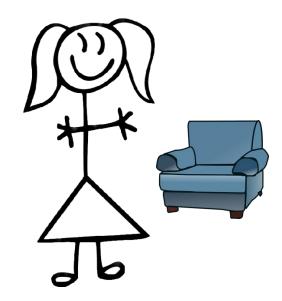


Selling to Multiple No Regret Buyers

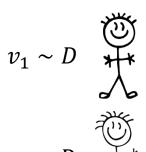
Linda Cai, S. Matthew Weinberg, Evan Wildenhain, Shirley Zhang

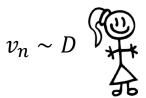
Princeton University

One seller with an item

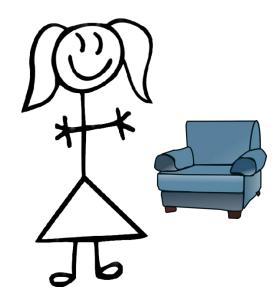


n independent and identical buyers

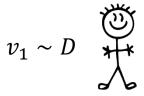


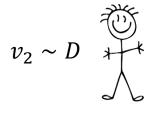


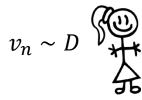
One seller with an item



n independent and identical buyers

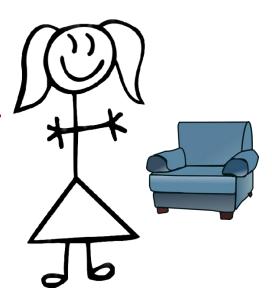




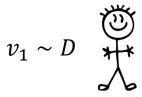


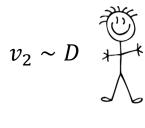
One seller with an item

Request Bid



n independent and identical buyers



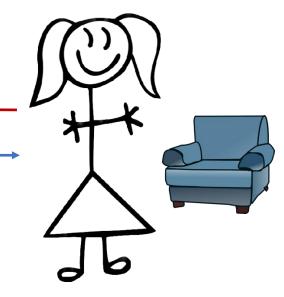


 $v_n \sim D$

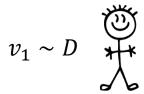
One seller with an item

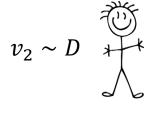


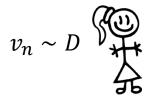
