

EECS 2011: Assignment 2

Due Date - Monday, Mar 27, in class!!!

Question 1 Property of Binary Trees

[15 points]

Assume a proper and complete binary tree (T) with $n > 1$ nodes. Let $E(T)$ represent the sum of the depths of all external nodes in T (see Figure 1), and $I(T)$ represent the sum of the depths of all internal nodes in T. Prove that:

- (a) $E(T) = O(n \cdot \log(n))$
- (b) $E(T) - I(T) = 2i$, where i represents the actual number of internal nodes in T.

(Hint: for part (b) use mathematical induction.)

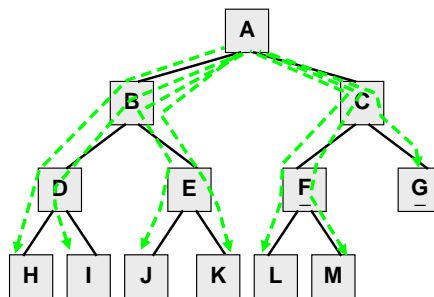


Figure 1 External Path Length

Solution:

Proof:

- (a) Based on the properties of proper binary trees, discussed in class:

$$E \leq e \cdot h \leq ((n+1)/2) \cdot \log(n) = O(n \cdot \log(n))$$

- (b) Assume $E - I = 2i$ holds for some arbitrary i .

Now consider case $i_{\text{new}}=i+1$. We should prove: $E_{\text{new}} - I_{\text{new}} = 2 \cdot i_{\text{new}} = 2 \cdot (i+1)$.

New tree with $i_{\text{new}}+1$ internal nodes is created by removing an external node from the tree, and replacing it with an internal node with two attached external nodes.

Hence $e_{\text{new}} = (e-1) + 2 = e+1 = i + 2$.

Suppose the removed external node is at depth d . Accordingly:

$$I_{\text{new}} = I + d$$

$$E_{\text{new}} = E - d + 2 \cdot (d+1) = E + d + 2.$$

Subsequently:

$$E_{\text{new}} - I_{\text{new}} = (E + d + 2) - (I + d) = (E - I) + 2 = (2i) + 2 = 2(i+1) = 2 \cdot i_{\text{new}}$$

Question 2 Binary Search Tree

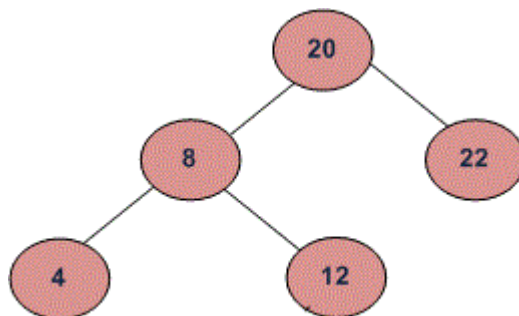
[15 points]

Given a BST and two numbers – a and b (where $a \leq b$) – propose an algorithm that prints all the elements k of the BST that satisfy: $k \geq a$ and $k \leq b$. Your algorithm should return the found keys in ascending order. You can assume that there are no duplicate keys in the tree. (You can also assume that a , b , and all the keys stored in the BST are integers.)

In addition to writing a brief description of the algorithm, you also need to implement the algorithm in Java (i.e., provide respective Java code), as well as justify the algorithm's running time.

Solution:

Suppose in this question, for the tree given below, if $a = 10$ and $b = 22$, your algorithm should print 12, 20 and 22.



One possible solution for this problem would be to do **inorder traversal** on the given tree (starting from the root). Continue on the left-subtree only if the current key $> a$, print the current node's key if the key is between a and b , continue on the right-subtree only if the current key $< b$.

The running time of this algorithm, in the worst case, would be $O(n)$ in the worst case – which is the running time of the inorder traversal on the tree

```
public void binaryInorderPrintKeys(Tree<E> T, Position<E> v, int a, int b) {  
    E e = v.element();  
    if (T.hasLeft(v) && e > a) binaryInorderPrintKeys(T, T.leftChild(v), a, b);  
    if (e >= a && e <= b) { System.out.println(e+" "); }  
    if (T.hasRight(v) && e < b) binaryInorderPrintKeys(T, T.rightChild(v), a, b);  
}
```