

MPC PERSTER

KEEPING DATA PRIVATE, TOGETHER

TEAM NGO

Yuening Yang
Chen Chang Lew
Junzhen Lou

INTRODUCTION

MPC PEERSTER
TEAM NGO



MOTIVATION

- Data is becoming increasingly valuable Nowadays.
- The need to preserve data privacy while maintaining availability is becoming a critical issue.
- Traditional methods for data sharing are no longer sufficient for protecting sensitive information.
- Provide a platform to raise incentives for people to use mpc

PROJECT OVERVIEW

A p2p community with access control, where participants can start an MPC, sell data for MPC to use, participate in MPC to earn fees, etc.



Junzhen Lou
MPC Protocols
Finite Field Arithmetic
Benchmark



Chen Chang Lew
MPC Progress
Benchmark
UI



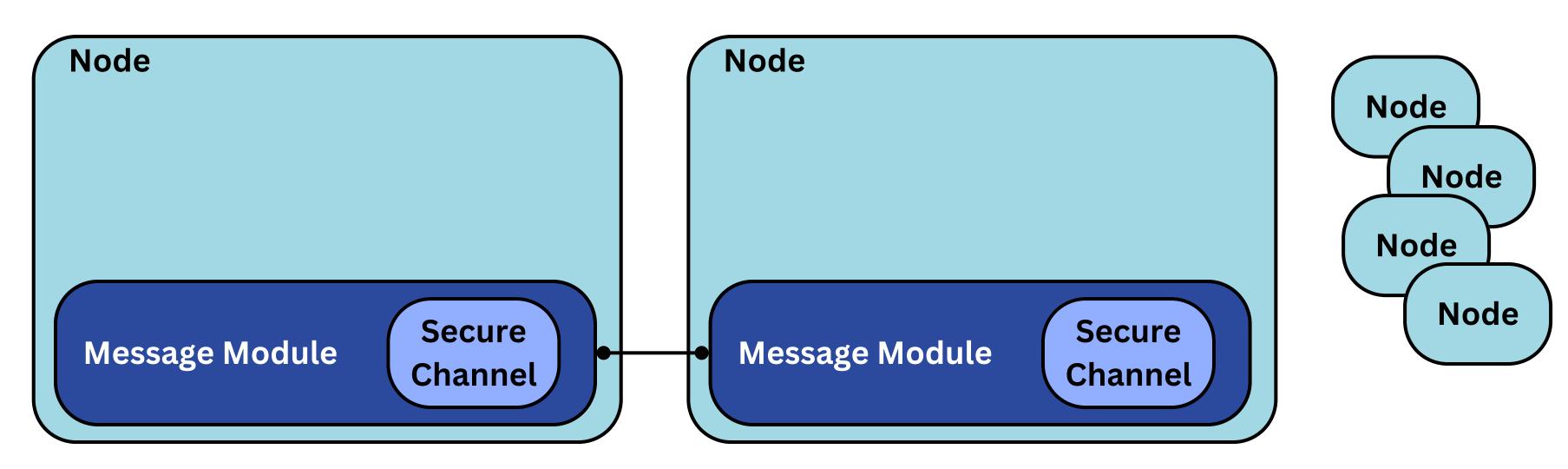
Yuening Yang
Secure channel
Blockchain
CLI

