

Grammer-2

COM S 319

Objectives

1. Learn formal and informal definitions of grammar.
2. Learn about types of grammar (Chomsky Hier..)
3. Learn about regular grammar
4. Learn about recognizer for regular grammar
5. Learn about LEX

REGULAR GRAMMER

Regular Grammar

- Production Rules have to have one of the forms

1. $A \rightarrow a$

2. $A \rightarrow aB$

3. $A \rightarrow \varepsilon$

where A and B stand for arbitrary variables and a stands for an arbitrary terminal. Epsilon is the empty string.

Note: There is an equivalent form, where middle rule is $A \rightarrow Ba$

Example1

- This was shown before. This is an example of a regular grammar
- $V = \{S, A, B, C\}$, $T = \{a, b, c\}$
- $S \rightarrow A$
- $A \rightarrow aA$
- $A \rightarrow B$
- $B \rightarrow bC$
- $C \rightarrow cC$
- $C \rightarrow \varepsilon$

Q1: What are strings in the language?

We have already learnt about regular expressions.

- Regular expressions express strings in regular language
- Regular grammar also expresses strings in regular language.
- Finite automaton is used to recognize regular expressions
- RE, RG, FA are equivalent!

REG EXP IMPLEMENTATION

Non-deterministic FA

- A NFA is a 5-tuple (Q, A, T, S, F)
 - Q is a FINITE set of states
 - A is the alphabet
 - T is the transition function
 - $Q \times A + \epsilon \rightarrow P(Q)$ (i.e. state & alphabet gives state)
 - S is the start state
 - F is set of final states
- NFA (non-deterministic finite automaton) can
 1. transition to multiple states on the same input and
 2. can also transition on epsilon

easier to express in NFA vs DFA (note DFA and NFA are equivalent)

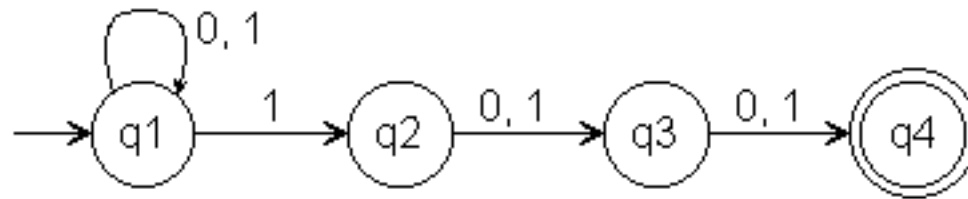
Accept a^*bc^*

- S_0 is start state
- state S_0 $\langle \text{epsilon}, S_1 \rangle$,
- state S_1 $\langle a, S_1 \rangle, \langle b, S_2 \rangle$,
- state S_2 $\langle c, S_2 \rangle$
- S_2 is final state

Transition Table

- can use transition table to represent finite automaton.
 - rows represent states
 - cols represent input
 - element represents next state
- can use simple traversal over input and keeping set of future states. when an input string has been fully processed – is any of the next states a final state? If so accept – else reject.

Example Transition table



	0	1
q1	{q1}	{q1, q2}
q2	{q3}	{q3}
q3	{q4}	{q4}
q4	\emptyset	\emptyset

LEX (LEXER OR LEXICAL ANALYSER)

lex

- describe rules for language.
- lex automatically creates lexical analyzer.

Format of lex rules file:

```
{definitions}
```

```
%%
```

```
{rules}
```

```
%%
```

```
{user subroutines}
```

Example Lex file

```
%{  
#include <stdio.h>  
%}  
  
%%  
[a-zA-Z][a-zA-Z0-9]*    printf("WORD ");  
[a-zA-Z0-9\\/.-]+      printf("FILENAME ");  
\"                printf("QUOTE ");  
\{                printf("OBRACE ");  
\}                printf("EBRACE ");  
;                printf("SEMICOLON ");  
\\n                printf("\\n");  
[ \\t]+           /* ignore whitespace */;  
%%
```

ANTLR

<http://www.antlr.org>

ANTLR (ANother Tool for Language Recognition) is a powerful parser generator for reading, processing, executing, or translating structured text or binary files. It's widely used to build languages, tools, and frameworks. From a grammar, ANTLR generates a parser that can build and walk parse trees.

ANTLR

- lex and yacc – standard unix utilities to build lexer and parser (to build compilers).
 - c code
 - lexer and parser rules kept in separate files
- ANTLR
 - completely Java code.
 - Both lexer and parser rules are specified in one file.
- Commands
 - antlr4 Expr.g4
 - javac Expr*.java
 - grun Expr prog -gui
 - 100+2*34 ^D

SUMMARY

- Basics definitions (terminal, alphabet, rules,..)
- Chomsky Hierarchy (or Types of grammars)
- Regular Grammar
- Equivalence of RE, RG, FA
- Implementation NFA using transition table
- Implementation using lex & ANTLR