# COMP9334 Capacity Planning for Computer Systems and Networks

Week 2A: Operational Analysis (2).

**Workload Characterisation** 

#### Last lecture

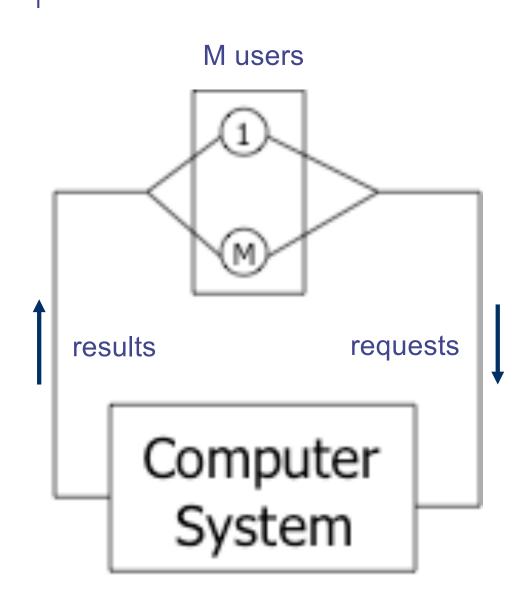
- Modelling a computer system as a queueing network
- Operational analysis on queueing networks
- We have derived these operational laws
  - Utilisation law U(j) = X(j) S(j)
  - Forced flow law X(j) = V(j) X(0)
  - Service demand law D(j) = V(j) S(j) = U(j) / X(0)
  - Little's law N = X R

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#### This lecture

- Operational analysis (Continued)
  - Using operational law for
    - Performance analysis
    - Bottleneck analysis
- Workload characterisation
  - Poisson process and its properties

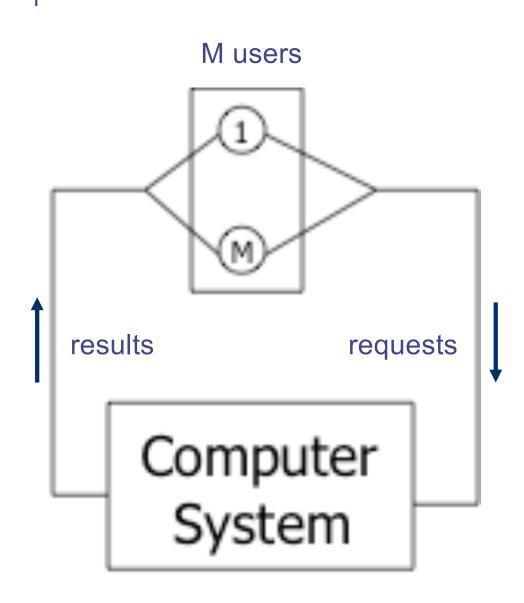
## Interactive systems



- An interactive system is used to model the interaction between humans (users) and computers
- The system consists of
  - A number of users
  - A computer system

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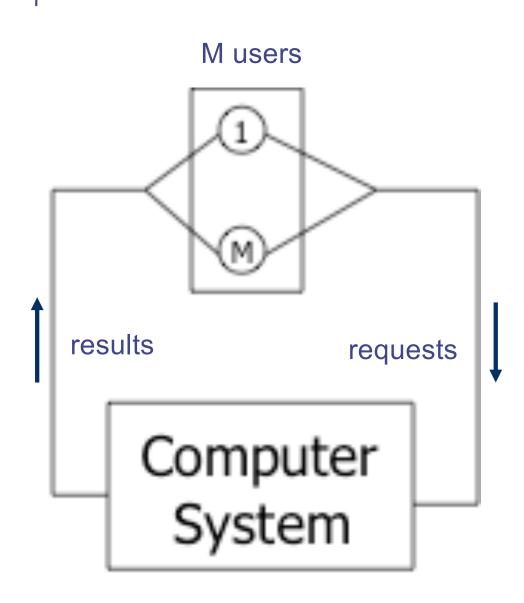
# Interactive systems (Cont'd)