SYSTEM OVERVIEW

MPC PEERSTER

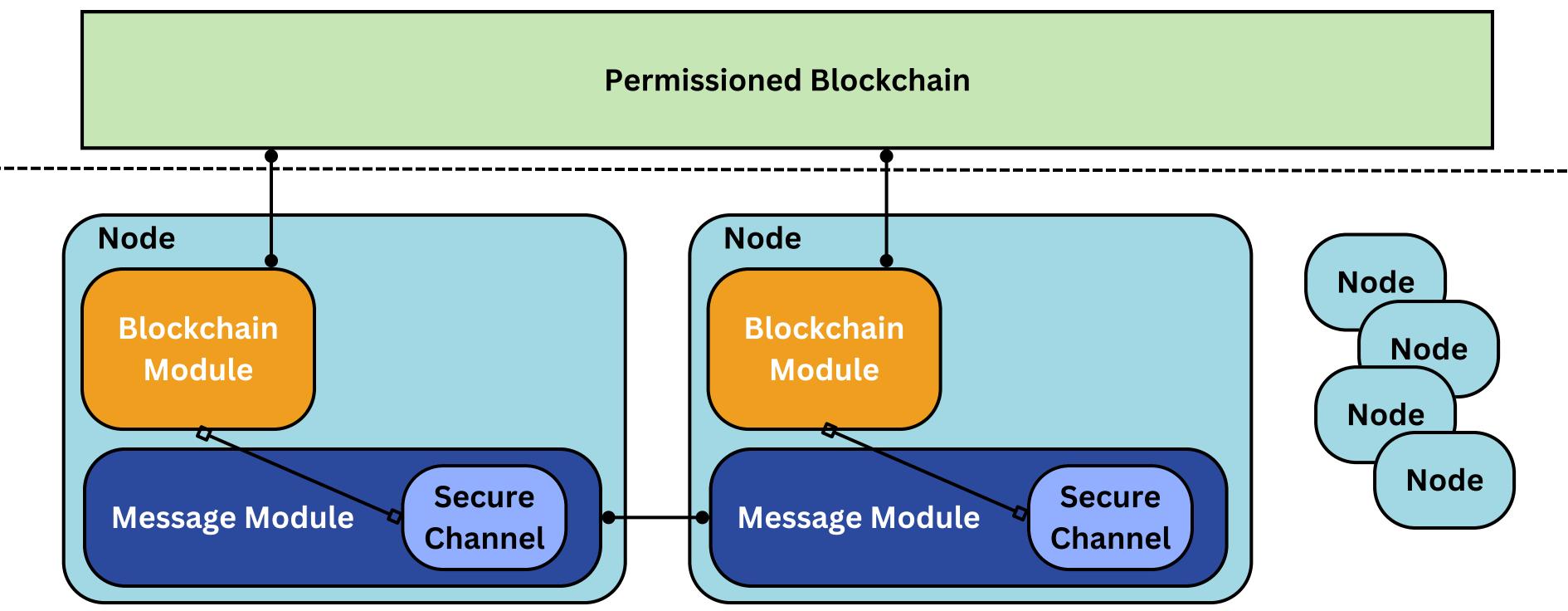


SYSTEM ARCHITECTURE



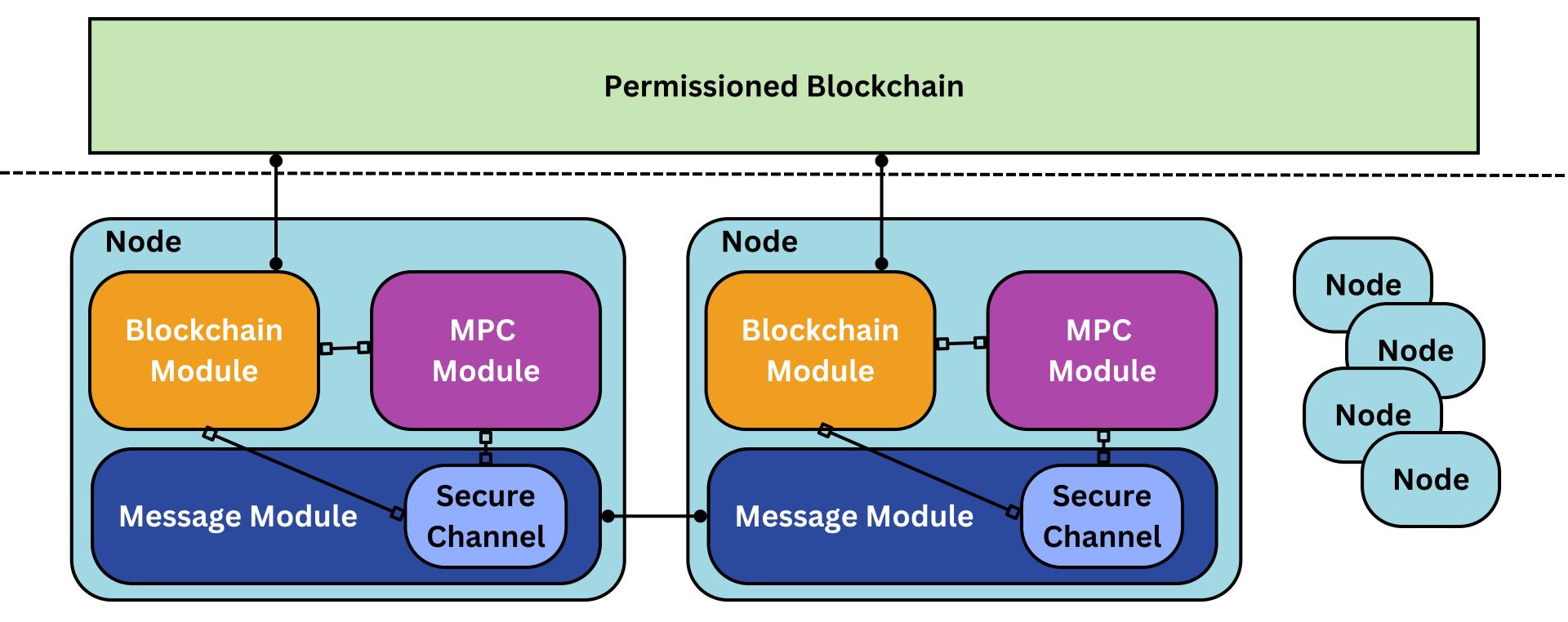
Secure channel supports end-to-end encryption and group private broadcast

SYSTEM ARCHITECTURE



• Blockchain Module is built on top of secure channel and form a permissioned blockchain

SYSTEM ARCHITECTURE



• MPC Module is built on top of secure channel

Overview

- Privacy-preserving distributed computation
- Secure and efficient data/information exchange
- allows parties to perform secure computation while keeping their own input private

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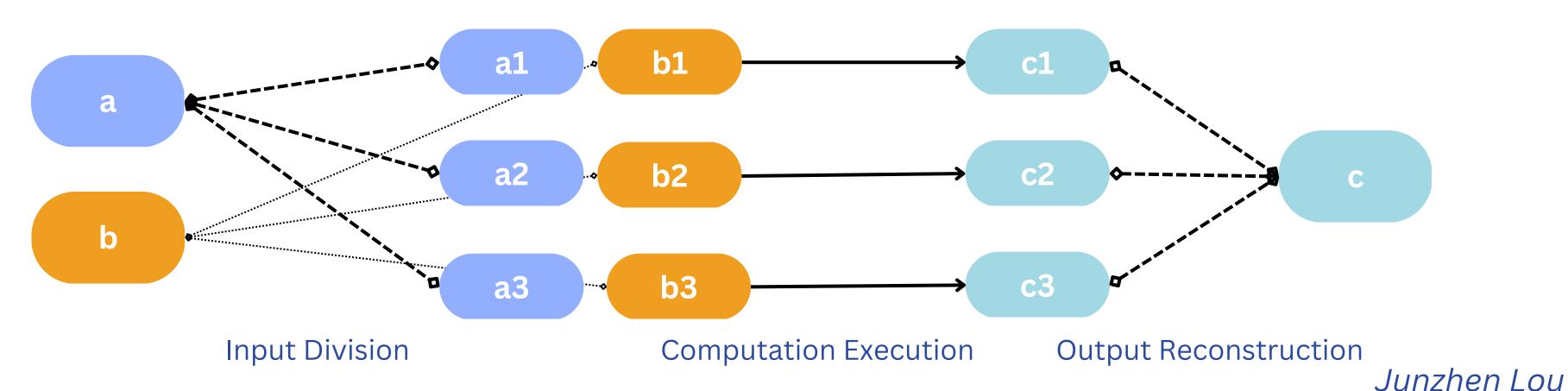
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Model

- Complete network
- Secure channels with broadcast
- Synchronous communication
- Passive adversary

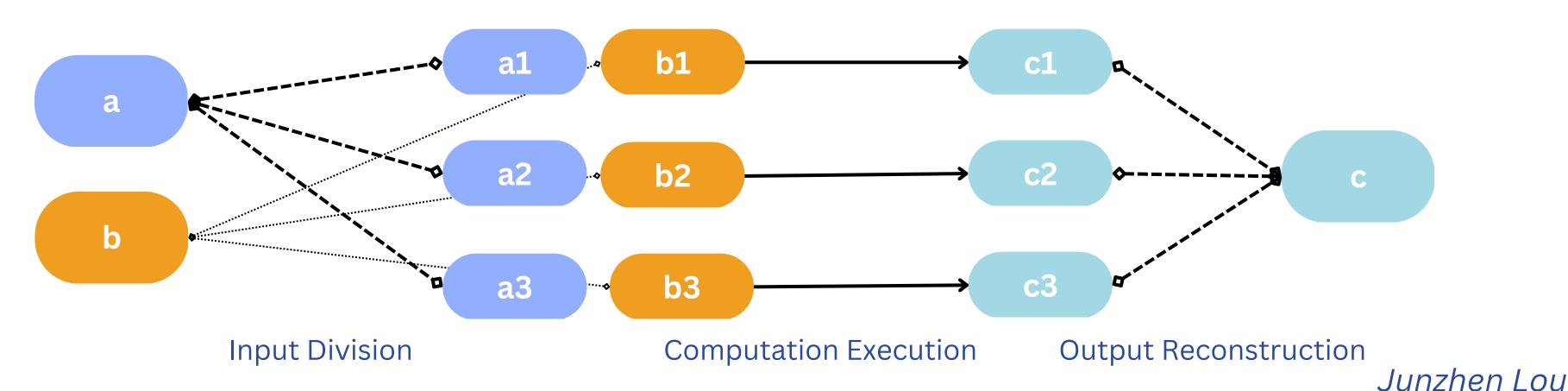
Steps

- Function Identification
- Input Division
- Computation Execution
- Output Reconstruction



Steps

- Function Identification
- Input Division Shamir Secret Sharing Scheme
- Computation Execution Secure Addition & Multiplicaiton Protocols
- Output Reconstruction Lagrange Interpolation

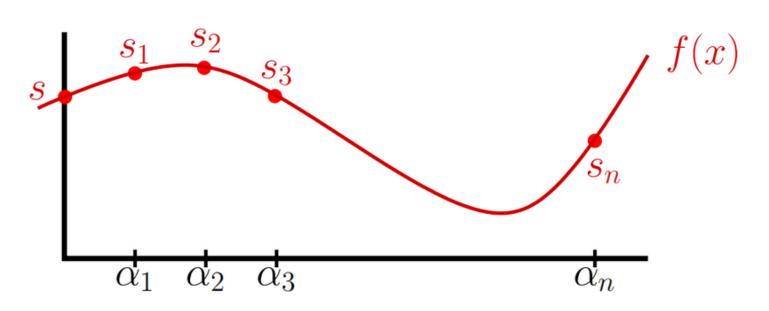


Input Division - Shamir Secret Sharing Scheme

- Threshold secret sharing scheme
- Based on polynomials
- Linear (addictive homomorphism)

