68k Software Model

ECE511: Digital System & Microprocessor

What we are going to learn in this session:

- Introduction to M68k microprocessor.
- Software model of M68k:
  - M68k internal architecture:
    - Registers in M68k.
    - Their functions.
  - Programmer-side view:
    - What you will use as a programmer.

Introduction to M68k Microprocessor

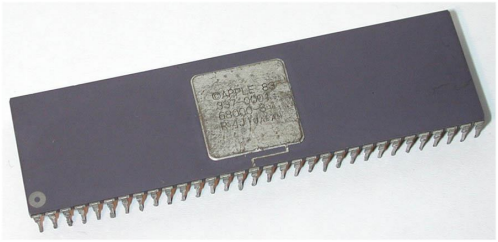
The M68000 Microprocessor

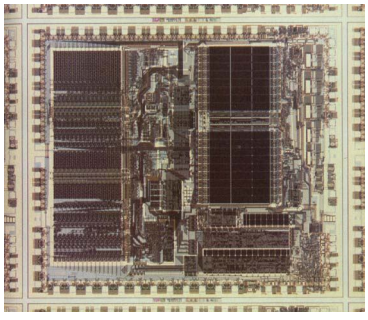
- M68000 (M68k) microprocessor.
- Manufactured by Motorola Semiconductors, 1979.
- 16-bit processor, but can perform 32-bit operations.
- Speed: 8-12 MHz.

The M68k Microprocessor

- Very advanced compared to 8-bit processors:
  - 16-bit data bus, 24-bit address bus.
  - Can execute instructions twice as fast.
- Still available today:
  - Simple, practical commands.
  - Robust: can be used for many applications.

The M68k Microprocessor

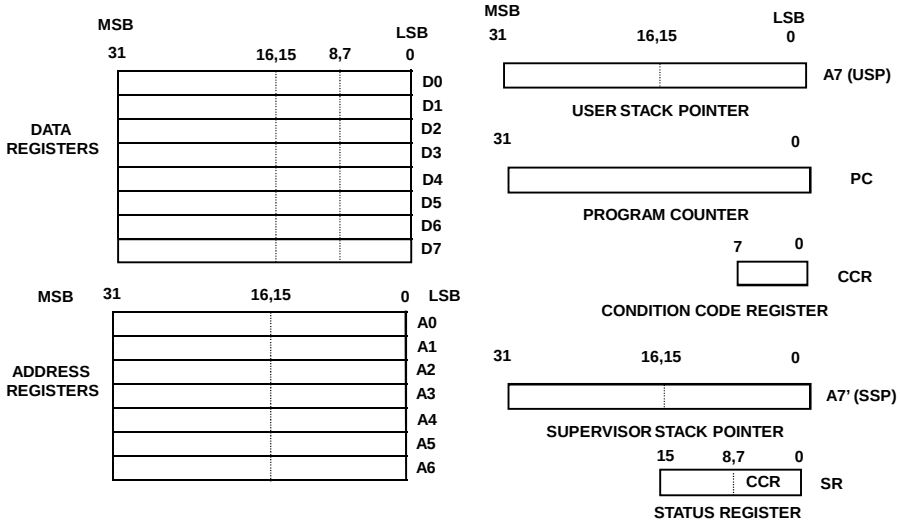




What's inside the M68k?

- Data Registers.
- Address Registers.
- Program Counter.
- Stack Pointer.
- Status Register.
- Arithmetic-Logic Unit (ALU).
- Instruction Register.

M68k Programmer Model



Data & Address Registers

Data Registers

- General-purpose registers:
  - Stores data/results for calculations.
  - High-speed “memory” inside M68k.
  - 8 registers (D0 to D7).
  - 32-bits long.

Data Registers

- Able to process:
  - 32-bits (Long word).
  - 16-bits (Word).
  - 8-bits (Byte).
  - 1-bit.

How Data Registers Work - Example

- The CPU wants to add together 2 numbers from memory locations A and C.
- Data stored into registers first, and added together.
- Results stored in register.

How Data Registers Work - Example

DATA REGISTERS (IN M68K)

	Contents
D0	
D1	
D2	
D3	
D4	
D5	
D6	
D7	

M68k wants to add together contents of A and C.

MEMORY

Add.	Contents
A	$12340000
B	…
C	$00005678
D	…
E	
F	
G	
H	

How Data Registers Work - Example

DATA REGISTERS (IN M68K)

	Contents
D0	$12340000
D1	$00005678
D2	
D3	
D4	
D5	
D6	
D7	

Step 1: M68k loads data from memory locations into data registers.

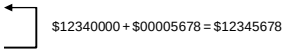
MEMORY

Add.	Contents
A	$12340000
B	…
C	$00005678
D	…
E	
F	
G	
H	

How Data Registers Work - Example

DATA REGISTERS (IN M68K)

	Contents
D0	$12345678
D1	$00005678
D2	
D3	
D4	
D5	
D6	
D7	



Step 2: M68k adds together the numbers and stores them in register. Result stored in D0.

MEMORY

Add.	Contents
A	$12340000
B	…
C	$00005678
D	…
E	
F	
G	
H	

Try It Yourself

```
START   ORG     $1000

A       EQU     $2000
C       EQU     $2008


        MOVE.L  #$12340000,A
        MOVE.L  #$00005678,C

        MOVE.L  A,D0
        MOVE.L  C,D1

        ADD.L   D1,D0

        END     START
```

What are the contents of addresses $2000 and $2008?

What are the contents of data registers D0 and D1?

What is the content of D0 after execution of ADD?

Address Registers

- Stores memory addresses of data and instructions.
- Eight registers in M68k.
- A0 – A7.
- A7 (x 2) is reserved as Stack Pointer (SP).
- 32-bits long (but only uses 24-bits).

How Address Registers Work - Example

M68k wants to add together contents of addresses $1000 and $1008.

DATA REGISTERS (IN M68K)

	Contents
D0	
D1	
D2	
D3	
D4	
D5	
D6	
D7	

ADDRESS REGISTERS (IN M68K)

	Contents
A0	
A1	
A2	
A3	
A4	
A5	
A6	
A7	

MEMORY

Add.	Contents
$1000	$12340000
$1004	…
$1008	$00005678
$100C	…
$1010	
$1014	
$1018	
$101C	

How Address Registers Work - Example

The addresses $1000 and $1008 are loaded into Address Registers.

