

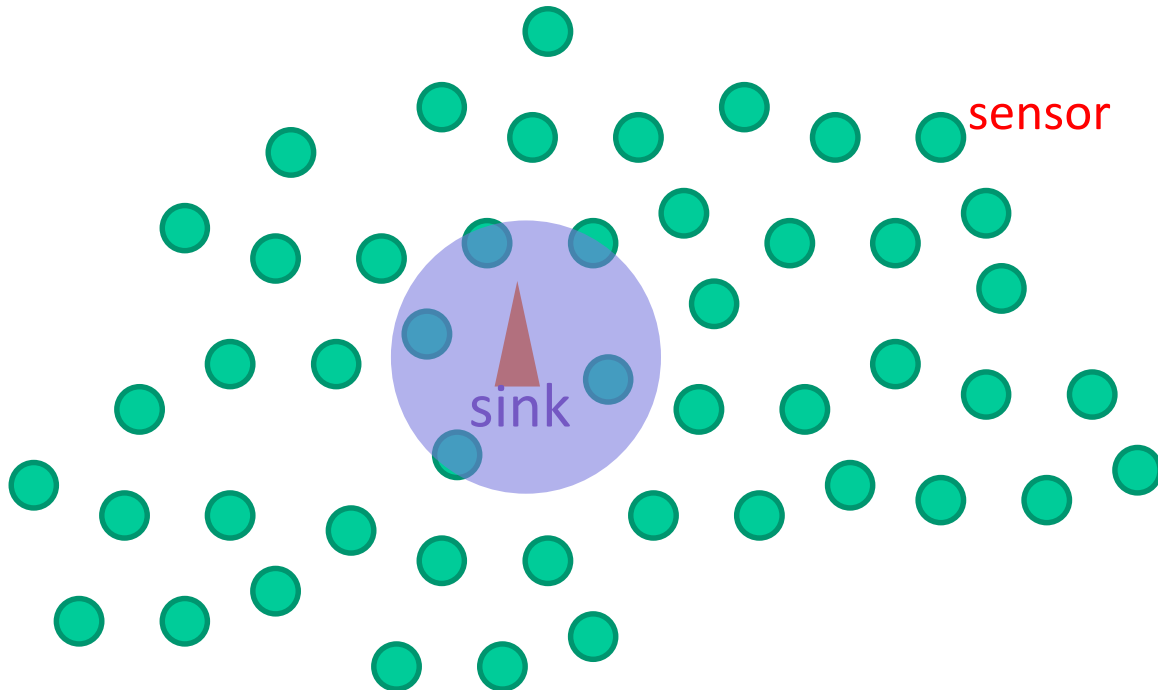
# Broadcasting Techniques for Mobile Ad Hoc Networks

**Broadcasting:** It is a process in which **one** node sends a packet to **all** other nodes in a network.

Broadcast: 1 → all communication

# Applications of broadcasting

- Broadcasting of net-wide **control information**
  - Ex.: The **sink** broadcasts its **location** and **ID** to all sensors.
- **Provides route establishment functionality**
  - The sink may locate a mobile sensor.
  - Serves as a **building block** in ad hoc networking

