Lab 0: Design Doc

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```

Preliminaries

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· Reference:

https://blog.csdn.net/u011011827/article/details/125656962?ops_request_misc=
{"request_id":"167772676616800188575949","scm":"20140713.130102334.pc_all."}&request_id=167772676616800188575949&
task-code-2~all~first_rank_ecpm_v1~rank_v31_ecpm-1-125656962-0-nullnull.142^v73^insert_down1,201^v4^add_ask,239^v2^insert_chatgpt&utm_term=shell中实现backspace

https://en.wikipedia.org/wiki/BIOS_interrupt_call#Interrupt_table

Booting Pintos

A1: Put the screenshot of Pintos running example here.

As the screenshot shows, the first booting is in QEMU and the second in Bochs.

Debugging

QUESTIONS: BIOS

B1: What is the first instruction that gets executed?

ljmp \$0x3630,\$0xf000e05b. As shown in the screenshot.

B2: At which physical address is this instruction located?

The PA is <code>0xfff0</code>.

As we know, PA = VA - PHYS_BASE, and PHYS_BASE is a multiple of 0x100000000. As the screenshot shows, VA of the first instruction is 0xf000fff0. So we can easily induce that PHYS_BASE = 0xf00000000, PA = 0xfff0

Ref:

- https://pkuflyingpig.gitbook.io/pintos/appendix/reference-guide/loading#physical-memory-map
- https://pkuflyingpig.gitbook.io/pintos/appendix/reference-guide/virtual-addresses#work-with-mapping-kernel-vm-one-to-one-to-pm

QUESTIONS: BOOTLOADER

B3: How does the bootloader read disk sectors? In particular, what BIOS interrupt is used?

First, the bootloader (beginning at 0x7c00) checks all the sectors on different disks one by one.

Relevant functions.

When checking a specific sector, it uses int 0x13 to call Interrupt 0x13 (Line 346, 0x7d30). The number in register %ah is 0x42 (Line 342, 0x7d2c). Looking up the table we know that the BIOS interrupt is "*Extended Read Sectors*".

```
read_sector:
        7d1f:
        7d20: 29 c0
                                        %eax,%eax
                          # LBA sector number [48:63]
        7d22: 50
                                  push %eax
        push %ax
                         # LBA sector number [32:47]
        push %ebx
                         # LBA sector number [0:31]
        7d24: 66 53
                          # Buffer segment
                         # Buffer offset (always 0)
        push %ax
7d27: 50
        push %ax
                         # Number of sectors to read
        push $1
        7d28: 6a 01
                                    push $0x1
                         # Packet size
        push $16
        7d2a: 6a 10
                                   push $0x10
                          # Extended read
        mov $0x42, %ah
                                  mov $0x42,%ah
        7d2c: b4 42
        mov %sp, %si
                           # DS:SI -> packet
        7d2e: 89 e6
                                   mov %esp,%esi
345
        int $0x13
                         # Error code in CF
        7d30: cd 13
                                   int $0x13
                          # Pop 16 bytes, preserve flags
        7d32:
```

	UCII	зеек то эреспіей паск
13h	0Dh	Reset Fixed Disk Controller
	15h	Get Drive Type
	16h	Get Floppy Drive Media Change Status
	17h	Set Disk Type
	18h	Set Floppy Drive Media Type
	41h	Extended Disk Drive (EDD) Installation Check
	42h	Extended Read Sectors
	43h	Extended Write Sectors
	44h	Extended Verify Sectors

https://en.wikipedia.org/wiki/BIOS_interrupt_call#Interrupt_table

B4: How does the bootloader decides whether it successfully finds the Pintos kernel?

The bootloader checks if a partition meets three criteria:

- · An used partition?
- · A Pintos kernel partition?
- · A bootable partition?

If yes, this is the bootable Pintos kernel. And we jump to load_kernel to load the kernel.

```
check partition:
        # Is it an unused partition?
        cmpl $0, %es:(%si)
        7c47: 26 66 83 3c 00 74
                                      cmpw $0x74,%es:(%eax,%eax,1)
         je next_partition
        7c4d: 10 e8
        # Print [1-4].
        call putc
        7c4f: ae
7c50: 00 26
                                      scas %es:(%edi),%al
                                             %ah,(%esi)
        # Is it a Pintos kernel partition?
        cmpb $0x20, %es:4(%si)
                                      cmpb $0x75,0x20(%esp,%eax,1)
        7c52: 80 7c 04 20 75
        jne next_partition
                                      push %es
        7c57: 06
        # Is it a bootable partition?
        cmpb $0x80, %es:(%si)
        7c58: 26 80 3c 80 74
                                            $0x74,%es:(%eax,%eax,4)
100
         je load kernel
         7c5d:
               20
                                       .byte 0x20
```

B5: What happens when the bootloader could not find the Pintos kernel?