Submit Bid



n independent and identical buyers





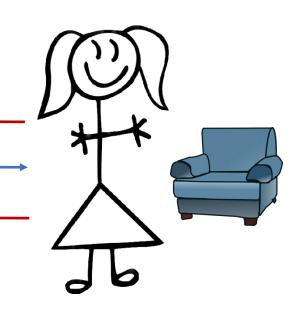


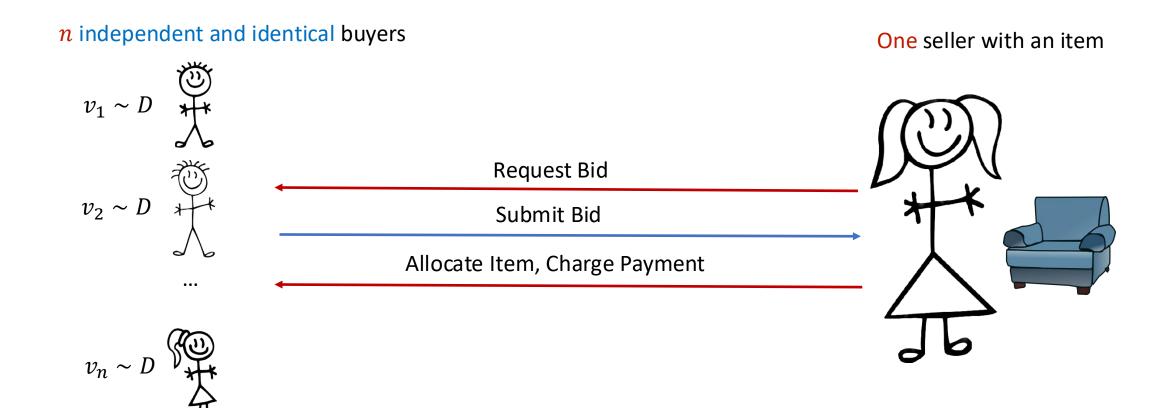
Request Bid

Submit Bid

Allocate Item, Charge Payment

One seller with an item

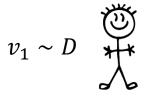


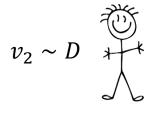


Goal: maximize their expected utility

= value of item * Pr[receiving the item] – expected payment

n independent and identical buyers



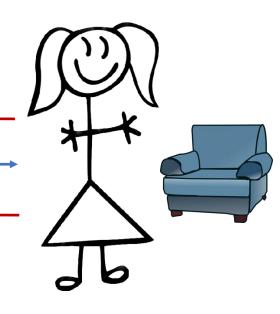


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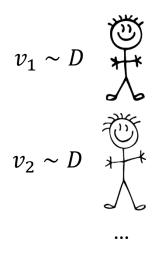


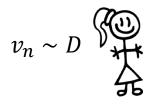
Goal: maximize expected revenue (payment received)

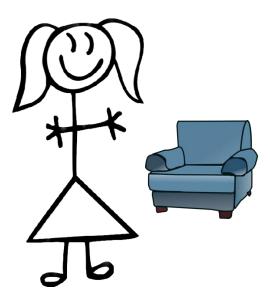
 $v_n \sim D$

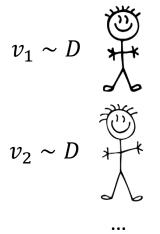
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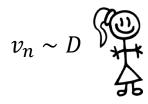
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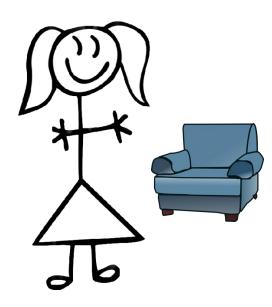


