

20201024 Homework



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1. (1) Prove $H_0(X) = \sum_{k \in L} q_k h_k$

Proof. ① If there is only one internal node, the lemma is trivially true.

② Assume it's true for all code trees with n internal nodes. When there are $n+1$ internal nodes, let k be an internal node such that k is the parent of a leaf c

$\begin{matrix} k \\ | \\ c \end{matrix}$

with maximum order. If the outcome of X is not a leaf descending from k , we

identify the outcome exactly, otherwise call this random variable V to be a

child of k . If we don't identify the outcome that is a leaf node of k , which

happens with q_k , we further identify which of child of k is the outcome

We call this W . $X = (V, W)$. The outcome of V is a situation " n ", then:

$$H(V) = \sum_{k' \in L \setminus \{k\}} q_{k'} h_{k'}. \text{ So } H(X) = H(V) + H(W|V) = \sum_{k' \in L \setminus \{k\}} q_{k'} h_{k'} + (1 - q_k) \cdot 0 +$$

$$q_k h_k = \sum_{k \in L} q_k h_k \quad \therefore H_0(X) = \sum_{k \in L} q_k h_k$$

(2) prove $L = \sum_{k \in L} q_k$.

Proof. Define $a_{ki} = \begin{cases} 1 & \text{if leaf } c_i \text{ is a descendent of internal node } k. \\ 0 & \text{otherwise.} \end{cases}$

$$h_k = \sum_{i \in L} a_{ki}, \quad q_k = \sum_i p_i a_{ki}$$

$$\therefore L = \sum_k p_k h_k = \sum_k p_k \sum_{i \in L} a_{ki} = \sum_{i \in L} p_i \sum_k a_{ki} = \sum_{i \in L} p_i = 1$$

2. ~~pro~~ prove the equivalence of 4 definitions of convergence in probability.

$$(1) \forall \epsilon > 0, \lim_{n \rightarrow \infty} P(|X_n - Y| < \epsilon) = 1. \quad (2) \forall \epsilon > 0, \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \epsilon) > 1 - \epsilon.$$

$$(3) \forall \epsilon > 0, \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \epsilon) = 1. \quad (4) \forall \epsilon > 0, \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \epsilon) > 1 - \epsilon.$$

$$(1) \Rightarrow (2): \forall \epsilon > 0, \lim_{n \rightarrow \infty} P(|X_n - Y| < \epsilon) = 1 > 1 - \epsilon.$$

$$(2) \Rightarrow (4): \forall \epsilon > 0, \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \epsilon) \geq \lim_{n \rightarrow \infty} P(|X_n - Y| < \epsilon) > 1 - \epsilon.$$

$$(1) \Rightarrow (3): \forall \epsilon > 0, 1 = \lim_{n \rightarrow \infty} P(|X_n - Y| < \epsilon) \leq \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \epsilon) \leq 1$$

$$\therefore \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \epsilon) = 1$$

$$(3) \Rightarrow (4): \forall \epsilon > 0, 1 = \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \epsilon) > 1 - \epsilon.$$

$$(4) \Rightarrow (1): \forall 0 < \delta < \epsilon, \lim_{n \rightarrow \infty} P(|X_n - Y| < \epsilon) \geq \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \delta) > 1 - \delta.$$

$$\therefore \lim_{n \rightarrow \infty} P(|X_n - Y| < \epsilon) \geq \lim_{\delta \rightarrow 0} \lim_{n \rightarrow \infty} P(|X_n - Y| \leq \delta) > \lim_{\delta \rightarrow 0} (1 - \delta) = 1.$$

$$\therefore \lim_{n \rightarrow \infty} P(|X_n - Y| < \epsilon) = 1$$

So the four definitions are equivalent.

3. Show examples such that (1) $x \in T_{[X]\epsilon}^n$ and $x \in W_{[X]\epsilon}^n$.

(2) $x \notin T_{[X]\epsilon}^n$ and $x \in W_{[X]\epsilon}^n$. (3) $x \in T_{[X]\epsilon}^n$ and $x \notin W_{[X]\epsilon}^n$. (4) $x \notin T_{[X]\epsilon}^n$ and $x \notin W_{[X]\epsilon}^n$.

$$(1) \begin{array}{c|c|c|c} X & 0 & 1 & 2 \\ \hline p & 0.5 & 0.5 & 0.25 \end{array} \quad x_1 = (1, 0, 2, 0). \quad H(X) = -\frac{1}{2} \log \frac{1}{2} - \frac{1}{4} \log \frac{1}{4} - \frac{1}{4} \log \frac{1}{4} = \frac{3}{2}.$$

$$(2) x_2 = (0, 1, 0, 1) \quad \frac{N(i; x)}{n} = \frac{2}{4} \neq p(i) \text{ so } x_2 \notin T_{[X]\epsilon}^n$$

$$(3) \text{ if } x \in T_{[X]\epsilon}^n \text{ then } x \in W_{[X]\epsilon}^n \text{ because } 2^{-n(H(x)+1)} \leq p(x) \leq 2^{-n(H(x)-1)}$$

$$\text{then } H(x) - 1 \leq -\frac{1}{n} \log p(x) \leq H(x) + 1, \quad y \rightarrow 0 \text{ when } \epsilon \rightarrow 0 \text{ and } \eta \rightarrow 0.$$

