

# Linear Programming Fundamentals

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# Outline

Linear Programming: A Historic and Vital Tool in OR

What is Linear Programming (LP) ?

Linear Programming applications

LP assumptions

Graphical solution of LP

Summary

# Linear Programming (LP): A Historic and Vital Tool in OR

- 1950: LP Discovery - A Landmark Moment in OR for Linear Programming.
- Linear Programming's Impact: Saving Money in Diverse Industries, such as education, forestry, petroleum, ...



- In a survey of Fortune 500 firms, 85% of the respondents said they had used LP

# Prototype example

You run a company that produces two products: **chairs and tables**.

- You have 500 hours of labor, 400  $m^2$  of wood, and 300  $m^2$  of fabric available.
- Each chair requires 3 hours of labor, 4  $m^2$  of wood, and 2  $m^2$  of fabric.
- Each table requires 4 hours of labor, 6  $m^2$  of wood, and 3  $m^2$  of fabric.
- You sell each chair for DZD 100 and each table for DZD 150.

# Prototype example

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- Each table requires 4 hours of labor, 6  $m^2$  of wood, and 3  $m^2$  of fabric.
- You sell each chair for DZD 100 and each table for DZD 150.

How many chairs and tables should you produce in order to maximize your profit?

# Prototype example

- Decision variables ?
- Objective function ?
- Constraints ?

# Prototype example

- **Decision variables:**
  - Let  $x_1$  be the number of chairs to produce.
  - Let  $x_2$  be the number of tables to produce.
- **Objective function:**
  - Maximize  $100x_1 + 150x_2$
- **Constraints:**
  - $3x_1 + 4x_2 \leq 500$  (labor constraint)
  - $4x_1 + 6x_2 \leq 400$  (wood constraint)
  - $2x_1 + 3x_2 \leq 300$  (fabric constraint)
  - $x_1, x_2 \geq 0$  (non-negativity constraint)

## Prototype example

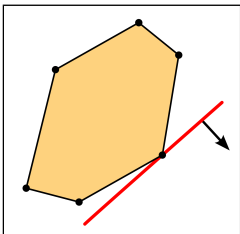
$$\begin{array}{ll}\text{Maximize} & 100x_1 + 150x_2 \\ \text{Subject to} & 3x_1 + 4x_2 \leq 500 \\ & 4x_1 + 6x_2 \leq 400 \\ & 2x_1 + 3x_2 \leq 300 \\ & x_1, x_2 \geq 0\end{array}$$

What is the solution approach for this problem on a two-dimensional plane?



# What is Linear Programming (LP) ?

An optimization model is a Linear Program (or LP) if it has **continuous variables**, a **single linear objective function**, and **all constraints are linear equalities or inequalities**.



# What is Linear Programming (LP) ?

- **Decisions variables** : that we seek to determine
  - $x_1, x_2, \dots, x_n$
- **Linear objective function**: (goal) that we need to optimize (maximize or minimize)
  - $z = c_1x_1 + c_2x_2 + \dots + c_nx_n$
  - $z = \sum_{i=1}^n c_i x_i$
- **Linear constraints**: that the solution must satisfy
  - $a_{1j}x_1 + a_{2j}x_2 + \dots + a_{nj}x_n \leq b_j$  for  $j = 1 \dots m$

# What is Linear Programming (LP)?

$$\begin{array}{ll}\textbf{Maximize} & z = c_1x_1 + c_2x_2 + \cdots + c_nx_n \\ \textbf{Subject to} & a_{11}x_1 + a_{21}x_2 + \cdots + a_{n1}x_n \leq b_1, \\ & a_{12}x_1 + a_{22}x_2 + \cdots + a_{n2}x_n \leq b_2, \\ & \vdots \\ & a_{1m}x_1 + a_{2m}x_2 + \cdots + a_{nm}x_n \leq b_m, \\ & x_i \geq 0 \quad (1 \leq i \leq n)\end{array}$$

- Here,  $x_i$  represents the decision variables,  $c_i$ ,  $a_{ij}$  and  $b_j$  are fixed constants extracted from the studied problem.  $x_i \geq 0$  represents non-negativity constraint.