Idea

- Random polynomial f of degree d is defined by d + 1 points
- s = f(0) = secret, party Pi gets share $si = f(\alpha i)$ for fixed αi
- Degree d = k-1 ⇒ k parties can reconstruct,
 k-1 cannot



n parties, k needed for reconstruction

Output Reconstruction - Lagrange Interpolation

Reconstruct a value from its shares

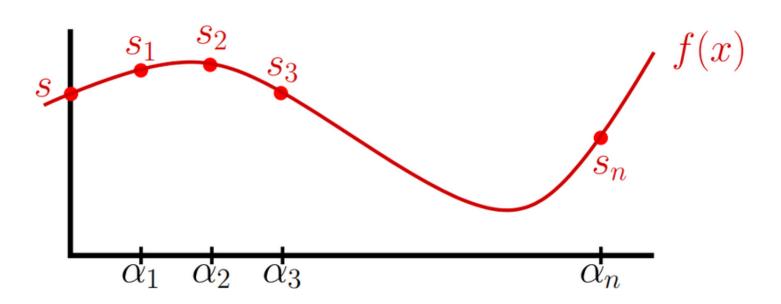
Protocol

- ∀Pi: send si to P
- P: compute s with Lagrange interpolation

$$s = \sum_{i=1}^{n} w_i s_i, \ w_i = \lambda_i(0) = \prod_{\substack{j=1 \ j \neq i}}^{n} \frac{-\alpha_j}{\alpha_i - \alpha_j}.$$

Privacy

 adversary with < k shares obtains no information about s



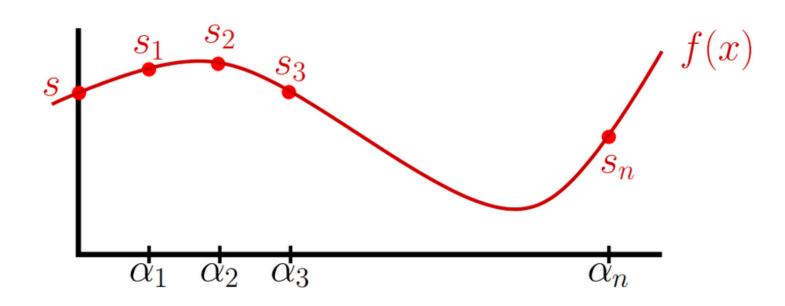
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Computation Execution - Secure Addition Protocol

- Shamir-Sharing is linear
- apply linear function on shares ⇒ apply the same function on original secret values

Protocol

- a, b, . . . shared by a1, ..., an, b1, ..., bn, etc.
- Every Pi computes ci = f(ai, bi,...)
- c1, ..., cn is a sharing of c = f(a, b, ...)



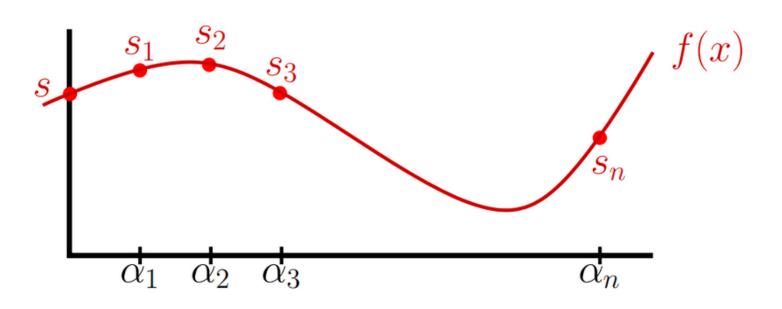
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Computation Execution - Secure Multiplication Protocol

- multiply two secrets ⇒ multiply two polynomials
- doubles the degree

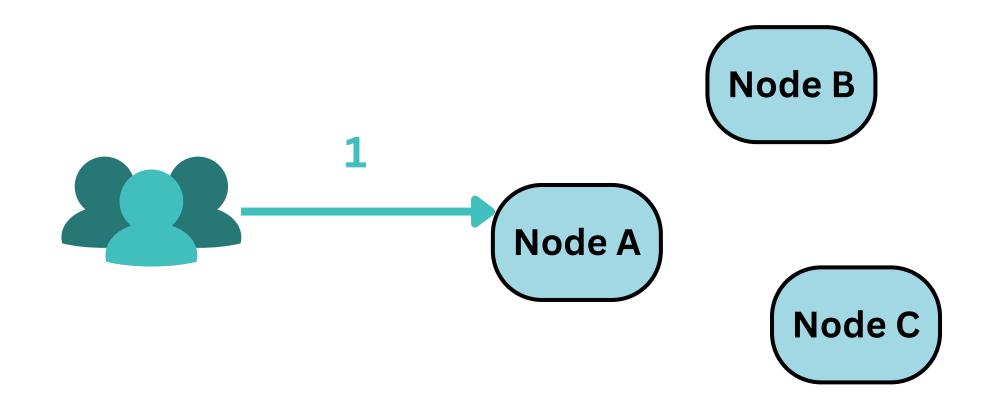
Protocol

- 0. a, b are shared by a1, ..., an, b1, ..., bn
- 1. $\forall Pi$: compute di = ai * bi
- 2. $\forall Pi$: share $di \rightarrow di1, \ldots, din$
- 3. $\forall Pj$: compute $cj = L(d1j, \ldots, dnj)$,
- L(): Lagrange Interpolation

