

Introduction to Database Systems CSE 414

Lecture 12: NoSQL

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Class Overview

- Unit 1: Intro
- Unit 2: Relational Data Models and Query Languages
- Unit 3: Non-relational data
 - NoSQL
 - Json
 - SQL++
- Unit 4: RDBMS internals and query optimization
- Unit 5: Parallel query processing
- Unit 6: DBMS usability, conceptual design
- Unit 7: Transactions
- Unit 8: Advanced topics (time permitting)

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Two Classes of Database Applications

- OLTP (Online Transaction Processing)
 - Queries are simple lookups: 0 or 1 join
E.g., find customer by ID and their orders
 - Many updates. E.g., insert order, update payment
 - **Consistency** is critical: **transactions** (more later)
- OLAP (Online Analytical Processing)
 - aka "Decision Support"
 - Queries have many joins, and group-by's
E.g., sum revenues by store, product, clerk, date
 - No updates

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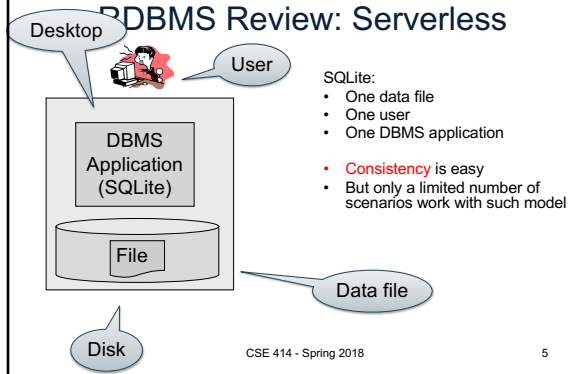
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NoSQL Motivation

- Originally motivated by Web 2.0 applications
 - E.g. Facebook, Amazon, Instagram, etc
 - Web startups need to scaleup from 10 to 100000 users very quickly
- Needed: very large scale OLTP workloads
- Give up on consistency
- Give up OLAP

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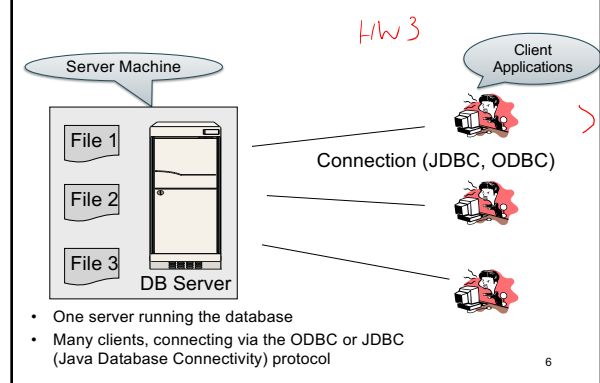
RDBMS Review: Serverless



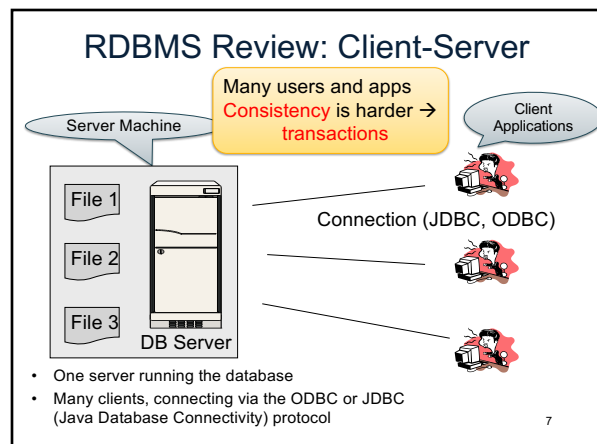
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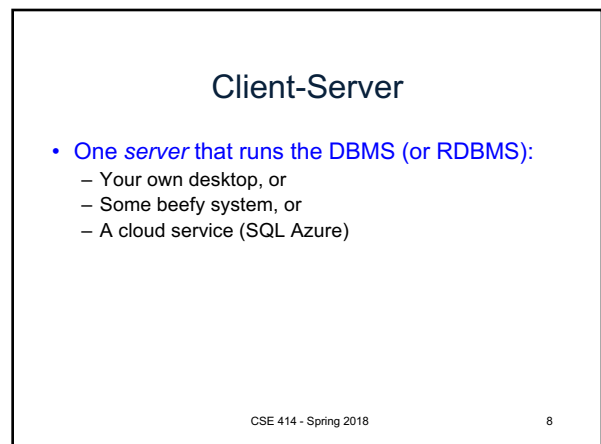
RDBMS Review: Client-Server



