

## 10.5 DESIGNING AN EFFICIENT COMPUTER DRUM

The position of a rotating drum is to be recognised by means of binary signals produced at a number of electrical contacts at the surface of the drum. The surface is divided into  $2^n$  sections, each consisting of either insulating or conducting material. An insulated section gives signal 0 (no current), whereas a conducting section gives signal 1 (current). For example, the position of the drum in figure 10.9 gives a reading 0010 at the four

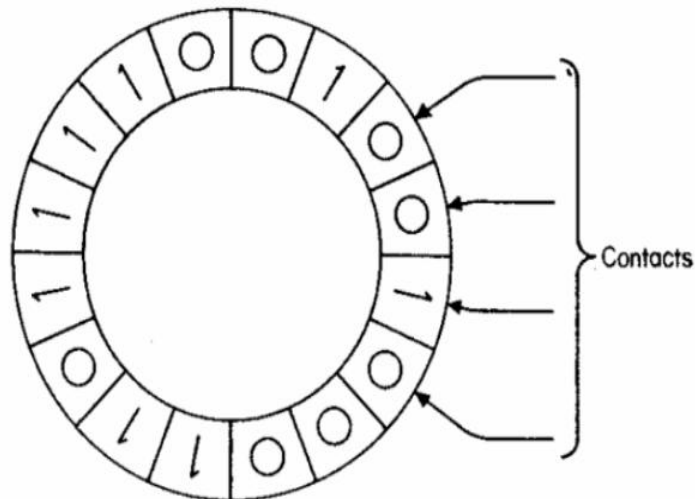


Figure 10.9. A computer drum

contacts. If the drum were rotated clockwise one section, the reading would be 1001. Thus these two positions can be distinguished, since they give different readings. However, a further rotation of two sections would result in another position with reading 0010, and therefore this latter position is indistinguishable from the initial one.

We wish to design the drum surface in such a way that the  $2^n$  different positions of the drum can be distinguished by  $k$  contacts placed consecutively around part of the drum, and we would like this number  $k$  to be as small as possible. How can this be accomplished?

First note that  $k$  contacts yield a  $k$ -digit binary number, and there are  $2^k$  such numbers. Therefore, if all  $2^n$  positions are to give different readings, we must have  $2^k \geq 2^n$ , that is,  $k \geq n$ . We shall show that the surface of the drum can be designed in such a way that  $n$  contacts suffice to distinguish all  $2^n$  positions.

We define a digraph  $D_n$  as follows: the vertices of  $D_n$  are the  $(n-1)$ -digit binary numbers  $p_1p_2 \dots p_{n-1}$  with  $p_i = 0$  or  $1$ . There is an arc with tail  $p_1p_2 \dots p_{n-1}$  and head  $q_1q_2 \dots q_{n-1}$  if and only if  $p_{i+1} = q_i$  for  $1 \leq i \leq n-2$ ; in other words, all arcs are of the form  $(p_1p_2 \dots p_{n-1}, p_2p_3 \dots p_n)$ . In addition,

each arc  $(p_1 p_2 \dots p_{n-1}, p_2 p_3 \dots p_n)$  of  $D_n$  is assigned the label  $p_1 p_2 \dots p_n$ .  $D_4$  is shown in figure 10.10.

Clearly,  $D_n$  is connected and each vertex of  $D_n$  has indegree two and outdegree two. Therefore (exercise 10.3.2)  $D_n$  has a directed Euler tour. This directed Euler tour, regarded as a sequence of arcs of  $D_n$ , yields a binary sequence of length  $2^n$  suitable for the design of the drum surface.

For example, the digraph  $D_4$  of figure 10.10 has a directed Euler tour  $(a_1, a_2, \dots, a_{16})$ , giving the 16-digit binary sequence 0000111100101101. (Just read off the first digits of the labels of the  $a_i$ .) A drum constructed from this sequence is shown in figure 10.11.

This application of directed Euler tours is due to Good (1946).