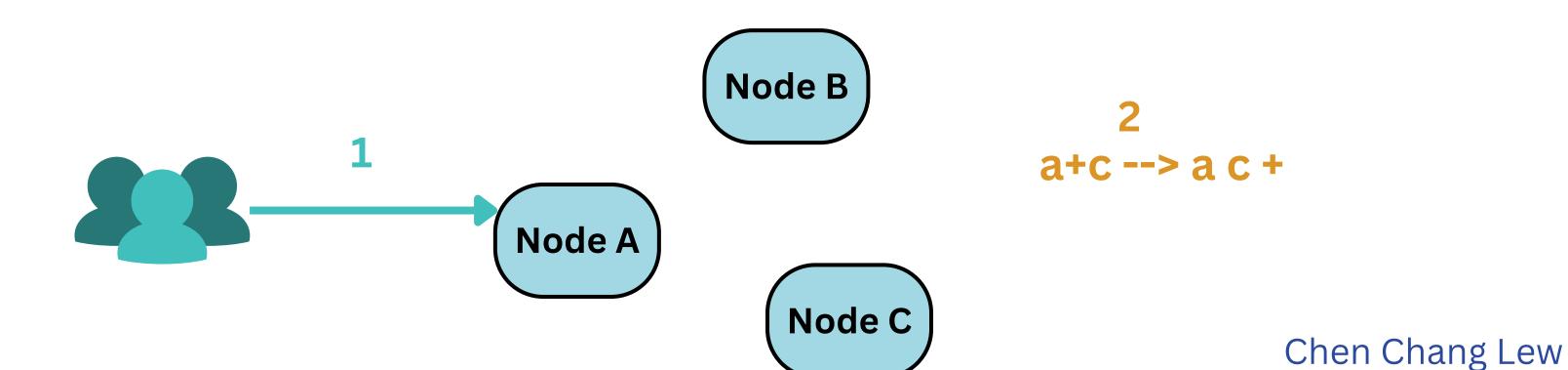


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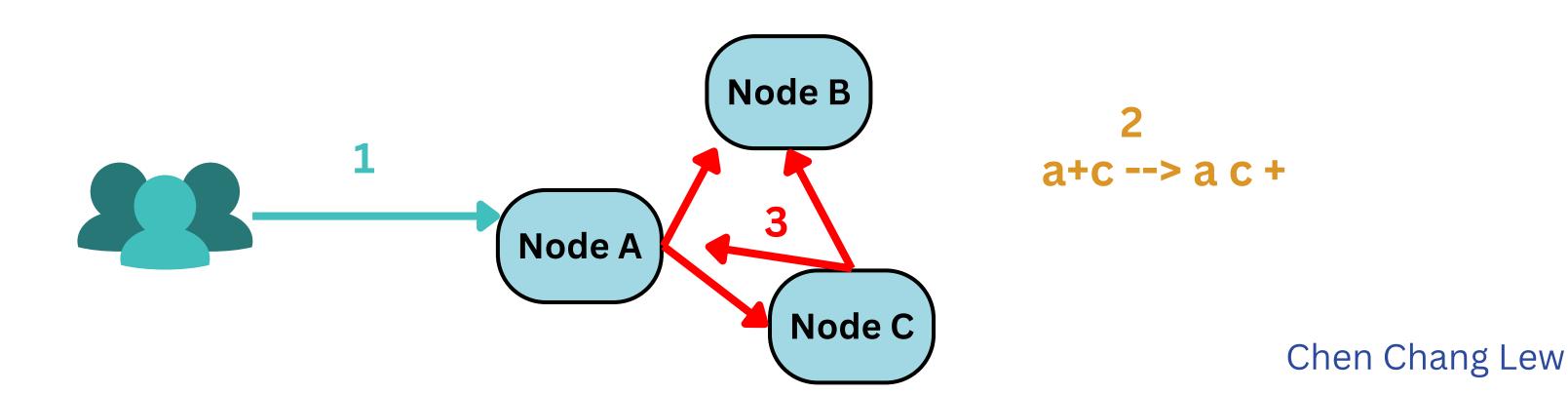
1. Users send an Expression for MPC to count. Ex: a+c



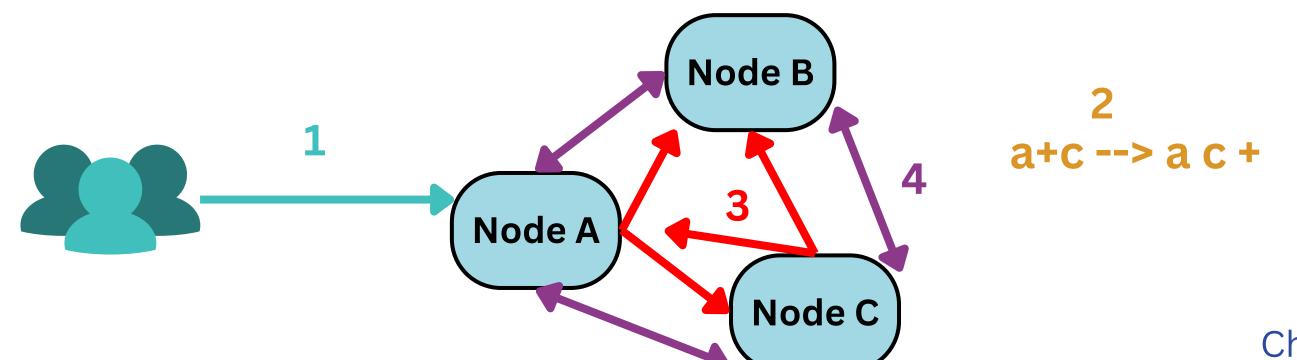
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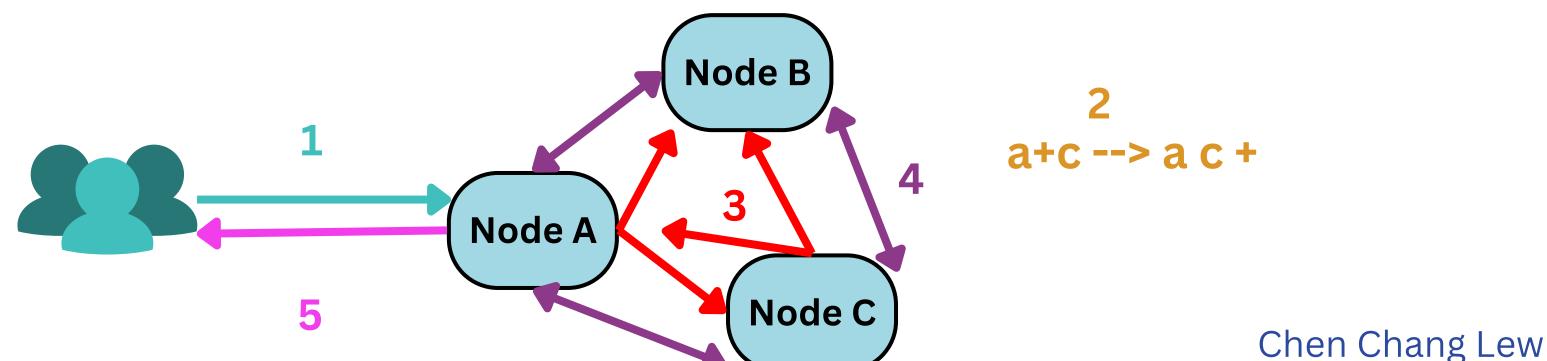
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- 5. Reconstruct the final value using Lagrange Interpolation and return result.



Overview

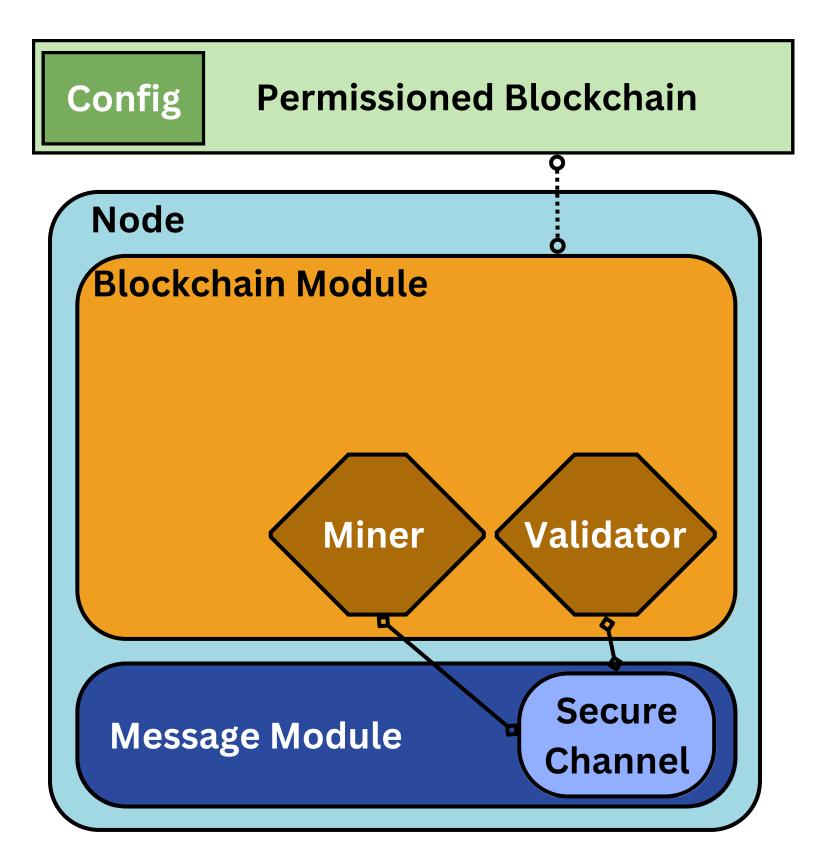
- A decentralized ledger. An Ethereum-like state machine
- In-Chain configuration specifies participants and other policies
- Support execution of specific types of transactions

