

*COMP9332 Network Routing and Switching*  
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# *Geographic Routing*

# *Lecture overview*

*This lecture examines the concept of using location (geography) to route packets in wireless adhoc (multi-hop) networks.*

# *Topics to be covered*

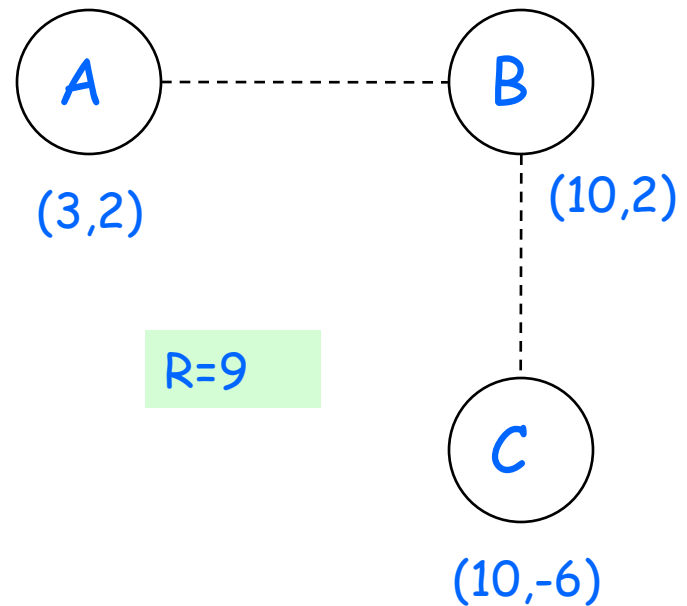
- Graph representation of multi-hop wireless networks
- Position beaconing and neighbourhood discovery
- Progress-based next-hop selection
- Greedy Perimeter Stateless Routing
  - Voids
  - Void traversal using right hand rule

# *Routing Options for Adhoc Networks*

- Table-driven (high overhead)
- No precomputing of routing tables
  - On-demand route computation (e.g. AODV protocol later)
  - Geographic or position-based routing (this lecture)

# Graph representation for multi-hop wireless networks - unit graphs

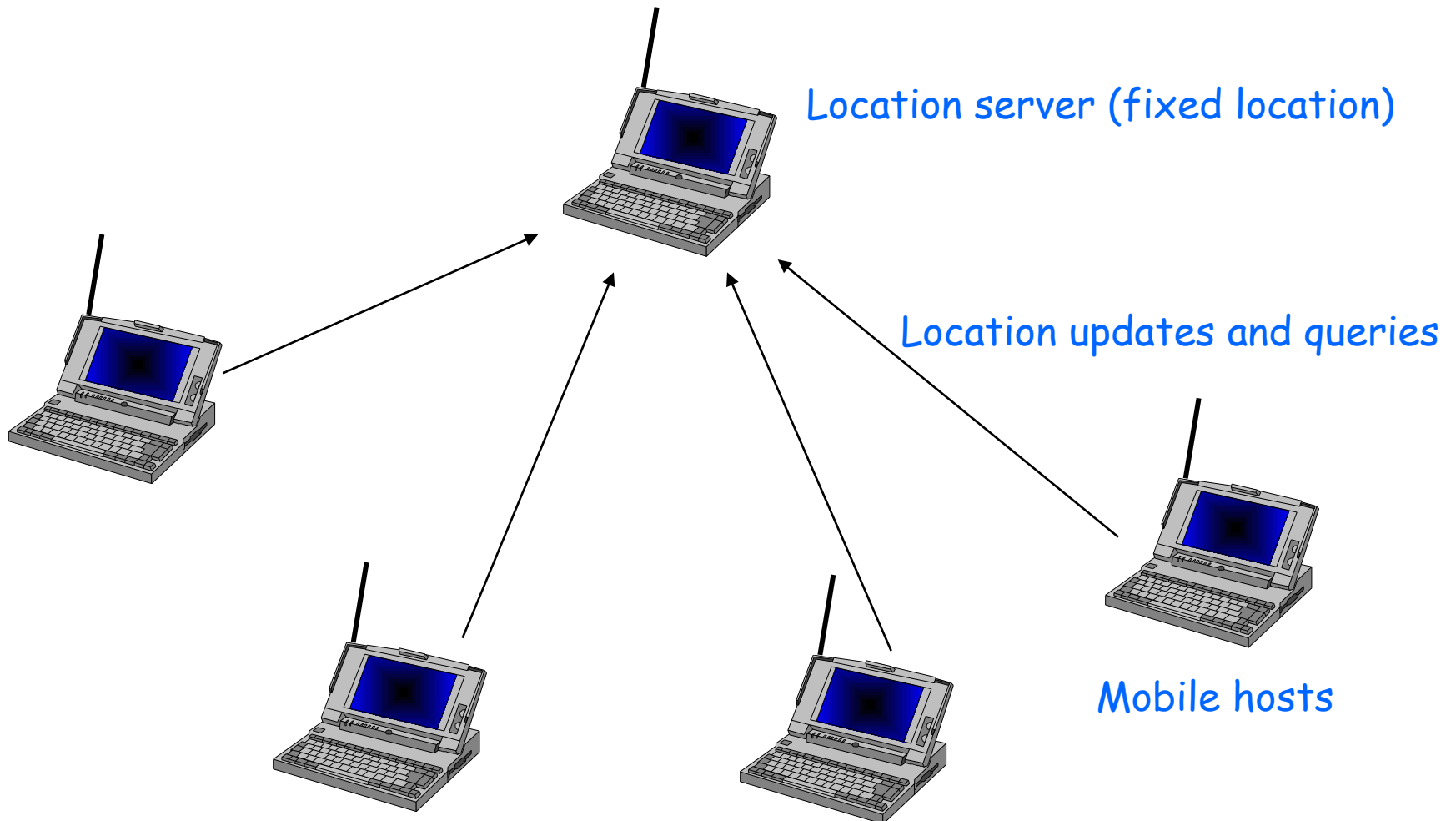
- All nodes have identical wireless communication range of  $R$
- An edge exists between two nodes iff Euclidean distance  $\leq R$