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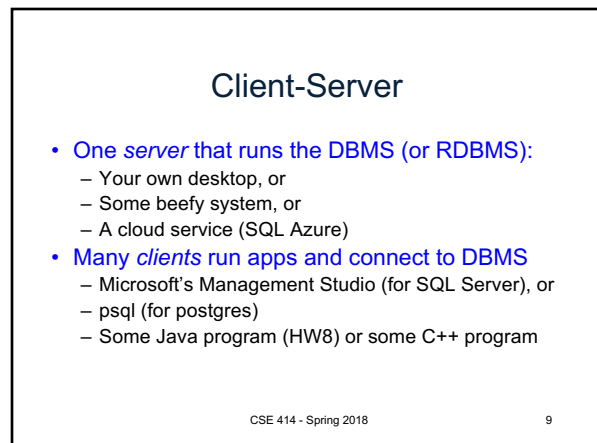


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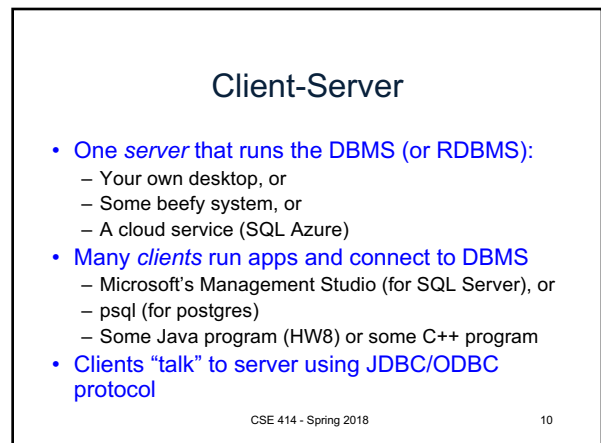
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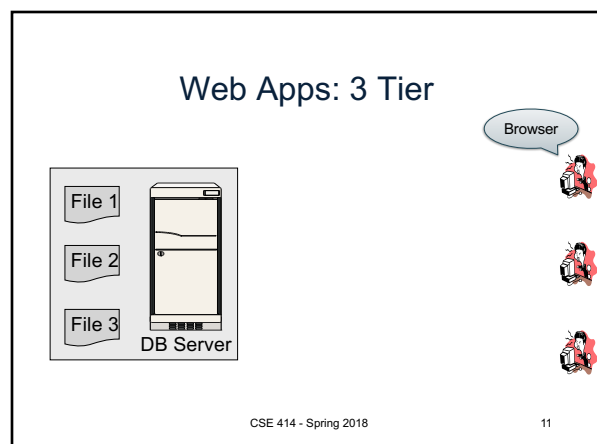
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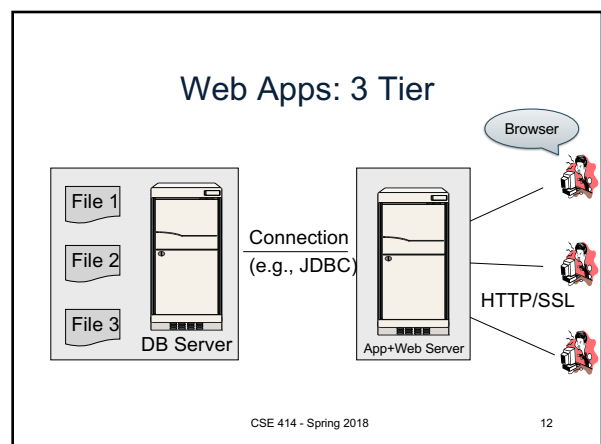
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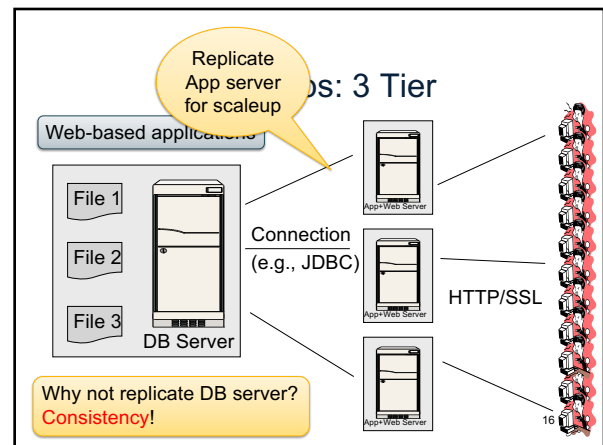
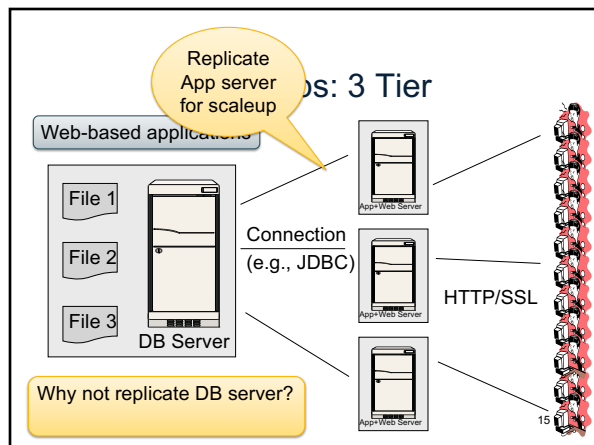
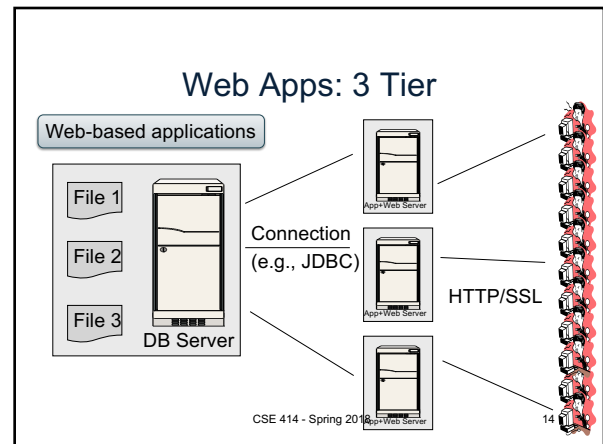
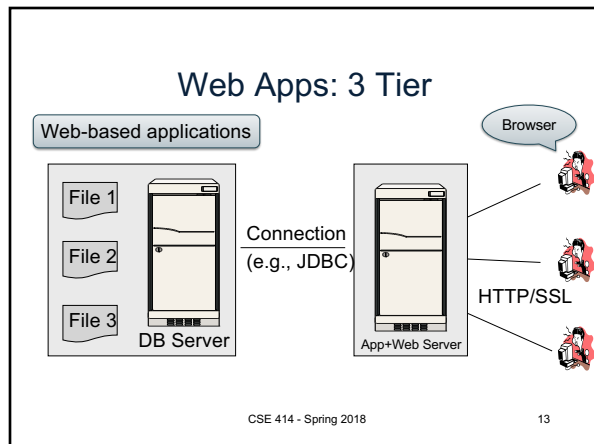
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Replicating the Database

- Two basic approaches:
 - Scale up through **partitioning**
 - Scale up through **replication**
- Consistency** is much harder to enforce

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Scale Through Partitioning

- Partition the database across many machines in a cluster
 - Database now fits in main memory
 - Queries spread across these machines
- Can increase throughput
- Easy for writes but reads become expensive!

Application updates here

Three partitions

May also update here

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Scale Through Replication

- Create multiple copies of each database partition
- Spread queries across these replicas
- Can increase throughput and lower latency
- Can also improve fault-tolerance
- Easy for reads but writes become expensive!