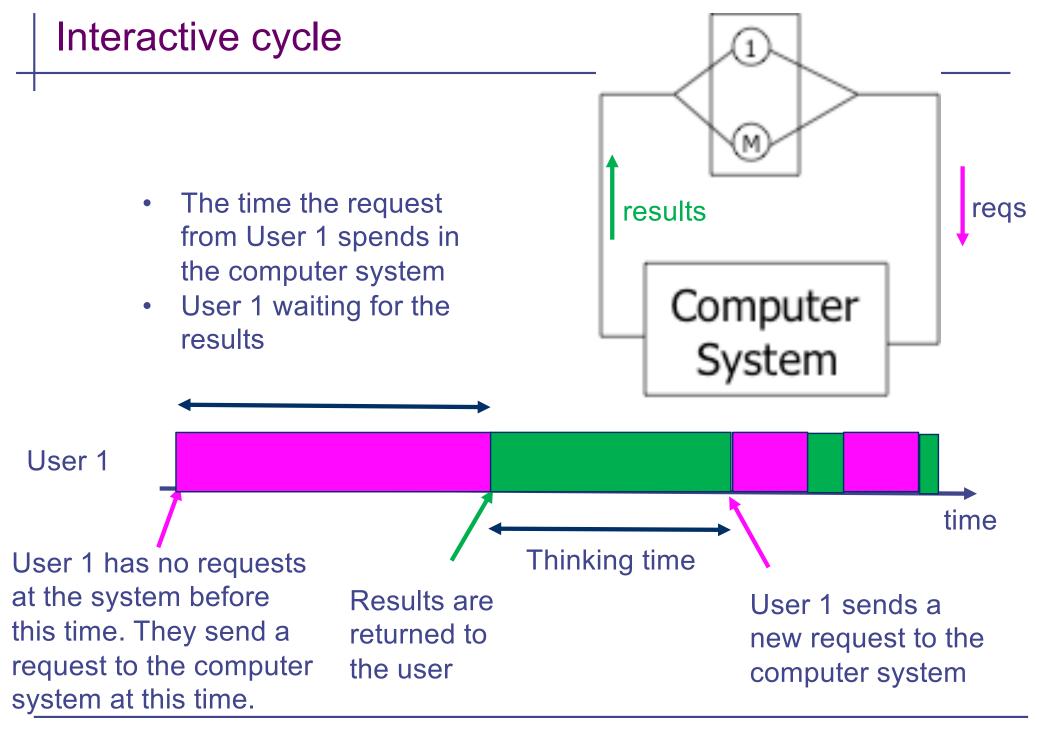
#### Interactions

- Users send requests to the computer system
- After finish processing a request, the computer system returns the result to the user
- A user, after inspecting the results from the computer system, will send another request to the system

# Interactive systems: Modelling assumptions

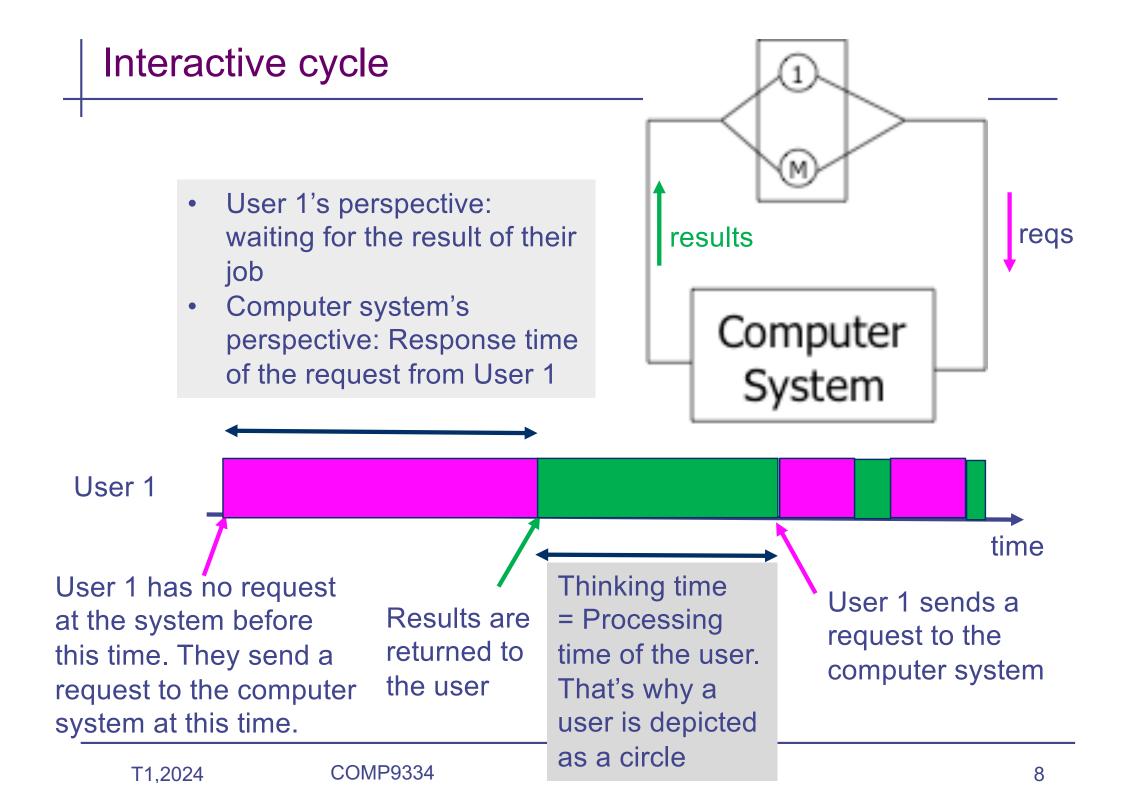


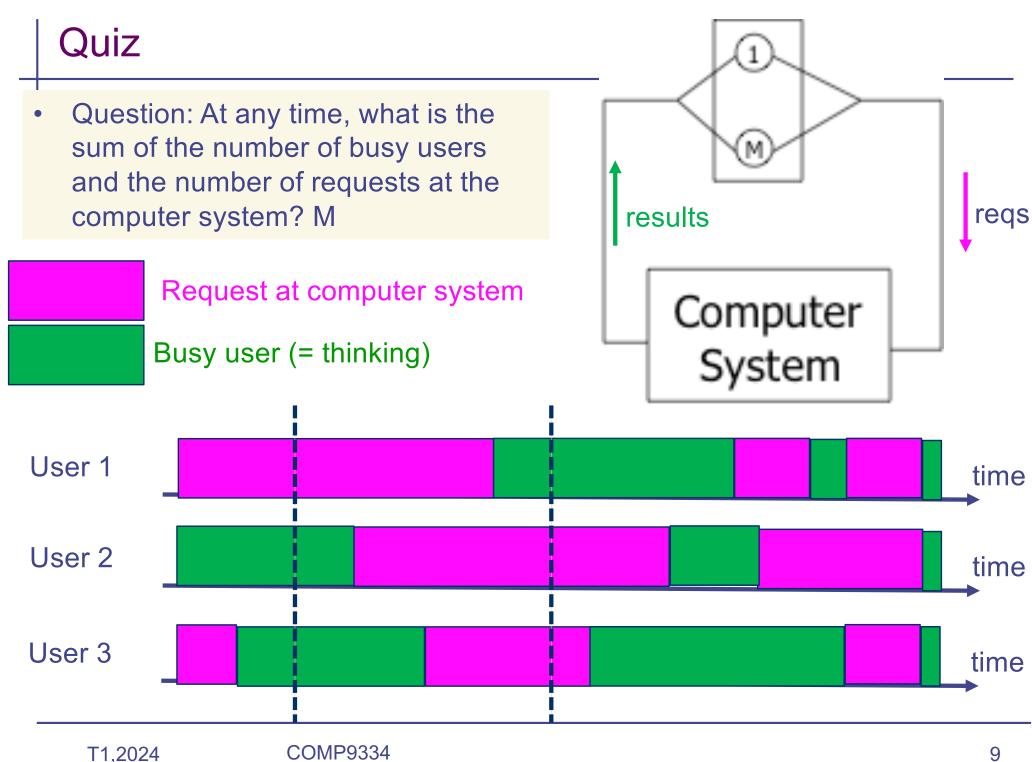
- Analyze interactive systems with specific assumptions
  - Fixed number of users denoted by M
  - Each user can have at most 1 request at the computer system
  - Each user goes through a cycle consisting of
    - Thinking time
    - Waiting for result time



