

# Programming Exercise 1

## Construction Heuristic

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# Approach

## 1. Preprocessing

- a. Battery constraints
- b. Time constraints
- c. Capacity constraints

## 2. Parallel Savings

- a. Initial route creation
- b. Calculate the savings
  - i. customer - customer
  - ii. charger - customer
- c. Delete charger only routes
- d. Combine routes

## 3. Verify/Visualize

# Used Ressources

- Provided paper (The Electric Vehicle-Routing Problem with Time Windows and Recharging Stations; Schneider, Stenger, Goeke) and slides
- Python 3.x
- Provided Verifier
- Graphviz/Dot for visualization

# Results

- Total Runtime: 206.125036 seconds
- Average Cost over all Instances: 1065.25
- Worst Cost:       instance: ../data/instances/rc101\_21.txt  
                  cost: 3115.0410853411863  
                  time: 3.442202091217041
- Worst Runtime:   instance: ../data/instances/r112\_21.txt  
                  cost: 1344.055276592745  
                  time: 9.748563051223755

# Computer

Prozessor	AMD Phenom(tm) II X4 955 Processor 3.20 GHz
Installiertes RAM	4,00 GB
Systemtyp	64-Bit-Betriebssystem, x64-basierter Prozessor

# Preprocessing

$q_i$  - demand of vertex  $i$

$C$  - vehicle load capacity

$e_i$  - beginning of a time window at vertex  $i$

$l_i$  - end of a time window at vertex  $i$

$Q$  - vehicle battery capacity

Arc  $(v,w)$  is blacklisted if:

$$v,w \in V: q_v + q_w > C$$

$$v \in V_d, w \in V_d: e_v + s_v + t_{vw} > l_w$$

$$v \in V_d, w \in V_d: e_v + s_v + t_{vw} + s_w + t_{wd} > l_d$$

$$v,w \in V, \forall j \in F_d, i \in F_d: h^*(d_{jv} + d_{vw} + d_{wi}) > Q$$

# Parallel Savings

Initial Routes:

- initial routes are created for every node
- customer might not be reachable without charger
- insert charger after customer and, if needed, before customer

# Parallel Savings

Calculate Savings:

- Calculate savings between:
  - a. Customer - Customer
  - b. Charger - Customer
  - c. Both not on Blacklist



# Parallel Savings

## Delete Charger only Routes

- delete routes that only contain charger
- adding chargers is not done by savings
- add instead nearest charger by necessity

