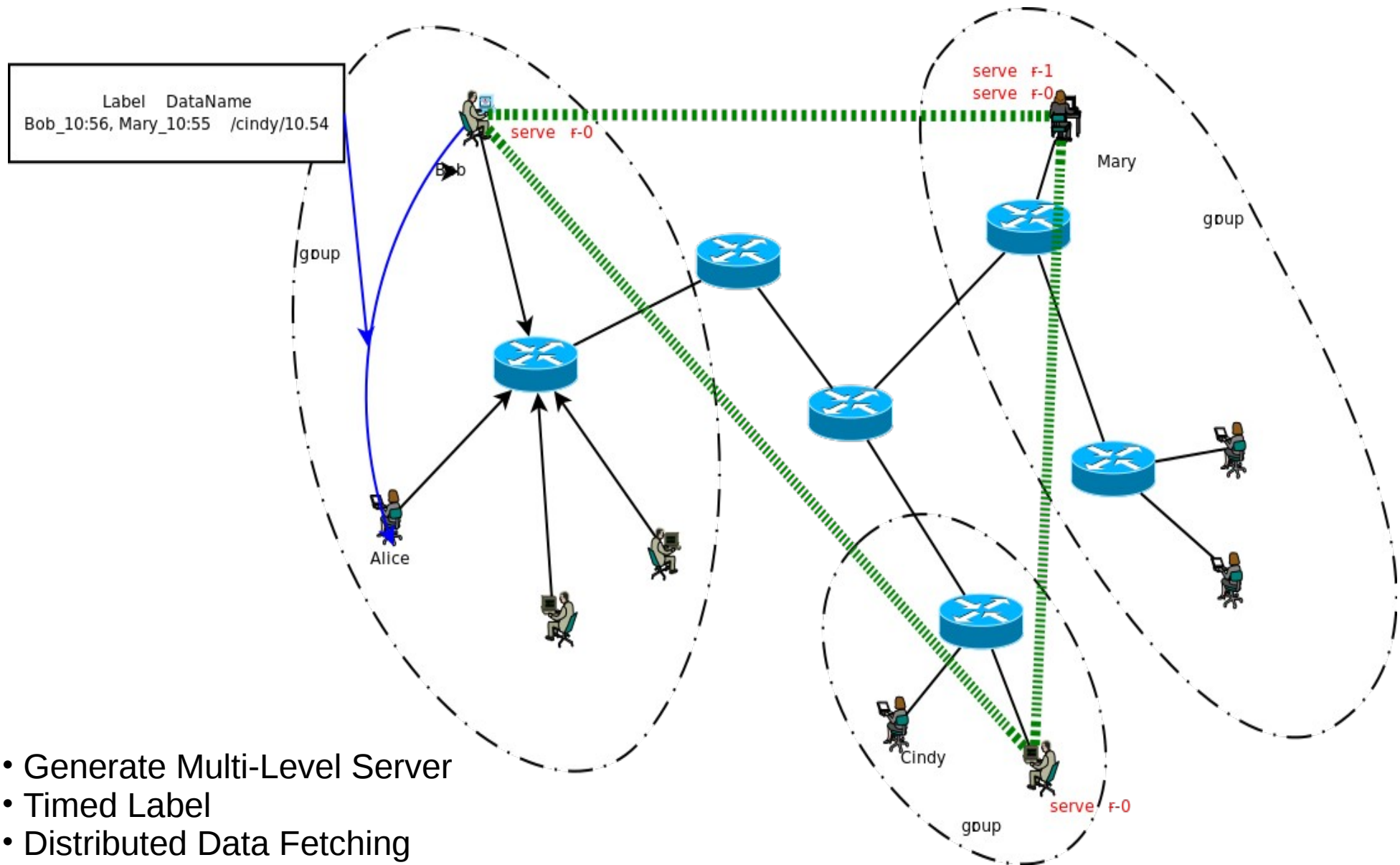


# 3. Design

- 3.1 Overview
- 3.2 Server Generation
  - Anything New Interest
- 3.3 Data Structure
  - True Message(for distributed Fetching)
  - Record(Timed Label)
  - Record Container(Index of Record)
- 3.4 Synchronization
  - Anything New Interest
  - Something New Interest
  - Data Fetching Interest