DATA REGISTERS (IN M68K)

	Contents
D0	
D1	
D2	
D3	
D4	
D5	
D6	
D7	

ADDRESS REGISTERS (IN M68K)

	Contents
A0	**$1000**
A1	**$1008**
A2	
A3	
A4	
A5	
A6	
A7	

MEMORY

Add.	Contents
$1000	$12340000
$1004	…
$1008	$00005678
$100C	…
$1010	
$1014	
$1018	
$101C	

How Address Registers Work - Example

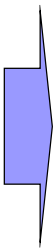
The values of A0 and A1 are sent to memory.

DATA REGISTERS (IN M68K)

	Contents
D0	
D1	
D2	
D3	
D4	
D5	
D6	
D7	

ADDRESS REGISTERS (IN M68K)

	Contents
A0	**$1000**
A1	**$1008**
A2	
A3	
A4	
A5	
A6	
A7	



MEMORY

Add.	Contents
$1000	$12340000
$1004	…
$1008	$00005678
$100C	…
$1010	
$1014	
$1018	
$101C	

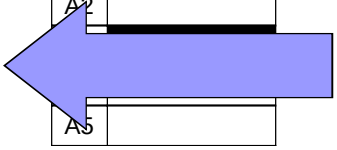
How Address Registers Work - Example

DATA REGISTERS (IN M68K)

	Contents
D0	$12340000
D1	$00005678
D2	
D3	
D4	
D5	
D6	
D7	

ADDRESS REGISTERS (IN M68K)

	Contents
A0	**$1000**
A1	**$1008**
A2	
A3	
A4	
A5	



M68k loads data from memory locations into data registers.

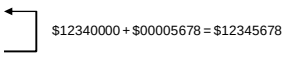
MEMORY

Add.	Contents
A	$12340000
B	…
C	$00005678
D	…
E	
F	
G	
H	

How Address Registers Work - Example

DATA REGISTERS (IN M68K)

	Contents
D0	$12345678
D1	$00005678
D2	
D3	
D4	
D5	
D6	
D7	



Step 2: M68k adds together the numbers and stores them in register. Result stored in D0.

MEMORY

Add.	Contents
A	$12340000
B	…
C	$00005678
D	…
E	
F	
G	
H	

Try It Yourself

```
START   ORG     $2000

        LEA     $1000,A0
        LEA     $1008,A1

        MOVE.L  #$12340000,(A0)
        MOVE.L  #$00005678,(A1)

        MOVE.L  (A0),D0
        MOVE.L  (A1),D1

        ADD.L   D1,D0

   END  START
```

What are the contents of addresses in A0 and A1?

What are the contents of data registers D0 and D1?

What is the content of D0 after execution of ADD?

Address Register vs. Data Register

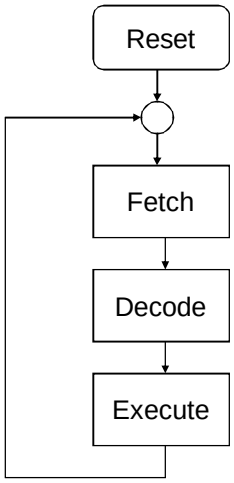
Comparison	Address Register	Data Register
Size	32-bit	32-bit
Total	7 (1 reserved)	8
Designation	A0 to A6 (A7 reserved)	D0 to D7
Data size	Word, Long	Byte, Word, Long
Instructions to access	Special instructions	General instructions
Purpose	Store addresses only	Store data only

Program Counter

Program Counter (PC)

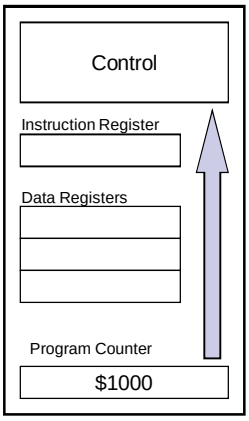
- Instructions must occur in correct sequence for program to run properly.
- PC makes sure this happens:
  - Special-purpose register inside CPU.
  - Keeps track of address of next instruction.
  - 32-bits: can point to any location in memory.

Remember this cycle?



Fetch – Step 1

CPU

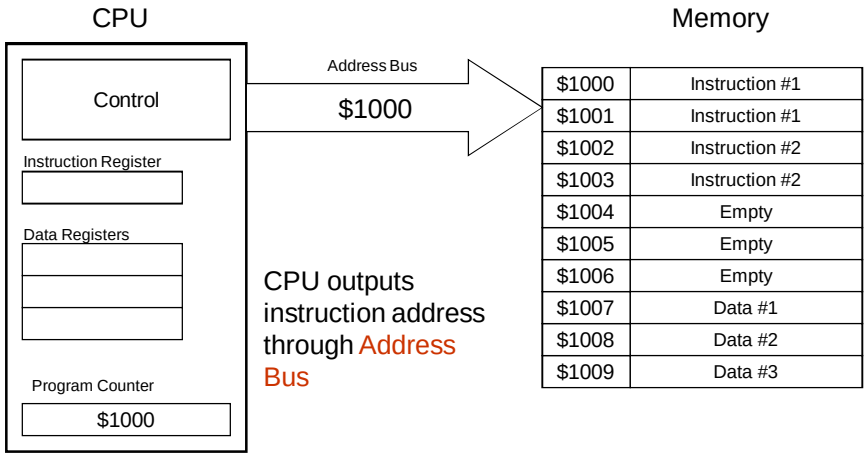


CPU gets instruction address from PC

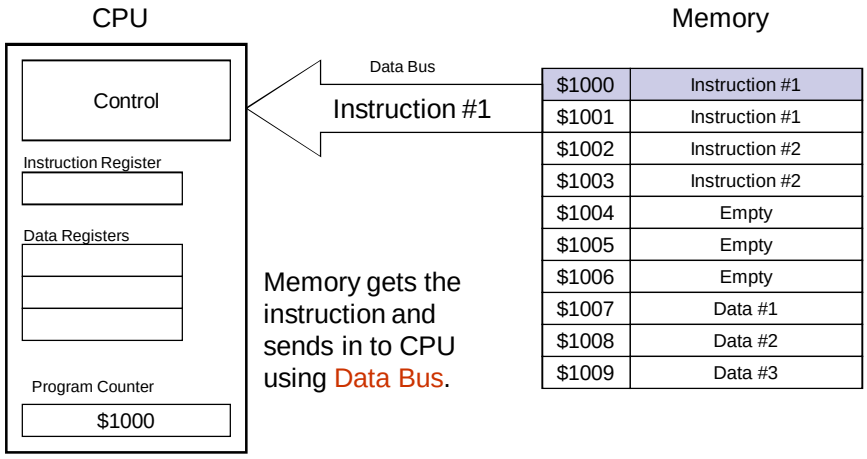
Memory

$1000	Instruction #1
$1001	Instruction #1
$1002	Instruction #2
$1003	Instruction #2
$1004	Empty
$1005	Empty
$1006	Empty
$1007	Data #1
$1008	Data #2
$1009	Data #3