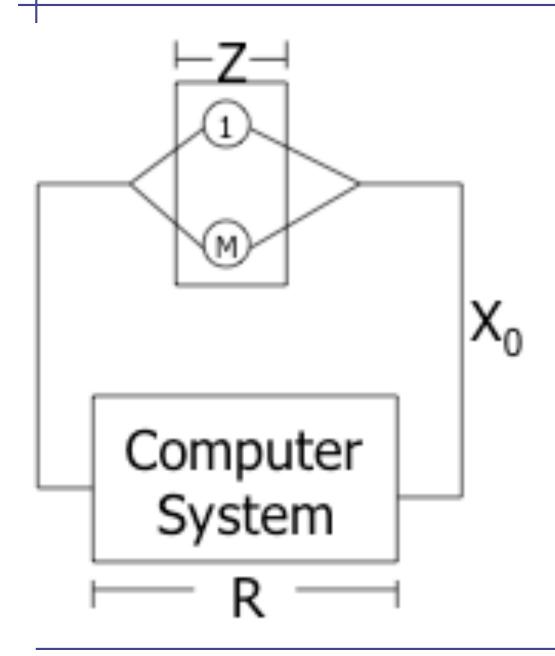
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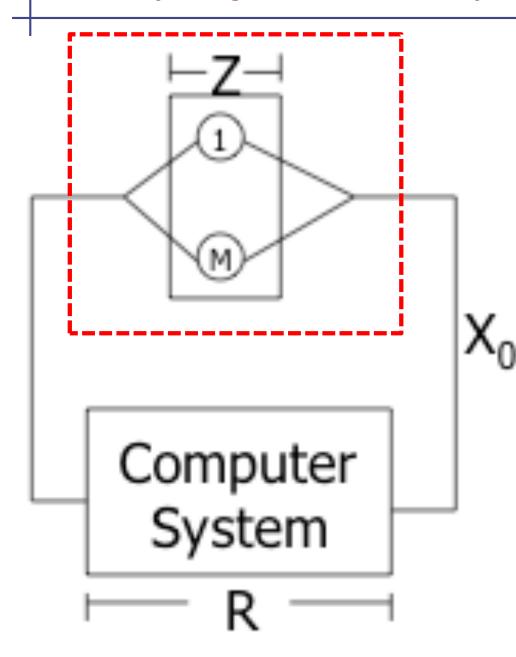


# Interactive system: Parameters



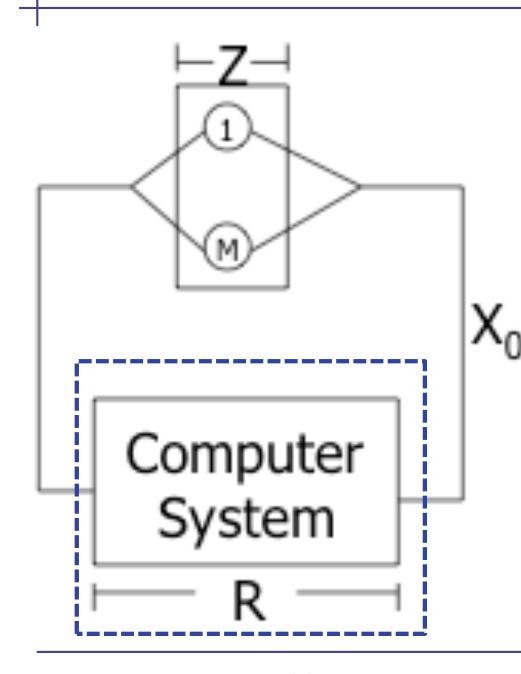
- M interactive users
- Z = mean thinking time
- R = mean response time of the computer system
- X0 = throughput

## Analyzing interactive system: Quiz 1



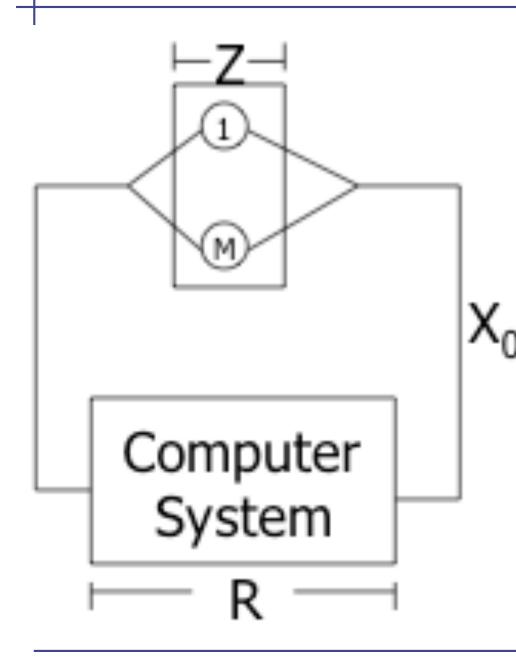
- Mavg = mean # busy users
- Z = mean thinking time
- X0 = throughput
- Apply Little's Law to the red box. What do you get?
  - Mavg = Z \* X0