Relational Model → NoSQL

- Relational DB: difficult to replicate/partition
- Given `Supplier(sno,...), Part(pno,...), Supply(sno,pno)`
 - Partition: we may be forced to join across servers
 - Replication: local copy has inconsistent versions
 - Consistency is hard in both cases (why?)
- NoSQL: simplified data model
 - Given up on functionality
 - Application must now handle joins and consistency

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Data Models

Taxonomy based on data models:

- **Key-value stores**
 - e.g., Project Voldemort, Memcached
- **Document stores**
 - e.g., SimpleDB, CouchDB, MongoDB

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Key-Value Stores Features

- **Data model:** (key,value) pairs
 - Key = string/integer, unique for the entire data
 - Value = can be anything (very complex object)

Key-Value Stores Features

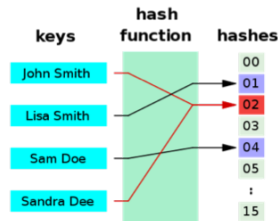
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 - Key = string/integer, unique for the entire data
 - Value = can be anything (very complex object)
- **Operations**
 - `get(key)`, `put(key,value)`
 - Operations on value not supported

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- **Distribution / Partitioning**

Aside: Hash Functions

- A function that maps any data to a “hash value” (e.g., an integer)



Aside: Hash Functions

- Example: data and hash value are integers
- Simple hash function:
 - $h(\text{key}) = \text{key} \% 42;$
 - $h(10) = 10$
 - $h(2) = 2$
 - $h(50) = 8$
- What does this have to do with data distribution?

Key-Value Stores Features

- Data model:** (key,value) pairs
 - Key = string/integer, unique for the entire data
 - Value = can be anything (very complex object)
- Operations**
 - get(key), put(key,value)
 - Operations on value not supported
- Distribution / Partitioning** – w/ hash function
 - No replication: key k is stored at server $h(k)$
 - 3-way replication: key k stored at $h1(k), h2(k), h3(k)$

How does get(k) work? How does put(k,v) work?

Flights(fid, date, carrier, flight_num, origin, dest, ...)
Carriers(cid, name)

Example

- How would you represent the Flights data as key, value pairs?

How does query processing work?

Flights(fid, date, carrier, flight_num, origin, dest, ...)
Carriers(cid, name)

Example

- How would you represent the Flights data as key, value pairs?
- Option 1: key=fid, value=entire flight record

How does query processing work?

Flights(fid, date, carrier, flight_num, origin, dest, ...)
Carriers(cid, name)

Example

- How would you represent the Flights data as key, value pairs?
- Option 1: key=fid, value=entire flight record
- Option 2: key=date, value=all flights that day

How does query processing work?

Flights(fid, date, carrier, flight_num, origin, dest, ...)
Carriers(cid, name)

Example

- How would you represent the Flights data as key, value pairs?
- Option 1: key=fid, value=entire flight record
- Option 2: key=date, value=all flights that day
- Option 3: key=(origin,dest), value=all flights between

How does query processing work?

Data Models

Taxonomy based on data models:

- **Key-value stores**
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- **Document stores**
 - e.g., SimpleDB, CouchDB, MongoDB

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Motivation

- In Key, Value stores, the Value is often a very complex object
 - Key = '2010/7/1', Value = [all flights that date]
- Better: allow DBMS to understand the *value*
 - Represent *value* as a JSON (or XML...) document
 - [all flights on that date] = a JSON file
 - May search for all flights on a given date

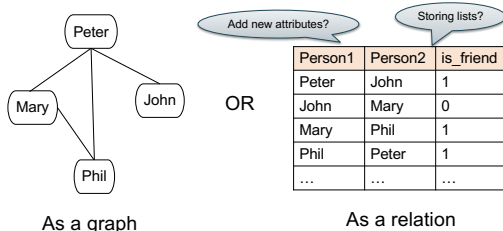
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Document Stores Features

- **Data model:** (key,document) pairs
 - Key = string/integer, unique for the entire data
 - Document = JSON, or XML
- **Operations**
 - Get/put document by key
 - Query language over JSON
- **Distribution / Partitioning**
 - Entire documents, as for key/value pairs

We will discuss JSON

Example: storing FB friends



We will learn the tradeoffs of different data models later this quarter

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JSON

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JSON - Overview

- JavaScript Object Notation = lightweight text-based open standard designed for human-readable data interchange. Interfaces in C, C++, Java, Python, Perl, etc.
- The filename extension is .json.