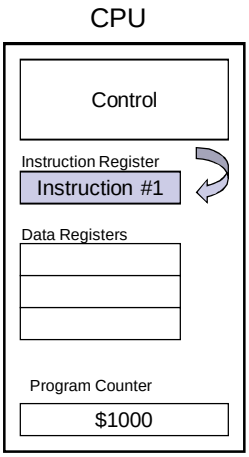
Fetch – Step 2



Fetch – Step 3



Fetch – Step 4

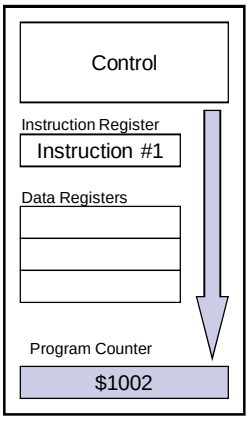


CPU stores instruction in Instruction Register

Memory

$1000	Instruction #1
$1001	Instruction #1
$1002	Instruction #2
$1003	Instruction #2
$1004	Empty
$1005	Empty
$1006	Empty
$1007	Data #1
$1008	Data #2
$1009	Data #3

Fetch – Step 5



After instruction has been loaded, CPU updates Program Counter.

Memory

$1000	Instruction #1
$1001	Instruction #1
$1002	Instruction #2
$1003	Instruction #2
$1004	Empty
$1005	Empty
$1006	Empty
$1007	Data #1
$1008	Data #2
$1009	Data #3

Program Counter (PC)

- Once M68k fetches instruction from memory, it automatically increments PC.
  - Just before M68k starts to execute the current instruction.
- PC always points to the next instruction during execution of current instruction.

Program Counter (PC)

- Before M68k can start executing a program, PC has to be loaded with address of its first instruction.
- During start-up, PC must be loaded with a number:
  - First instruction the M68k has to execute.
  - Usually at simple location – $0000000.

Program Counter (PC)

- PC can be loaded with new value to "jump" over instructions.
- When the PC is modified, new instruction is taken from new PC value, not the address in sequence.
- Can use this technique in subroutines, exceptions and branches.

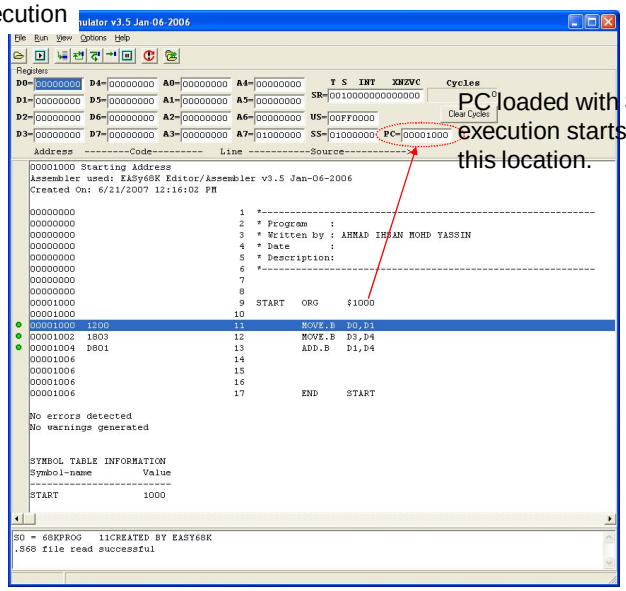
Try It Yourself

```
START    ORG      $1000

         MOVE.B   D0,D1
         MOVE.B   D3,D4
         ADD.B    D1,D4


         END      START
```

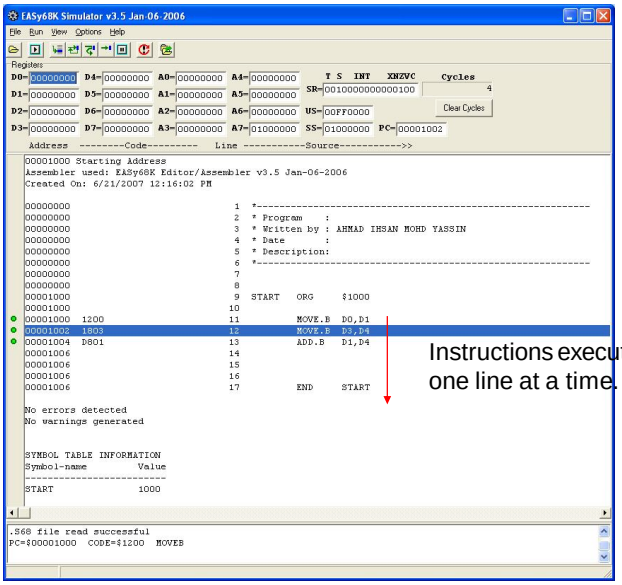
What are the contents of memory address starting from $1000?

Execute the program step by step. Notice that the value of PC changes automatically after each instruction.