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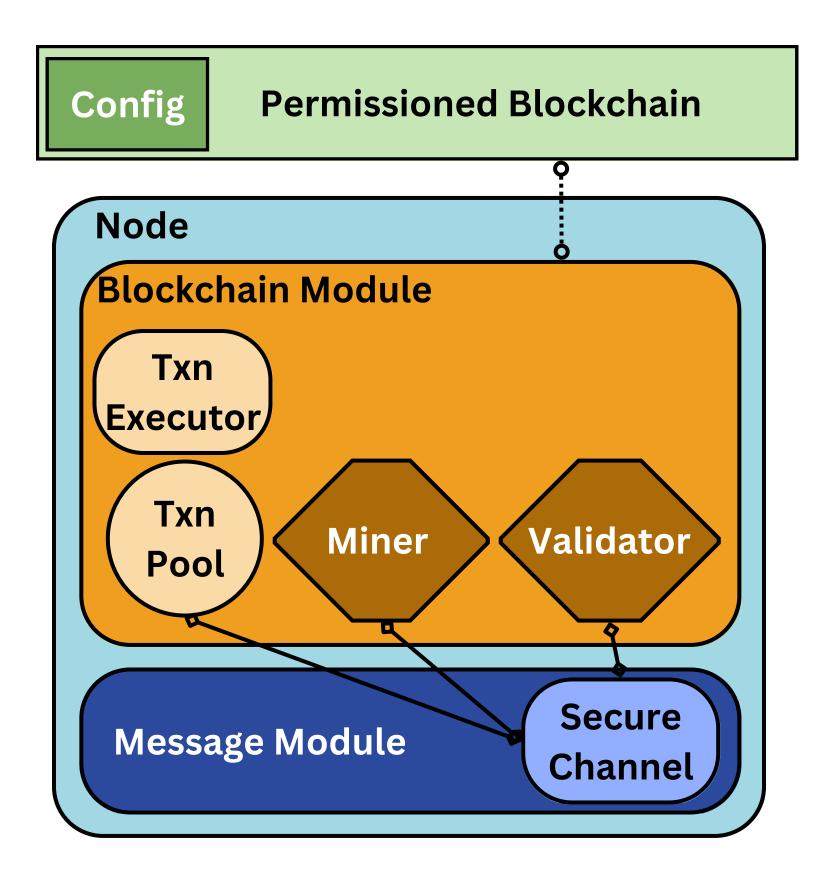
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Purpose

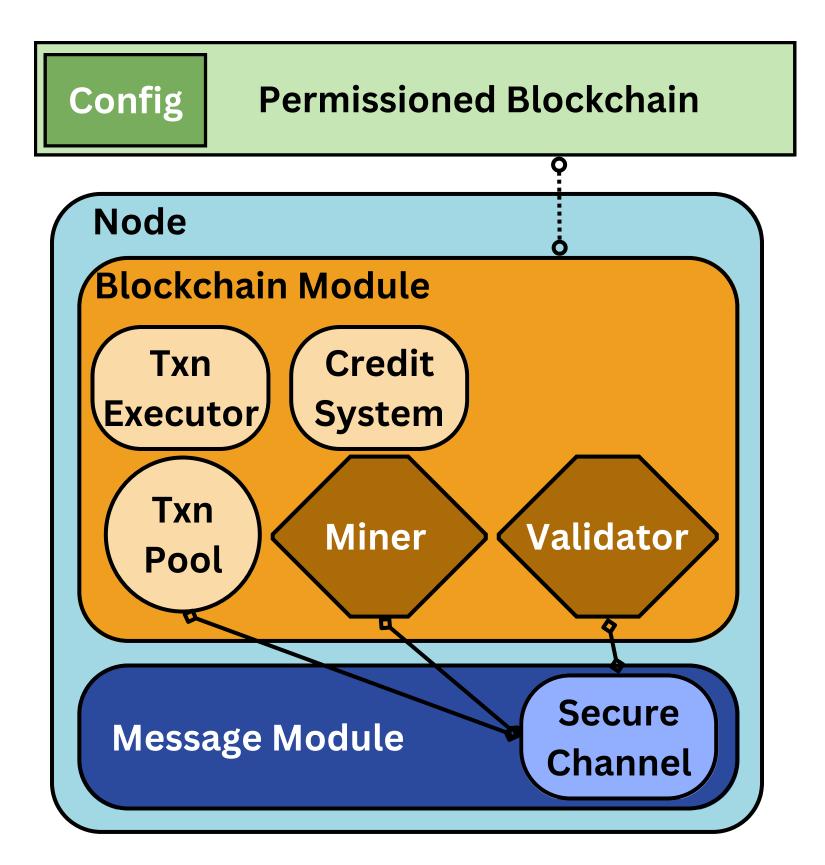
- Create an incentive for parties to provide their data in MPC
- Filter adversaries for MPC
- Facilitate consensus on network assets and MPC participants as well as distribution of encryption keys



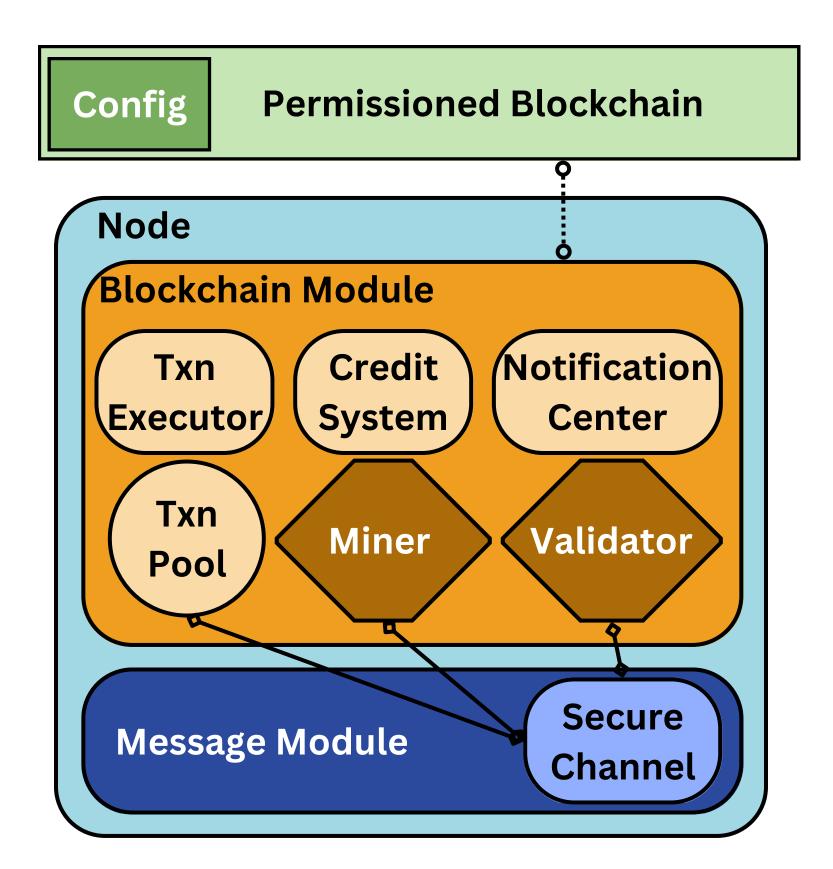
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- Txn Executor: validate and execute txn.
 Only support dedicated txn types