# 3.1 Overview

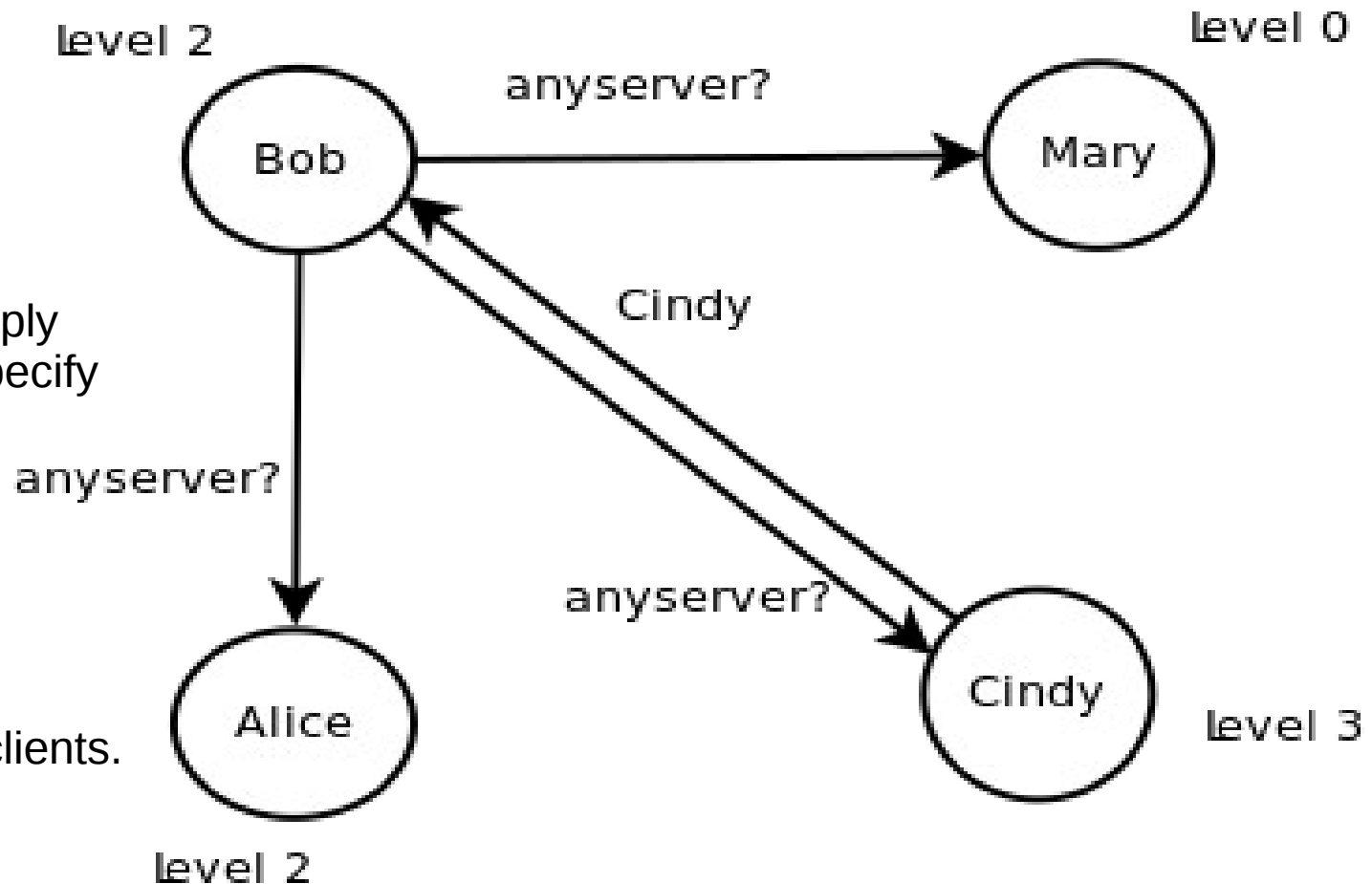




## 3.2 Server Generation

### Algorithm:

- Wait for random time
- Send Anyserver Request
  - Positive: set server
  - Negative: ready to reply
- Use different time-out to specify different levels.
- When a server fails:
  - Select one from his clients.



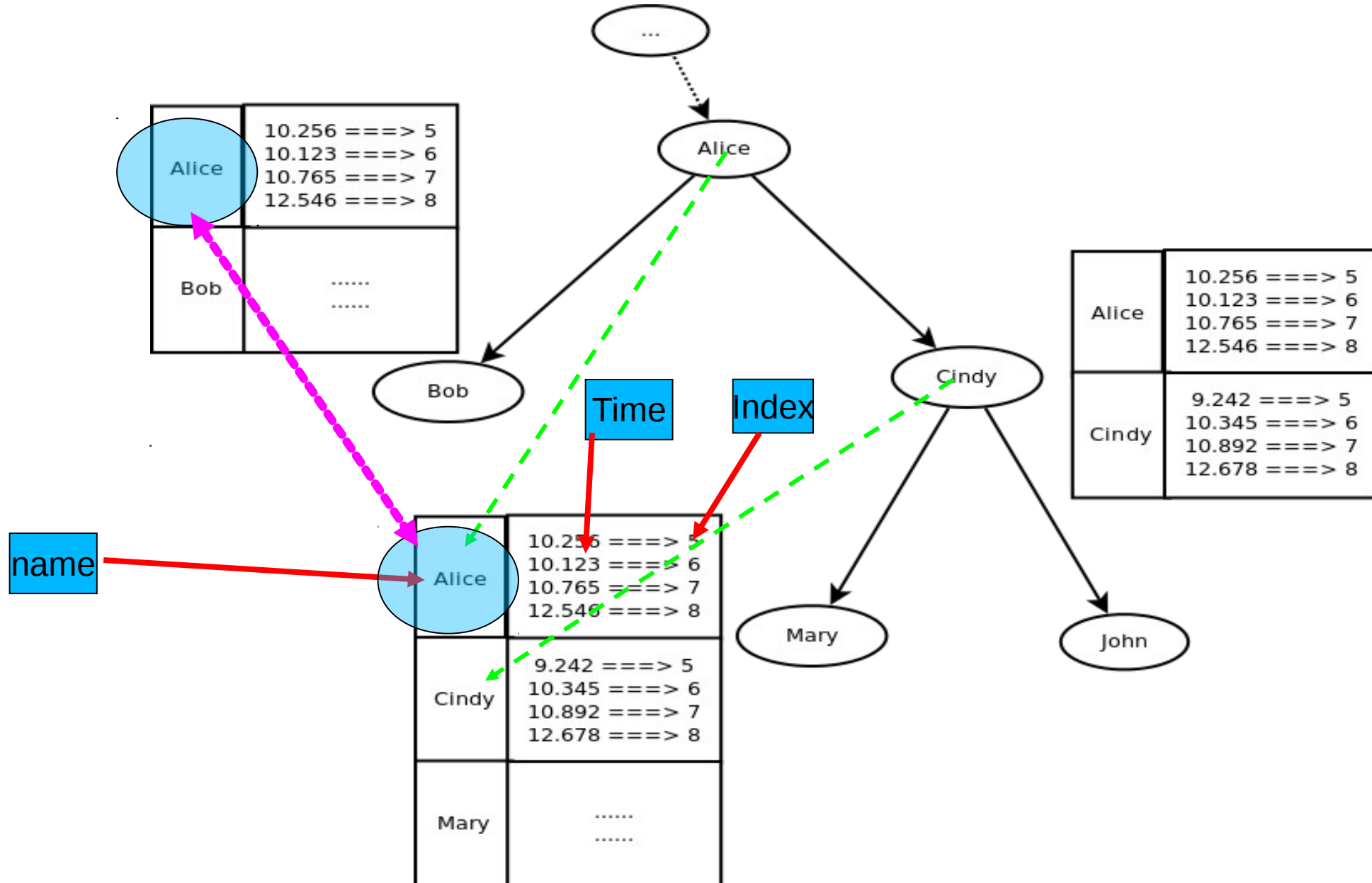


# 3.3 Data Structure

## Timed Labels



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# 3.3 Data Structure



Actual Data

Messages	
double time	std::string content
5.6783	Hello, everyone!
13.1870	Another Msg!

