

CO142 - Discrete Structures

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Recommended Books

- K.H. Rosen. *Discrete Mathematics and its Applications*
- J.L. Gersting. *Mathematical Structures for Computer Science*
- J.K. Truss. *Discrete Mathematics for Computer Science*
- R. Johnsonbaugh. *Discrete Mathematics*
- C. Schumacher. *Fundamental Notions of Abstract Mathematics*

However, these books don't cover the same content. Learn his notation.

Logical Formula, and Notation

This notation will be shared with **CO140**.

| | |
|-----------------------|---|
| $A \wedge B$ | A and B both hold |
| $A \vee B$ | A or B holds (or both) |
| $\neg A$ | A does not hold |
| $A \Rightarrow B$ | if A holds, then so does B |
| $A \Leftrightarrow B$ | A holds if and only if B holds |
| $\forall x(A)$ | the predicate A holds for all x |
| $\exists x(A)$ | the predicate A holds for some x |
| $a \in A$ | the object a is in the set A (a is an element of A) |
| $a \notin A$ | the object a is not in the set A |
| $=_A$ | tests whether two elements of A are the same |

Sets

Sets are like data types in Haskell: Haskell data type declaration;

```
data Bool = False | True
{false, true}                                set of boolean values
[true, false, true, false]                   list of boolean values
{false, true} = {true, false}                 set equality (note that order doesn't matter)
```

A set is a collection of objects from a pool of objects. Each object is an *element*, or a *member* of the set. A set *contains* its elements. Sets can be defined in the following ways;

| | |
|-------------------------|--|
| $\{a_1, \dots, a_2\}$ | as a collection of n distinct elements |
| $\{x \in A \mid P(x)\}$ | for all the elements in A , where P holds |
| $\{x \mid P(x)\}$ | for all elements, where P holds (dangerous - Russel's paradox) |

Use of "triangleq"

The use of \triangleq is for "is defined by". Hence the empty set, $\emptyset \triangleq \{\}$. The difference between \triangleq and $=$, is that the former cannot be proven, it is fact, whereas the latter takes work to prove.

Russel's paradox

Not everything we write as $\{x \mid P(x)\}$ is automatically a set. Assume $R = \{X \mid X \notin X\}$ is a set, the set of all sets which don't contain themselves. As R is a set, then $R \in R$, or $R \notin R$ (law of excluded middle), and thus we can do a case by case analysis.

- Assume $R \in R$. By the definition of R , it then follows that $R \notin R$ (if $R \in R$, then it doesn't satisfy the definition of R) - which is a contradiction.
- Assume $R \notin R$. It then follows that $R \in R$, as it follows the definition of R , hence it is another contradiction.

As both assumptions lead to contradictions, it's possible to write sets which aren't defined. We should only select from a set that we know is defined; $\{x \in A \mid P(x)\}$ - where A is a well-defined set.