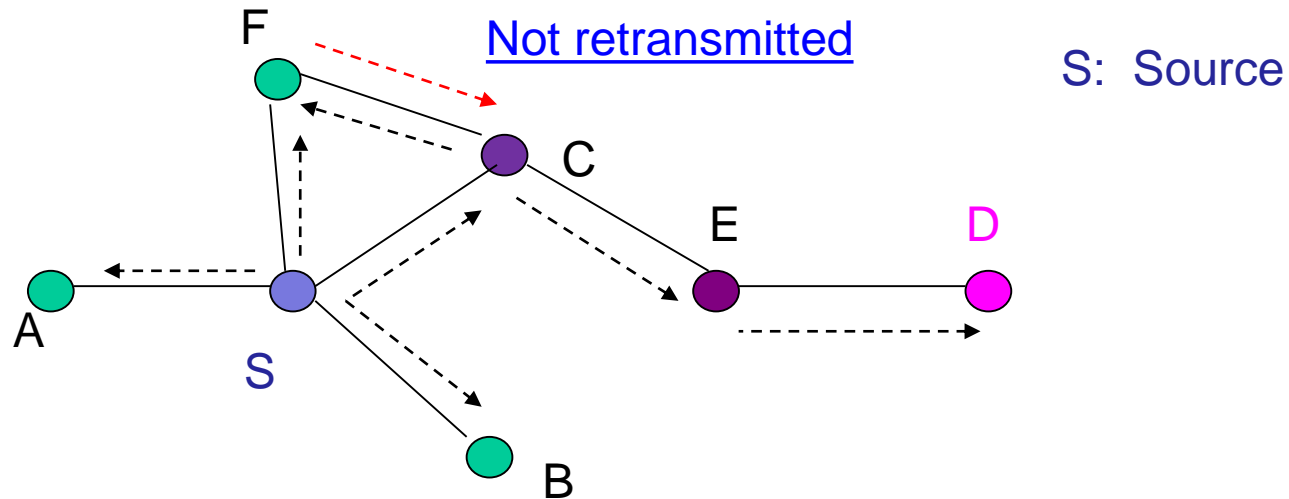


# A simple form of broadcasting: flooding

- Flooding:
  - Protocol:
    - The source **transmits** its packet to all its neighbors.
    - Each receiving node **retransmits** each received unique packet exactly once.
  - **Data structure:** Each node maintains and uses a **table** with the following structure:

Source node	Most recent <b>sequence number</b> of packet from the source
-------------	--

# Flooding



In a network with  $N$  nodes, there are **about**  $N$  (re)transmissions.

# Efficient broadcasting

- **Goal:** Minimize the number of retransmissions.
  - ➔ Some nodes do **not** retransmit.
- (Attempt to) ensure that a broadcast packet is delivered to each node in the network.

$$\text{Delivery ratio} = \frac{\text{\# of nodes receiving a copy of the packet}}{\text{Total \# of nodes in the network}}$$

Delay (aka Latency) = How long does it take for the last node to receive a copy of the packet?

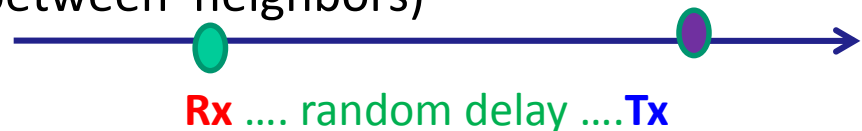
# Common attributes

- MAC layer assumption

- Example: 802.11 MAC with no RTS/CTS

- Jitter the scheduling of a broadcast packet:

- Schedule it for retransmission after a random delay (reduces the probability of collision between neighbors)



- Random Access Delay

- After a node receives a new packet, it **waits** for a while to receive the same packet from other neighbors. Based on this info, the node decides to broadcast or not.



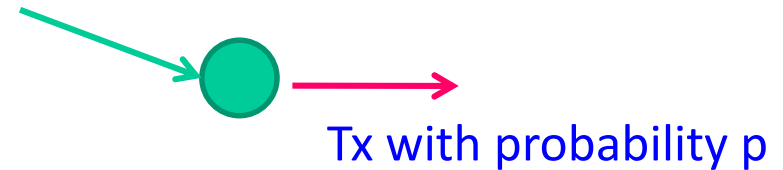
- A packet is not transmitted more than once: prevent loop

# Classes of efficient flooding algorithms

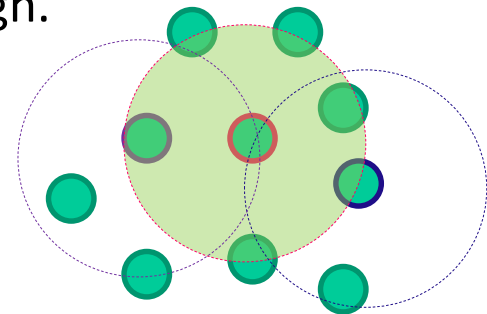
- Probability based methods
  - » Probabilistic scheme, counter-based scheme
- Area based methods
  - » Distance based
- Neighbor knowledge methods
  - » Flooding with self pruning
  - » Scalable broadcast
  - » Dominant pruning
  - » Multipoint relaying

# Probabilistic scheme (Prob. Based)

- This works similar to **flooding**.
- Nodes **re**broadcast with a **predetermined probability**.



- Impacts of network **density**
  - **Dense networks:** multiple nodes share similar transmission coverage areas. (→ some nodes need not retransmit.)
  - **Sparse networks:** there is much less shared coverage. The **probability** parameter needs to be high.