# Analyzing interactive system: Quiz 2



- Navg = average # requests in the computer system
- R = mean response time at the computer system
- X0 = throughput
- Apply Little's Law to the computer system (i.e. the blue box), what do you get?
  - Navg = R \* X0

## Analyzing interactive system: Quiz 3

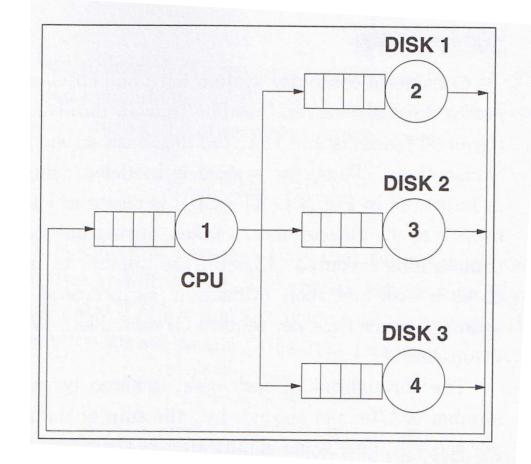


- Quiz 1: Mavg = X0 \* Z
- Quiz 2: Navg = X0 \* R
- What is Mavg + Navg?
  - M = Mavg + Navg
- Interactive response time law
  - M = X0 \* (Z+R)

## The operational laws

- These are the operational laws
  - Utilisation law U(j) = X(j) S(j)
  - Forced flow law X(j) = V(j) X(0)
  - Service demand law D(j) = V(j) S(j) = U(j) / X(0)
  - Little's law N = X R
  - Interactive response time M = X(0) (R+Z)
- Applications
  - Mean value analysis (later in the course)
  - Bottleneck analysis
  - Modification analysis

# Bottleneck analysis - motivation

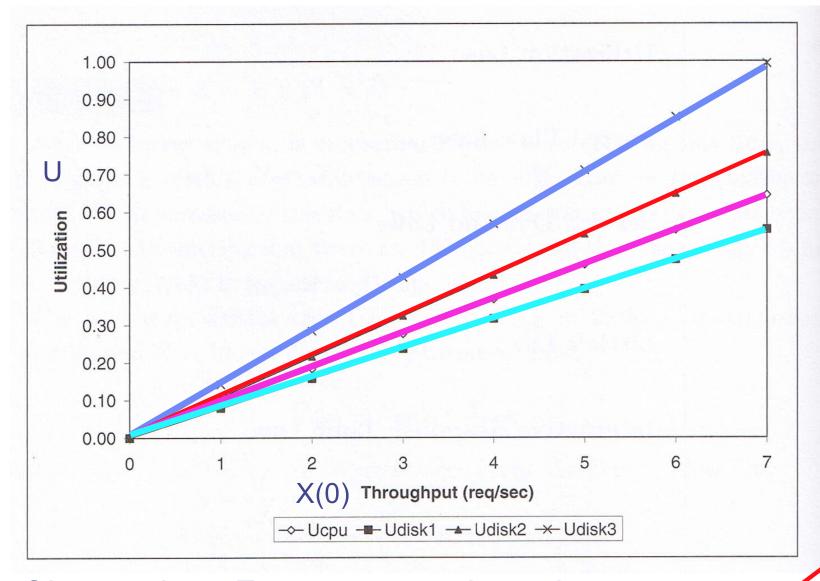


	D(j)	Utilisation
Disk 1	79ms	0.30
Disk 2	108ms	0.41
Disk 3	142ms	0.54
CPU	92ms	0.35

Service demand law: D(j) = U(j) / X(0)==> U(j) = D(j) X(0)

Utilisation increases with increasing throughput and service demand

#### Utilisation vs. throughput plot U(j) = D(j) X(0)



Disk 3

Disk 2 CPU

Disk 1

What determines this order?

Observation: For all system throughput: Utilisation of Disk 3 > Utilisation of Disk 2 > Utilisation of CPU compatibles at ion of Disk 1

## Bottleneck analysis

- Recall that utilisation is the busy time of a device divided by measurement time
  - What is the maximum value of utilisation?
- Based on the example on the previous slide, which device will reach the maximum utilisation first?

# Bottleneck (1)

- Disk 3 has the highest service demand
- It is the bottleneck of the whole system

Operational law: 
$$X(0) = \frac{U(j)}{D(j)}$$
 Utilisation limit: 
$$U(j) \leq 1$$
 
$$X(0) \leq \frac{1}{D(j)}$$

# Bottleneck (2)

$$X(0) \leq \frac{1}{D(j)}$$
 Should hold for all K devices in the system

$$i.e.X(0) \le \frac{1}{D(1)}, ..., X(0) \le \frac{1}{D(K)}$$

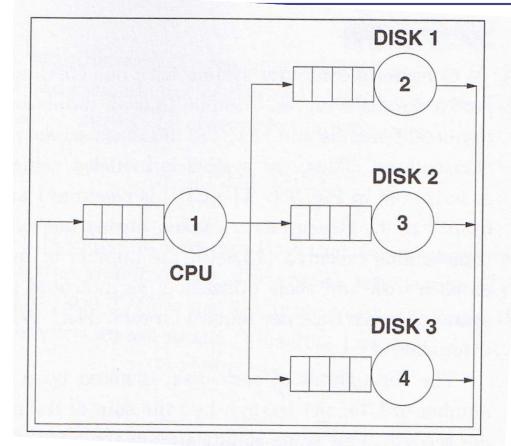
$$\Rightarrow X(0) \le \min \frac{1}{D(j)}$$

$$\Rightarrow X(0) \le \frac{1}{\max D(j)}$$

Bottleneck throughput is limited by the maximum service demand

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#### Bottleneck exercise



	D(j)	Utilisation
Disk 1	79ms	0.30
Disk 2	108ms	0.41
Disk 3	142ms	0.54
CPU	92ms	0.35