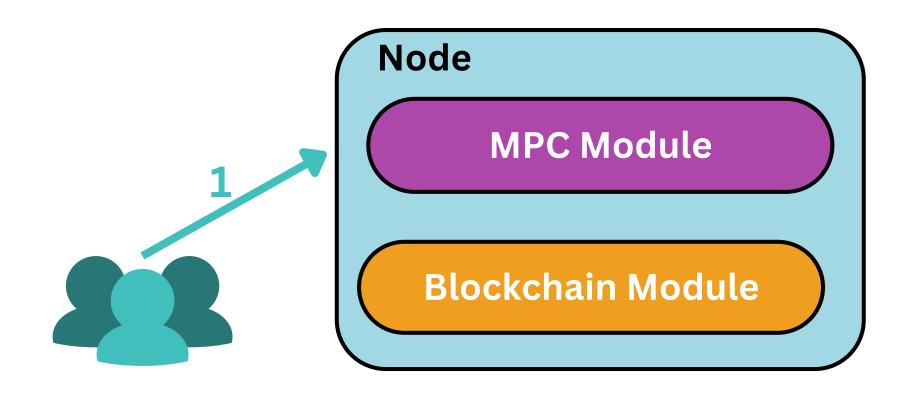


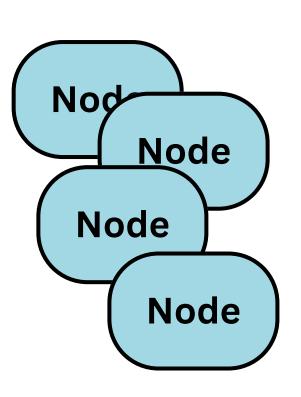
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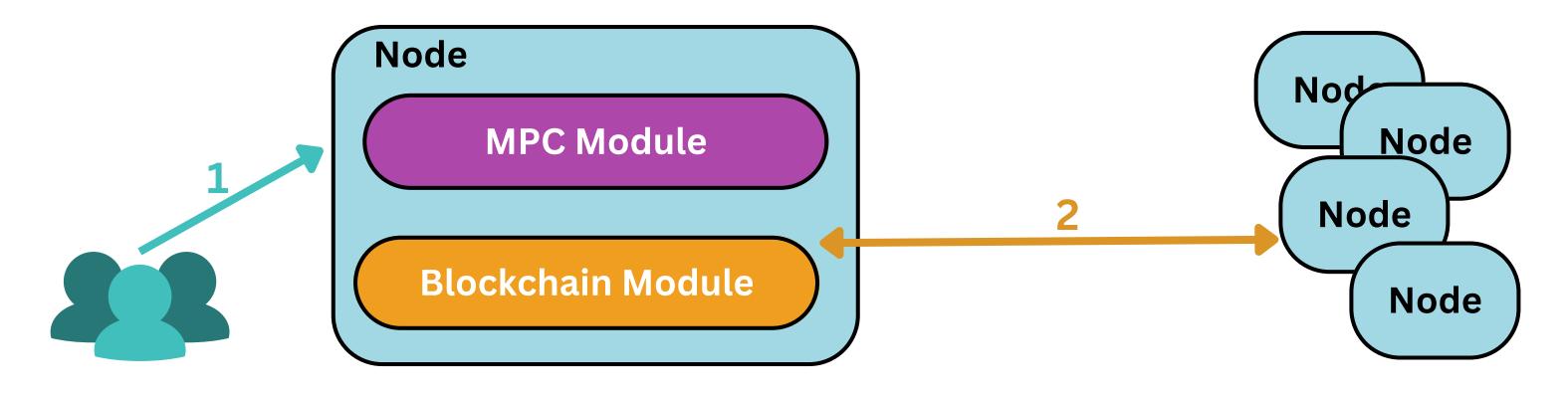
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- Notification Center: notifies other modules of new transaction commits

1. User sends MPC request with a budget to be paid

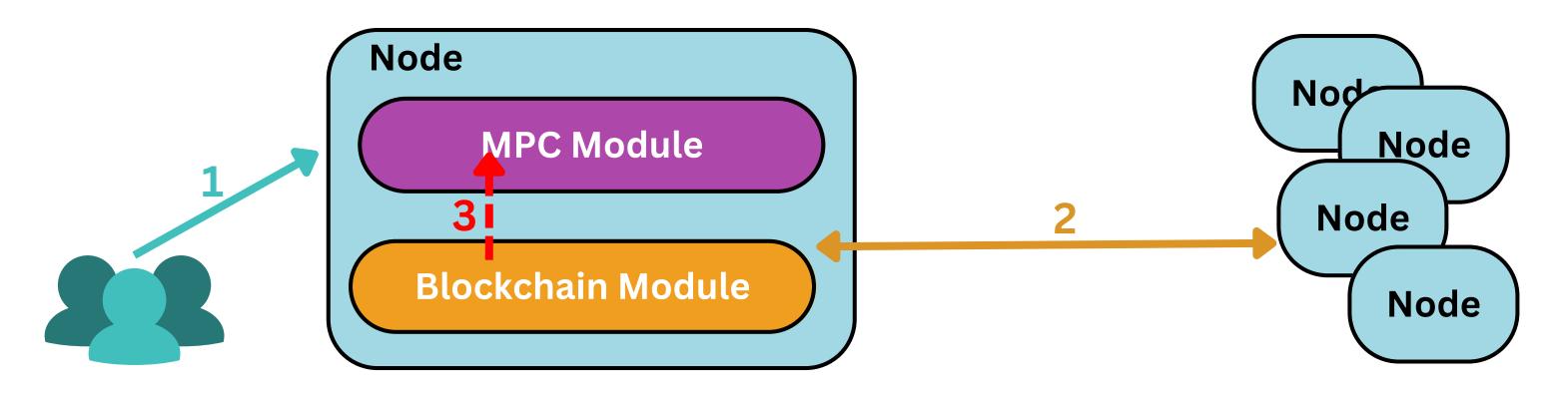




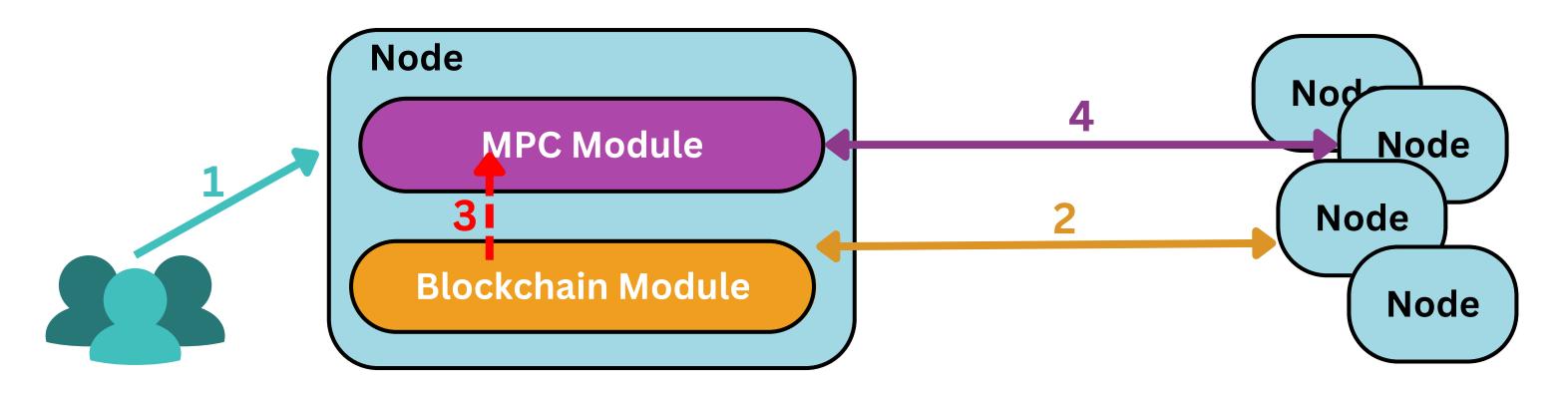
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- 2. Blockchain Module includes the PreMPC txn into the ledger



