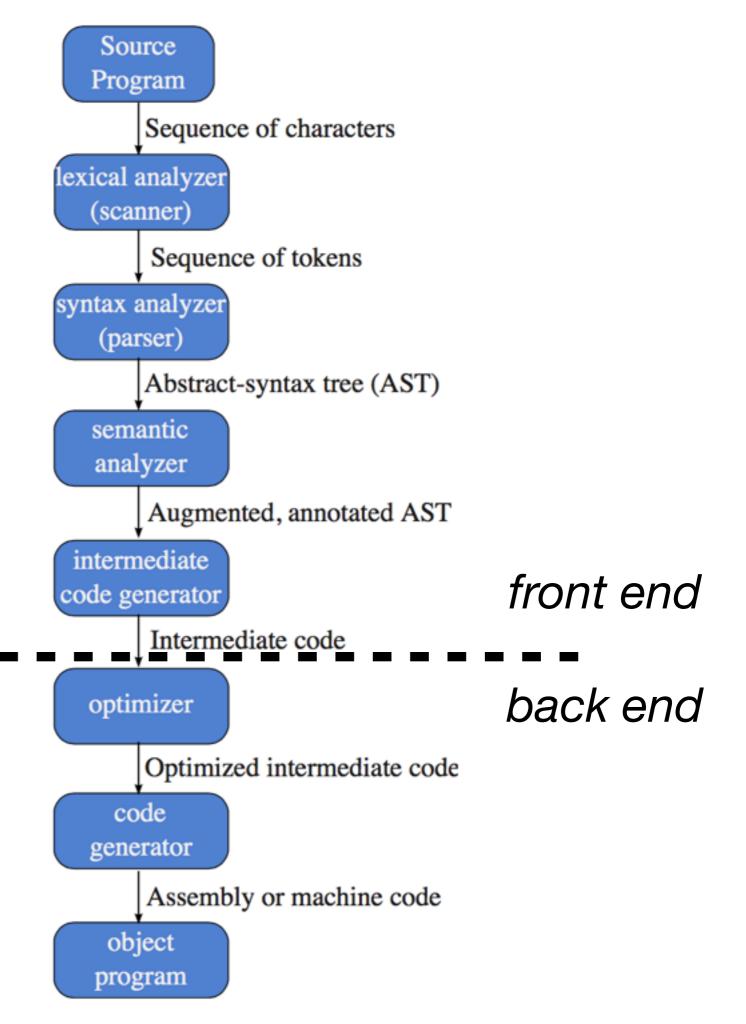
Finite-state machines

CS 536

Last time



The scanner

Translates sequence of chars into sequence of tokens

Each time scanner is called it should:

find longest sequence of chars corresponding to a token

return that token

Scanner generator

Generates a scanner!!!

Needs one regular expression for each token

Needs regular expressions for things to ignore comments, whitespace, etc.

To understand how it works, we need FSMs finite state machines

FSMs: Finite State Machines

Aka finite automata

Input: string (seq of chars)

