Conditional Loop Instructions

- LOOPZ and LOOPE
- LOOPNZ and LOOPNE

LOOPZ and **LOOPE**

Syntax:

LOOPE destination
LOOPZ destination

- Logic:
 - ECX ← ECX 1
 - if ECX > 0 and ZF=1, jump to destination
- Useful when scanning an array for the first element that does not match a given value.

LOOPNZ and **LOOPNE**

- LOOPNZ (LOOPNE) is a conditional loop instruction
- Syntax:

LOOPNE destination

- Logic:
 - ECX ← ECX 1;
 - if ECX > 0 and ZF=0, jump to destination
- Useful when scanning an array for the first element that matches a given value.

LOOPNZ Example

The following code finds the first positive value in an array:

```
.data
array SWORD -3,-6,-1,-10,10,30,40,4
sentinel SWORD 0
.code
  mov esi, OFFSET array
  mov ecx, LENGTHOF array
next:
  test WORD PTR [esi],8000h ; test sign bit
  pushfd
                          ; push flags on stack
   add esi, TYPE array
                          ; pop flags from stack
  popfd
                          ; continue loop
  loopnz next
   jnz quit
                        ; none found
   quit:
```

Your turn . . .

Locate the first nonzero value in the array. If none is found, let ESI point to the sentinel value:

```
.data
array SWORD 50 DUP(?)
sentinel SWORD OFFFFh
.code
   mov esi, OFFSET array
   mov ecx, LENGTHOF array
L1: cmp WORD PTR [esi],0
                                 ; check for zero
   (fill in your code here)
quit:
```

... (solution)

```
.data
array SWORD 50 DUP(?)
sentinel SWORD OFFFFh
.code
  mov esi, OFFSET array
  mov ecx, LENGTHOF array
L1: cmp WORD PTR [esi],0 ; check for zero
                          ; push flags on stack
  pushfd
  add esi, TYPE array
                          ; pop flags from stack
  popfd
  loope next
                          ; continue loop
  jz quit
                          ; none found
  quit:
```

Conditional Structures

- Block-Structured IF Statements
- Compound Expressions with AND
- Compound Expressions with OR
- WHILE Loops
- Table-Driven Selection

Block-Structured IF Statements

Assembly language programmers can easily translate logical statements written in C++/Java into assembly language. For example:

```
if( op1 == op2 )
  X = 1;
else
  X = 2;
```

```
mov eax,op1
cmp eax,op2
jne L1
mov X,1
jmp L2
L1: mov X,2
L2:
```

Your turn . . .

Implement the following pseudocode in assembly language. All values are unsigned:

```
if( ebx <= ecx )
{
  eax = 5;
  edx = 6;
}</pre>
```

Your turn . . .

Implement the following pseudocode in assembly language. All values are 32-bit signed integers:

```
if( var1 <= var2 )
  var3 = 10;
else
{
  var3 = 6;
  var4 = 7;
}</pre>
```

Compound Expression with AND [1/3]

- When implementing the logical AND operator, consider that HLLs use short-circuit evaluation
- In the following example, if the first expression is false, the second expression is skipped:

```
if (al > bl) AND (bl > cl)
   X = 1;
```

Compound Expression with AND [2/3]

```
if (al > bl) AND (bl > cl)
X = 1;
```

This is one possible implementation . . .

Compound Expression with AND [3/3]

```
if (al > bl) AND (bl > cl)
    X = 1;
```

But the following implementation uses 29% less code by reversing the first relational operator.

Your turn . . .

Implement the following pseudocode in assembly language. All values are unsigned:

```
if(ebx <= ecx) AND
  (ecx > edx )
{
  eax = 5;
  edx = 6;
}
```

Compound Expression with OR [1/2]

- When implementing the logical OR operator, consider that HLLs use short-circuit evaluation
- In the following example, if the first expression is true, the second expression is skipped:

```
if (al > bl) OR (bl > cl)
  X = 1;
```

Compound Expression with OR [2/2]

```
if (al > bl) OR (bl > cl)
  X = 1;
```

We can use "fall-through" logic to keep the code as short as possible:

WHILE Loops

A WHILE loop is really an IF statement followed by the body of the loop, followed by an unconditional jump to the top of the loop. Consider the following example:

```
while( eax < ebx)
  eax = eax + 1;</pre>
```

This is a possible implementation:

Your turn . . .

Implement the following loop, using unsigned 32-bit integers:

```
while( ebx <= val1)
{
    ebx = ebx + 5;
    val1 = val1 - 1
}</pre>
```

Table-Driven Selection [1/3]

- Create a table containing lookup values and the offsets of labels or procedures
- Use a loop to search the table

Table-Driven Selection [2/3]

Step 1: Create a table containing lookup values and procedure offsets.

```
.data
CaseTable BYTE 'A' ; lookup value

DWORD Process_A ; address of procedure
EntrySize = ($ - CaseTable)
BYTE 'B'
DWORD Process_B
BYTE 'C'
DWORD Process_C
BYTE 'D'
DWORD Process_D

NumberOfEntries = 4
```

Table-Driven Selection [3/3]

Step 2: Use a loop to search the table. When a match is found, we call the procedure offset stored in the current table entry.

```
mov ebx,OFFSET CaseTable
                                ; point EBX to the table
   mov ecx, NumberOfEntries
                                 ; loop counter
                                 : match found?
L1: cmp al,[ebx]
                                 : no: continue
   jne L2
   call NEAR PTR [ebx + 1]
                                ; yes: call the procedure
   imp L3
                                ; and exit the loop
L2: add ebx, EntrySize
                                 ; point to next entry
   loop L1
                                 ; repeat until ECX = 0
L3:
```

required for procedure pointers

Application: Finite-State Machines

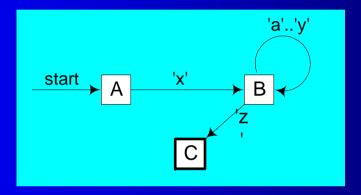
- A finite-state machine (FSM) is a graph structure that changes state based on some input. Also called a state-transition diagram.
- We use a graph to represent an FSM, with squares or circles called nodes, and lines with arrows between the circles called edges (or arcs).
- A FSM is a specific instance of a more general structure called a directed graph (or digraph).
- Three basic states, represented by nodes:
 - Start state
 - Terminal state(s)
 - Nonterminal state(s)