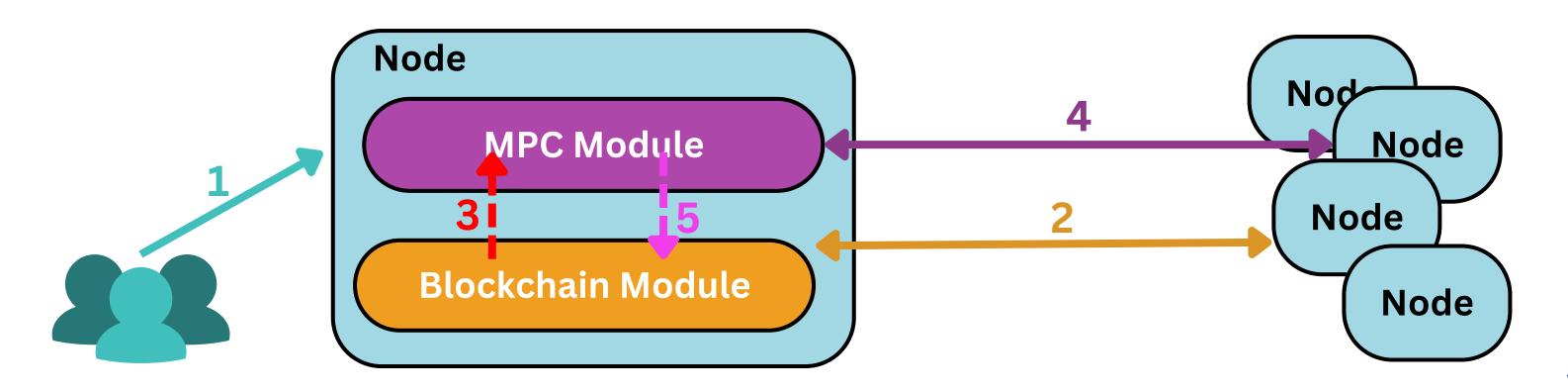
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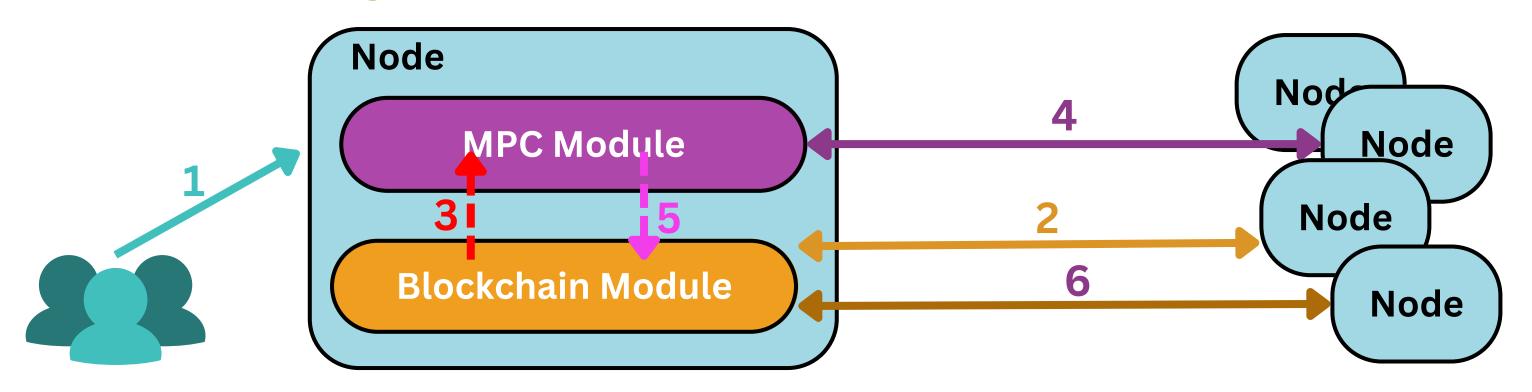
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- 4. MPC Module starts MPC



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- 4. MPC Module starts MPC
- 5. MPC Module notifies Blockchain Module who sends a PostMPC txn



- 1. User sends MPC request with a budget to be paid
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- 3. Blockchain Module notifies MPC Module of PreMPC txn commit
- 4. MPC Module starts MPC
- 5. MPC Module notifies Blockchain Module who sends a PostMPC txn
- 6. Locked budget is released to participants (#PostMPC Txn > 1/2)