Output: accept / reject

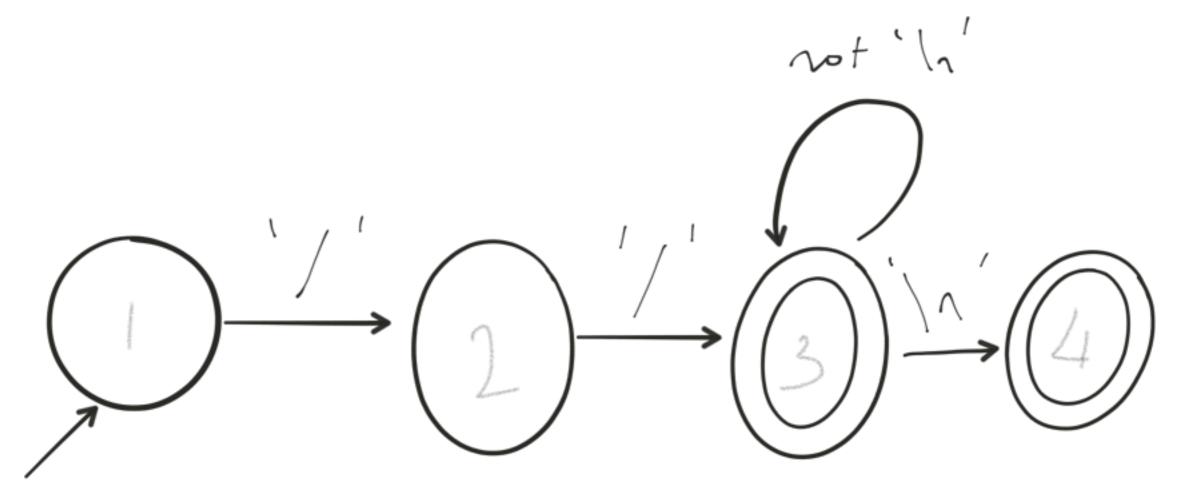
i.e., input is legal in language

FSMs

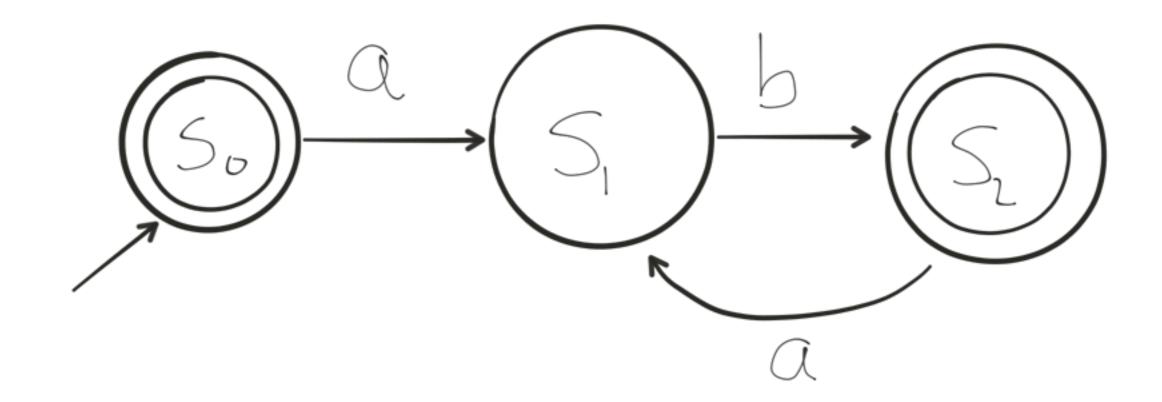
Represent regular languages
Good enough for tokens in PLs

Example 1

single line comments with //



Example 2



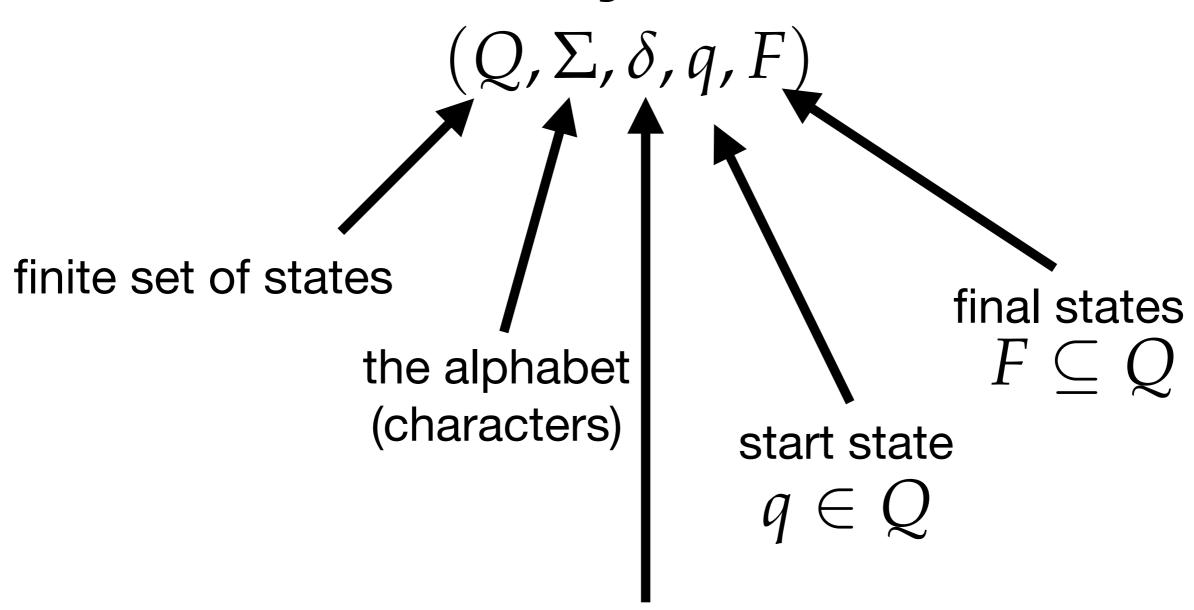
What language does this accept?

Can you find an equivalent, but smaller, FSM?