# What is Linear Programming (LP)?

- Any LP problem can be converted to the previous form 11.
- An equality ( $=$ ) can be represented by two inequalities ( $\leq$  and  $\geq$ )
- An inequality  $\geq$  can be converted to  $\leq$  by multiplying by  $-1$ .
- A negative variable  $x_i$  can be replaced with  $x_i = -x'_i$  and  $x'_i \geq 0$ .
- An unrestricted variables (i.e., it can take both positive and negative values) can be replaced by two variables  $x_i = x_i^+ - x_i^-$  and  $x_i^+, x_i^- \geq 0$ .

# LP Applications: Resource Allocation

- LP commonly used for allocating resources to activities with limited resources
- The objective is to **choose levels of activities** that achieve the **best overall measure of performance**.
- E.g. Optimizing the scheduling of personnel and equipment in a hospital or allocating budgets across different departments in a company
- Organizations can make the most efficient use of their resources and achieve better performance outcomes using LP.

Resource	Resource Usage per Unit of Activity				Amount of Resource Available
	Activity				
	1	2	...	$n$	
1	$a_{11}$	$a_{12}$	...	$a_{1n}$	$b_1$
2	$a_{21}$	$a_{22}$	...	$a_{2n}$	$b_2$
.					.
.	...	...	...	...	.
.					.
$m$	$a_{m1}$	$a_{m2}$	...	$a_{mn}$	$b_m$
Contribution to Z per unit of activity	$c_1$	$c_2$	...	$c_n$	

## Example: Resource Allocation

- **Context:** manufacturing company has two factories and three product lines.
- **The company's objective:** maximize the profit.
- **Constraint 1:** each factory has a limited capacity for producing each product line.
- **Constraint 2:** limited demand for each product line.

## Example: Resource Allocation

- **Context:** manufacturing company has two factories and three product lines.
- **The company's objective:** maximize the profit.
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- **Constraint 2:** limited demand for each product line.

**Maximize**  $15x_1 + 20x_2 + 10x_3$

**Subject to:**

$$2x_1 + 3x_2 + x_3 \leq 30 \quad (\text{Factory 1 capacity constraint})$$

$$x_1 + 2x_2 + 4x_3 \leq 40 \quad (\text{Factory 2 capacity constraint})$$

$$x_1 + x_2 + x_3 \leq 20 \quad (\text{Product 1 demand constraint})$$

$$2x_2 + 3x_3 \leq 25 \quad (\text{Product 2 demand constraint})$$

$$x_1 + 2x_3 \leq 15 \quad (\text{Product 3 demand constraint})$$

$$x_1, x_2, x_3 \geq 0 \quad (\text{Non-negativity constraint})$$



# LP Applications

1. Production planning:
  - Determine optimal mix of products to produce
  - Determine quantity of each product to manufacture
  - Utilize resources such as labor, materials, and machines efficiently
2. Transportation and distribution:
  - Minimize transportation costs
  - Satisfy demand and supply constraints
3. Portfolio optimization:
  - Select best mix of assets
  - Maximize returns while minimizing risks

# LP Applications

## 4. Marketing optimization:

- Optimize advertising campaigns and promotions
- Determine best media channels, advertising frequency, and promotion offers
- Maximize sales and profits

## 5. Energy and environmental management:

- Optimize use of energy resources
- Reduce environmental impacts
- Determine optimal mix of energy sources and production technologies

## 6. Other applications ...

# LP assumptions

- LP assumptions are implicit in the mathematical formulation.
- In order to use LP from a modeling perspective, it is necessary for the assumptions about the data to be fulfilled..

Understanding LP assumptions can help evaluate how well linear programming applies to a problem.

## LP assumptions: Proportionality

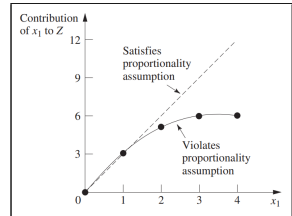
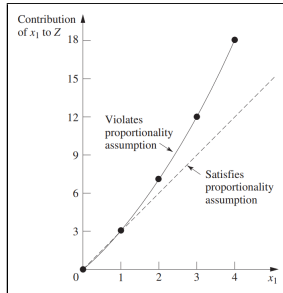
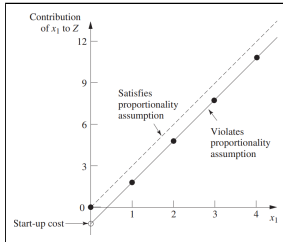
- The contribution of each variable to the objective function  $z$  and functional constraints is proportional to  $x_j$ .
- This is represented by the terms  $c_j x_j$  in the objective function and  $a_{ij} x_j$  in the functional constraint.