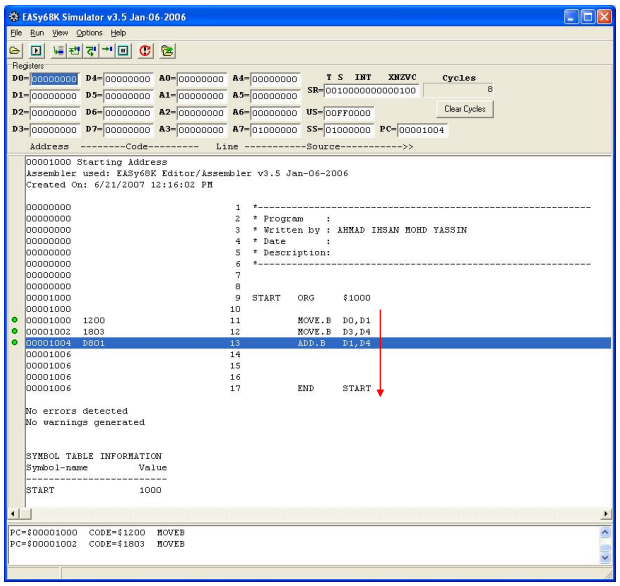
* Normal execution

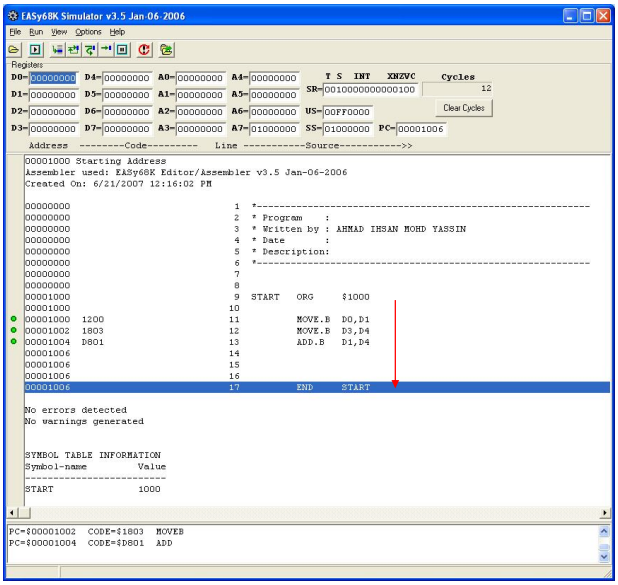


PC loaded with $1000, execution starts with this location.

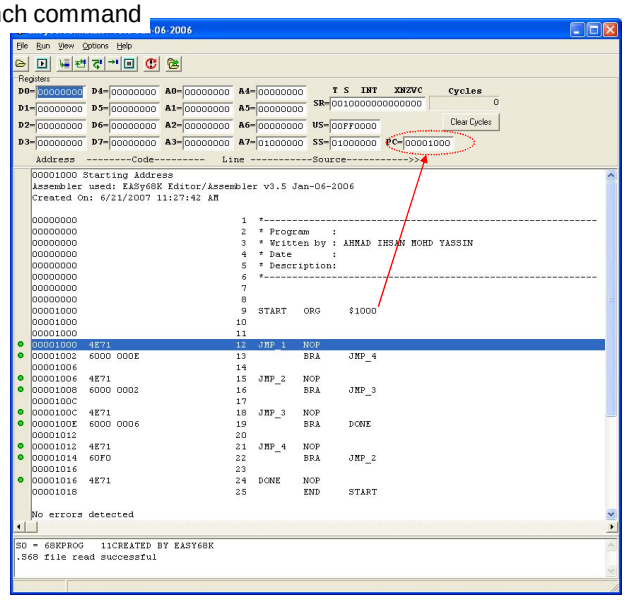


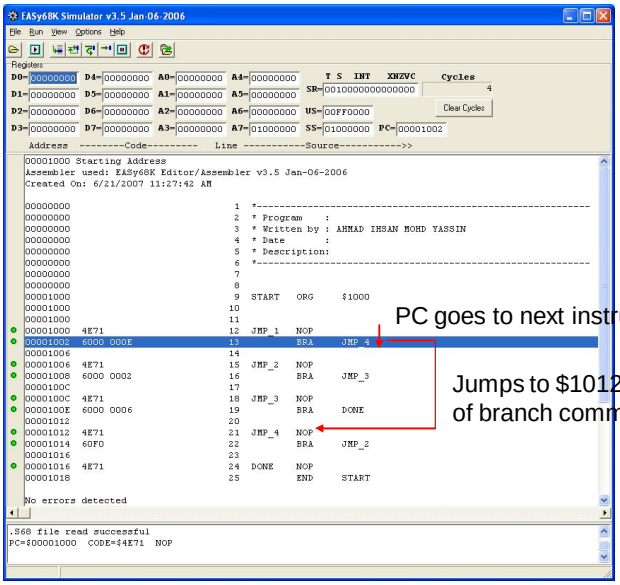
Instructions executed one line at a time.


EASy68K Simulator v3.5 Jan-06-2006
File Run View Options Help
Registers
D0=00000000 D4=00000000 A0=00000000 A4=00000000 T S INT XNZVC Cycles
D1=00000000 D5=00000000 A1=00000000 A5=00000000 SR=0010000000000100 8
D2=00000000 D6=00000000 A2=00000000 A6=00000000 US=00FF0000 Clear Cycles
D3=00000000 D7=00000000 A3=00000000 A7=01000000 SS=01000000 PC=00001004
Address --------Code--------- Line ----------Source---------->>
00001000 Starting Address
Assembler used: EASy68K Editor/Assembler v3.5 Jan-06-2006
Created On: 6/21/2007 12:16:02 PM
00000000 1 *---
00000000 2 * Program :
00000000 3 * Written by : AHMAD IHSAN MOHD YASSIN
00000000 4 * Date :
00000000 5 * Description:
00000000 6 *---
00000000 7
00000000 8
00001000 9 START ORG $1000
00001000 10
00001000 1200 11 MOVE.B D0,D1
00001002 1803 12 MOVE.B D3,D4
00001004 D801 13 ADD.B D1,D4
00001006 14
00001006 15
00001006 16
00001006 17 END START
No errors detected
No warnings generated
SYMBOL TABLE INFORMATION
Symbol-name Value
START 1000
PC=$00001000 CODE=$1200 MOVEB
PC=$00001002 CODE=$1803 MOVEB


EASy68K Simulator v3.5 Jan-06-2006
File Run View Options Help
Registers
D0=00000000 D4=00000000 A0=00000000 A4=00000000 T S INT XNZVC Cycles
D1=00000000 D5=00000000 A1=00000000 A5=00000000 SR=0010000000000100 12
D2=00000000 D6=00000000 A2=00000000 A6=00000000 US=00FF0000 Clear Cycles
D3=00000000 D7=00000000 A3=00000000 A7=01000000 SS=01000000 PC=00001006
Address --------Code--------- Line ----------Source---------->>
00001000 Starting Address
Assembler used: EASy68K Editor/Assembler v3.5 Jan-06-2006
Created On: 6/21/2007 12:16:02 PM
00000000 1 *---
00000000 2 * Program :
00000000 3 * Written by : AHMAD IHSAN MOHD YASSIN
00000000 4 * Date :
00000000 5 * Description:
00000000 6 *---
00000000 7
00000000 8
00001000 9 START ORG $1000
00001000 10
00001000 1200 11 MOVE.B D0,D1
00001002 1803 12 MOVE.B D3,D4
00001004 D801 13 ADD.B D1,D4
00001006 14
00001006 15
00001006 16
00001006 17 END START
No errors detected
No warnings generated
SYMBOL TABLE INFORMATION
Symbol-name Value
START 1000
PC=$00001002 CODE=$1803 MOVEB
PC=$00001004 CODE=$D801 ADD