# *Geographic routing*

- Nodes know their location (e.g. using GPS)
- Nodes know location of one-hop neighbours
  - neighbours within radio range are one-hop away
  - every node periodically broadcast their positions
- Nodes know the location of the destination
  - using some location service, or
  - publicly known fixed servers
  - when they receive a packet from that node
- Packet header contains location information
- Routing is based on location, not IP address
  - next hop is decided based on nodes own position, position of one-hop neighbours, and position of the destination

# *Location Service*



- S wants to send to D
- S must select a next hop from A,B,C,E,F
- Which one should node S select as next hop?

## *Geographic routing selecting next hop*

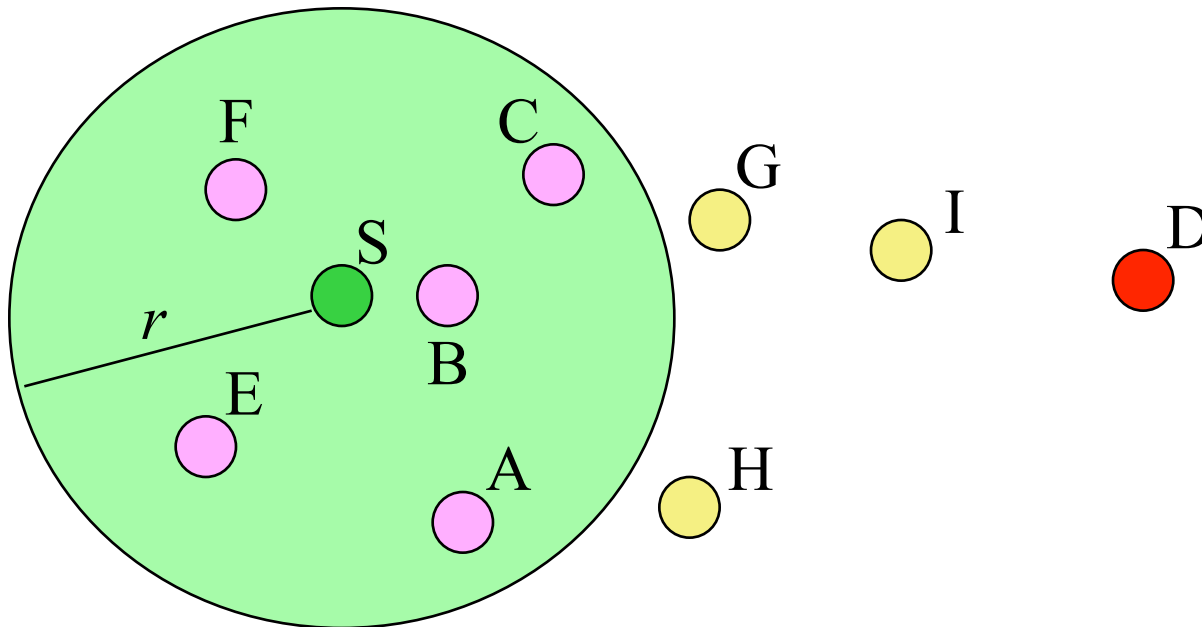


Fig 1: S (source) has a radio range of  $r$ . A, B, C, E and F are within radio range of S