In a LP model, only variables with an exponent of 1 are allowed in the terms of any function, including the objective function and constraints.

## LP assumptions: Proportionality

$x_1$	Proportionality Satisfied	Proportionality Violated		
		Case 1	Case 2	Case 3
0	0	0	0	0
1	3	2	3	3
2	6	5	7	5
3	9	8	12	6
4	12	11	18	6

# LP assumptions: Proportionality



## LP assumptions: Additivity

- Proportionality assumption in LP model restricts exponents to 1.
- It does not forbid cross-product terms like  $x_i x_j$ .
- Additivity assumption in LP model rules out cross-product terms.

Every function in the model can be represented as the sum of the separate contributions of the individual activities.

## LP assumptions: Additivity

Value of the objective function $z$			
$(x_1, x_2)$	Additivity Satisfied	Additivity Violated	
		Case 1	Case 2
(1,0)	3	3	3
(0,1)	5	5	5
(1,1)	8	9	7



# LP assumptions: Divisibility

- The divisibility assumption allows decision variables to have values that are not limited to integers.
- Integer Programming can be used if the problem requires integer decision variables.

LP allows non-integer decision variable values, implying the possibility to run at fractional levels, which can result in more efficient solutions in some scenarios.

# LP assumptions: Certainty

- The certainty assumption in LP models assumes that all the LP coefficients are **known fixed constants**.
- In practical applications, satisfying the certainty assumption in can be challenging as parameter values are often **estimated or predicted values**.
- Therefore, **sensitivity analysis** is crucial to identify sensitive parameters that may affect the optimal solution.

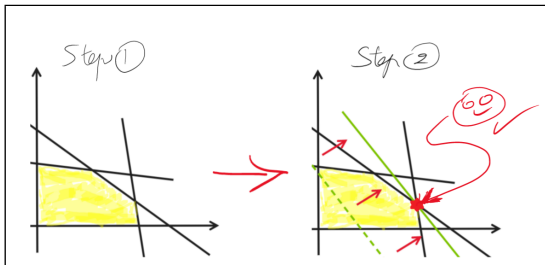
## LP assumptions: Conclusion

- Mathematical models are just approximations and simplifying assumptions.
- Reasonably high correlation between the model's prediction and the actual problem is sufficient for useful analysis.
- If any of the assumptions are violated significantly, alternative models are available.
- LP outperforms other complex methods in terms of its effectiveness and efficiency.

# Graphical solution of LP

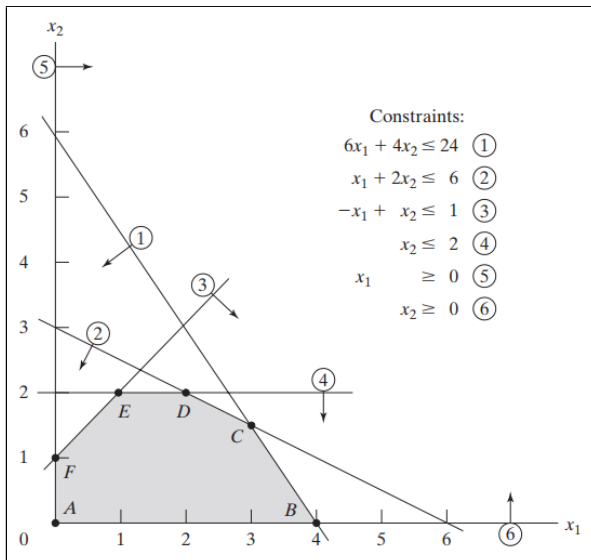
The graphical solution includes two steps:

1. Determination of the feasible solution space (feasible region).
2. Determination of the optimum solution among all points of the feasible region.



