MASTER THEOREM (Abdul Bari) T(n) = aT(n/b) + f(n) $a \ge 1$ $b \ge 1$ $b \ge 1$ Case 1: If loga > K then O (nloga) Case 2: If $\log a = k$ then: if $p > -1 = O(n^k \log^{p+1} n)$ if $P=-1=\theta(n^{k}\log\log n)$ if $p < -1 = \theta(n^{k})$

Case 3: If loga < k Hen:

Ef P > 0 O(nt log Pn)

if P < 0 O(nt)

$$q: T(n) = 2T(n/2)+1$$

$$a=2$$

$$b=2$$

$$f(n) = \theta(1)$$

$$= \theta(n^{0} \log^{0} n)$$

$$\log a = 1$$

$$(ase 1: \theta(n))$$

$$case 1: \theta(n)$$

$$q = 4T(n/2) + n$$

$$a = 9$$

60 a = 2

(ase 1 0 (n2)

 $f(n) = \theta(n)$ $= \theta(n) \log^{n} n$

2 > 1

T(n) - 8 + (n/2) + n $(ase 1 = o(n^3)$ log 8 = 3 3>1 T(n)= 9T(n/3)+ (099=2 > K=0 Case 1 $\theta(n^2)$ 8T(1/2) + n log n log 8 = 3 4(n) = 0(n/log/n) case $1 = \Theta(n^3)$

eg 5: 2T(1/2) + n 09 2 = 1 = 0 (nt log P+1 n) = O (nlogn) eg b: T(n) = 2T(n/2) + 1 10gn k=1 $\log a=1$ Case 2: case p=-1to (n log (logn)) > T(n) = 2T(n/2) + n