The system throughput is upper bounded by  $\frac{1}{0.142}$  = 7.04 jobs/s If we upgrade Disk 3 by a new disk which is 2 times faster, which device will be the bottleneck after the upgrade? You can assume that service time is inversely proportional to disk speed.

#### Another throughput bound

Little's law

$$N = R \times X(0) \ge \left(\sum_{i=1}^{K} D_i\right) \times X(0)$$

$$\Rightarrow X(0) \le \frac{N}{\sum_{i=1}^{K} D_i}$$

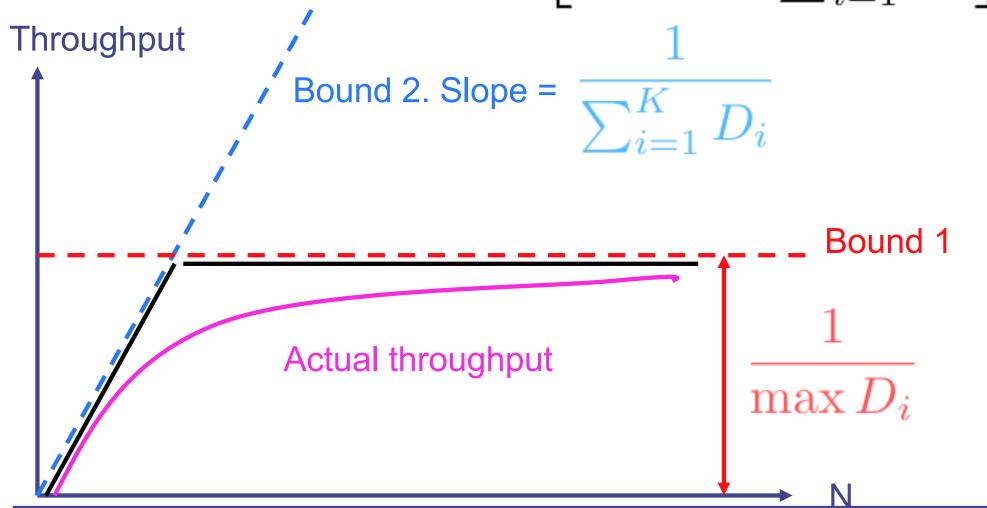
Previously, we have

Therefore:

$$X(0) \le \min \left[ \frac{1}{\max D_i}, \frac{N}{\sum_{i=1}^K D_i} \right]$$

# Throughput bounds

$$X(0) \le \min \left[ \frac{1}{\max D_i}, \frac{N}{\sum_{i=1}^K D_i} \right]$$



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# Bottleneck analysis

- Simple to use
  - Needs only utilisation of various components
- Assumes service demand is load independent

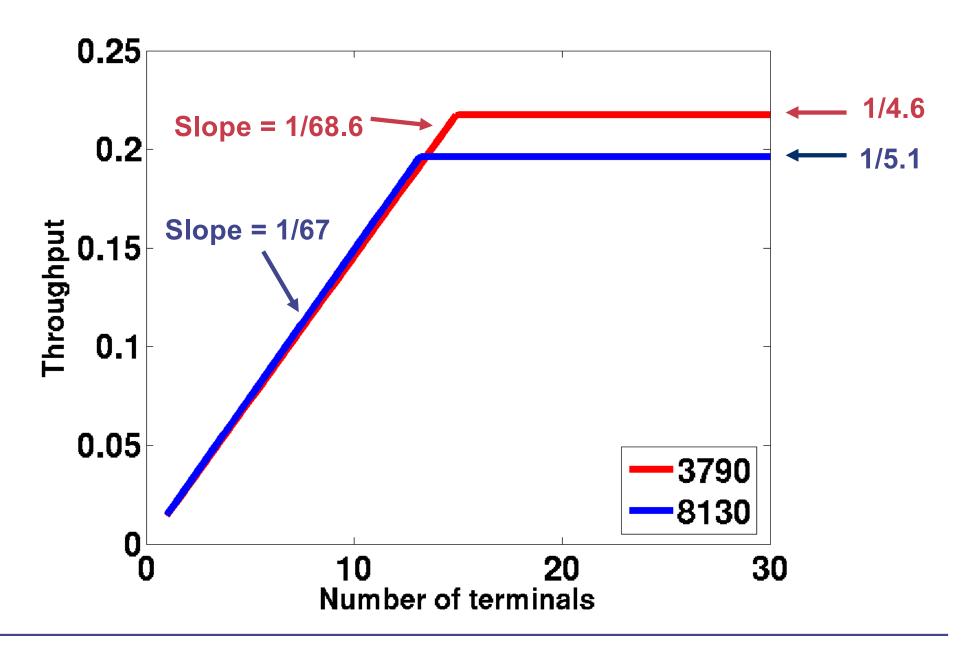
# Modification analysis (1)

- (Reference: Lazowska Section 5.3.1)
- A company currently has a system (3790) and is considering switching to a new system (8130). The service demands for these two systems are given below:

	Service demand (seconds)	
System	CPU	Disk
3790	4.6	4.0
8130	5.1	1.9

- The company uses the system for interactive application with a think time of 60s.
- Given the same workload, should the company switch to the new system?
- Exercise: Answer this question by using bottleneck analysis. For each system, plot the upper bound of throughput as a function of the number of interactive users.