If bootloader could not find the kernel, it calls *interrupt 18h* by the instruction INT 6x18. This interrupt outputs an error message telling BIOS that the bootloader fails to find the kernel.

18h

Execute Cassette BASIC: On IBM machines up to the early PS/2 line, this interrupt would start the ROM Cassette BASIC. Clones did not have this feature and different machines/BIOSes would perform a variety of different actions if INT 18h was executed, most commonly an error message stating that no bootable disk was present. Modern machines would attempt to boot from a network through this interrupt. On modern machines this interrupt will be treated by the BIOS as a signal from the bootloader that it failed to complete its task. The BIOS can then take appropriate next steps.^[3]

https://en.wikipedia.org/wiki/BIOS interrupt call#Interrupt table

B6: At what point and how exactly does the bootloader transfer control to the Pintos kernel?

Ljmp *start (Line 225), the bootloader jumps to the start point of the kernel.

The bootloader saves the start address in position "start" (which is initially part of the loader's code).

```
mov $0x2000, %ax
7cbf: b8 00 20 8e c0
                                    $0xc08e2000, %eax
mov %ax, %es
mov %es:0x18, %dx
7cc4: 26 8b 16
                                    %es:(%esi),%edx
7cc7: 18 00
                                    %al,(%eax)
mov %dx, start
7cc9: 89 16
                                    %edx,(%esi)
7ccb: d7
                              xlat %ds:(%ebx)
7ccc: 7c c7
                                    7c95 <load_kernel+0x17>
movw $0x2000, start + 2
7cce: 06
                              push %es
7ccf: d9 7c 00 20
                              fnstcw 0x20(%eax,%eax,1)
ljmp *start
7cd3: ff 2e
                              1jmp
                                    *(%esi)
       d7
7cd5:
                                    %ds:(%ebx)
```

QUESTIONS: KERNEL

B7: At the entry of pintos_init(), what is the value of expression init_page_dir[pd_no(ptov(0))] in hexadecimal format?

The answer is 0.

Use gdb to set a breakpoint in pintos_init. And use print to print the value of expression.

```
(gdb) b pintos_init
Breakpoint 1 at 0xc00202b6: file ../../threads/init.c, line 82.
(gdb) c
Continuing.
The target architecture is assumed to be i386
=> 0xc00202b6 <pintos_init>: push %ebp

Breakpoint 1, pintos_init () at ../../threads/init.c:82
(gdb) p init_page_dir[pd_no(ptov(0))]
=> 0xc000efef: int3
=> 0xc000efef: int3
$1 = 0
(gdb) [
```

▼ to delete

ptov(0) transfers PA 0 to VA, which is VA = PA + PHYS_BASE = 0xf00000000 . pd_no = VA>>PDSHIFT = VA>>(10+12) = d

```
static inline void *ptov (uintptr_t paddr)
{
   ASSERT ((void *) paddr < PHYS_BASE);
   return (void *) (paddr + PHYS_BASE);
} // transfer PA to VA

static inline uintptr_t pd_no (const void *va) {
   return (uintptr_t) va >> PDSHIFT;
} // Obtains page directory index from a virtual address
```

- B8: When palloc_get_page() is called for the first time,
- B8.1 what does the call stack look like?

Set a breakpoint in ${\tt palloc_get_page}$ and run until the breakpoint.

Issue bt to see the call stack, which includes start(), pintos_init(), paging_init() and current palloc_get_page().

```
Breakpoint 1, palloc_get_page (flags=(PAL_ASSERT | PAL_ZERO)) at ../../threads/palloc.com/
(gdb) bt

#0 palloc_get_page (flags=(PAL_ASSERT | PAL_ZERO)) at ../../threads/palloc.c:112

#1 0xc0020522 in paging_init () at ../../threads/init.c:219

#2 0xc002031b in pintos_init () at ../../threads/init.c:104

#3 0xc002013d in start () at ../../threads/start.S:180
```

B8.2 what is the return value in hexadecimal format?

Issue finish to run until the function finishes. The return value is <code>excololooo</code>, which is a void pointer.

B8.3 what is the value of expression init_page_dir[pd_no(ptov(0))] in hexadecimal format?