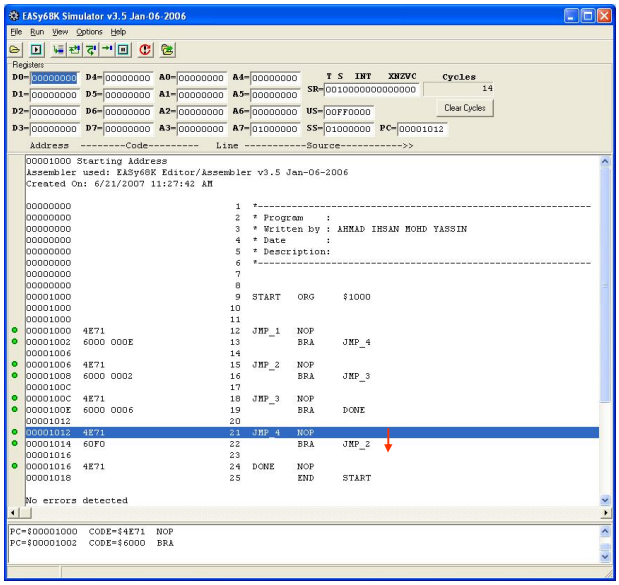
* Using Branch command

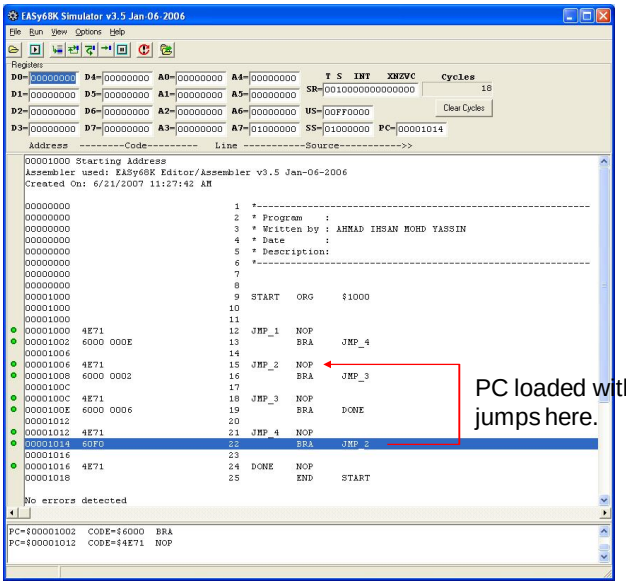




PC goes to next instruction

Jumps to $1012 because of branch command.





PC loaded with $1006, jumps here.

Stack Pointer

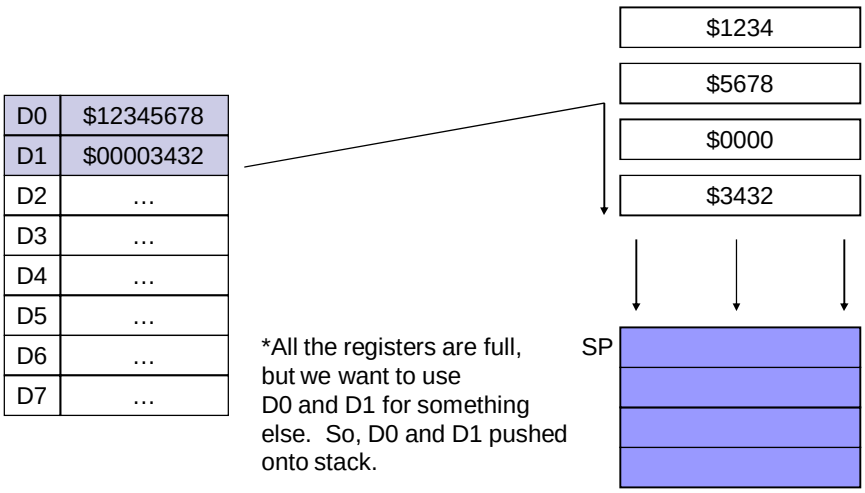
Software Stack

- Sometimes, registers not enough to store data:
  - Only 8 data registers available.
- Solution: reserve part of memory for storing data – software stack.
- Data can be stored/retrieved by pushing/popping.

Software Stack

- Stacks are usually used to store:
  - Data register values.
  - Program Counter.
  - Status Register.
  - Other information.
- Stacks operate on LIFO basis (Last In First Out).
- Pushes (puts) items on stack and pops (takes) in reverse order.

The Software Stack - Push



The Software Stack – After Push

D0	
D1	
D2	
D3	
D4	
D5	
D6	
D7	

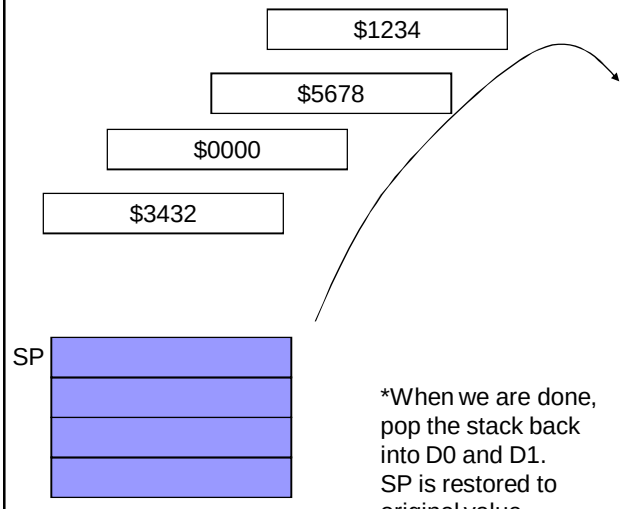
*D0 and D1 can now be cleared and used for other data. SP set as address to uppermost layer.

Stack

SP

$1234
$5678
$0000
$3432

The Software Stack - Pop



D0	$12345678
D1	$00003432
D2	
D3	
D4	
D5	
D6	
D7	

*When we are done, pop the stack back into D0 and D1. SP is restored to original value.

Stack Pointer (SP)

- Stack pointer: special register to keep record of stack use:
  - Keeps track of top-most stack position.
- A7 used to store SP.
- 32-bits: large enough to address all memory locations.