nameTimeIndex	
name	map<double, int>
Alice	
Bob	
Cindy	

When Receive Anythingnew Interest

- Compare time
- Get Index
- Get Labels

map<double, int>	
time	index
10.6477	5
12.8576	14
15.6743	17

Record Container

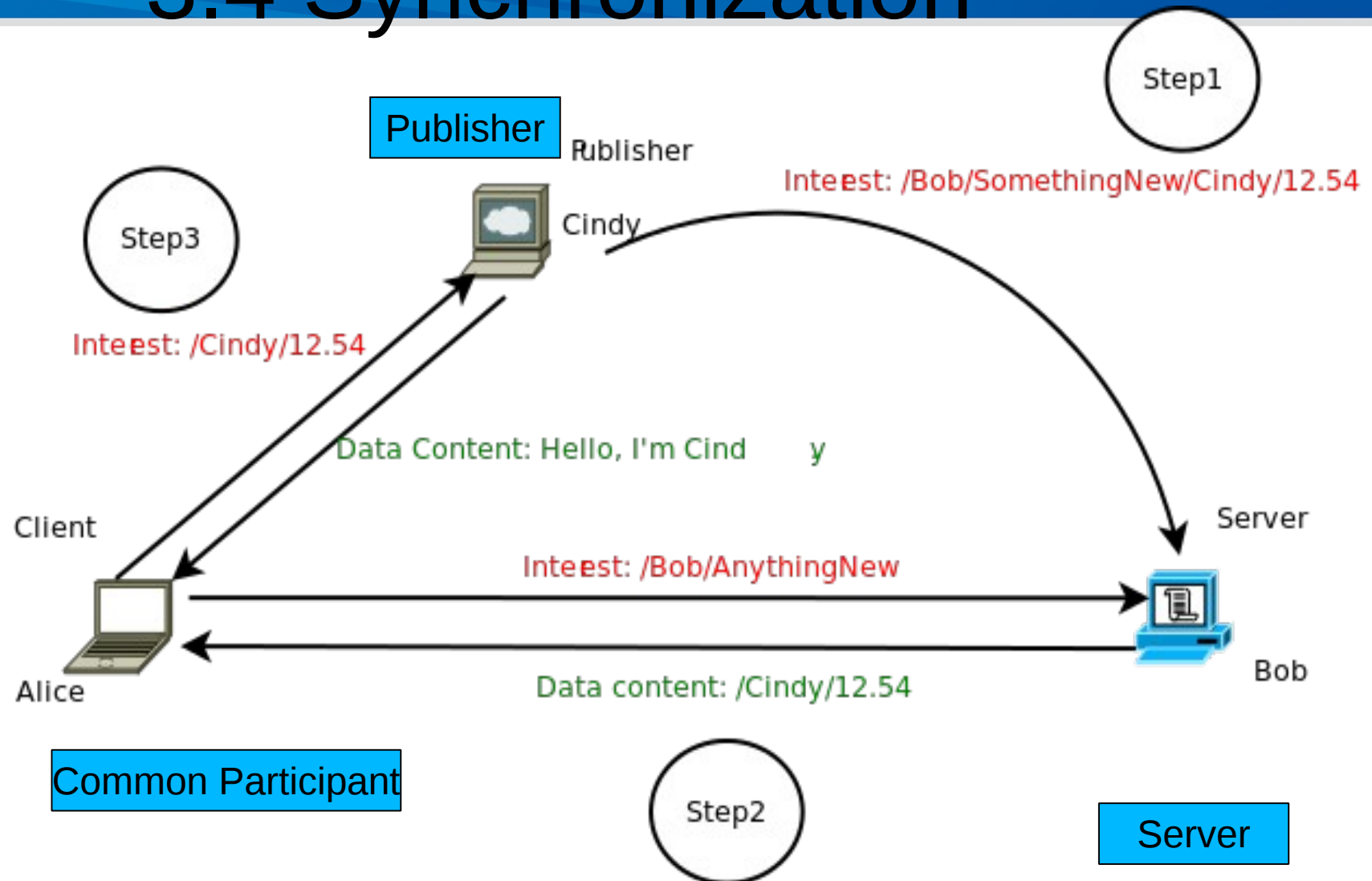
RecordContainer	
nameTimeIndex	
Records	

Records(Timed Labels)

dataName	/Alice/drsapp/5.6783
Record	
name	time
Alice	5.6783
Bob	5.7854
Cindy	5.9905
Labels	