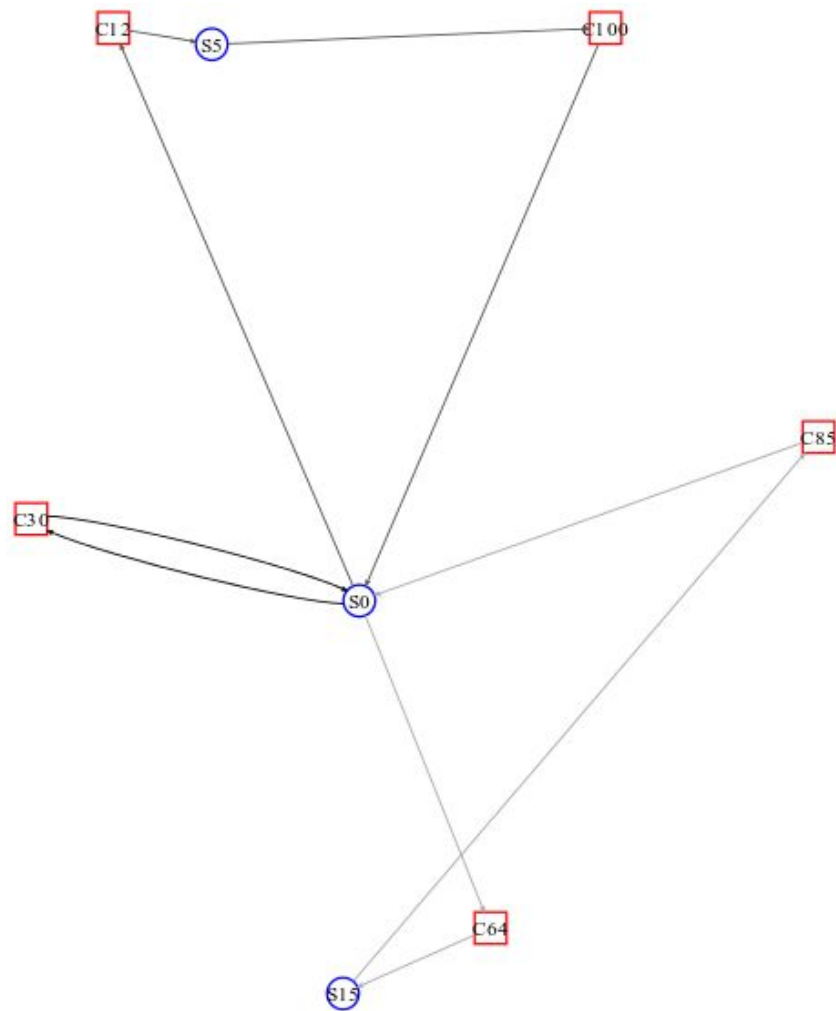
# Parallel Savings

## Combine Routes

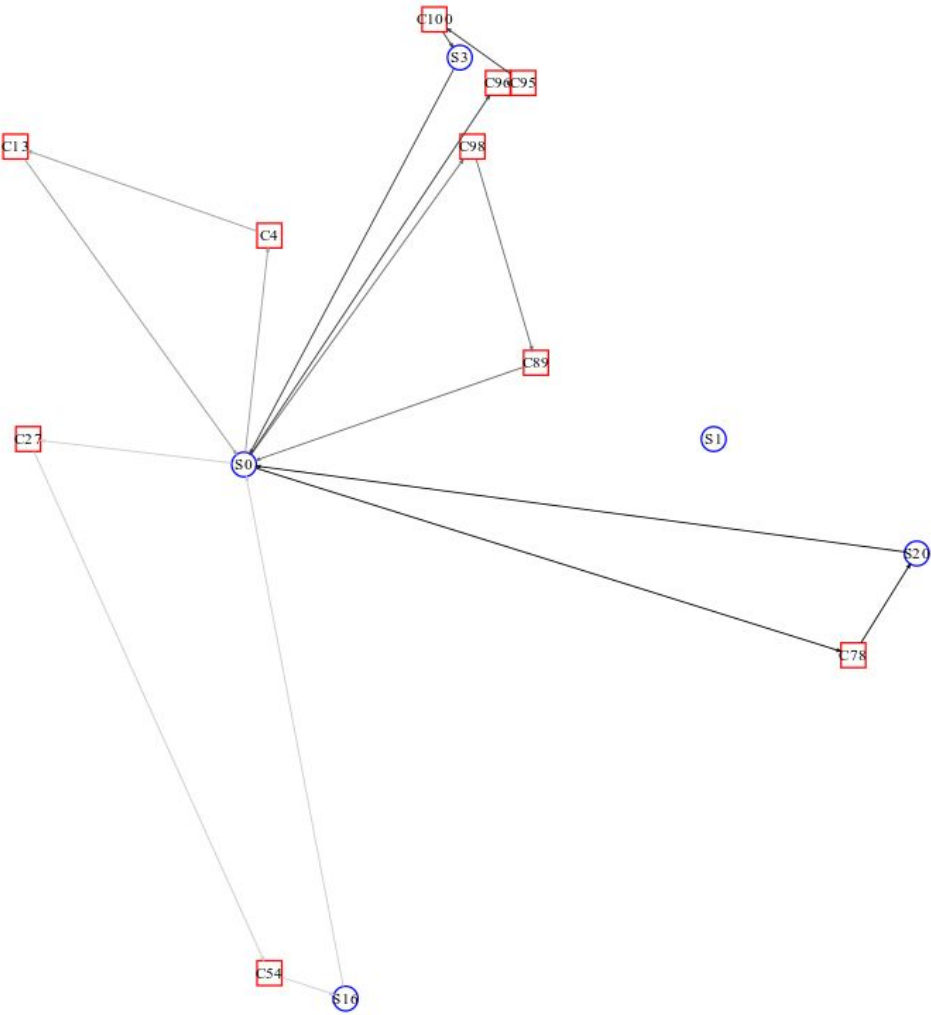
- Create new combined Route
- Check if and where charge is exceeded
  - iterate from there backwards and insert nearest charger if possible
  - if no previous insertion makes combined routes feasible then insert the closest charger between two farthest nodes
  - since we insert chargers until we make our route feasible, we have to decline routes that have more than some number of chargers, because we can create an infinite loop
- Check capacity, time-windows and charge capacity
- If newly created Route is feasible, add it and remove the old ones
- At the end, check for one edge case