$$\therefore x \in W_{[X]\eta}^n, x \in W_{[X]\epsilon}^n.$$

$$(4) x_4 = (0, 0, 0) \quad \frac{N(i; x)}{n} = \frac{1}{3} \neq p(i) \text{ so } x_4 \notin T_{[X]\epsilon}^n.$$

$$-\frac{1}{n} \log p(x_4) = -\frac{1}{3} \log (0.5 \times 0.5 \times 0.5) = -\log 0.5 = 1 \neq H(X) = \frac{3}{2}. \text{ so } x_4 \notin W_{[X]\epsilon}^n.$$

$$(3) \text{ let } \begin{array}{c|c|c} X & 0 & 1 \\ \hline p & 0.1 & 0.9 \end{array} \quad H(X) = -0.1 \log 0.1 - 0.9 \log 0.9 \approx 0.467$$

$$n = 100, \delta = 0.1, \epsilon = 0.1$$

$$\text{For } T_{[X]\epsilon}^n: |N(0; x) - np(0)| + |N(1; x) - np(1)| \leq n\epsilon \Rightarrow 85 \leq N(1; x) \leq 95.$$

$$\text{For } W_{[X]\epsilon}^n: |-N(0; x) \log p(0) - N(1; x) \log p(1) - nH(X)| \leq n\epsilon \Rightarrow 87 \leq N(1; x) \leq 93.$$

\therefore When $N(1; x) = 94$ or 95 , $x \in T_{[X]\epsilon}^n$ and $x \notin W_{[X]\epsilon}^n$.

$$\text{for example: } x = (0, 0, 0, 0, 0, \underbrace{1, 1, \dots, 1}_{95 \uparrow 1})$$

1. Prove the following two lemmas of Local Redundancy Theorem

* **Lemma 4.19.** $H_D(X) = \sum_{k \in L} q_k h_k$. (maybe following example will show you an easier description than the textbook in the proof.)

Lemma 4.19 $H_D(X) = \sum_k q_k h_k$

1st Page

A. $H(X) = \sum -\frac{1}{5} \log \frac{1}{5} = \log 5$.

B. 四个内节点:

$$\sum_{k=1}^4 q_k h_k = \log 5$$

内部节点:

$$\begin{cases} h_4: -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} & q_4 = \frac{2}{5} \\ h_3: -\frac{2}{5} \log \frac{2}{5} - \frac{1}{5} \log \frac{1}{5} & q_3 = \frac{3}{5} \\ h_2: -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} & q_2 = \frac{2}{5} \\ h_1: -\frac{2}{5} \log \frac{2}{5} - \frac{3}{5} \log \frac{3}{5} & q_1 = 1 \end{cases}$$

总结:

$$\sum_{k=1}^4 q_k h_k = \begin{cases} -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} & ④ \\ -\frac{2}{5} \log \frac{2}{5} - \frac{1}{5} \log \frac{1}{5} & ③ \\ -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} & ② \\ -\frac{2}{5} \log \frac{2}{5} - \frac{3}{5} \log \frac{3}{5} & ① \end{cases} = \log 5$$

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A. $H(X) = (-\frac{1}{5} \log \frac{1}{5}) \times 3 + (-\frac{2}{5} \log \frac{2}{5})$

B. 三个内节点:

$$\sum_{k=1}^3 q_k h_k$$

内部节点:

$$\begin{cases} h_3: -\frac{2}{5} \log \frac{2}{5} - \frac{1}{5} \log \frac{1}{5} & q_3 = \frac{3}{5} \\ h_2: -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} & q_2 = \frac{2}{5} \\ h_1: -\frac{2}{5} \log \frac{2}{5} - \frac{3}{5} \log \frac{3}{5} & q_1 = 1 \end{cases}$$

总结:

$$\sum_{k=1}^3 q_k h_k = \begin{cases} -\frac{2}{5} \log \frac{2}{5} - \frac{1}{5} \log \frac{1}{5} & ③ \\ -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} & ② \\ -\frac{2}{5} \log \frac{2}{5} - \frac{3}{5} \log \frac{3}{5} & ① \end{cases}$$

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总结:

$$\begin{cases} \sum q_k h_k \text{ 端改变值: } q_4 h_4 \\ = -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} \\ H(X) \text{ 端改变值: } -\frac{1}{5} \log \frac{1}{5} - \frac{1}{5} \log \frac{1}{5} \\ - (-\frac{2}{5} \log \frac{2}{5}) \\ = -q_{4a} \log q_{4a} - q_{4b} \log q_{4b} \\ - (-q_4 \log q_4) \\ = -q_{4a} \log \frac{q_{4a}}{q_4} - q_{4b} \log \frac{q_{4b}}{q_4} \end{cases}$$

* **Lemma 4.20.** $L = \sum_{k \in L} q_k$.

2. Prove the equivalence of four definitions of convergence in probability.

3. Show examples such that

1) $x \in T_{[X]_\varepsilon}^n$ and $x \in W_{[X]_\varepsilon}^n$

$$2) \quad \boldsymbol{x} \notin T_{[X]\varepsilon}^n \text{ and } \boldsymbol{x} \in W_{[X]\varepsilon}^n$$

$$3) \quad \boldsymbol{x} \in T_{[X]\varepsilon}^n \text{ and } \boldsymbol{x} \notin W_{[X]\varepsilon}^n$$

$$4) \quad \boldsymbol{x} \notin T_{[X]\varepsilon}^n \text{ and } \boldsymbol{x} \notin W_{[X]\varepsilon}^n$$