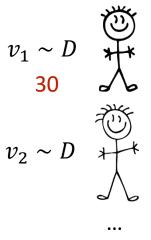


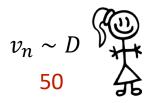


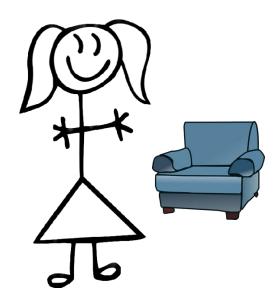


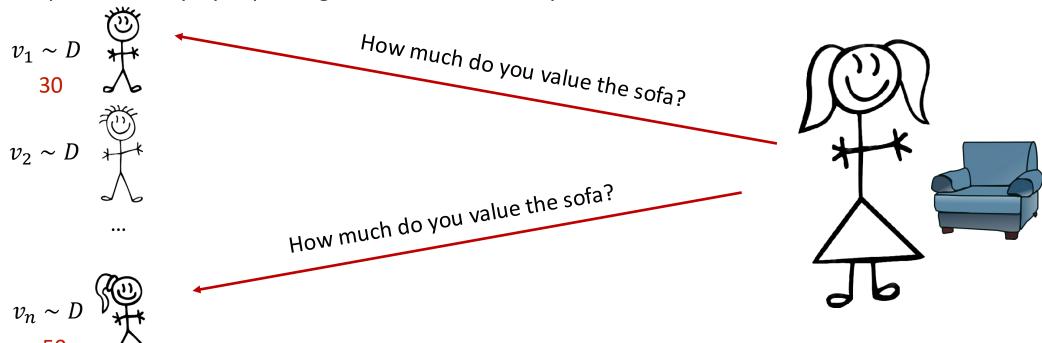


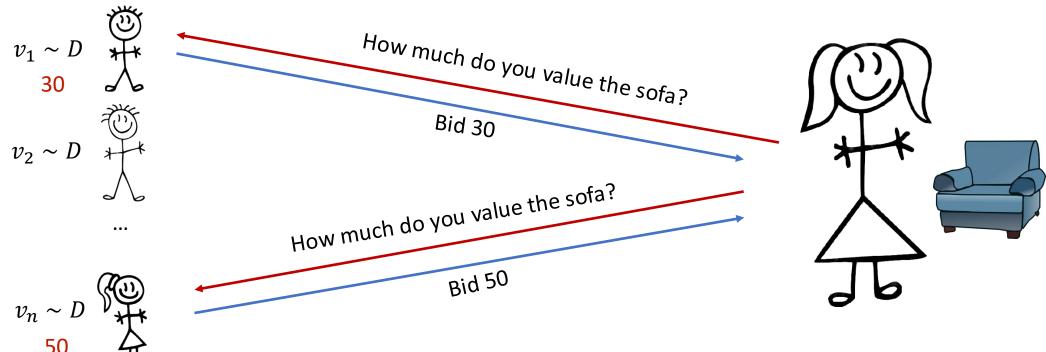




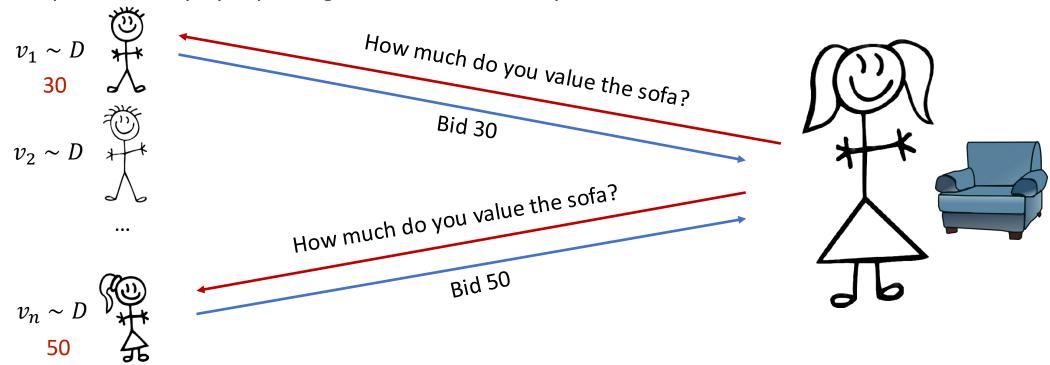








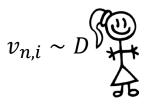
One Shot Auction: classical Bayesian auction design study *truthful* auctions: all buyers can maximize their expected utility by reporting their value truthfully

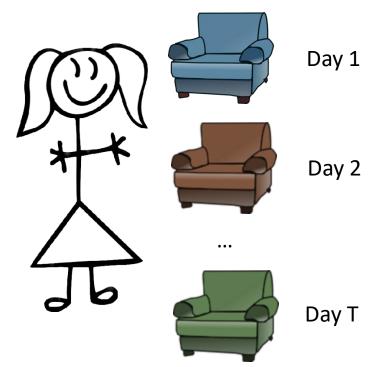


Myerson gives a complete characterization of revenue optimal truthful auction

Repeated Auction: there are T rounds, in each round the buyers' values are drawn independently from the prior distribution

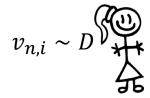
$$v_{1,i} \sim D$$





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Day 1

Day 1

Day 2

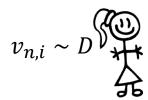
Day T

Goal: maximize their expected total utility

= $\sum_{i=1}^{T}$ (value of item in round i)* Pr[receiving the item in round i] – expected payment in round i