Stack Pointer in Action – Moving Data into Stack

```
MOVE.L    #$11112222,D0
MOVE.L    D0,-(A7)
```

*initial SP at $1000

SP (inside M68k):

$0FFC

	$0FF8	
	$0FFA	
NEW SP VALUE ➔	$0FFC	1111
	$0FFE	2222
OLD SP VALUE ➔	$1000	Not Filled

Stack Pointer in Action – Moving Data into Stack

```
MOVE.L      #$33334444,D0
MOVE.L      D0,-(A7)
```

*initial SP at $1000

SP (inside M68k):

$0FF8

*SP always keeps track of upper-most location in stack

NEW SP VALUE ➔

OLD SP VALUE ➔

$0FF8	3333
$0FFA	4444
$0FFC	1111
$0FFE	2222
$1000	Not Filled

Stack Pointer in Action – Getting Data Back from Stack

```
MOVE.L      (A7)+,D0
```

SP (inside M68k):

$0FFC

D0:

$33334444

OLD SP VALUE ➔

NEW SP VALUE ➔

$0FF8	Popped back to D0
$0FFA	Popped back to D0
$0FFC	1111
$0FFE	2222
$1000	Not Filled

Try It Yourself

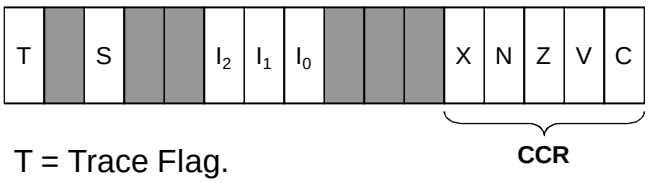
```
START   ORG     $1000

        MOVE.L  #$11112222,D0
        MOVE.L  #$33334444,D1

        MOVE.L  D0,-(A7)        * PUSH
        MOVE.L  D1,-(A7)

        MOVE.L  #$AAAABBBB,D0
        MOVE.L  #$CCCCDDDD,D1

        MOVE.L  (A7)+,D1        *POP
        MOVE.L  (A7)+,D0

    END  START
```

Status Register

Status Register

- Special purpose register, 16-bits wide.
- Stores important control and status bits of arithmetic calculations.
- Consists of:
  - Trace flag.
  - Supervisor flags.
  - Interrupt mask flags.
  - Condition Code Register (CCR) flags.

The Status Register



T = Trace Flag.
S = Supervisor Flag.
I_2, I_1, I_0 = Interrupt Mask Flags.
X = Extend Flag.
N = Negative Flag.
Z = Zero Flag.
V = Overflow Flag.
C = Carry Flag.

Status Register – Trace Bit

- Trace bit (T) used for program tracing:
  - Used to debug programs.
  - When trace is on, the program is executed one line at a time.
  - Waits for user input before going to next line.
  - T = 1, enabled. T = 0, disabled.

Status Register – Supervisor Bit

- Used to store privilege state of M68k.
- Two states defined:
  - S = 1, Supervisor mode.
  - S = 0, User mode.
- Controls access to critical system instructions and data.

Why do we need User/Supervisor modes?

- In any M68k system, there are:
  - Critical system data that must not be changed by inexperienced users. These data are stored in special memory locations.
  - Potentially destructive instructions that cannot be executed by inexperienced users.
    - Reset system.
    - Halt system.
    - Edit Status Register, Stack Pointer values.
- User/SV modes limits access by not allowing certain actions if users does not have SV privileges.

Supervisor vs. User Mode

Property	User Mode	Supervisor Mode
Activated when	S = 0	S = 1
Memory Access	Access to only user memory locations.	Access to both user and supervisor memory locations.
Instruction Set	Limited instruction set.	Unlimited instruction set.
Can modify SR/SP?	False.	True.
Stack pointer	User Stack Pointer*.	Supervisor Stack Pointer*.

*A7 (SP) actually has two registers, SSP and USP.

Status Register – Interrupt Mask Bits (IMB)

- 3-bits (I_2, I_1, I_0) as interrupt masks:
  - Stores interrupt levels from 000 to 111.
  - 000 (0) is lowest priority, least important.
  - 111 (7) is highest priority, most important.
  - Interrupt request compared to IMB to determine whether it should be executed.

What are interrupts?

- Allows M68k to prioritize processing:
  - More important tasks executed first.
  - Less important tasks executed later.
- Requested by external device:
  - Asks to be serviced.
  - Compare priority to IMB.
  - If higher than IMB, update IMB and service interrupt.
  - If request lower than IMB, request stored, executed later.

Interrupt Example

M68k

External Peripheral

1. M68k is executing instructions normally.
2. External peripheral has important task for M68k.
3. External peripheral asks for attention from M68k.
4. M68k compares interrupt level to SR.

T		S			I_2	I_1	I_0				X	N	Z	V	C

M68k

5. If external interrupt higher than current.
6. Update interrupt bits, save important data into stack.
7. Handle interrupt, restore interrupt bits.
8. Pop stack, resume normal execution.

5. If external interrupt lower than current.
6. Wait for higher-level interrupt being handled.

Condition Code Register (CCR)

Condition Code Register (CCR)

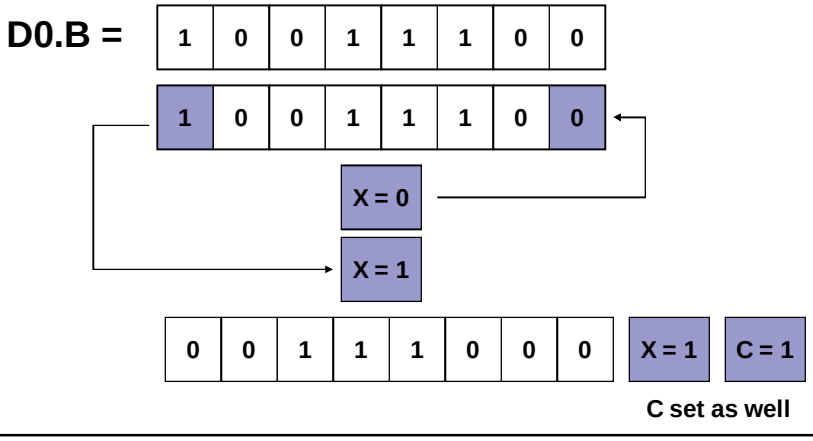
- Used to store status of evaluated conditions.
- Final 5-bits of SR.
- Bits in CCR:
  - X: Stores extra bit for arithmetic shifts.
  - N: Whether the result is negative.
  - Z: Whether the result is zero.
  - V: Whether an arithmetic overflow has occurred.
  - C: Whether a carry/borrow has occurred.