# Modification analysis (2)



# Operational analysis

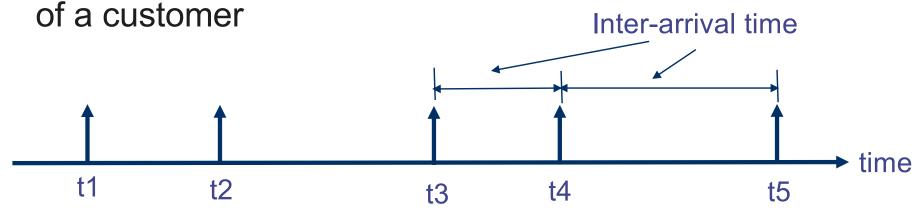
- Operational analysis allows you to bound the system performance but it does NOT allow you to find the throughput and response time of a system
- To order to find the throughput and response time, we need to use queueing analysis
- To order to use queueing analysis, we need to specify the workload in a probabilistic manner

## Workload analysis

- Performance depends on workload
  - When we look at the performance bound earlier, the bounds depend on number of users and service demand
  - Queue response time depends on the job arrival probability distribution and job service time distribution
    - Recall from Lecture 1A:
      - Uniform arrival times and uniform processing times result in zero waiting time
      - But non-uniform distributions give non-zero waiting time
- Need to specify workload by using probability distribution.
- We will look at a well-known arrival process called Poisson process today.

#### Arrival process

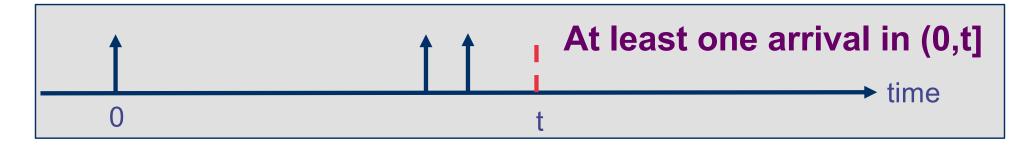
Each vertical arrow in the time line below depicts the arrival

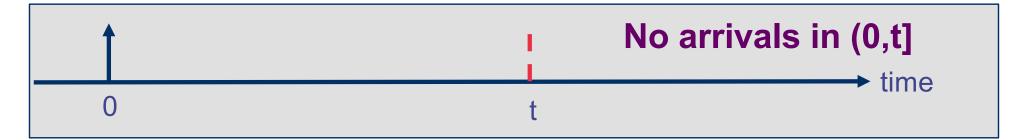


- An arrival can mean
  - A telephone call arriving at a call centre
  - A transaction arriving at a computer system
  - A customer arriving at a checkout counter
  - An HTTP request arriving at a web server
- The inter-arrival time distribution will impact on the response time.
- We will study an inter-arrival time distribution that results from a large number of independent customers.

# Describing arrivals probabilistically

Assume a customer arrives at time 0

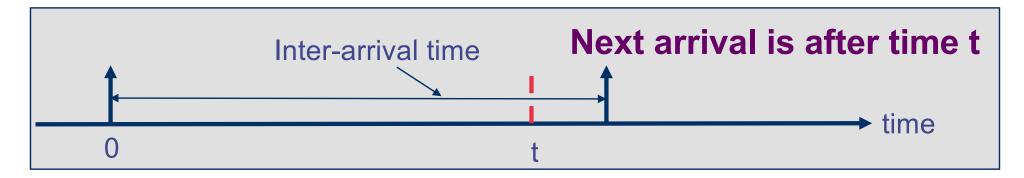


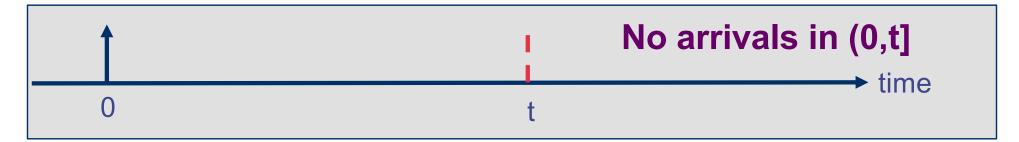


- Quiz: What is the relation between the following two probabilities?
  - Prob[at least one arrival in (0,t]]
  - Prob[no arrivals in (0,t]]
- Answer: They add up to 1
- Moral: "No arrivals" is not boring, it tells you something

# Inter-arrival time probability

Assume a customer arrives at time 0





- Quiz: What is the relation between the following two probabilities?
  - Prob[Inter-arrival time is > t]
  - Prob[no arrivals in (0,t]]
- Answer: Equal
- Next step: Find Prob[no arrivals in (0,t]] for independent customers

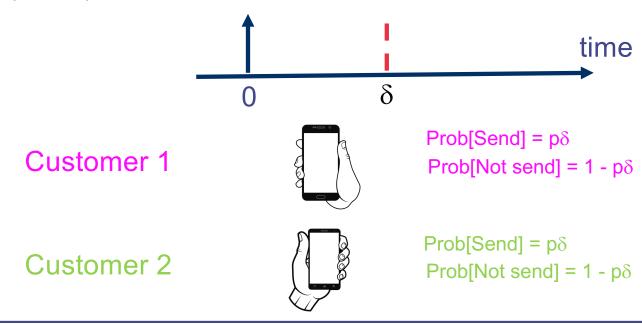
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# Many independent arrivals (1)

- Problem set up:
  - An arrival at time 0
  - A large pool of N independent customers