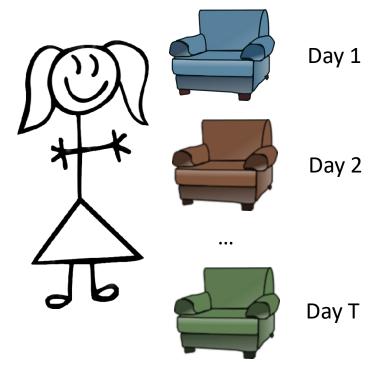
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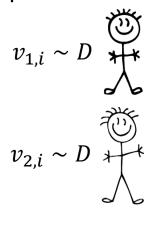
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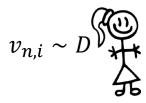
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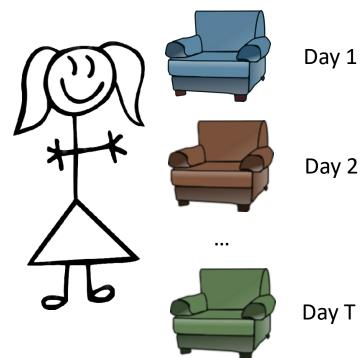


Goal: maximize expected **total** revenue (payment received)

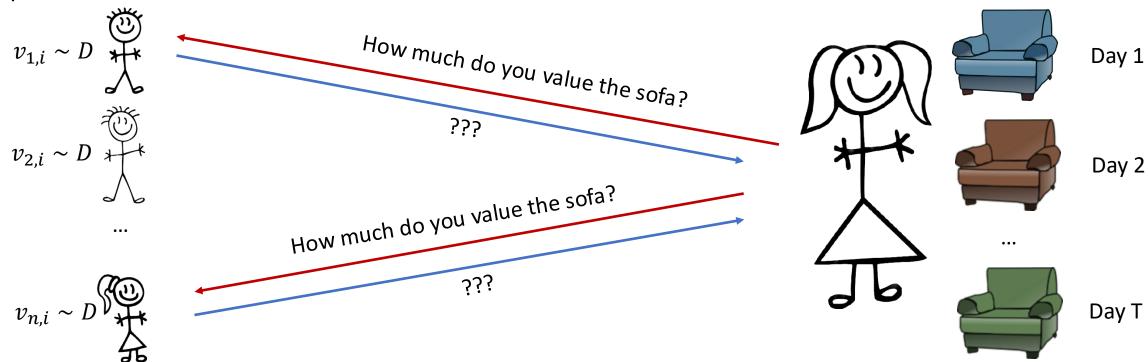
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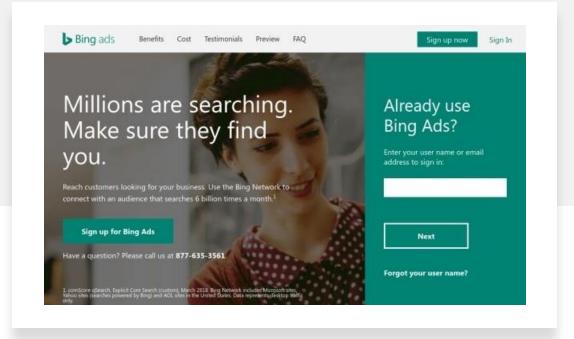




Repeated Auction: there are T rounds, in each round the buyers' values are drawn independently from the prior distribution

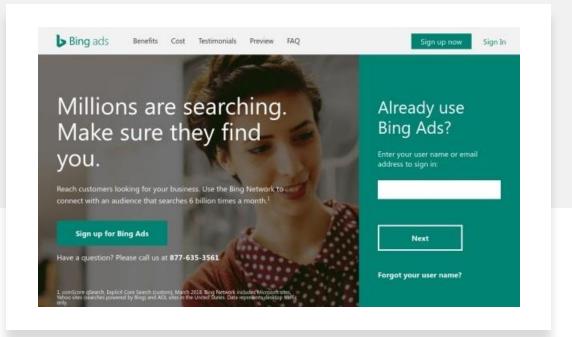


Question: not clear how to incentivize truthfulness, how do we model buyer behavior?