X: Extend Bit

- Purpose:
  - To allow rotate and shift operations:
    - Stores the extra shifted bit during rotation.
  - For multi-precision arithmetic in BCD operations.
  - Usually set according to C.

X Example – Rotate

MOVE.B #$9C,D0 * Assuming X = 0 before ROXL
ROXL.B #1,D0



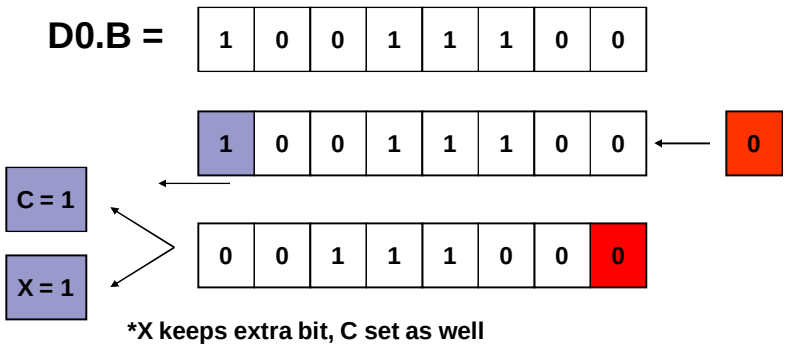
C set as well

Try It Yourself

```
START   ORG     $1000

        MOVE.B  #$9C,D0
        ROXL.B  #1,D0

END     START
```

X Example – Arithmetic Shift

ASL.B #1,D0



*X keeps extra bit, C set as well

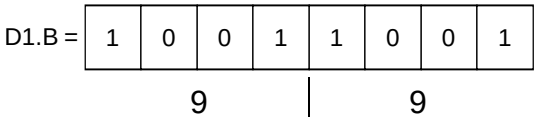
Try It Yourself

```
START   ORG     $1000

        MOVE.B  #$9C,D0
        ASL.B   #1,D0

        END     START
```

X Example – BCD Operation

D0 = $00000099
D1 = $00000099
X = 1 before execution
ABCD D0,D1

```
    99
+   99
+    1
  199
```



X/C = 1 (result > 99).
Z = clear since result non-zero.
N, V = undefined

Try It Yourself

```
START    ORG      $1000

         MOVE.B   #$99,D0
         MOVE.B   #$99,D1

         ABCD     D0,D1

         END      START
```

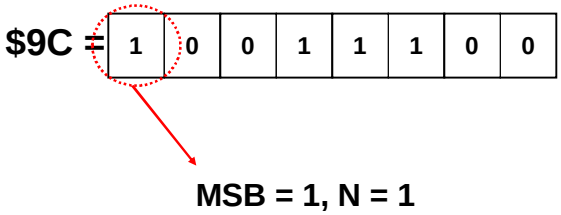
N – Negative Bit

- Purpose: to test whether the result is negative.
- Does this by examining the MSB:
  - In 2's complement, MSB is sign bit.
  - If MSB = 0, N = 0 (not negative).
  - If MSB = 1, N = 1 (negative).

N Example

- MOVE.B #0,D0
- MOVE.B #100,D1
- SUB.B D1,D0

D0 = $00 (0)
D1 = $64 (100)
D0 = 0 – 100 = $9C (-100)



Try It Yourself

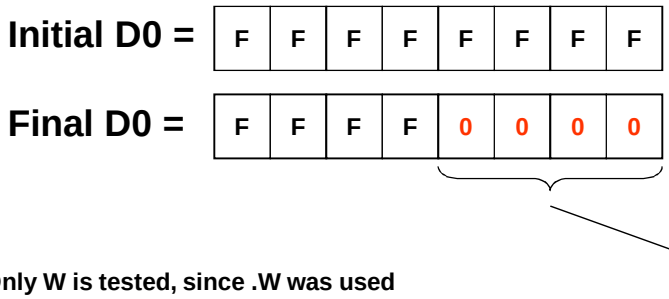
```
START   ORG     $1000

        MOVE.B  #0,D0
        MOVE.B  #100,D1
        SUB.B   D1,D0

        END     START
```

Z – Zero Bit

- Purpose: to test whether the result is zero.
- Does this by examining all active bits.
  - If all active bits = 0, Z = 1 (is zero).
  - If not all active bits = 0, Z = 0 (not zero).

Z Example

- MOVE.L #$FFFFFFFF,D0
- MOVE.W #$0000,D0



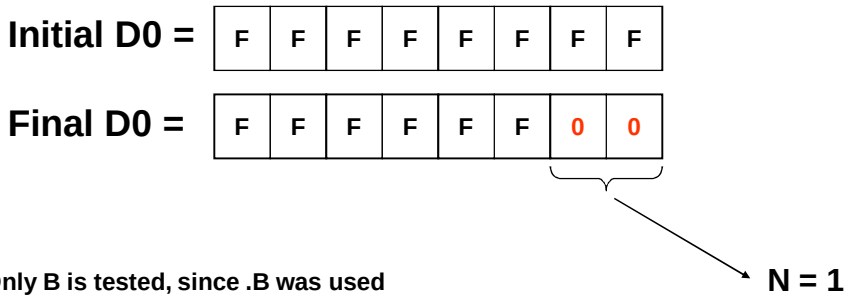
*** Only W is tested, since .W was used**

Try It Yourself

```
START   ORG     $1000

        MOVE.L  #$FFFFFFFF,D0
        MOVE.W  #$0000,D0

END     START
```

Z Example

- MOVE.L #$FFFF,D0
- SUB.B #$FF,D0



*** Only B is tested, since .B was used**

Try It Yourself

```
START   ORG     $1000

        MOVE.L  #$FFFFFFFF,D0
        SUB.B   #$FF,D0

        END     START
```

V – Overflow Bit

- Checks whether sign bit is changed as a result of arithmetic operation.
  - Set when arithmetic operation changes MSB.
  - Only few instructions can change V.
  - Check with M68k instruction reference.

V- Overflow Bit

- In ADD, SUB and CMP, V is modified differently.
- For ADD, SUB & CMP:
  - P + P = N
  - N + N = P
  - P – N = N
  - N – P = P

* P = Positive Number (MSB = 0)
  N = Negative Number (MSB = 1)

V Example

- MOVE.B #$41,D0
- MOVE.B #$46,D1
- ADD.B D1,D0

D0 = 65 + 70

		2's complement
D0.B =	0 1 0 0 0 0 0 1	+65
+ D1.B =	0 1 0 0 0 1 1 0	+70
New D0.B =	1 0 0 0 0 1 1 1	-121?