Use gdb print, and we see the value is 0.

```
(gdb) p init_page_dir[pd_no(ptov(0))]
=> 0xc000ef7f: int3
=> 0xc000ef7f: int3
$2 = 0
```

B9: When palloc_get_page() is called for the third time,

B9.1 what does the call stack look like?

```
After the breakpoint is set, issue c three times to get to the thrid calling of palloc_get_page().

Issue bt to see the call stack, which includes start(), pintos_init(), thread_start(), thread_create() and current palloc_get_page().
```

B9.2 what is the return value in hexadecimal format?

Issue fin, and we see the return value is <code>oxc01030000</code>, which is a void pointer.

```
(gdb) fin

Run till exit from #0 palloc_get_page (flags=PAL_ZERO) at ../../threads/palloc.c:112

=> 0xc0020bf9 <thread_create+55>: add $0x10,%esp

0xc0020bf9 in thread_create (name=0xc002ea51 "idle", priority=0, function=0xc0021028 <idle>, aux=0xc000efac) at ../../threads/thread.c:178

Value returned is $3 = (void *) 0xc0103000
```

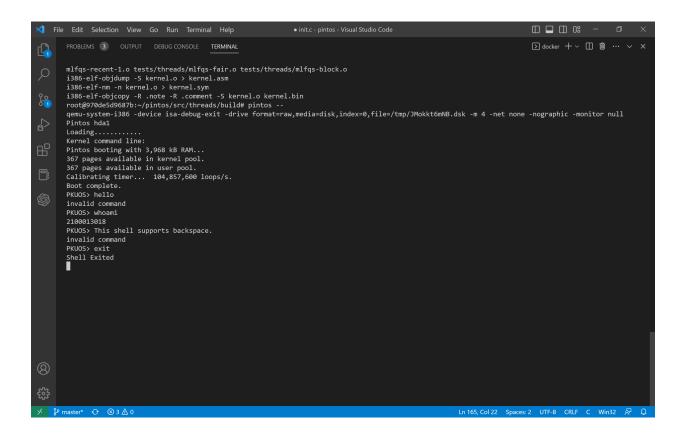
B9.3 what is the value of expression init_page_dir[pd_no(ptov(0))] in hexadecimal format?

Use gdb print to get the value, which is 1056807 in decimal format. Use p/x to get the hexadecimal value 0x102027.

```
(gdb) p init_page_dir[pd_no(ptov(0))]
=> 0xc000ef3f: int3
=> 0xc000ef3f: int3
$4 = 1056807
(gdb) p/x init_page_dir[pd_no(ptov(0))]
=> 0xc000ef3f: int3
=> 0xc000ef3f: int3
$5 = 0x102027
```

Kernel Monitor

C1: Put the screenshot of your kernel monitor running example here. (It should show how your kernel shell respond to whoami, exit, and other input.)



C2: Explain how you read and write to the console for the kernel monitor.

I use an infinite loop to print "PKUOS" and get command line input. If a new line is detected, we enter the next loop.

```
while(print_prompt){
    printf("%s", prompt);
    int idx=0;
    memset(cmdline, 0, MAXLINE);

// get cmdline
// get cmdline
// get cmdline
// parse cmdline
// parse cmdline
if(strcmp(cmdline, "whoami") == 0){...
else if(strcmp(cmdline, "exit") == 0){...
else{...
}
```

For each line, I write a while loop to get the chars one by one and:

- print it if printable
- · delete a char if it is a backspace
- overlook it if the length of a line has exceeded MAXLINE (a macro of max line length)
- break current loop and begin a new line if it is a vr (vn results in error)

Note: As for the backspace, I refer to the article below.

https://blog.csdn.net/u011011827/article/details/125656962?ops_request_misc=

 $\label{lem:code_2} $$ \{ $\ensuremath{"request_id"}: $\ensuremath{"167772676616800188575949"}, $\ensuremath{"scm"}: $\ensuremath{"20140713.130102334.pc_all."} \} $$ $$ $$ $\ensuremath{"request_id=167772676616800188575949\&biz$} $$ $$ $$ $\ensuremath{task-code-2-all-first_rank_ecpm_v1-rank_v31_ecpm-1-125656962-0-null-1-12566962-0-null-1-1256696860-0-null-1-125669660-0-null-1-125669660-0-null-1-125666660-0-null-1-125666660-0-null-1-125666$

null.142^v73^insert_down1,201^v4^add_ask,239^v2^insert_chatgpt&utm_term=shell中实现backspace

(Note: We wrote a similar shell in ICS Shelllab, so I refer to my own code ($\stackrel{\centerdot}{\bullet}\,\omega\,\stackrel{\centerdot}{\bullet}\,)y$)