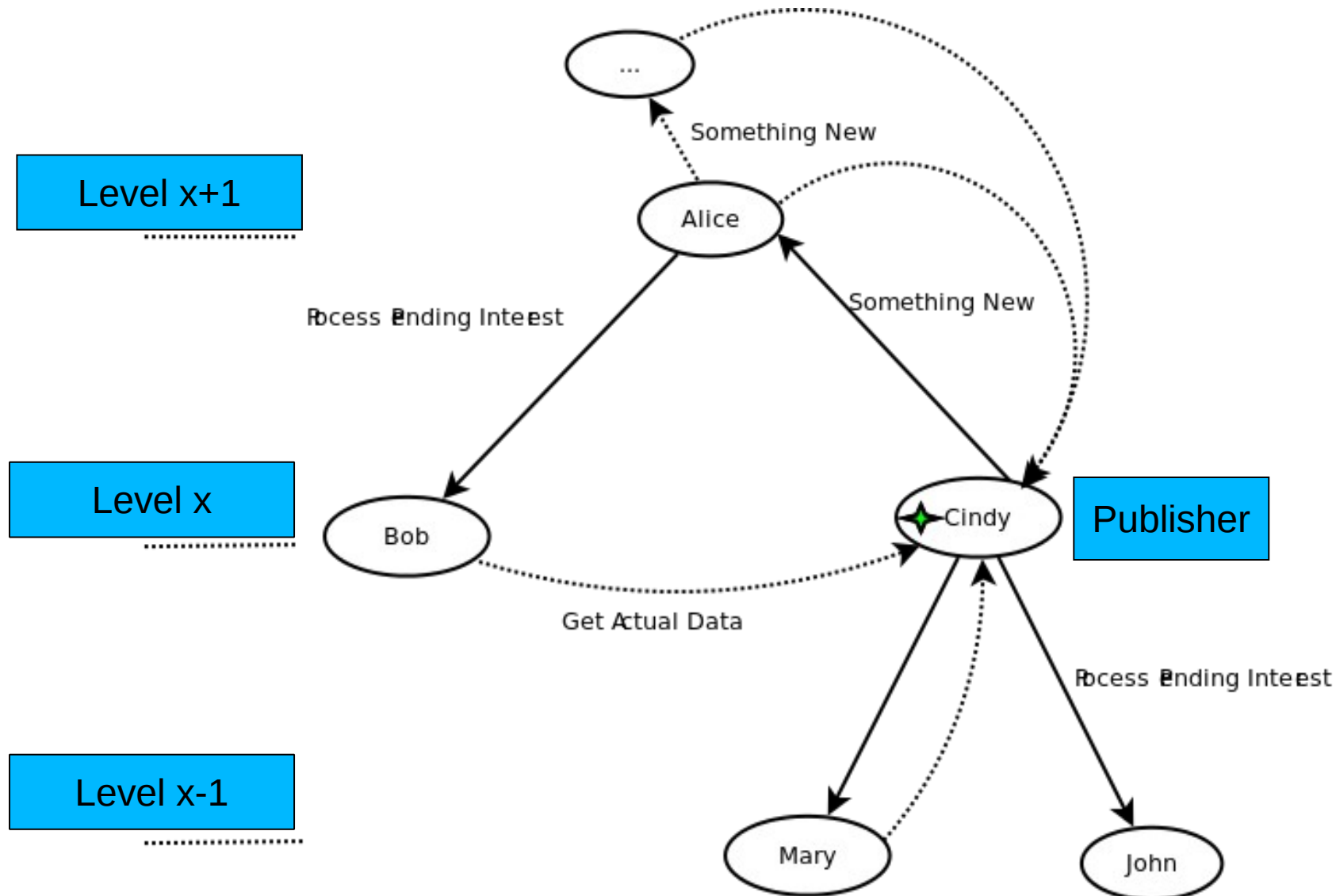
## 3.4 Synchronization



- STEP1: Publisher to Server: Something New
- STEP2: Server to Client: Satisfy Pending Interest
- STEP3: Client to Publisher: Distributed Data Fetching



## 3.4. Synchronization--Multi-Level





# 4. Evaluation

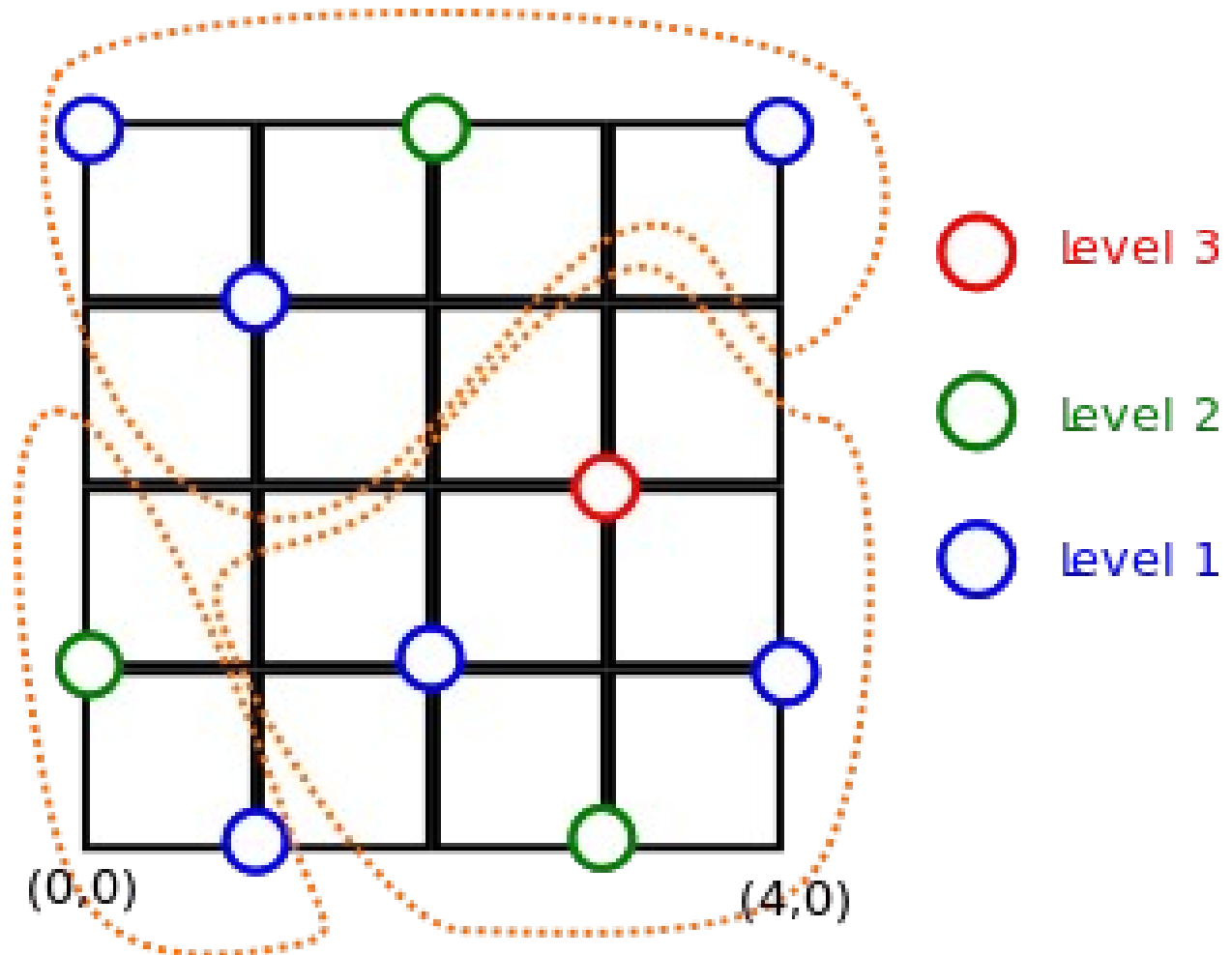
## Topology & Server Generation

### TOPOLOGY:

- 5x5 grid topology.
- Every node contains a participant.