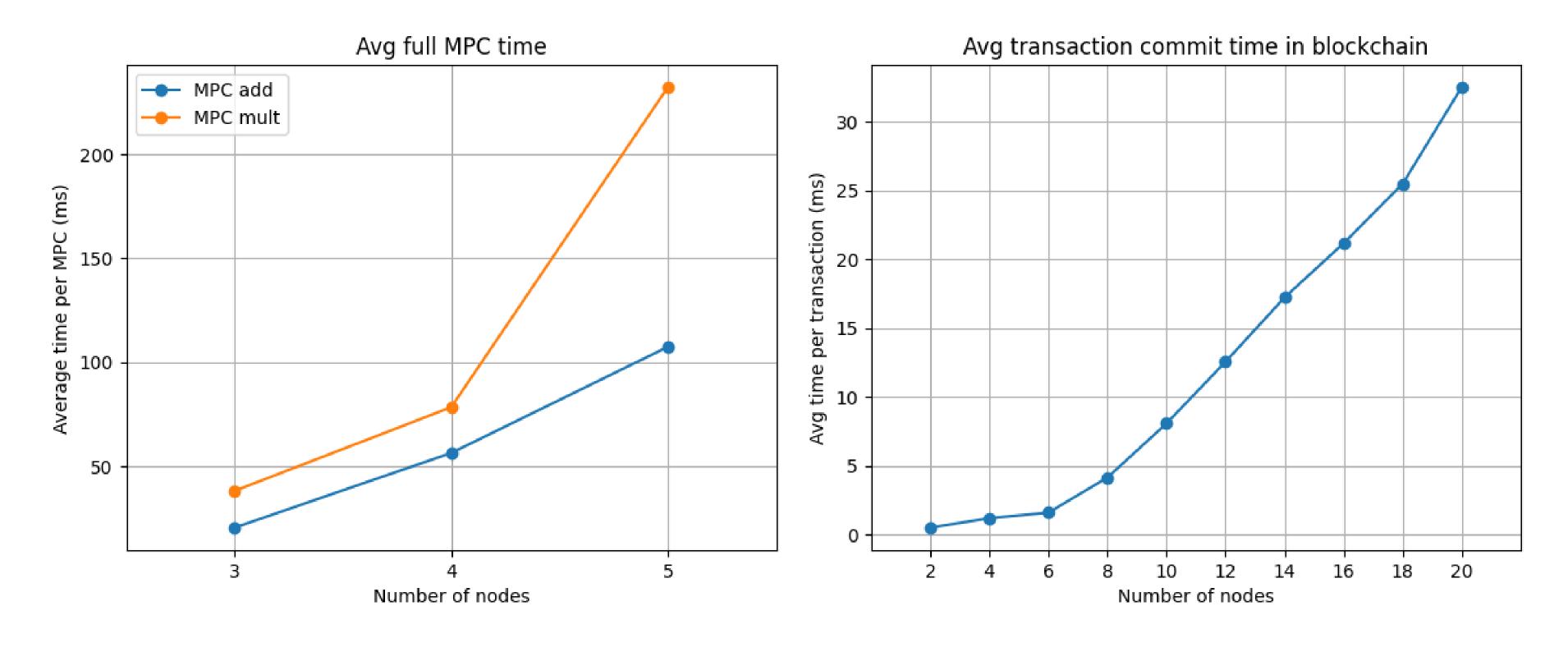


EVALUATION

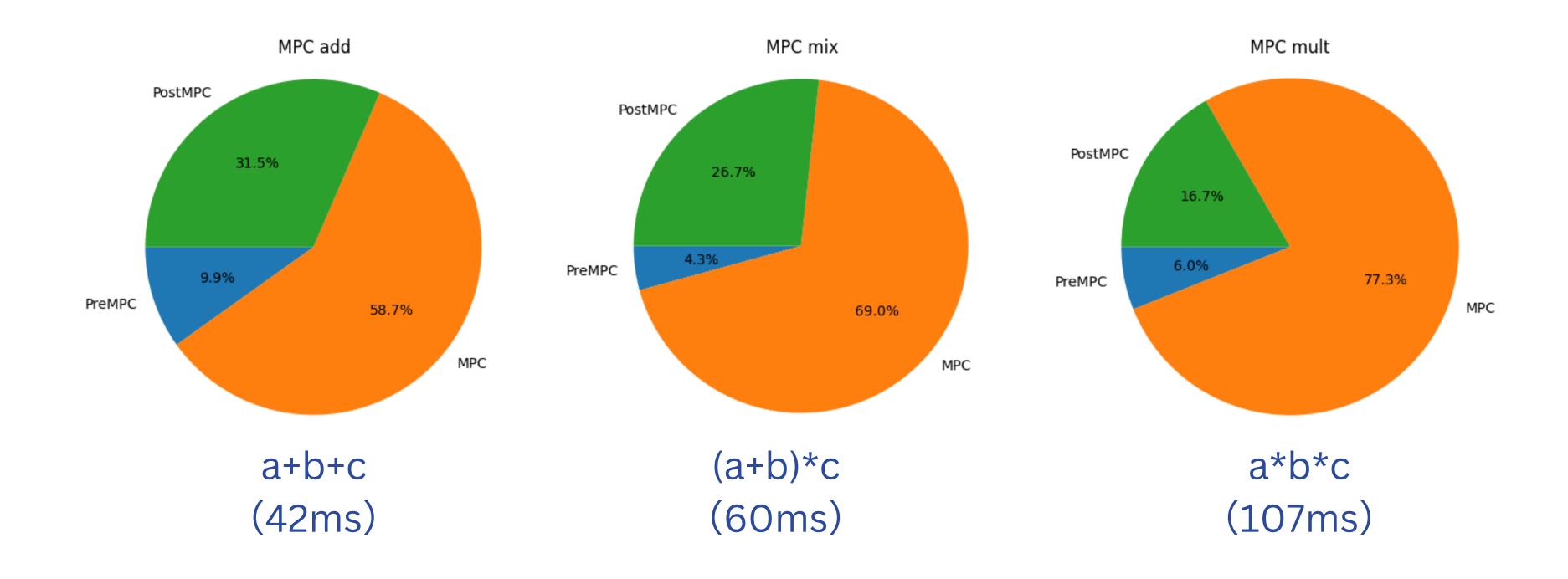
MPC PEERSTER



PERFORMANCE



PERFORMANCE



SECURITY

MPC

Blockchain

Things Provided

- Secure channels
- Synchronous communication
- Passive Adversary
- Information-theoretic level security
 & privacy against <1/2 corruptions

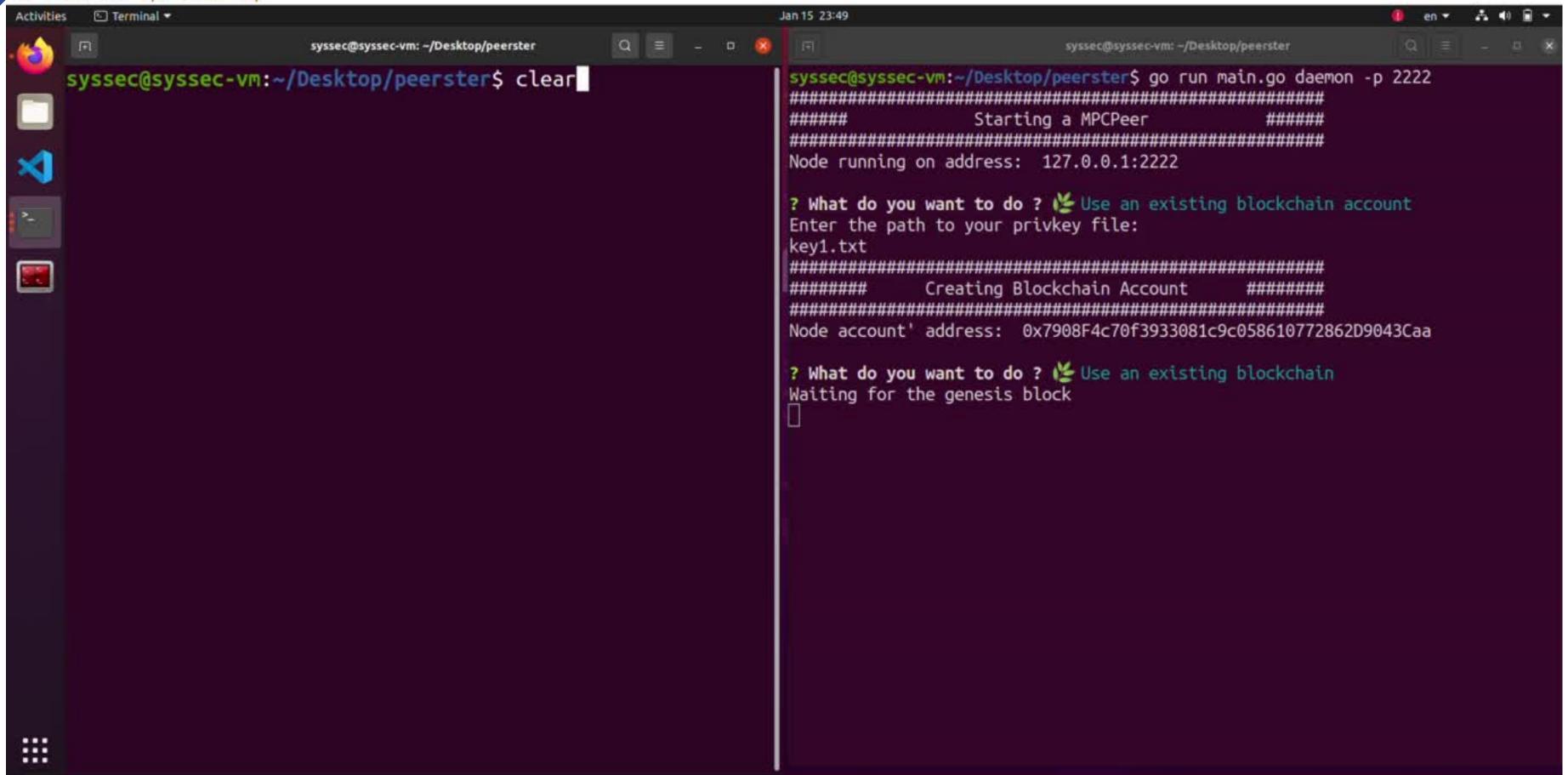
- Access control & Authentication
- Accountability
- Correctness against Active Adversaries who can
 - produce fake transactions
 - replay authenticated transactions
 - produce faulty blocks

Future Improvments

- Asynchronous communicaiton
- Privacy against Active Adversary (crytographic security: Commitment)
- Asynchronous consensus algorithm
- Protection against DoS attack
- Real blockchain privacy



UI available at: https://chenlew.wixsite.com/mpcpeerster *Must be connected to the localhost:7122 node



THANK YOU!

Code available at: https://github.com/lilyyangyn/peerster

Code Contribution			
Chen Chang Lew	MPC Integration Flow, Benchmark UI	Total	
		+2833	-857
Junzhen Lou	MPC Module and Protocols, Finite Field Arithmetic, Benchmark	Total	
		+2126	-297
Yuening Yang	Secure channel, Blockchain, Integration MPC on Blockchain, Benchmark CLI tool	Total	
		+15010	-4774