Finite-State Machine

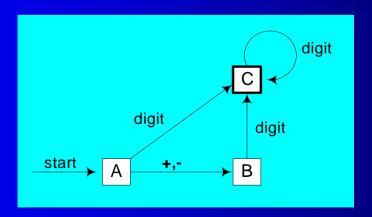
- Accepts any sequence of symbols that puts it into an accepting (final) state
- Can be used to recognize, or validate a sequence of characters that is governed by language rules (called a regular expression)
- Advantages:
 - Provides visual tracking of program's flow of control
 - Easy to modify
 - Easily implemented in assembly language

FSM Examples

 FSM that recognizes strings beginning with 'x', followed by letters 'a'..'y', ending with 'z':

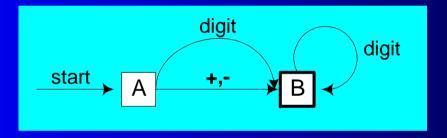


FSM that recognizes signed integers:



Your turn . . .

 Explain why the following FSM does not work as well for signed integers as the one shown on the previous slide:



Implementing an FSM

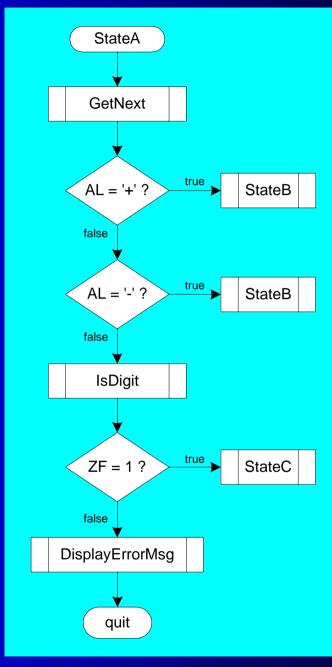
The following is code from State A in the Integer FSM:

```
StateA:
   call Getnext
                        ; read next char into AL
   cmp al, '+'
                        ; leading + sign?
   ie StateB
                        ; go to State B
   cmp al,'-'
                        ; leading - sign?
  ie StateB
                        ; go to State B
  call IsDigit ; ZF = 1 if AL = digit
   iz StateC
                  ; go to State C
   call DisplayErrorMsg ; invalid input found
   jmp Quit
```

View the Finite.asm source code.

Flowchart of State A

State A accepts a plus or minus sign, or a decimal digit.



Runtime Expressions

- .IF, .ELSE, .ELSEIF, and .ENDIF can be used to create block-structured IF statements.
- Examples:

```
.IF eax > ebx
    mov edx,1
.ELSE
    mov edx,2
.ENDIF
```

```
.IF eax > ebx && eax > ecx
    mov edx,1
.ELSE
    mov edx,2
.ENDIF
```

 MASM generates "hidden" code for you, consisting of code labels, CMP and conditional jump instructions.

Relational and Logical Operators

Operator	Description
expr1 == expr2	Returns true when expression1 is equal to expr2.
expr1 != expr2	Returns true when expr1 is not equal to expr2.
expr1 > expr2	Returns true when expr1 is greater than expr2.
expr1 >= expr2	Returns true when <i>expr1</i> is greater than or equal to <i>expr2</i> .
expr1 < expr2	Returns true when expr1 is less than expr2.
expr1 <= expr2	Returns true when <i>expr1</i> is less than or equal to <i>expr2</i> .
! expr	Returns true when expr is false.
expr1 && expr2	Performs logical AND between expr1 and expr2.
expr1 expr2	Performs logical OR between expr1 and expr2.
expr1 & expr2	Performs bitwise AND between expr1 and expr2.
CARRY?	Returns true if the Carry flag is set.
OVERFLOW?	Returns true if the Overflow flag is set.
PARITY?	Returns true if the Parity flag is set.
SIGN?	Returns true if the Sign flag is set.
ZERO?	Returns true if the Zero flag is set.

MASM-Generated Code

```
.data
val1   DWORD 5
result DWORD ?
.code
   mov eax,6
mov eax,6
.IF eax > val1
   jbe @C0001
mov result,1
.ENDIF
Generated code:

mov eax,6
cmp eax,val1
jbe @C0001
mov result,1
@C0001:
```

MASM automatically generates an unsigned jump (JBE).

MASM-Generated Code

```
.data
val1    SDWORD 5
result SDWORD ?
.code
    mov eax,6
mov eax,6
.IF eax > val1
    jle @C0001
mov result,1
.ENDIF
Generated code:

mov eax,6
cmp eax,val1
jle @C0001
mov result,1
@C0001:
```

MASM automatically generates a signed jump (JLE).

.REPEAT Directive

Executes the loop body before testing the loop condition associated with the .UNTIL directive.

Example:

```
; Display integers 1 - 10:

mov eax,0
.REPEAT
   inc eax
   call WriteDec
   call Crlf
.UNTIL eax == 10
```

.WHILE Directive

Tests the loop condition before executing the loop body The .ENDW directive marks the end of the loop.

Example:

```
; Display integers 1 - 10:

mov eax,0
.WHILE eax < 10
  inc eax
  call WriteDec
  call Crlf
.ENDW</pre>
```