# *Concept of Progress*

- *S selects neighbour A as next hop to send a packet to D*
- *Progress =  $d(S,D) - d(A,D)$*
- *Progress can be positive or negative*

# *Progress-based next-hop selection*

- Three approaches
  - Random progress
  - Most forward within radius
  - Greedy forwarding

# *Progress-based next-hop selection*

## *random progress*

- Equal probability for all positive-progress neighbours
- K neighbours that would make positive progress
- Probability of any of these k neighbours to be selected as the next hop =  $1/k$
- Exercise: In Fig 1, what are the probabilities of nodes A,B,C,E, and F to be selected as next hop?
- Packet is dropped if no neighbour with positive progress, even if a route exists!

# *Progress-based next-hop selection*

*most forward within radius*

- Pick the neighbour with maximum progress
  - Don't worry about positive or negative progress
- Exercise: In Fig 1, what are the probabilities of nodes A,B,C,E, and F to be selected as next hop?
- Packet is not necessarily dropped if no neighbour with positive progress (improves chances of reaching the destination at the risk of creating loops)

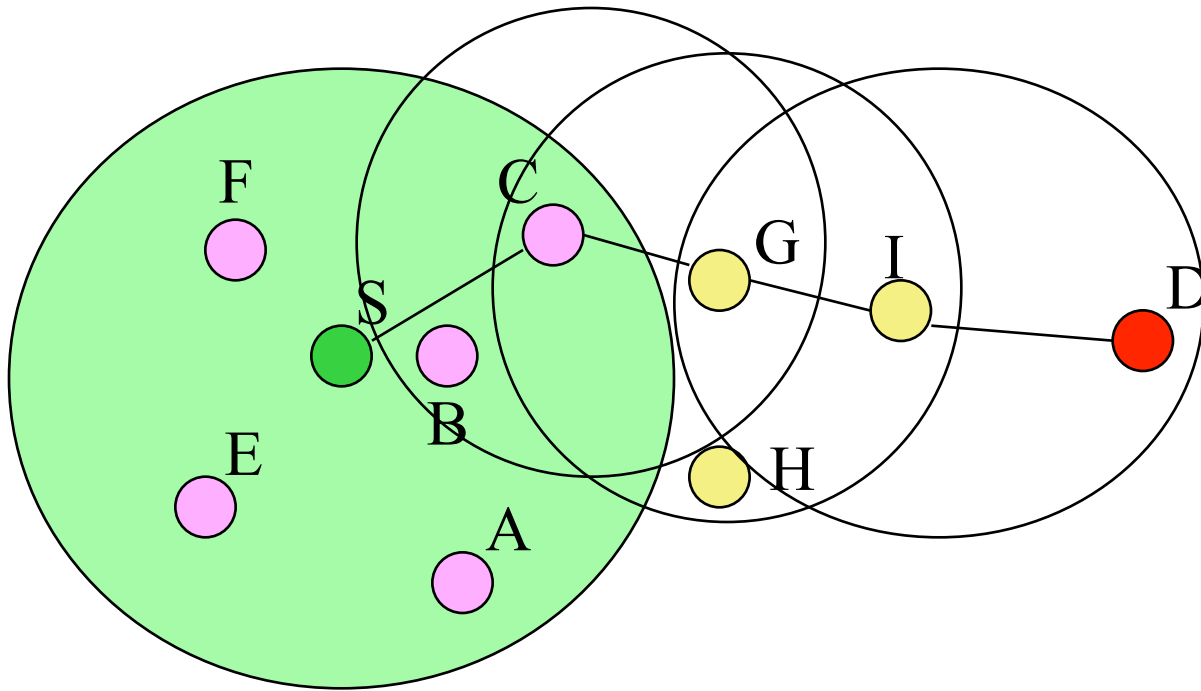
# *Progress-based next-hop selection*

## *greedy forwarding*

- Pick the neighbour with maximum positive progress
- Exercise: In Fig 1, what are the probabilities of nodes A,B,C,E, and F to be selected as next hop?
- Packet would be dropped if no neighbour with positive progress
- Greedy because it wants to reach the destination as quickly as possible (contrast it with *random progress*)