### Function Correctness:

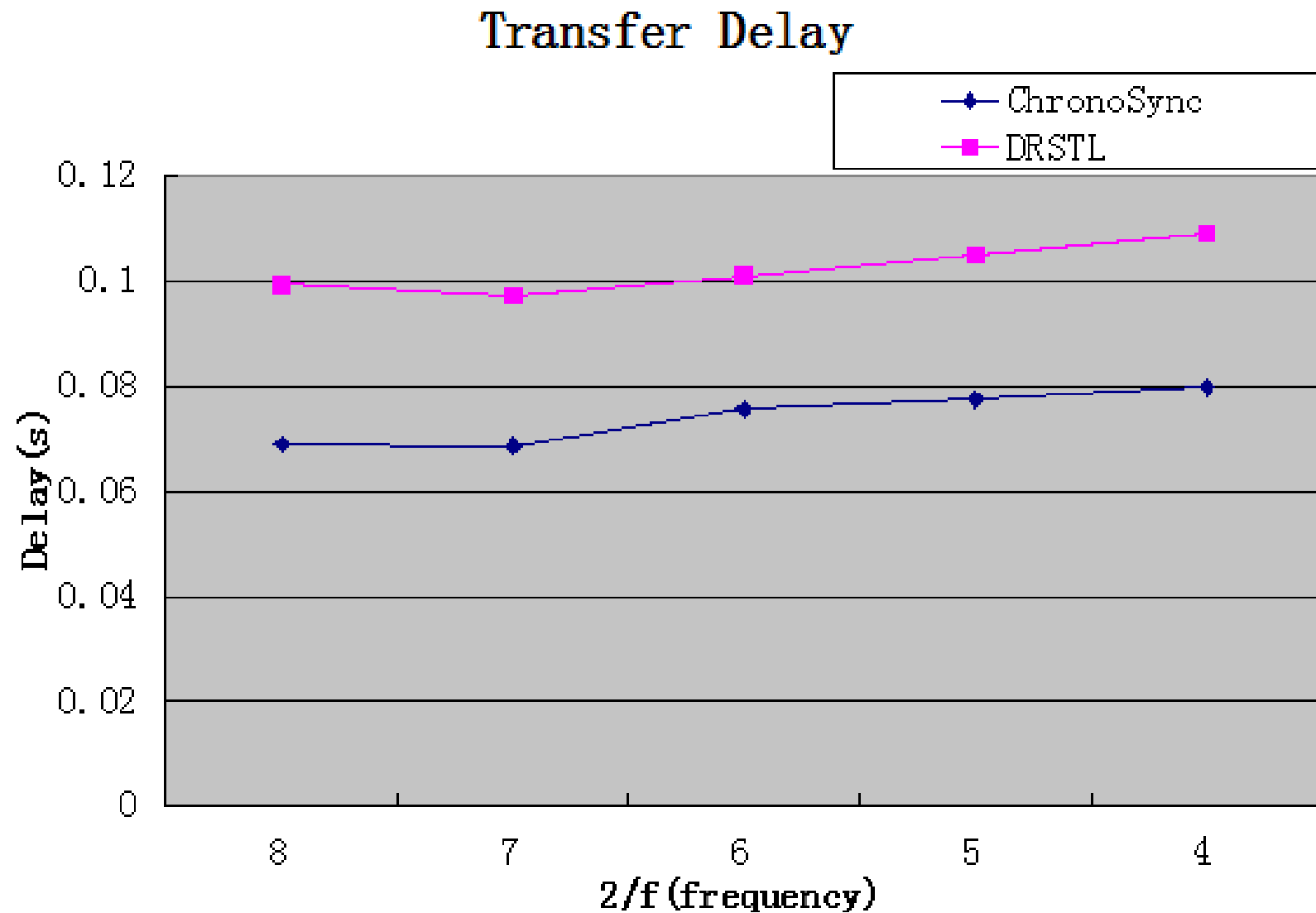
- Server Generated as expect.
- All participant received all Data.



# 4. Evaluation



Zhu, Z., & Afanasyev, A. (2013). Lets ChronoSync: Decentralized Dataset State Synchronization in Named Data Networking. In Proceedings of the 21st IEEE International Conference on Network Protocols(ICNP 2013)