V = 1

Try It Yourself

```
START    ORG      $1000

         MOVE.B   #$41,D0
         MOVE.B   #$46,D1

         ADD.B    D1,D0

         END      START
```

V Example

- MOVE.B #$DE,D0
- MOVE.B #$9F,D1

} D0 = -34 + (-97)

- ADD.B D1,D0

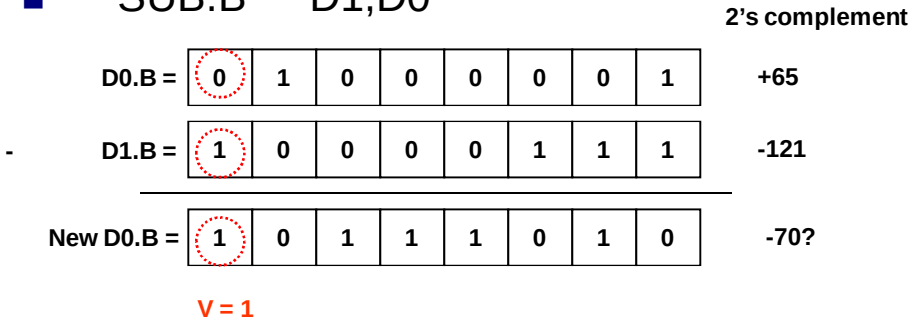
Try It Yourself

```
START    ORG      $1000

         MOVE.B   #$DE,D0
         MOVE.B   #$9F,D1

         ADD.B    D1,D0

         END      START
```

V Example

- MOVE.B #$41,D0
- MOVE.B #$87,D1
- SUB.B D1,D0

D0 = +65 – (-121)



Try It Yourself

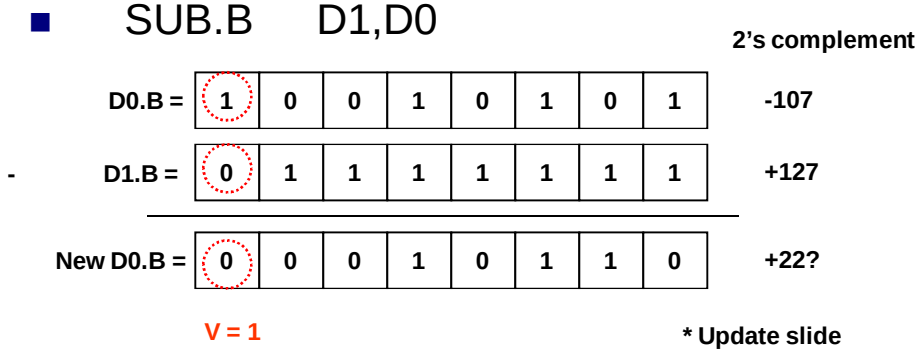
```
START    ORG      $1000

         MOVE.B   #$41,D0
         MOVE.B   #$87,D1

         SUB.B    D1,D0

         END      START
```

V Example

- MOVE.B #$95,D0
- MOVE.B #$7F,D1

} D0 = (-107) – 127

- SUB.B D1,D0



*** Update slide**

Try It Yourself

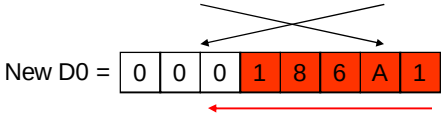
```
START    ORG      $1000

         MOVE.B   #$95,D0
         MOVE.B   #$7F,D1

         SUB.B    D1,D0

         END      START
```

DIVU Example - Overflow

D0 = $00186A1 (100,001 decimal)
D1 = $0000001 (1 decimal)
DIVU D1,D0

100,001 / 1 = 100,001 = $186A1

	Quotient	Remainder
Decimal	100,001	0
Hex	$186A1	$0000



Quotient overflowed to remainder.

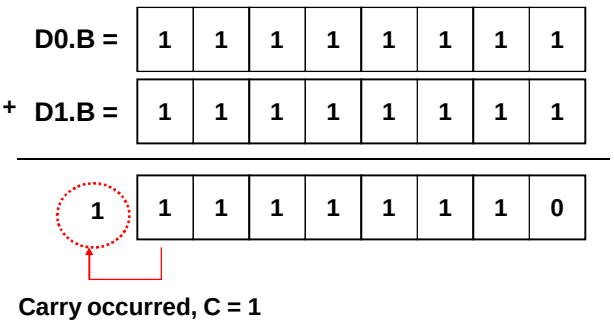
V = 1 (divide overflow)

C – Carry Bit

- Purpose: to test whether a borrow/carry has occurred:
  - Carry occurs when addition results are "carried forward".
  - Borrow occurs when subtraction results are "borrowed".
- Does this by examining the results of addition/subtraction operations.

C Example - Carry

- MOVE.B #$FF,D0
- MOVE.B #$FF,D1
- ADD.B D1,D0

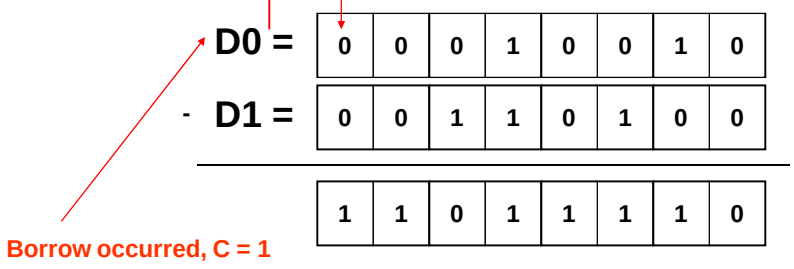


Try It Yourself

```
START    ORG      $1000

         MOVE.B   #$FF,D0
         MOVE.B   #$FF,D1

         ADD.B    D1,D0

         END      START
```

C Example - Borrow

- MOVE.B #$12,D0
- MOVE.B #$34,D1
- SUB.B D1,D0



Try It Yourself

```
START    ORG      $1000

         MOVE.B   #$12,D0
         MOVE.B   #$34,D1

         SUB.B    D1,D0

         END      START
```

Conclusion

Conclusion

- Data registers store data for manipulation.
- Address registers can only store address.
- Status register hold various status and control bits during M68k operations:
  - T actives trace mode.
  - S activates supervisor mode.
  - IMB stores level of serviced interrupt.
  - CCR store status each instruction performed.
  - Remember how conditions tested, results.

The End

Please read:

Antonakos, pg. 59 – 61.

Gilmore, 71 – 98.