We will emphasize JSON as semi-structured data

JSON Syntax

```
{ "book": [
  { "id": "01",
    "language": "Java",
    "author": "H. Javerson",
    "year": 2015
  },
  { "id": "07",
    "language": "C++",
    "edition": "second",
    "author": "E. Sepp",
    "price": 22.25
  }
]
```

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JSON vs Relational

- Relational data model
 - Rigid flat structure (tables)
 - Schema must be fixed in advance
 - Binary representation: good for performance, bad for exchange
 - Query language based on Relational Calculus
- Semistructured data model / JSON
 - Flexible, nested structure (trees)
 - Does not require predefined schema ("self describing")
 - Text representation: good for exchange, bad for performance
 - Most common use: Language API; query languages emerging

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JSON Terminology

- Data is represented in name/value pairs.
- Curly braces hold objects
 - Each object is a list of name/value pairs separated by , (comma)
 - Each pair is a name followed by ':' (colon) followed by the value
- Square brackets hold arrays and values are separated by , (comma).

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JSON Data Structures

- Collections of name-value pairs:
 - { "name1": value1, "name2": value2, ... }
 - The "name" is also called a "key"
- Ordered lists of values:
 - [obj1, obj2, obj3, ...]


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Avoid Using Duplicate Keys

The standard allows them, but many implementations don't

```
{ "id": "07",
  "title": "Databases",
  "author": "Garcia-Molina",
  "author": "Ullman",
  "author": "Widom"
}
```



```
{ "id": "07",
  "title": "Databases",
  "author": [ "Garcia-Molina",
              "Ullman",
              "Widom" ]
}
```

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JSON Datatypes

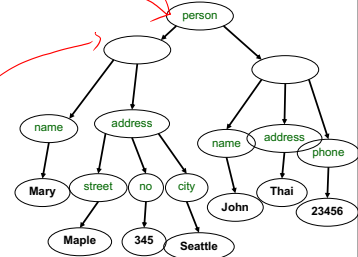
- Number
- String = double-quoted
- Boolean = true or false
- nullempty

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JSON Semantics: a Tree !

```
{ "person":
  [ { "name": "Mary",
      "address":
        { "street": "Maple",
          "no": 345,
          "city": "Seattle" },
        { "name": "John",
          "address":
            { "street": "Thailand",
              "no": 2345678 } }
    ]
}
```



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JSON Data

- JSON is **self-describing**
- Schema elements become part of the data
 - Relational schema: **person(name, phone)**
 - In JSON "person", "name", "phone" are part of the data, and are repeated many times
- Consequence: JSON is much more flexible
- JSON = **semistructured** data

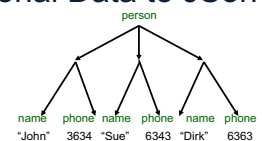
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Mapping Relational Data to JSON

Person

name	phone
John	3634
Sue	6343
Dirk	6363



```
{ "person":
  [ { "name": "John", "phone": 3634 },
    { "name": "Sue", "phone": 6343 },
    { "name": "Dirk", "phone": 6363 }
  ]
}
```

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Mapping Relational Data to JSON

May inline foreign keys

Person

name	phone
John	3634
Sue	6343

Orders

personName	date	product
John	2002	Gizmo
John	2004	Gadget
Sue	2002	Gadget

```
{ "Person":
  [ { "name": "John",
      "phone": 3634,
      "Orders": [ { "date": 2002,
                    "product": "Gizmo" },
                  { "date": 2004,
                    "product": "Gadget" }
                ]
    },
    { "name": "Sue",
      "phone": 6343,
      "Orders": [ { "date": 2002,
                    "product": "Gadget" }
                ]
    }
  ]
}
```