How an FSM works

```
curr_state = start_state
let in_ch = current input char
repeat
 if there is edge out of curr_state with
 label in_ch into next state
   cur_state = next_state
   in_ch = next char of input
 o/w stuck // error condition
until stuck or input string is consumed
string is accepted iff entire string is
consumed and cur_state = final_state
```

FSMs, formally



transition relation

$$\delta: Q \times \Sigma \to Q$$

FSMs, formally

$$(Q, \Sigma, \delta, q, F)$$

FSM accepts string

$$\chi_1 \chi_2 \chi_3 \dots \chi_n$$

$$\iff$$

$$\delta(\dots \delta(\delta(\delta(q, \chi_1), \chi_2), \chi_3) \dots, \chi_n) \in F$$

The language of FSM M is the set of all words it accepts, denoted L(M)

FSM example, formally

$$(Q, \Sigma, \delta, q, F)$$

$$Q=\{s_0,s_1\}$$
 a b c $\Sigma=\{a,b,c\}$ so s_1 $g=s_0$ so s_1

$$\delta = s_0, a \rightarrow s_1$$

 $s_1, b \rightarrow s_0$

anything else, machine is stuck

Coding an FSM

```
curr_state = start_state
done = false
while (!done)
 ch = nextChar()
 next = transition[curr_state][ch]
 if (next == error || ch == E0F)
   done = true
 else
   curr state = next
return curr_state == final_state
```

FSM types: DFA & NFA

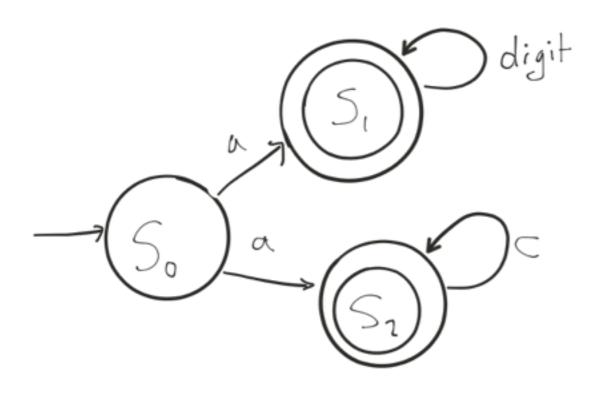
Deterministic

no state has > 1 outgoing edge with same label

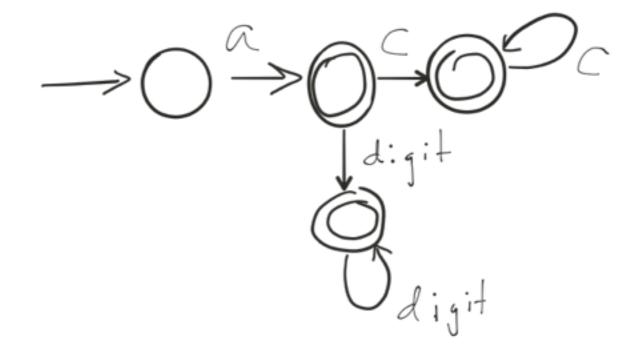
Nondeterministic

states may have multiple outgoing edges with same label edges may be labelled with special symbol ϵ (empty string) ϵ -transitions can happen without reading input

NFA example

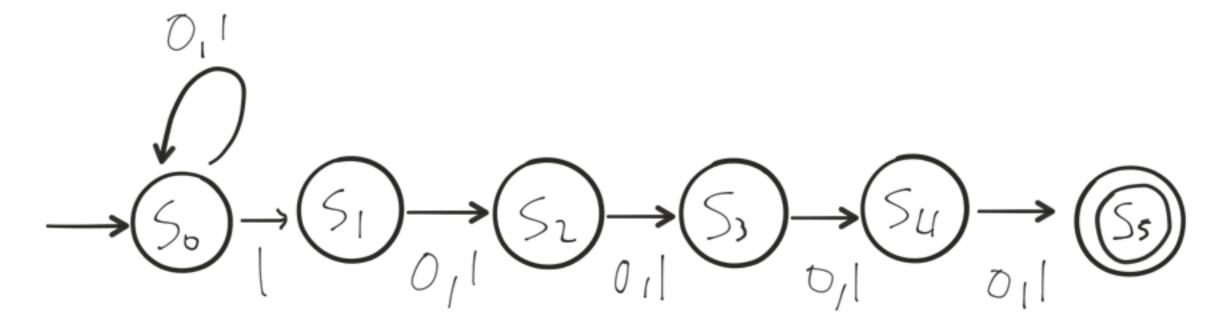


Equivalent DFA



Why NFA?

Much more compact



What does this accept?

An equivalent DFA needs 2^5 states

Extra example

Hex literals

must start with 0x or 0X followed by at least one hex digit (0-9,a-f,A-F) can optionally have long specifier (I,L) at the end

Extra example

A C/C++ identifier is a sequence of one or more letters, digits, or underscores. It cannot start with a digit.

What if you wanted to add the restriction that it can't end with an underscore?

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Recap

The scanner reads stream of characters and finds tokens

Tokens are defined using regular expressions, which are finite-state machines

Finite-state machines can be non-deterministic

Next time: understand connection between deterministic and non-deterministic FSMs