# Greedy Forwarding

*try to minimise number of hops*



- **S computes distance from each neighbour to the destination** (Note: S knows its position, position of D, and position of its neighbours, A,B,C,E,F)
- **Chooses the neighbour which is closest to the destination as the next hop** (next hop must make positive *geographic progress* towards destination to avoid loop)
- **S forwards to C. C forwards to G. G forwards to I.**
- **Finally I forwards to destination D.**

# *Benefits of Geography*

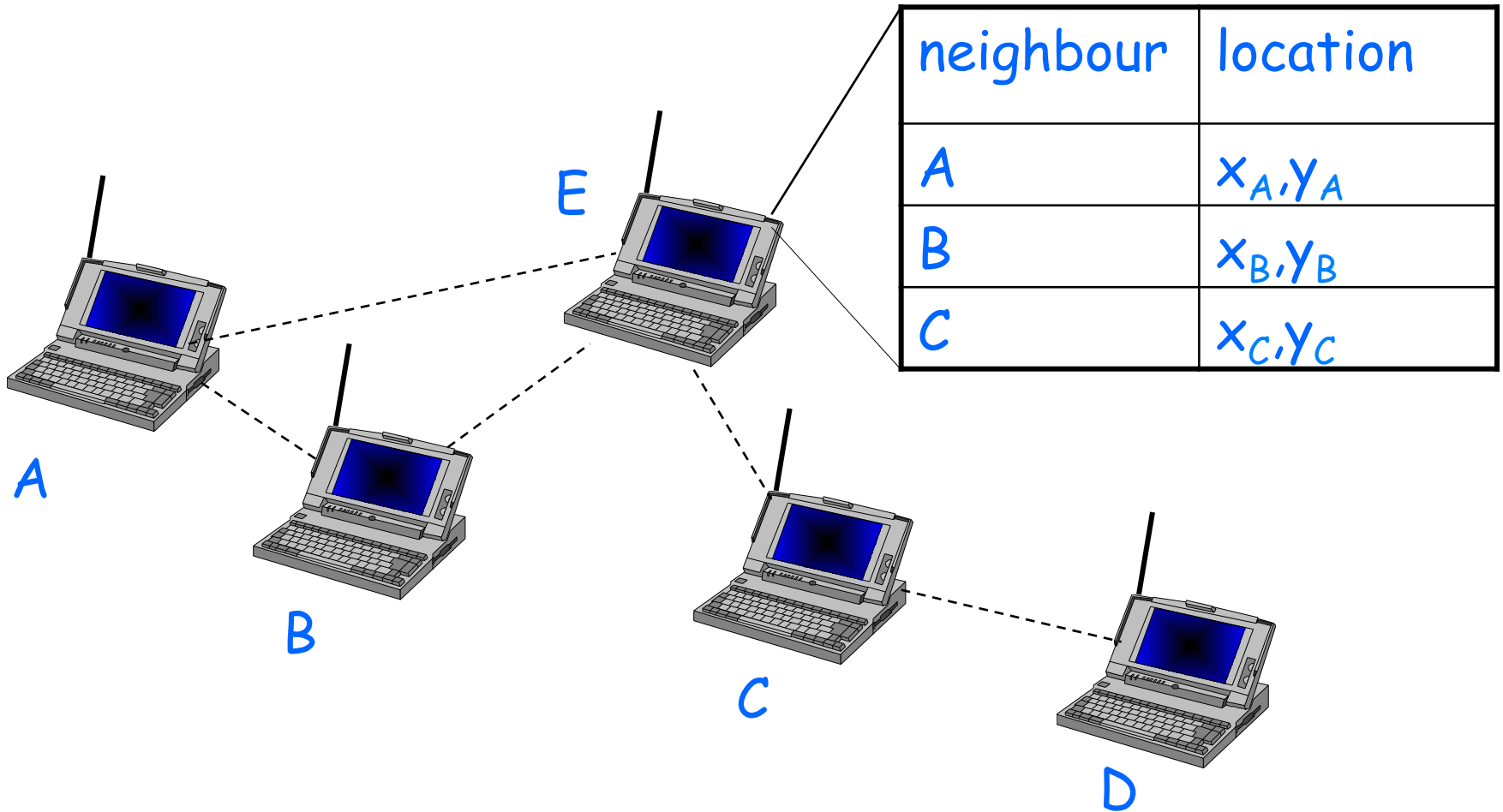
- Self-describing
- As node density increases, shortest path tends toward Euclidean straight line between source and destination
- Node's state concerns only one-hop neighbors:
  - Low per-node state:  $O(\text{density})$
  - Low routing protocol overhead: state (location) pushed only one hop