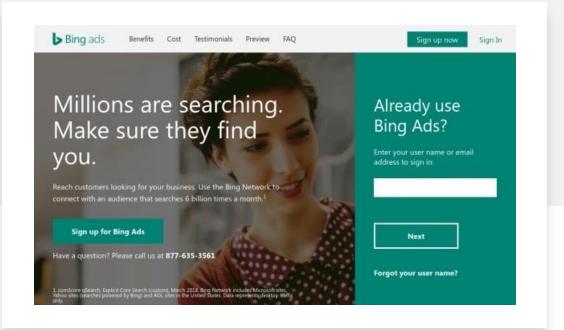


 In classical auction theory, agents are typically assumed to be rational/best responding





- In classical auction theory, agents are typically assumed to be rational/best responding
- [Nekipelov, Syrgkanis, and Tardos 2015] Bidder behavior is more consistent with no-regret learning in repeated auctions (e.g. in Bing ad auction)









Action m





In a repeated game with T rounds with context space V and action space A:

• In each round t, player i observes a context $v_{i,t} \in V$, then takes an action $a_{i,t} \in A$ and observe a reward $r_t(a_{i,t})$ (the reward function may be correlated with the context)



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Action 1 Action 2 Action 3 Action m

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$$\sum_{t=1}^{T} r_t(a_{i,t}) \ge \max_{fixed\ action\ strategy\ s} \sum_{t=1}^{T} r_t(\text{action}\ s\ \text{takes}\ \text{given}\ \text{context}\ v_{i,t}) - o(T)$$

Given a repeated auction designed by the seller:

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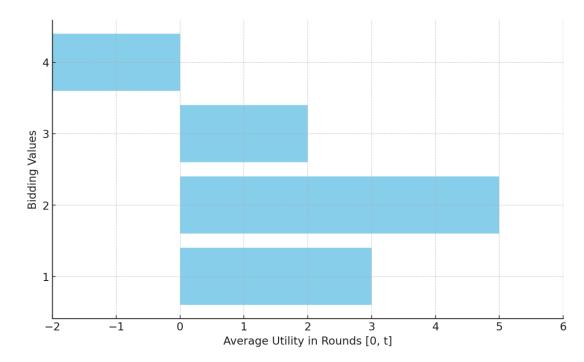
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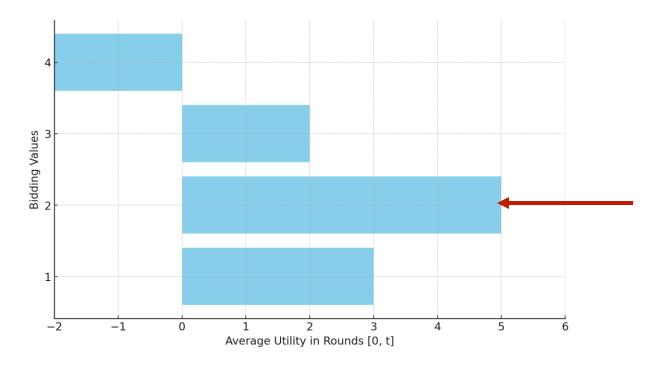
 Intuition: if an action seems much better historically it would be played with much higher probability

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A mean-based no regret algorithm will choose to bid 2 with high probability

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- 1) Allocates random costs to each expert at the initial time;
- 2) Subsequently, at each time step, it selects the expert with the lowest cumulative cost, inclusive of the initial step.

Once the difference between cumulative utility of two actions exceed the initial cost, the inferior action will never be chosen.