# Result Examples

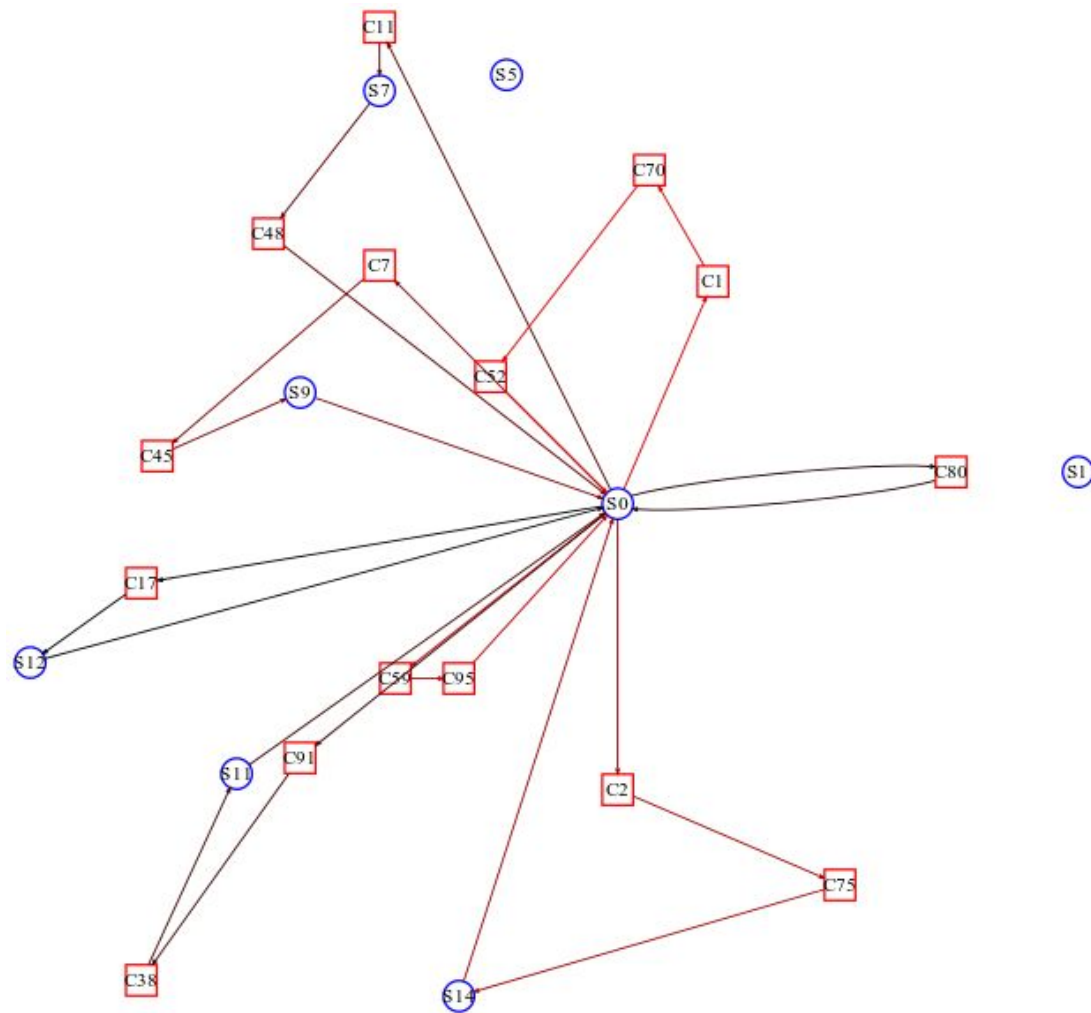
c101C5



