

Ex 4.6

(let $\underbrace{((v_1 e_1) (v_2 e_2) \dots)}_{\text{ve}}$ $\underbrace{(\dots)}_{\text{body}})$
 $((\lambda (v_1 v_2 \dots) (\dots)) e_1 e_2 \dots)$

(define (let → combination exp)
 (cons (list 'lambda (map car (ve exp)) (body exp))
 (map cadr (ve exp))))

(define (ve exp) (cadr exp))

(define (body exp) (caddr exp))

(define (let? exp) (tagged-list? exp 'let))
 ↓
 if list beg with let

→ Add to eval:

((let? exp) (eval (let → combination exp) env))

$(\text{let } ((x \ 3)) (\text{let } ((y \ 4)) (\text{let } (= 2) \ (\text{body})))$

Ex 4.7

$(\text{define } (\text{let}^* \rightarrow \text{nested-lets}) \ \text{exp})$
 $\quad (\text{define } (\text{letloop}' \ \text{ve}) (\text{if} (\text{null?} \ \text{ve}) (\text{car} \ (\text{body} \ \text{exp})))$
 $\quad \quad (\text{cons} \ (\text{list}' \ \text{let} \ (\text{car} \ \text{ve}))$
 $\quad \quad \quad (\text{letloop} \ (\text{cdr} \ \text{ve}))))))$
 $\quad (\text{letloop} \ (\text{ve} \ \text{exp})))$

$(\text{let}^* ((v_1 \ e_1) (v_2 \ e_2) \dots) \ (\text{body}))$
 $\Rightarrow (\text{let } ((v_1 \ e_1)) (\text{let } ((v_2 \ e_2) \ \dots \ (\text{body})))) \dots$

yes it is enough since it will convert to a
let expression and be evaluated accordingly.

Ex 4.13

```
(define (make-unbound! var env)
  (let ((vars (cdr (assoc var env))))
    (cond ((null? vars) #t)
          ((eq? var (car vars)) (begin
                                   (set-car! vars null) (set-car! vals null))
          (else (make-unbound! (cdr vars) (cdr vals)))))
    (let ((frame (first-frame env)))
      (make-unbound! (frame-variables frame))))
```

Ex 4.15

$(p \ a) \xrightarrow{\text{halts}} \text{value}$
 $\xrightarrow{\text{not}} \text{error / forever}$

$(\text{halts? } p \ a)$

$(\text{define } (\text{try } p))$

$(\text{if } (\text{halts? } p \ p) \ (\text{run forever}) \rightarrow \text{it halts } p \text{ on } p$
 $'\text{halted}))$

try does the inverse behavior of the procedure.

Assume $(\text{halts? } \text{try } \text{try})$ halts then $(\text{try } \text{try})$ will run forever.

Assume $(\text{halts? } \text{try } \text{try})$ runs forever then $(\text{try } \text{try})$ will halt.

These 2 assumption + outcome contradict each other.

They also assume existence of halts? , So halts? is not possible.

- if / cond

```
(if (= n 1) 1 (* n (fact (- n 1))))  
(cond ((= n 1) 1)  
      (else (* n (fact (- n 1))))))
```

(factorial 3)

```
(if (= 3 1) 1 (* 3 (factorial 2)))  
    it (= 2 1) 1 (factorial 1)  
    it (= 1 1) 1 (factorial 0))
```

since if is a special form, it evaluates only one of its parts. if we redefine it and send the function the 2 parts, it will try to evaluate both parts before calling on the function if. we will get to an infinite loop.

forever