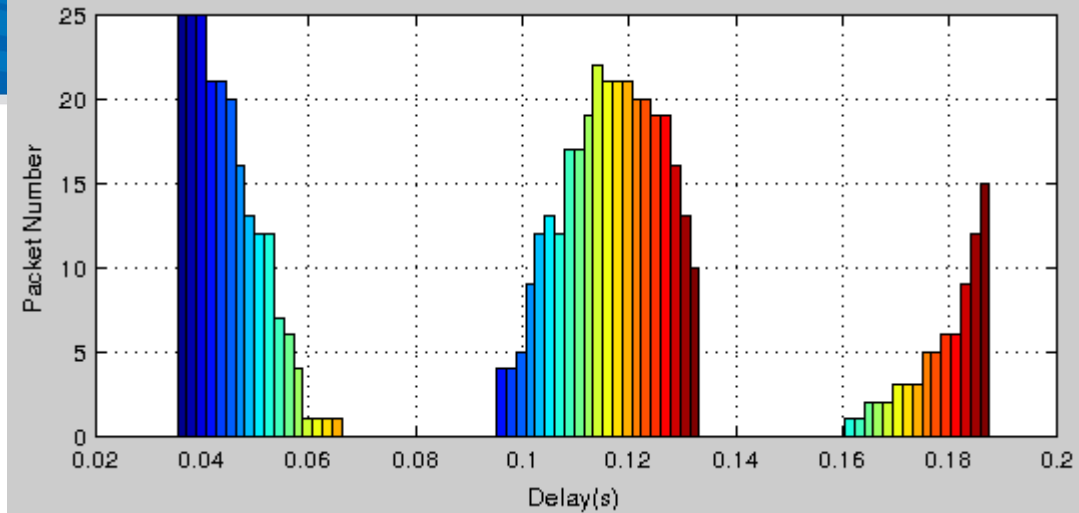
# 4. Evaluation



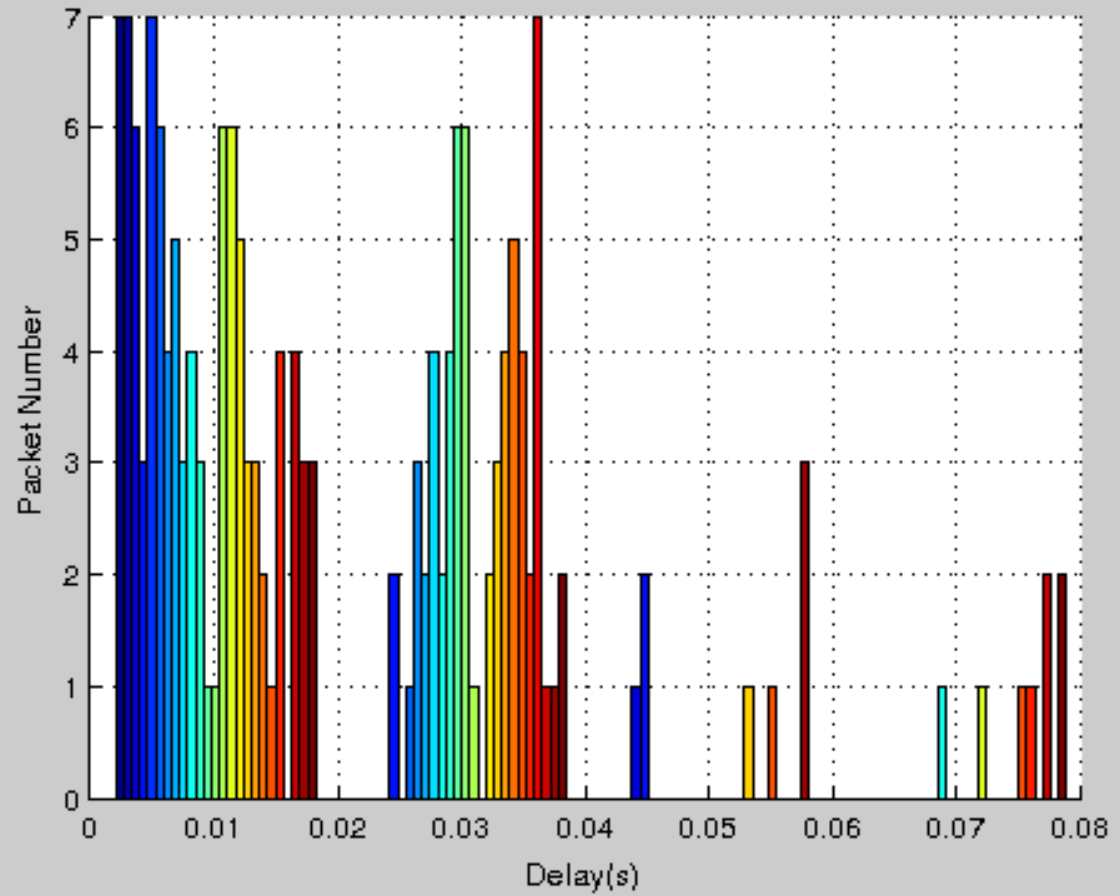
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Total Delay  
(Control Message +  
Data Fetching)



Data Fetching Delay



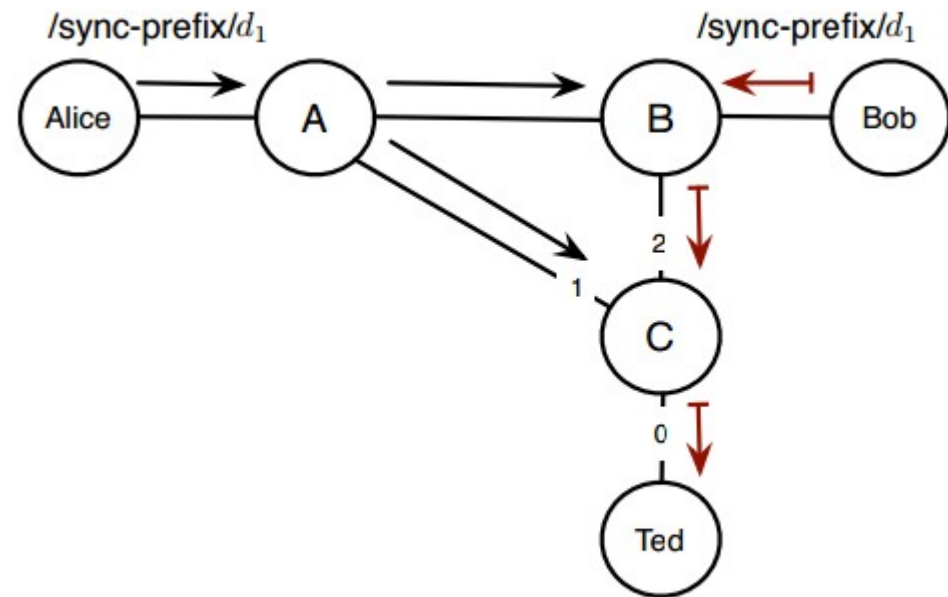




# ChronoSync's Fatal Downside

When Simultaneous Happens:

- Only one will get to the receiver!
- They are divided to different groups
- Recovery Interest has to be sent