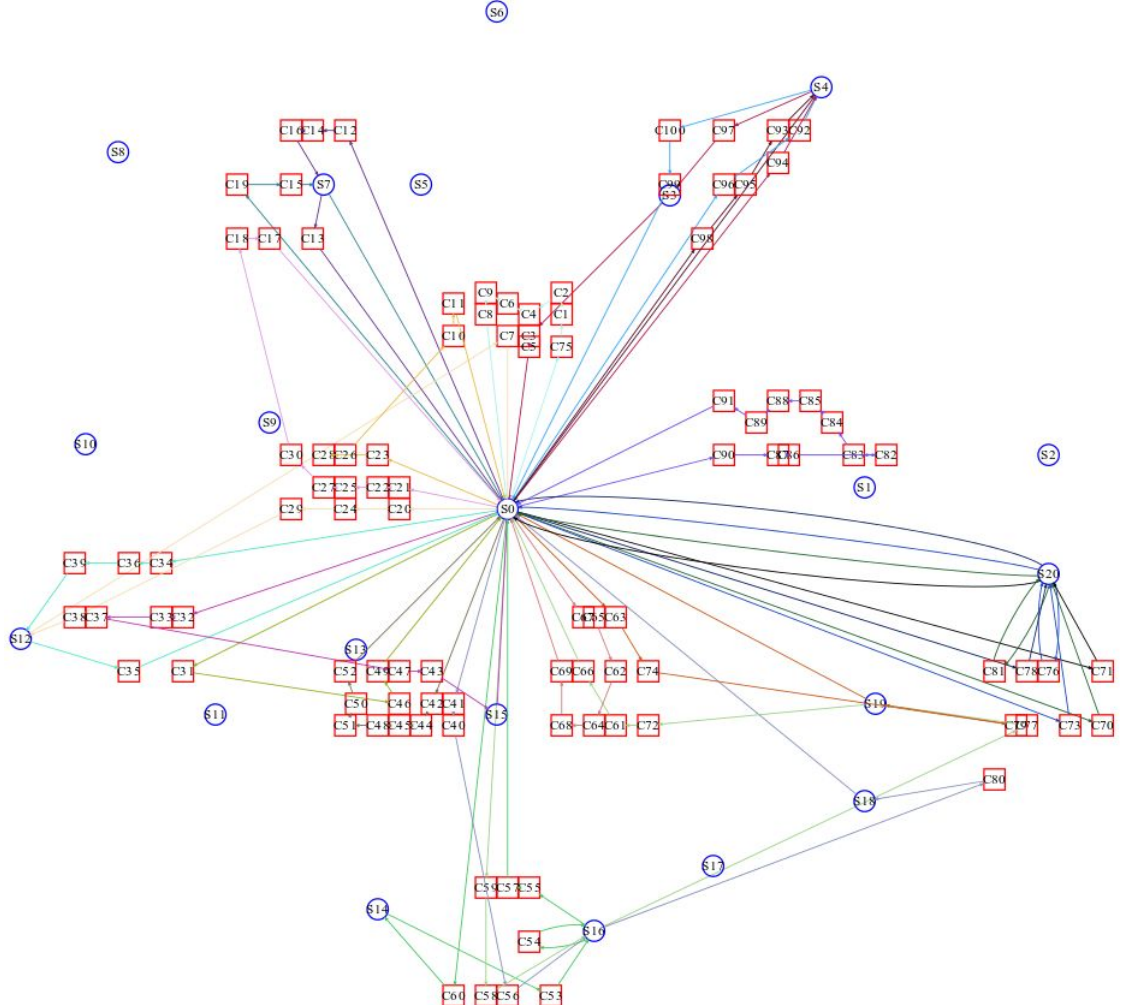
c101C10