# Counter-based scheme (Prob. Based)

- **Observation**

- There exists an **inverse relationship** between the number of times a packet is received at a node and the probability of the node reaching additional area.

$$\text{Prob}(\text{reaching additional area}) = k \cdot 1 / \# \text{of copies received}$$

- **Protocol:** (when a **new** packet is received)

- **Initialize:** Counter = 1 and RAD = random(0, Tmax)
- **Decision based on RAD**
  - » RAD not expired: **Counter++** for each redundant packet received.
  - » RAD expires: **If Counter < Threshold**, **rebroadcast** the packet;  
**else**, drop the packet

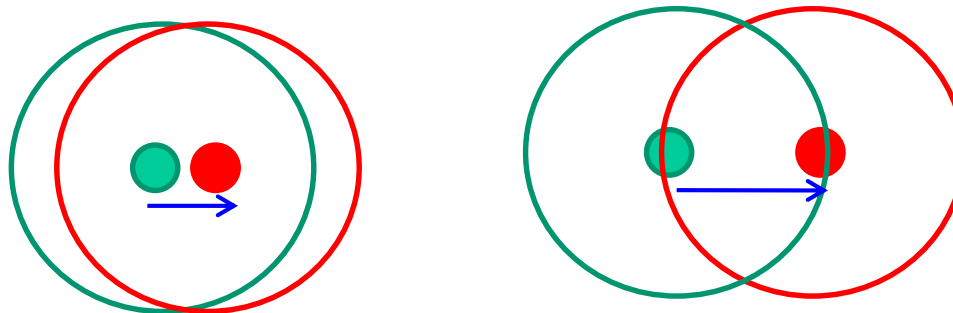
- **Advantages:** Simplicity, adaptability to local topology

# Area based methods

## – Observation

- If a receiving node is **very close** to the sender
  - The **additional area** covered by the receiver is **very small**.
- If a receiving node is **farthest** from the sender, the **additional area** covered by the receiver is **large** (61%).

- **Idea:** A node can evaluate **additional coverage area** based on all received redundant packets.



## Distance-based scheme

- Protocol (when a new packet is received)
  - Initiate a RAD.
  - Decisions based on RAD
    - RAD has not expired: Cache all redundant packets.
    - RAD expires: examine the locations of all transmitters of the redundant packets:
      - » If the distance between this node and any node (from above) is less than a threshold, drop the packet.
      - » Else, retransmit the packet.

# Neighbor knowledge methods

- Flooding with self-pruning
- Scalable broadcast algorithm (SBA)
- Dominant pruning
- Multipoint relaying
- Ad hoc broadcast protocol

# Flooding with self pruning

- Self pruning → a receiving node decides whether or not to further send the packet ...
- Protocol (Each node knows its 1-hop neighbors)
  - Sender: Includes its 1-hop neighbors in packet header before (re)transmitting.
  - Receiver: (when a new packet is received)
    - » Compare own neighbor list with the sender's.
    - » If the receiving node would not reach any additional nodes, drop the packet. Else, retransmit.

# Scalable Broadcast Algorithm (SBA)

- Protocol

- **Sender (say, A):** (Re)transmits a packet (no neighbor info)

- **Receiver (say, B):** A **new** packet is received

- » **B knows** all its neighbors common to **A** that have received the packet. **← via “hello” packets.**

- » **B starts** a RAD.

- For each redundant packet, **find** if B can reach more nodes by retransmitting the packet.

