

# EL9343 Homework 1

Due: Sept. 14th 8:00 a.m.

1. Prove the following properties of asymptotic notation:

(a)  $n = \omega(\sqrt{n})$

$$\forall c > 0, \exists n_0 = c^2 + 1 > 0, \text{ such that } n > c\sqrt{n}, \forall n \geq n_0$$

$$\therefore n = \omega(\sqrt{n})$$

(b) If  $f(n) = \Omega(g(n))$ , and  $h(n) = \Theta(g(n))$ , then  $f(n) = \Omega(h(n))$

$$f(n) = \Omega(g(n)) \Leftrightarrow \exists c_1 > 0, n_1 > 0, \text{ such that } f(n) \geq c_1 g(n), \forall n \geq n_1$$

$$h(n) = \Theta(g(n)) \Leftrightarrow \exists c_2 > 0, n_2 > 0, \text{ such that } h(n) \leq c_2 g(n), \forall n \geq n_2$$

$$\therefore g(n) \geq \frac{1}{c_2} h(n), \forall n \geq n_2$$

$$\therefore \exists c_3 = \frac{c_1}{c_2} > 0, n_3 = \max(n_1, n_2) > 0, \text{ such that } \forall n \geq n_3, f(n) \geq c_1 g(n) \geq c_1 \frac{1}{c_2} h(n) = c_3 h(n)$$

$$\therefore f(n) = \Omega(h(n))$$

(c)  $f(n) = O(g(n))$  if and only if  $g(n) = \Omega(f(n))$  (*Transpose Symmetry* property)

$$\text{Part1: } f(n) = O(g(n)) \implies g(n) = \Omega(f(n))$$

$$f(n) = O(g(n)) \Leftrightarrow \exists c_1 > 0, n_1 > 0, \text{ such that } f(n) \leq c_1 g(n), \forall n \geq n_1$$

$$\therefore c'_1 = \frac{1}{c_1}, \forall n \geq n_1, g(n) \geq c'_1 f(n) \Leftrightarrow g(n) = \Omega(f(n))$$

$$\text{Part2: } f(n) = O(g(n)) \Leftarrow g(n) = \Omega(f(n))$$

$$g(n) = \Omega(f(n)) \Leftrightarrow \exists c_2 > 0, n_2 > 0, \text{ such that } g(n) \geq c_2 f(n), \forall n \geq n_2$$

$$\therefore c'_2 = \frac{1}{c_2}, \forall n \geq n_2, f(n) \leq c'_2 g(n) \Leftrightarrow f(n) = O(g(n))$$

2. Indicate, for each pair of expressions  $(A, B)$  in the table below, whether  $A$  is  $O$ ,  $o$ ,  $\Omega$ ,  $\omega$ , or  $\Theta$  of  $B$ . Assume that  $k \geq 1, \epsilon > 0$ , and  $c > 1$  are constants. Your answer should be in the form of the table with “yes” or “no” written in each box.

	$A$	$B$	$O$	$o$	$\Omega$	$\omega$	$\Theta$
a	$\lg^k n$	$n^\epsilon$	yes	yes	no	no	no
b	$n^k$	$c^n$	yes	yes	no	no	no
c	$\sqrt{n}$	$n^{\sin n}$	no	no	no	no	no
d	$2^n$	$2^{n/2}$	no	no	yes	yes	no
e	$n^{\lg c}$	$c^{\lg n}$	yes	no	yes	no	yes
f	$\lg(n!)$	$\lg(n^n)$	yes	no	yes	no	yes

3. You have 5 algorithms, A1 took  $O(n)$  steps, A2 took  $\Theta(n \log n)$  steps, and A3 took  $\Omega(n^2)$  steps, A4 took  $o(n^3)$  steps, A5 took  $\omega(n^{3/2})$  steps. You had been given the exact running time of each algorithm, but unfortunately you lost the record. In your messy desk you found the following formulas:

(a)  $4(5^{3 \log_5 n}) + 12n + 9527$

(b)  $\sqrt[5]{3n!}$

(c)  $\frac{5^{\log_{16} n} 2}{6} + 4n + 17$

(d)  $3n \log_3 n + (\log_2 n)^3$

(e)  $\log_4 \log_2 n + 61$

(f)  $2^{5 \log_4 n}$

(g)  $(\log_2 n)^2 + \log_3 \log_3 n$

For each algorithm write down all the possible formulas that could be associated with it.

(a)  $4(5^{3 \log_5 n}) + 12n + 9527$

$$4(5^{3 \log_5 n}) + 12n + 9527 = \Theta(5^{3 \log_5 n} + n) = \Theta(n^3 + n) = \Theta(n^3)$$

(b)  $\sqrt[5]{3n!}$

By Stirling's approximation,  $n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$

$$\sqrt[5]{3n!} \approx \sqrt[5]{3} (2\pi n)^{1/10} \left(\frac{n}{e}\right)^{n/5} = \Omega(n^n)$$

(c)  $\frac{5^{\log_{16} n^2}}{6} + 4n + 17$

$$\frac{5^{\log_{16} n^2}}{6} + 4n + 17 = \Theta(n^{\log_{16} 25} + n) = \Theta(n^{\log_{16} 25}) \approx \Theta(n^{1.20})$$

(d)  $3n \log_3 n + (\log_2 n)^3$

$$3n \log_3 n + (\log_2 n)^3 = \Theta(n \log n)$$

(e)  $\log_4 \log_2 n + 61$

$$\log_4 \log_2 n + 61 = \Theta(\log \log n)$$

(f)  $2^{5 \log_4 n}$

$$2^{5 \log_4 n} = \Theta(n^{2.5})$$

(g)  $(\log_2 n)^2 + \log_3 \log_3 n$

$$(\log_2 n)^2 + \log_3 \log_3 n = \Theta((\log_2 n)^2)$$

Thus,

A1: e, g

A2: d

A3: a, b, f

A4: c, d, e, f, g

A5: a, b, f

4. Show what is printed by the following algorithm when called with **MAXIMUM**( $A, 1, 5$ ) where  $A = [7, 3, 8, 5, 9]$ ? Where the function **PRINT** simply prints its arguments in some appropriate manner.

0: **MAXIMUM**( $A, l, r$ )

1: **if** ( $r - l == 0$ ) **then**

2:     **return**  $A[r]$

3: **end if**

4:

5:  $lmax = \mathbf{MAXIMUM}(A, l, \lfloor (l+r)/2 \rfloor)$

6:  $rmax = \mathbf{MAXIMUM}(A, \lfloor (l+r)/2 \rfloor + 1, r)$

7: **PRINT**( $rmax, lmax$ )

8: **if**  $rmax < lmax$  **then**

9:     **return**  $lmax$

10: **else**

11:     **return**  $rmax$

12: **end if**

The printed result is,

3, 7

8, 7

9, 5

9, 8