It is a fully distributed protocol.  
Lack of Control Ability

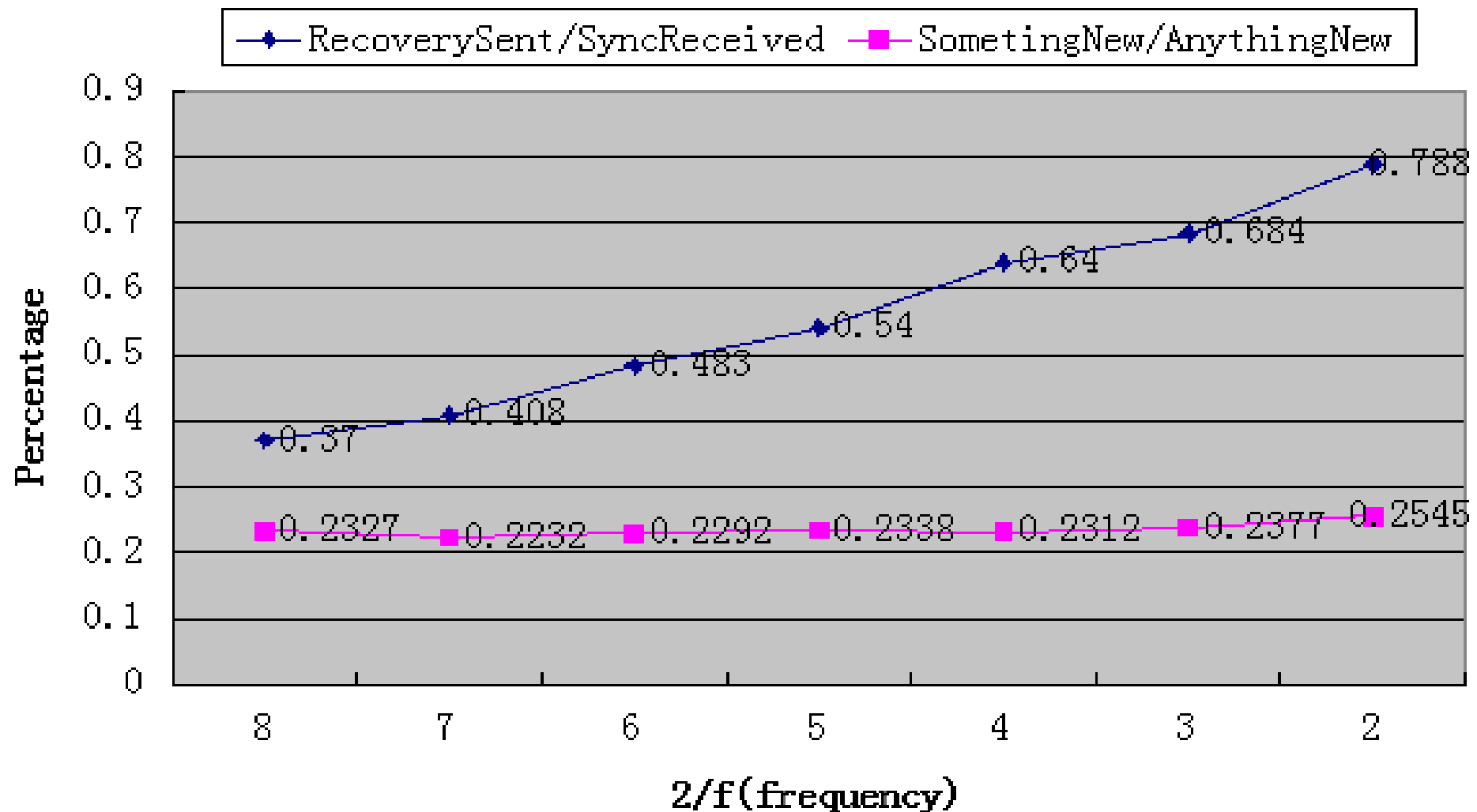
Fig. 6: An example of simultaneous data generation