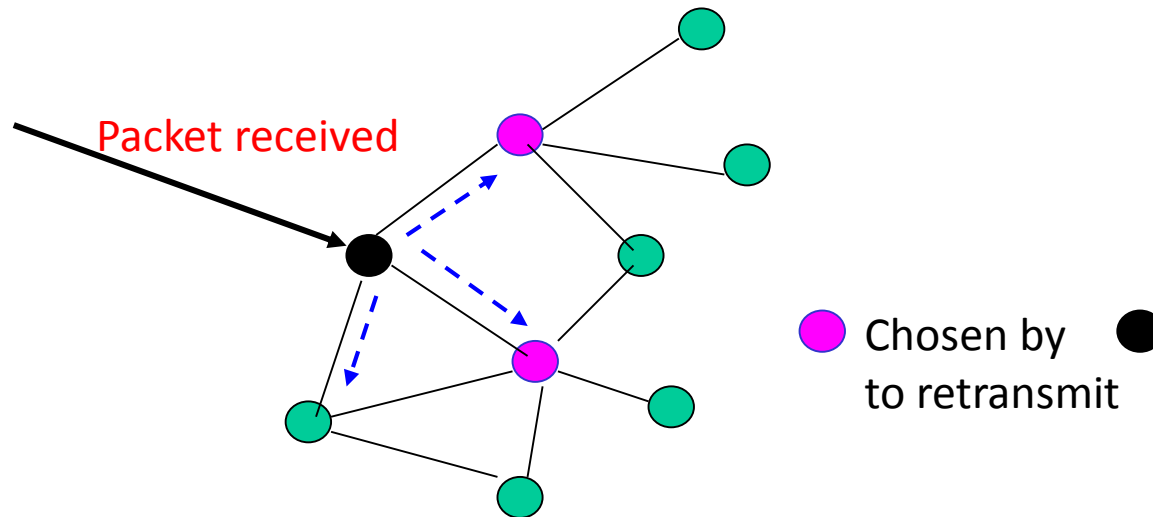
- **RAD expires:** Retransmit, if B can reach more nodes.

# SBA (contd.)

- Dynamic adjustment of RAD
  - RAD is proportional to  $(d_{N_{\max}}/d_{me})$ , where
    - »  $d_{N_{\max}}$  is the  $\max$  neighbor degree among all neighbors.
    - »  $d_{me}$  is the node's neighbor degree.
- Idea: Nodes with most neighbors usually broadcast before the others.

# Dominant Pruning

- **Assumption:** Nodes know their 2-hop neighbors.
- (Unlike SBA) **Rebroadcasting nodes proactively choose some/all** of their 1-hop neighbors as **rebroadcasting nodes**.





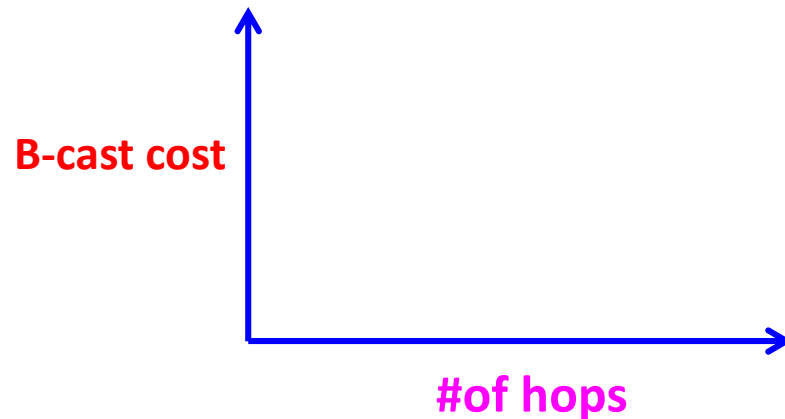
# Dominant Pruning (Contd.)

- Protocol: when a (new) packet is received
  - Check if the header contains the address of this node:
    - » No: Do not retransmit
    - » Yes: - Compute the subset of neighbors that should retransmit.
      - Include those nodes in the packet header and retransmit the packet.
- Identification of neighbors who will retransmit (Greedy Set Cover alg.)
  - Identify some 1-hop neighbors such that all (not just remaining) 2-hop neighbors will be covered.

## Parameterized neighborhood based flooding for ad hoc wireless networks

Vijay Dheap, M. A. Munawar, K. Naik, P. A. S. Ward

Military Communications Conference, 2003. MILCOM 2003. IEEE, Volume: 2



# Multipoint Relaying

- **Similar** to Dominant Pruning
  - Forwarding nodes are chosen by **upstream** senders. Forwarding nodes are called *multipoint relays* (MRPs).
- **Different** from Dominant Pruning
  - MRPs are included in “Hello” packets.
  - “Hello” packets trigger packet retransmissions.
- Computation of MRPs
  1. **Initialize**: MRP = {one 1-hop neighbor}.
  2. **Determine** the 2-hop neighbors that are reachable via the MRP set.
  3. **Identify** a **new 1-hop neighbor** that will cover **most** of the uncovered 2-hop neighbor, and add **it** to MRP.
  4. **Repeat** steps 2-3 **until** all 2-hop neighbors are covered.