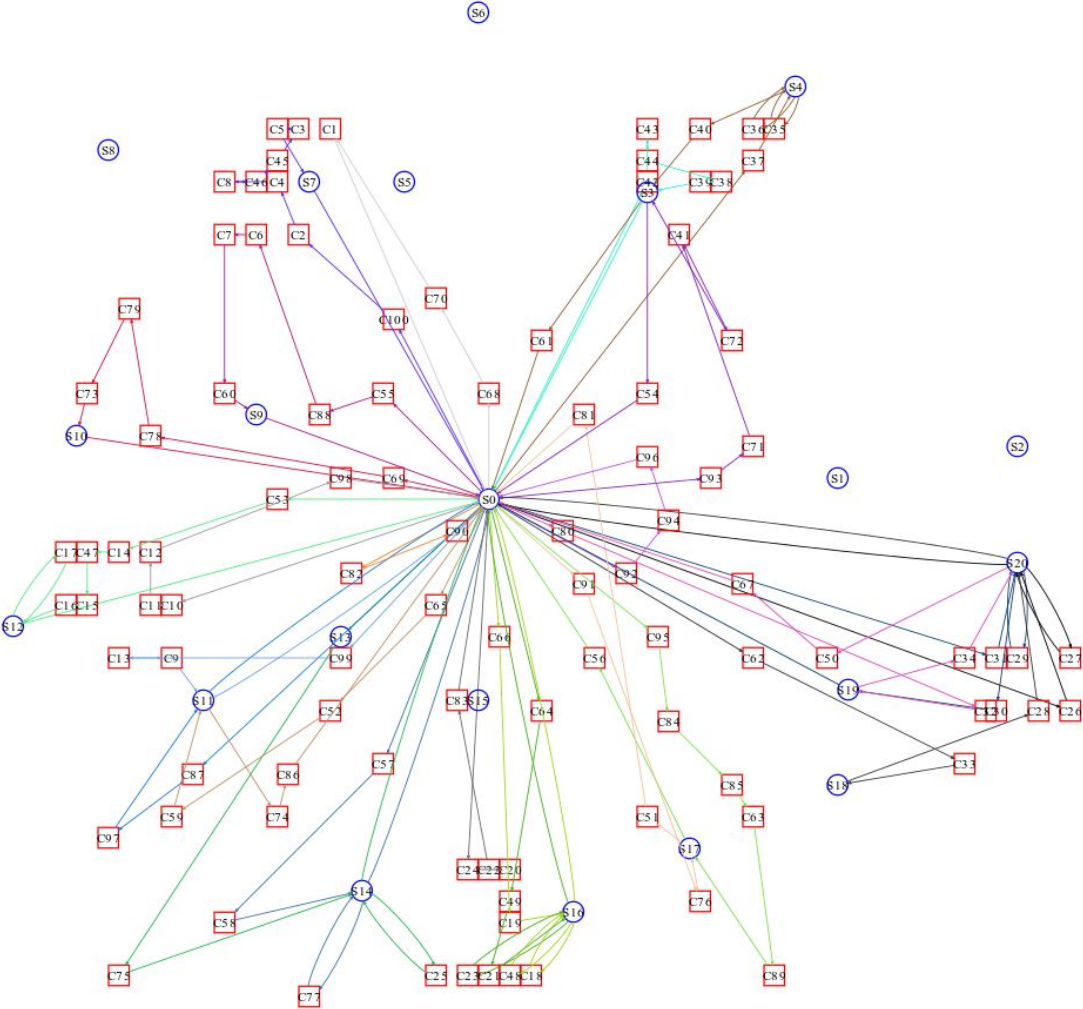
rc102C15



c101\_21



rc105\_21





rc102\_21