# 4. Evaluation

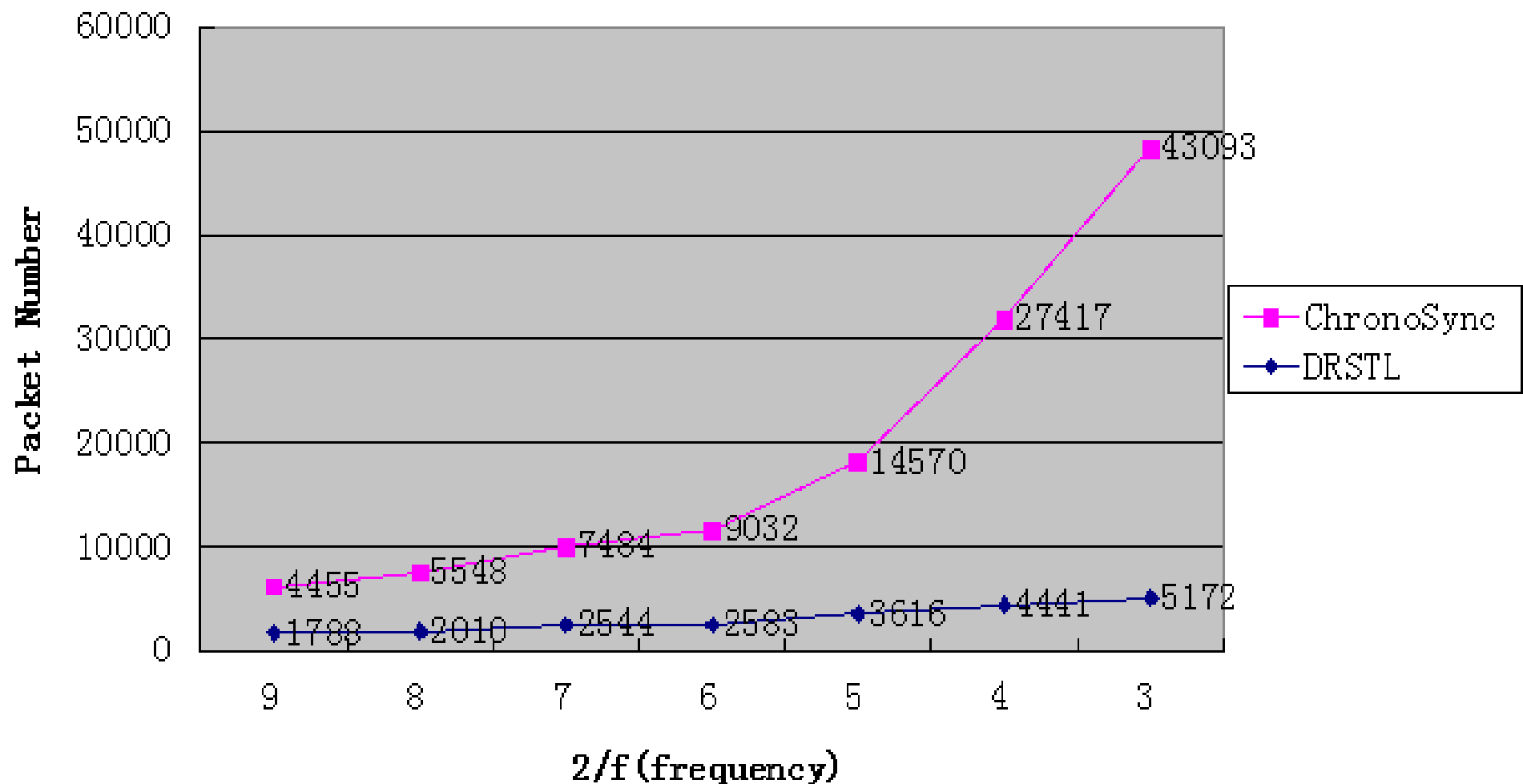
## Portion In Interests





## 4. Evaluation

Cumulative Incoming Interest Count For All  
Nodes In 5 Seconds





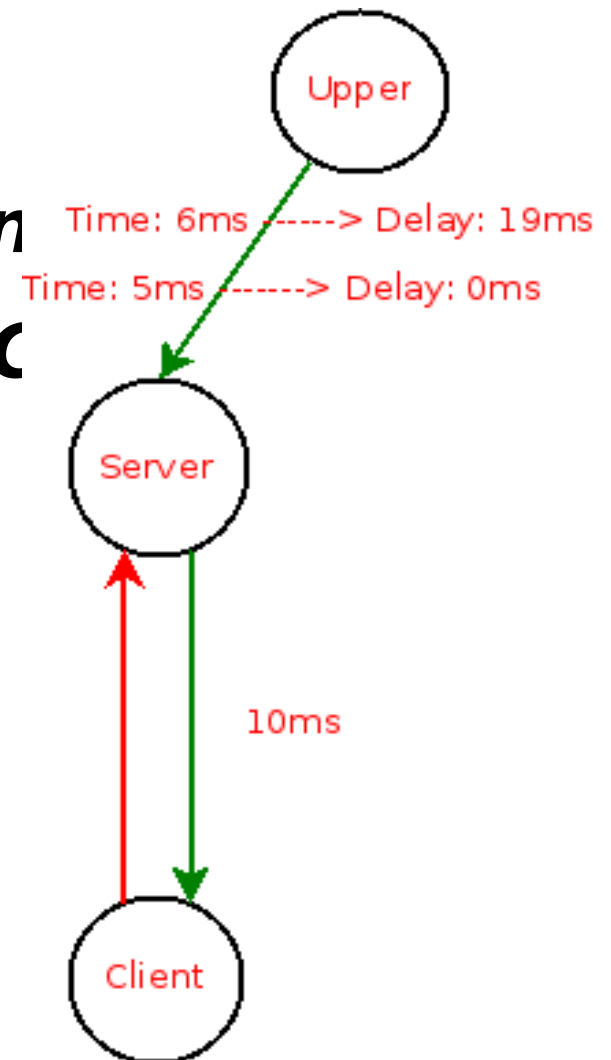
# 5. Future Work

## --- Inter-Level Delay

- Why delay increases when frequency increases

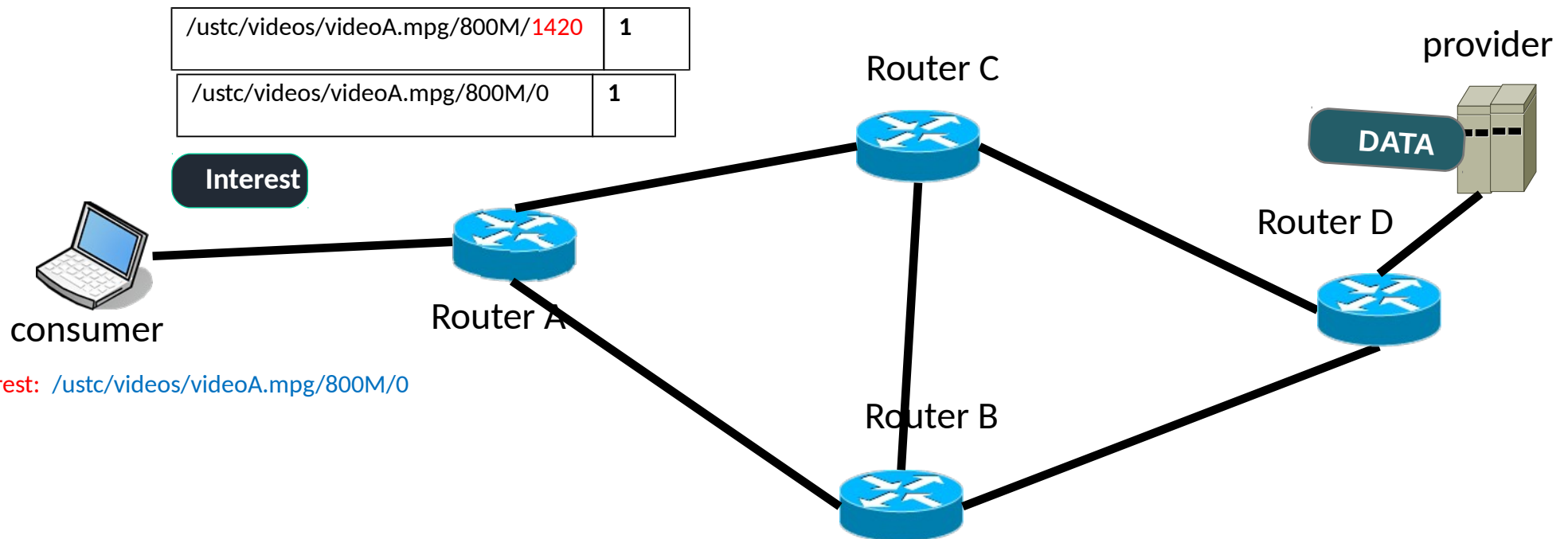
- Desired Solution:

❖ ***A Novel Flow-based Contin  
Data Replying Mechanism in IC***





# A Novel Flow-based Continuous Data Replying Mechanism in ICN



After the first Interest

- the consumer is **ready for receiving many data**
- The Provider is will **continuously put data**
- Routers will hold on the record, and **forward every data back**



# 5. Future Work

## --- Server Generation Algorithm

- Best fit for every given topology
- Robust
- Infrastructure
- ❖ A controller who holds information of the whole topology. Like SDN.