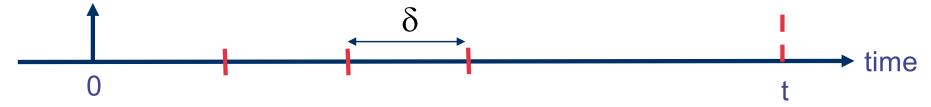
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- Behaviour of each customer: Within a small time interval of  $\delta$ , a customer sends a request (or arrives) with a probability of p $\delta$ 
  - p is a constant
- Quiz: If there are 2 (= N) customers, what is the probability that both of them do not send any request in the time interval  $\delta$ 
  - Answer:  $(1 p \delta)^2$



# Many independent arrivals (2)

- Aim: Want to find the probability of no arrivals in (0,t]
- Divide the time t into intervals of width δ



- No arrivals in (0,t] = no arrivals in each interval  $\delta$  from N users
- Probability of no arrivals in  $\delta = (1 p \delta)^N \approx 1 Np\delta$
- There are t / δ intervals
- Probability of no arrivals in (0,t] is

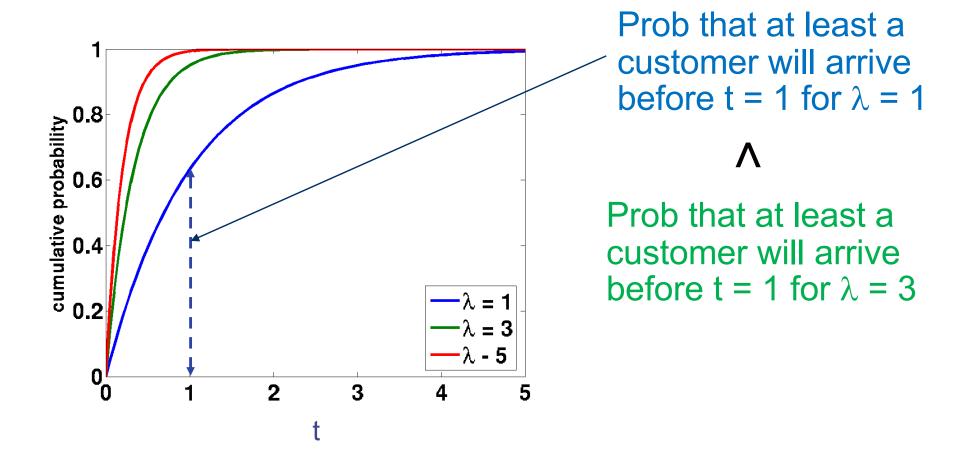
$$(1 - Np\delta)^{\frac{t}{\delta}} \rightarrow e^{-Npt} \text{ as } \delta \rightarrow 0$$

## Exponential inter-arrival time

- We have showed Probability(no arrivals in (0,t]) =  $\exp(-Npt)$
- Probability(inter-arrival time > t) =  $\exp(-Npt)$
- This means Probability(inter-arrival time  $\leq$  t) =  $1 \exp(-Npt)$
- What this shows is the inter-arrival time distribution for independent arrival is exponentially distributed
- Define:  $\lambda = Np$ 
  - λ is the mean arrival rate of customers

## Exponential distribution - cumulative distribution

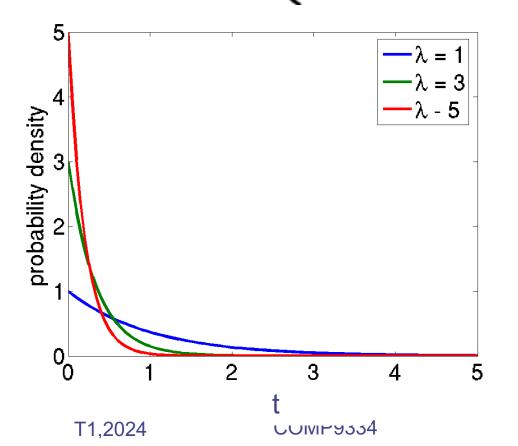
- Cumulative distribution of inter-arrival time with customer arrival rate  $\lambda$ 
  - Prob(inter-arrival time  $\leq$  t) = 1 exp(-  $\lambda$  t)

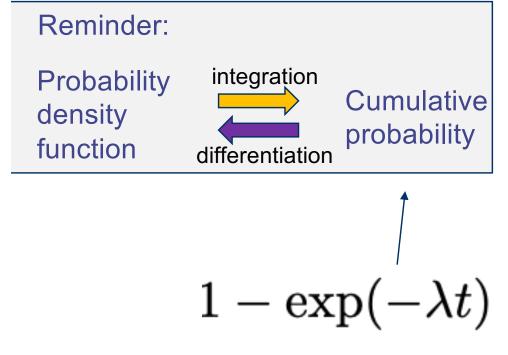


#### **Exponential distribution**

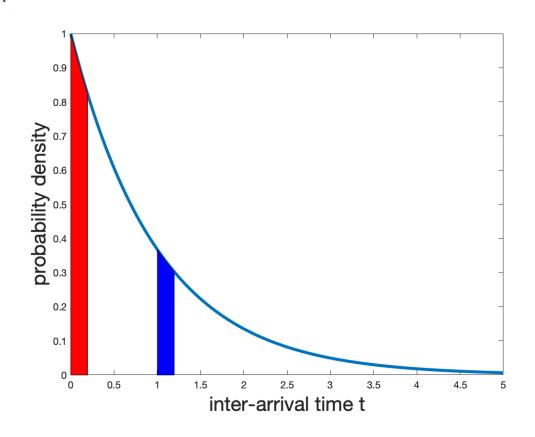
• A continuous random variable is exponentially distributed with rate  $\lambda$  if it has probability density function

$$f(t) = \begin{cases} \lambda \exp(-\lambda t) & \text{for } t \ge 0\\ 0 & \text{for } t < 0 \end{cases}$$