# *Position beaconing*

- **Control packet:** Nodes periodically broadcast (single-hop) their positions in special beacon (*control*) packets
- **Piggybacking:** Nodes piggyback position information in usual *data* packets
  - Other nodes hearing the data packet learns the position of the sender



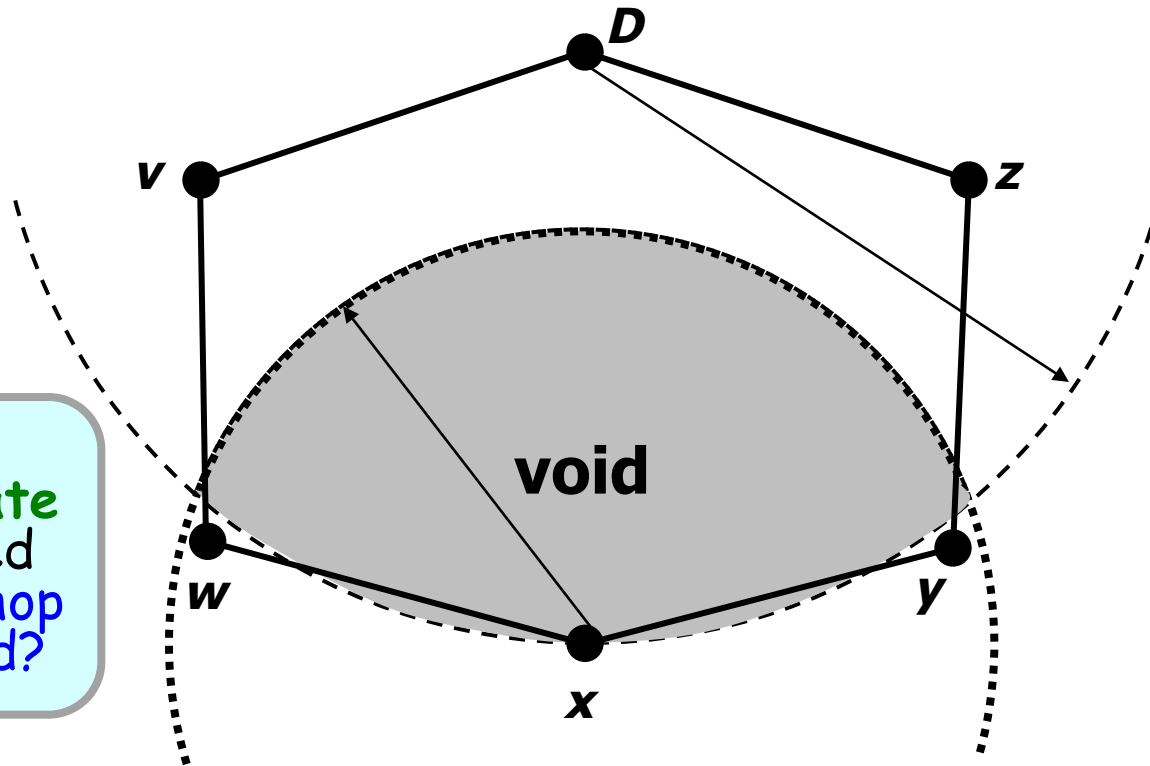
# Neighbour table



# Voids

## Failure of Greedy Forwarding

Greedy forwarding not always possible! Consider:



How can we  
**circumnavigate**  
voids? ...based  
only on one-hop  
neighborhood?