Exploitability of Learning Agents

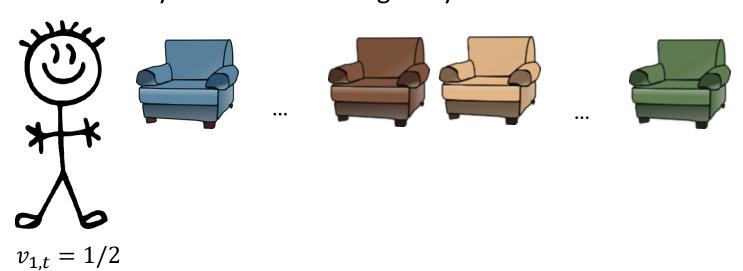
- Question: How much revenue can the seller extract knowing that the buyers are using mean-based no regret learning algorithms (but not the specific algorithm they are using)?
- [Braverman, Mao, Schneider, and Weinberg'18] For the one seller one buyer setting:
 - For naïve mean-based no regret buyer: Full surplus extraction (seller gets revenue ≈ optimal social welfare and buyer gets ≈ zero utility)
 - For non overbidding mean-based no regret buyer: Optimal revenue can be unboundedly worse than for naïve learners, yet still unboundedly better than a seller who runs the same auction in each round.

Exploitability of Learning Agents

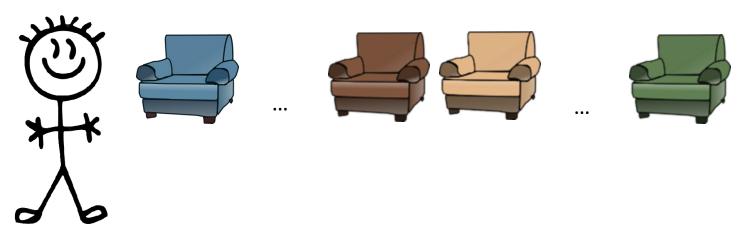
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 - For naïve mean-based no regret buyers: Full surplus extraction (seller gets revenue
 ≈ optimal social welfare and buyer gets ≈ zero utility)
 - For non-overbidding mean-based no regret buyers: much harder to understand seller's optimal revenue than the single buyer setting (formal barriers)

Idea: lure the buyer into bidding high by offering the item for cheap at first, then start charging the item with a high price. It will take a while for the mean-based no regret buyer to realize they should not bid high anymore.

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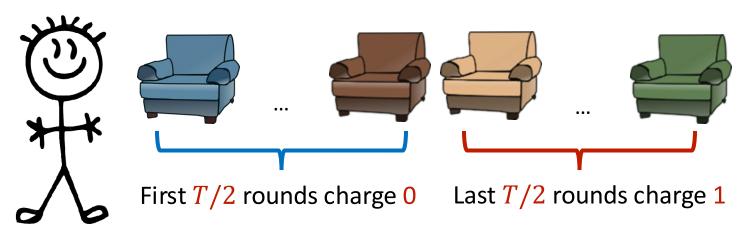