# Probability density function (PDF)



Reminder: PDF f(t)

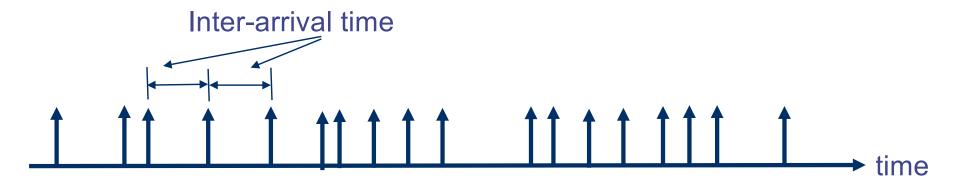
Probability( $t \le T \le t + \delta t$ )

 $= f(t) \delta t$ 

Red area = probability that inter-arrival time is in the interval [0,0.2] Blue area = probability that the inter-arrival time is in the interval [1,1.2]

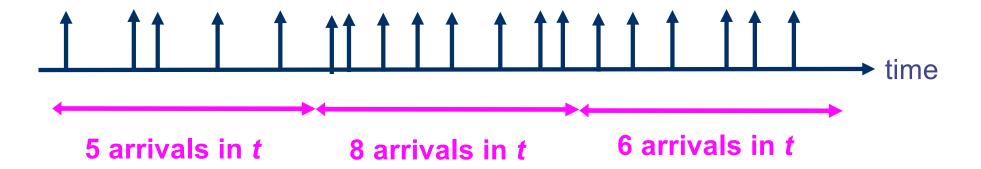
#### Two different methods to describe arrivals

Method 1: Continuous probability distribution of inter-arrival time



#### Two different methods to describe arrivals

Method 2: Use a fixed time interval (say *t*), and count the number of arrivals within *t*.

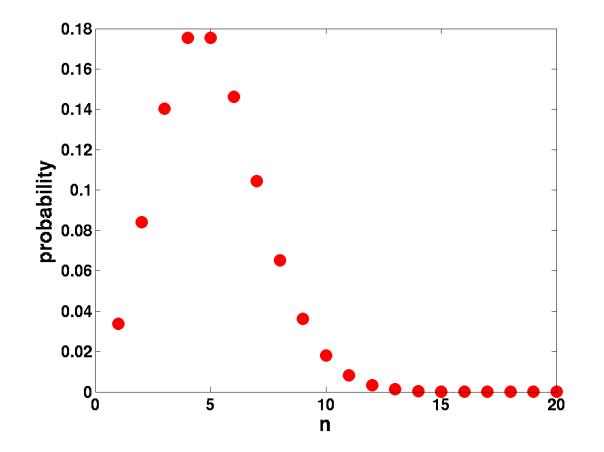


- The number of arrivals in t is random
- The number of arrivals must be a non-negative integer
- We need a discrete probability distribution:
  - Prob[#arrivals in t = 0]
  - Prob[#arrivals in t = 1]
  - etc.

# Poisson process (1)

• Definition: An arrival process is Poisson with parameter  $\lambda$  if the probability that n customers arrive in any time interval t is

 $\frac{(\lambda t)^n e^{-\lambda t}}{n!}$ 



#### Example:

Example:

 $\lambda$ = 5 and t = 1

Note: Poisson is a discrete probability distribution.

# Poisson process (2)

- Theorem: An exponential inter-arrival time distribution with parameter  $\lambda$  gives rise to a Poisson arrival process with parameter  $\lambda$
- How can you prove this theorem?
  - A possible method is to divide an interval t into small time intervals of width  $\delta$ . A finite  $\delta$  will give a binomial distribution and with  $\delta \rightarrow 0$ , we get a Poisson distribution.

# Customer arriving rate

• Given a Poisson process with parameter  $\lambda$ , we know that the probability of n customers arriving in a time interval of t is given by:

 $\frac{(\lambda t)^n e^{-\lambda t}}{n!}$ 

 What is the mean number of customers arriving in a time interval of t?

$$\sum_{n=0}^{\infty} n \frac{(\lambda t)^n e^{-\lambda t}}{n!} = \lambda t$$

• That's why  $\lambda$  is called the arrival rate.

#### Customer inter-arrival time

- You can also show that if the inter-arrival time distribution is exponential with parameter  $\lambda$ , then the mean inter-arrival time is  $1/\lambda$
- Quite nicely, we have
   Mean arrival rate = 1 / mean inter-arrival time

# Application of Poisson process

- Poisson process has been used to model the arrival of telephone calls to a telephone exchange successfully
- Queueing networks with Poisson arrival is tractable
  - We will see that in the next few weeks.
- Beware that not all arrival processes are Poisson! Many arrival processes we see in the Internet today are not Poisson. We will see that later.

#### References

- Operational analysis
  - Lazowska et al, Quantitative System Performance, Prentice Hall, 1984.
     (Classic text on performance analysis. Now out of print but can be downloaded from <a href="http://www.cs.washington.edu/homes/lazowska/qsp/">http://www.cs.washington.edu/homes/lazowska/qsp/</a>
    - Chapters 3 and 5 (For Chapter 5, up to Section 5.3 only)
  - Alternative 1: You can read Menasce et al, "Performance by design", Chapter 3. Note that Menasce doesn't cover certain aspects of performance bounds.
     So, you will also need to read Sections 5.1-5.3 of Lazowska.
  - Alternative 2: You can read Harcol-Balter, Chapters 6 and 7. The treatment is more rigorous. You can gross over the discussion on ergodicity.
- Poisson process: Harcol-Balter Chapter 11