$x$  is not allowed to forward to  $y$  or  $w$ , because it would not make geographic progress towards  $D$

# *Greedy Perimeter Stateless Routing (GPSR)*

# GPSR

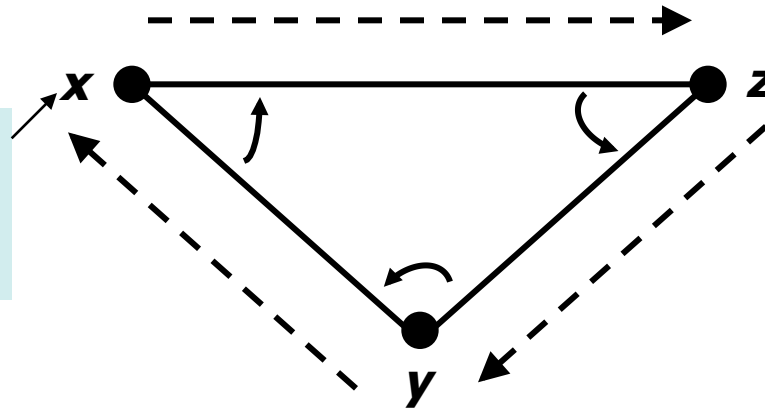
- A greedy-forwarding-based routing protocol
- Ability to traverse voids when possible (perimeter routing mode)
- Switches between two modes
  - Greedy mode (when greedy forwarding is possible)
  - Perimeter mode (when faced with a void)

# Void Traversal

## The Right-hand Rule

- Well-known graph traversal: *right-hand rule*
- Requires only neighbors' positions

X receives a packet from y, and forwards it to z

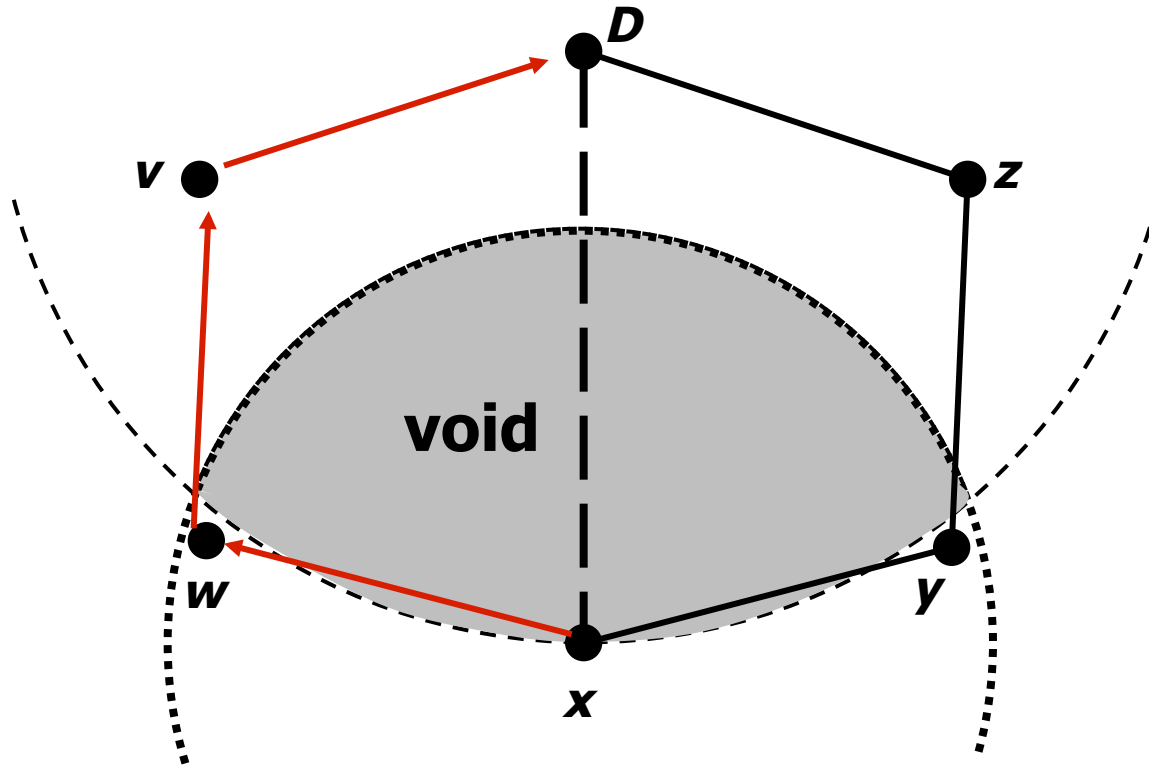


**Right-hand rule:** when arriving at  $x$  from  $y$ , the next edge traversed is the next one *sequentially counterclockwise* about  $x$  from edge  $x-y$ .