# Multipoint Relaying (Contd.)

- Protocol

- Nodes transmit “Hello(MRP)” packets from time to time.
- When a node receives a “Hello(MRP)” packet, it checks if it is a member of the MRP set.
  - » **Yes:** the node retransmits all data packets received from that source.
  - » **No:** Do not forward packets.
- **Note:** frequency of “Hello” packets affect delay performance.

# Performance Comparison

- Effect of a “null” (perfect) MAC
  - To know the inherent strength of the protocol
- Effect of network congestion
  - Packet size
  - Network size
  - Source rate

(With 802.11, static network, and fixed # of nodes.)
- Effect of mobility (null MAC, fixed # of nodes)
  - Random waypoint mobility model with zero pause
    - » Nodes choose a random location and move there with a given mean speed. Next, choose another location and another speed to reach there, ...
  - Constrained mobility ← more natural

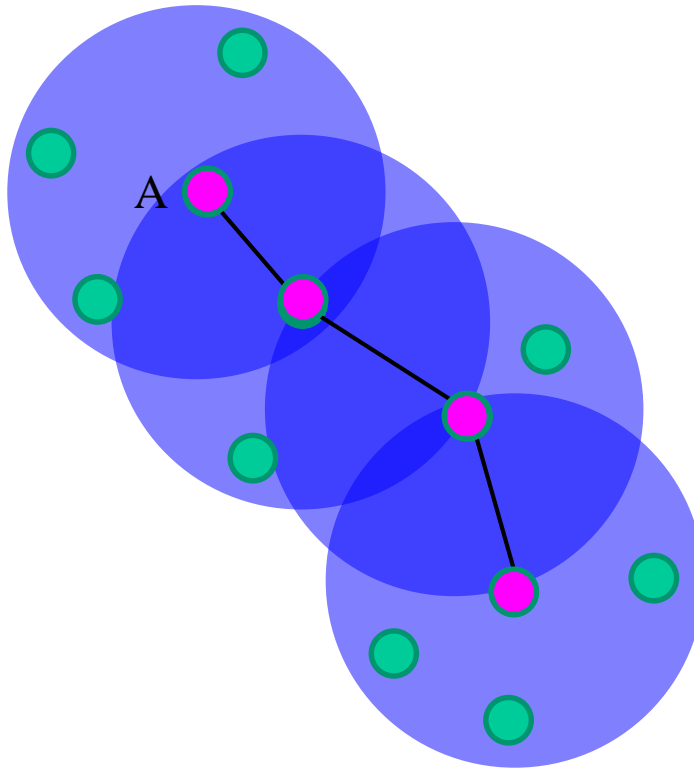
# Algorithm efficiency (Null MAC)

- **Delivery ratio:** The % of network nodes who receive a given broadcast packet.
  - » Dense network: All achieve 100%.
  - » Sparse networks: Simple flooding and protocols that utilize neighbor knowledge perform better.
- **Number of retransmitting nodes vs. Node density**
  - Benchmarks
    - » Worst: Flooding (every node retransmits)
    - » Best (theoretical): Min. Connected Dominating Set (MCDS)

An MCDS is the smallest set of retransmitting nodes such that the set of nodes are connected and non-MCDS nodes are within 1-hop of at least one MCDS node.

# Min Connected Dominating Set

Compute MCSD for A



MCSD for A = {  }

# Algorithm efficiency (Contd.)

- Number of **retransmitting nodes** vs. **Node density**
  - Benchmarks (Contd.)

**MCDS** < Neighbor knowledge < Area based < Prob. based < **Flood**



# Effect of congestion (802.11)

- **Delivery ratio** degrades with higher packet rate.
  - Increased interference and no ACK → more packet loss
- Protocols that minimize the number of retransmissions have better delivery ratio.
- **Delay:** The time it takes for the *last node* to receive a broadcast packet. ← **Latency**
  - Delay increases (exponentially?) with broadcast rate.
    - » Worst: Flooding
    - » Best: MRP
  - Proposal: The concept of RAD → higher delivery ratio

# Effect of mobility (Null MAC)

- **Delivery ratio**
  - **MRP:** Worst
  - **Others:** Near 100%