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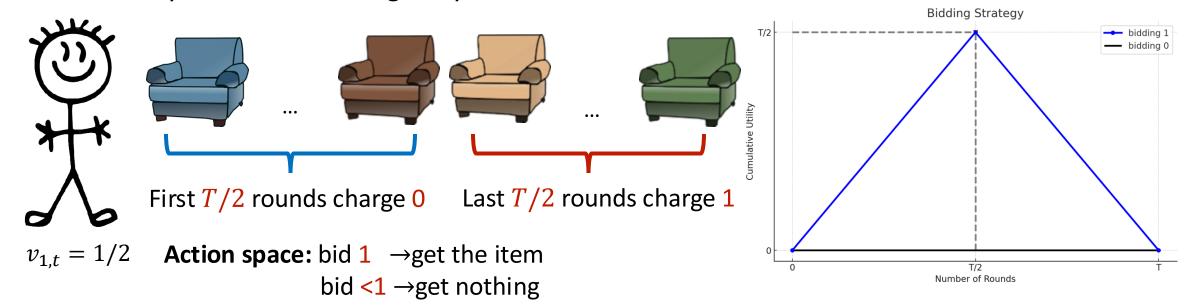
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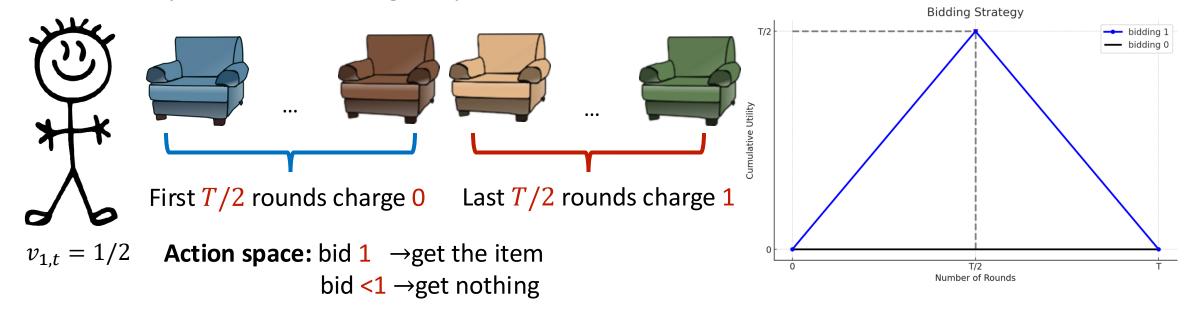


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- Naïve mean-based no regret buyer: bid 1 for all but o(T) rounds, average utility ≈ 0
- Seller revenue : 1 o(1)

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Our Idea:

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- In each phase, the seller lures the buyer with the highest value into bidding high an intended value by offering the item for a fair price during first half of the phase, then start charging surcharge in the second half of the phase.
- Once the no-regret buyers learn to switch to other actions than their intended value at the end of the phase, repeat above process for a different set of intended values.

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In a Phase with R rounds:

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In a Phase with R rounds:

Buyer value to bid mapping:

$$v_1 \rightarrow b_1$$

$$v_2 \rightarrow b_2$$

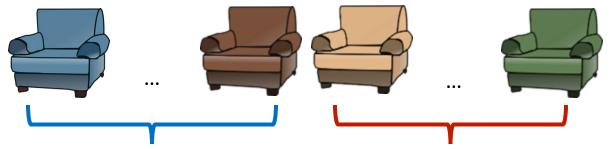
• • •

$$v_m \to b_m$$

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In a Phase with R rounds:



First R/2 rounds, run second price auction on underlying value based on bids

Last R/2 rounds, run second price auction with surcharge

Buyer value to bid mapping:

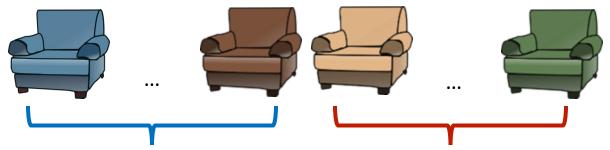
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$$v_m \to b_m$$

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Challenge for the seller: when buyers with different value for the item choose the same bid, the seller cannot treat them differently

In a Phase with R rounds:



First R/2 rounds, run second price auction on underlying value based on bids

bids $(b_1, b_2) \Rightarrow$ charge v_1 buyer with v_2

Last R/2 rounds, run second price auction with surcharge

Buyer value to bid mapping:

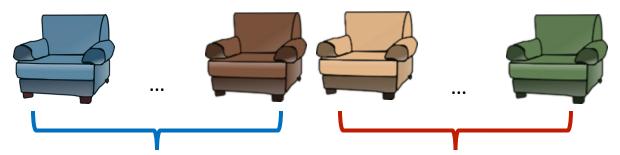
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- In the Ads auction setting, the sellers are typically bidding for multiple different keywords simultaneously. What would happen in the multiple item multiple buyer case?

THANK YOU!