# *More on Right-hand Rule*

- Traverses a closed polygon in clock-wise order
- If the void is an interior of a polygon, the packet will move through the perimeter of the void to the nodes that are closer to the destination

# Traversing void using right-hand-rule



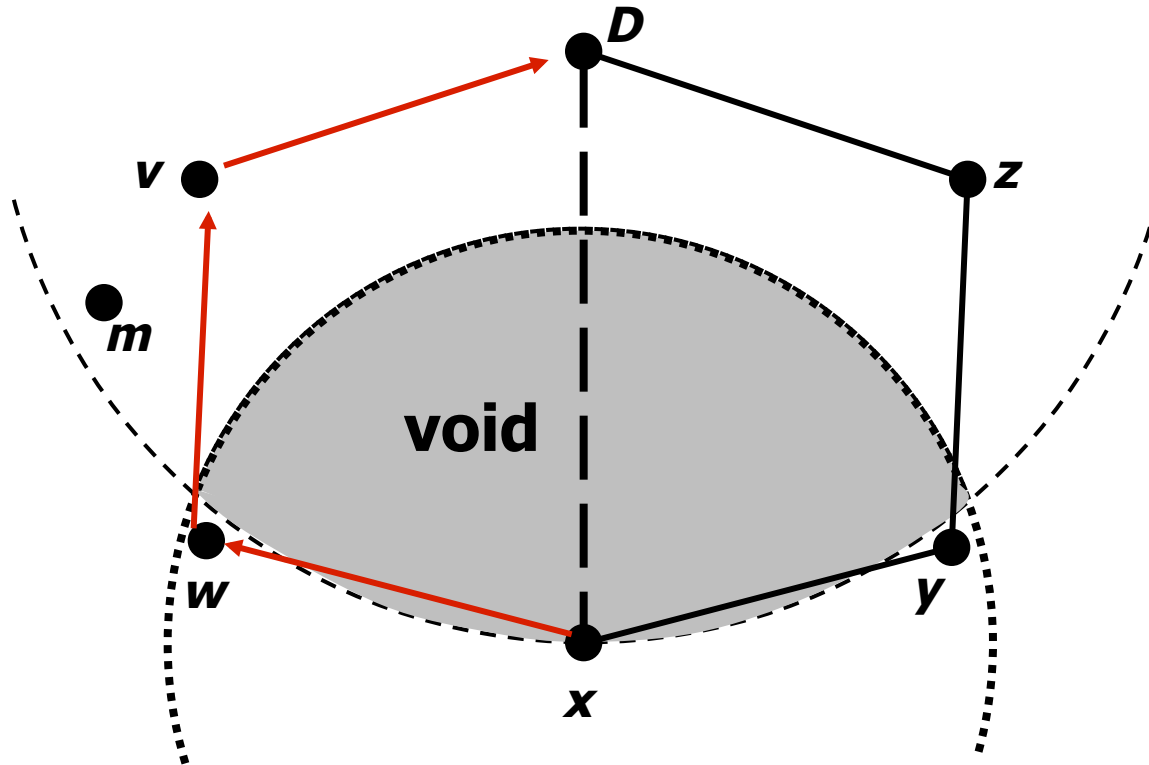
- First step: forward to first edge counterclockwise about  $x$  from line  $xD$  ( $x$  forwards to  $w$ )
- Then after apply right-hand-rule ( $w \rightarrow v \rightarrow D$ )

# *GPSR Packet Header*

| Field | Function                                |
|-------|---|
| D     | Destination Location                    |
| Lp    | Location packet entered perimeter mode  |
| Lf    | Point on xV packet entered current face |
| e0    | First edge traversed on current face    |
| M     | Packet Mode: Greedy or Perimeter        |



# Traversing void using right-hand-rule



- Q. Could  $w$  forward to  $m$ ?
- A. No.

# *Which mode to select?*

## *Perimeter or Greedy?*

- Default is *greedy*
- Select *perimeter* when no positive progress towards destination is possible
  - Stay in perimeter mode until positive progress is possible again
- Switch back to *greedy* if positive progress is possible again

# *Bibliography*

- Karp and Kung, “GPSR: Greedy perimeter stateless routing for wireless networks,” ACM MOBICOM 2000
- Position Based Routing Algorithms For Ad Hoc Networks: A Taxonomy (2001)by Silvia Giordano , Ivan Stojmenovic , Ljubica Blazevic, Ad Hoc Wireless Networking