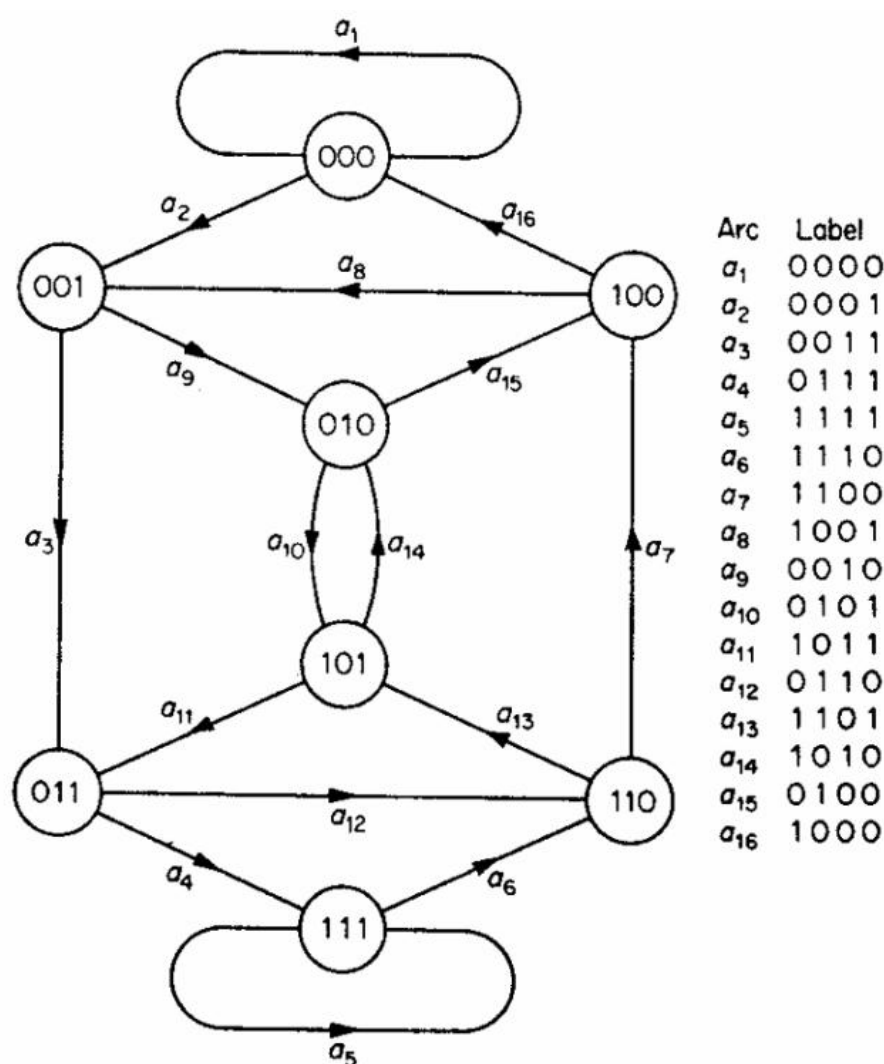


Figure 10.10

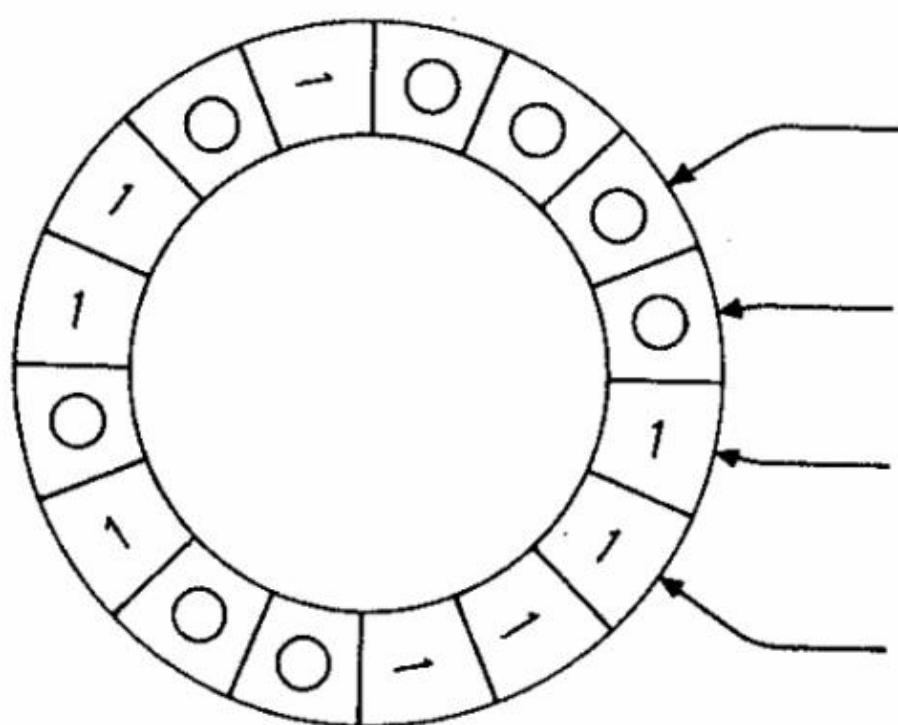


Figure 10.11