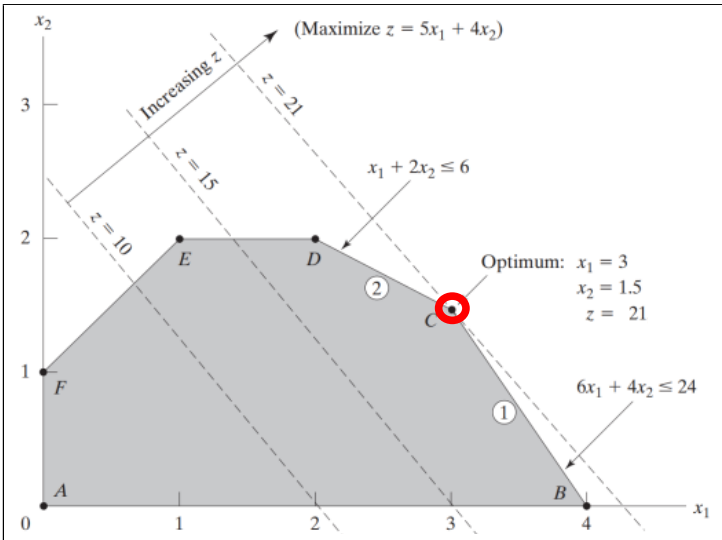
# Determination of the feasible region

1. **Feasible region:** set of possible solutions that meet the problem constraints.
2. **Graph inequalities:** represent constraints on a coordinate plane.
3. **Shade region:** satisfy all constraints to determine feasible solution space.



# Determination of the optimum solution

1. **Optimum solution:** determine best solution from feasible solution space.
2. **Graph the objective function:** Fix  $z$  and draw the line called **isoprofit/ iso-cost**
3. **Shift the line:** shift the iso-profit line in a parallel manner from its original position in a direction that increases of  $z$ .





# Why study the graphical solution of LP ?

- In practice, LPs can have hundreds or even thousands of variables and constraints.
- Two-variable LP may seem limited, but the graphical solution offers a key insight.
- The optimal solution of an LP is always at a **corner point** of the feasible region.

## Key result

The graphical solution of a two-variable LP problem provides a fundamental insight into LP optimization, which is that **the optimal solution is always associated with a corner point of the feasible region.**

# Why Corner Points help to solve LP ?

- This result limits the search for the optimal solution from an **infinite number of feasible points** to a **finite number of corner points**.
- Simplex method is based on this powerful result to solve complex LP problems.
- In two dimension case, the upper bound of the number of corner points is

**Number of corner points**  $\leq \binom{n}{2} = \frac{n(n-1)}{2}$  : where  $n$  is the number of constraints.

## Key result

Intelligently navigating the corner points of a LP problem can simplify solving very large LP problems using optimization algorithms.

# Convex Sets, Extreme Points, and LP

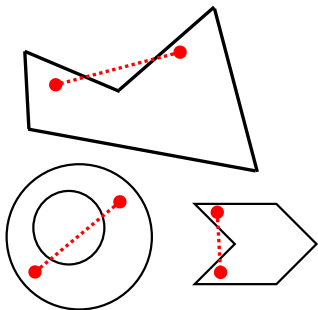
## Convex set

A set of points  $S$  is a convex set if the line segment joining any pair of points in  $S$  is wholly contained in  $S$ .

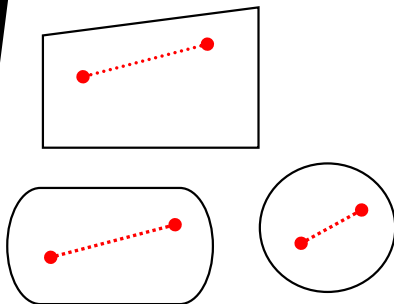
$S$  is convex  $\Rightarrow \forall A, B \in S, \beta \in [0, 1], \beta A + (1 - \beta)B \in S$

## Extreme Point (Corner point)

For any convex set  $S$ , a point  $P$  in  $S$  is an extreme point if each line segment that lies completely in  $S$  and contains the point  $P$  has  $P$  as an endpoint of the line segment.



**Non-convex Sets**



**Convex Sets**

# Convex Sets, Extreme Points, and LP

- Feasible region **cannot comprise several disconnected feasible regions.**
- Convexity guarantees **the absence of local optima.**
- **No local optimum** that could trap the solver.

## Key result

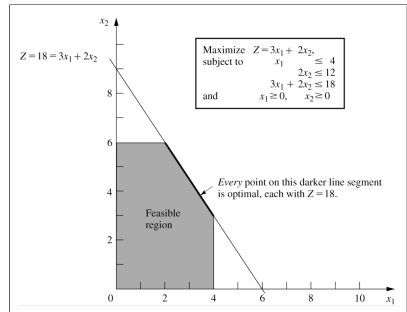
Feasible region convexity is a critical property and it assures that if the LP problem has an optimal solution, then **at least one corner point of the feasible region is a global optimum.**

# LP Special Cases

1. **Alternative or multiple optimal solutions:** LPs have an infinite number of optimal solutions.
2. **Infeasible LPs:** LPs have no feasible solutions.
3. **Unbounded LPs:** There are points in the feasible region with arbitrarily large (in a max problem) z-values.

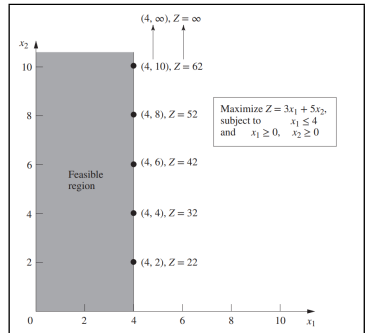
# Alternative or multiple optimal solutions

- The isoprofit intersects an **entire line segment** corresponding to the constraint.
- The decision maker can use a **secondary criterion** to choose between optimal solutions.



# Unbounded LPs

- Objective function can **be arbitrarily large**.
- Unbounded optimal solution **should not occur in a correctly formulated LP**.
- **Modify the problem formulation** to remove the unboundedness to make the problem solvable.





# Infeasible LPs

- **Feasible region to be empty** (contain no points), resulting in an infeasible LP.
- Because the optimal solution to an LP is the best point in the feasible region, an **infeasible LP has no optimal solution**.
- **Over-constraining the problem**, which makes the constraints are too tight to allow any feasible solution.
- **Introduction of incompatible constraints**, such as when one constraint requires a variable to be positive while another requires it to be negative.
- **Modify the problem formulation** to remove the contradictory constraint or constraints and make the problem feasible.

# Summary

- LP problem has three parts:
  1. **Decision variables:** variable that we can control and aim to optimize by finding the best values.
  2. **Objective function:** a linear function of decision variables that we want to maximize or minimize.
  3. **Constraints:** linear inequalities or equalities that restrict the values of decision variables.
- Optimal solution to the LP is a point in the feasible region that maximizes / minimizes the objective function.
- Verifying LP assumptions with data assesses the feasibility of using LP.

# Summary

- The feasible region for any LP is a convex set, and an optimal solution to an LP is an extreme point of the feasible region.
- To solve a max LP with two decision variables graphically, follow these steps:
  1. **Step 1:** Graph the feasible region.
  2. **Step 2:** Draw an isoprofit line.
  3. **Step 3:** Move parallel to the isoprofit line in the direction of increasing  $z$ . The last point in the feasible region that contacts an isoprofit line is an optimal solution to the LP.
- The optimal solution to a max LP with two decision variables is the extreme point of the feasible region that lies on the highest isoprofit line.

# Summary

Four cases that can occur when solving an LP:

1. **Case 1:** The LP has a unique solution.
  - Only one point in the feasible region that maximizes or minimizes the objective function.
2. **Case 2:** The LP has more than one optimal solution.
  - Only one point in the feasible region that maximizes or minimizes the objective function.
  - The feasible region has multiple points that optimize the objective function.
  - Graphically, the isoprofit line last hits an entire line segment before leaving the feasible region.

# Summary

5. **Case 3:** The LP is infeasible.

- No points in the feasible region.
- The LP has no feasible solution.

6. **Case 4:** The LP is unbounded.

- There are points in the feasible region with arbitrarily large objective function values.
- Graphically, when we move parallel to an isoprofit line in the direction of increasing  $z$ , we never lose contact with the LP's feasible region.

# References

1. Winston, Wayne & Goldberg, Jeffrey. **Operations research: applications and algorithms.**
2. Frederick S. Hillier and Gerald J. Lieberman. **Introduction to Operations Research.**
3. Hamdy, A. Taha. **Operations Research: An Introduction**
4. Robin J. Wilson. **Introduction to graph theory.**