





MARX – Uncovering Class Hierarchies in C++ Programs

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Introduction

Introduction

- We present Marx, a tool to extract class hierarchies from legacy binaries.
- This eases **static analysis** of C++ applications.
- We use the results for **security-related** use cases:
 - 1. Control-Flow Integrity
 - 2. Type-safe Object Reuse

MOTIVATION

- C++ applications are hard to analyze **statically**.
 - · Class inheritance.
 - · Objects allocated on the heap.
 - · Indirect function calls (polymorphism).

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- · Solving these problems aids in the reverse engineering process.
- Exploit mitigations would profit from solutions to these problems.
 - · Counterfeit Object-oriented Programming.
 - · Fine-grained control over forward edge.

Control-Flow Integrity

Code

```
Function X

[...]

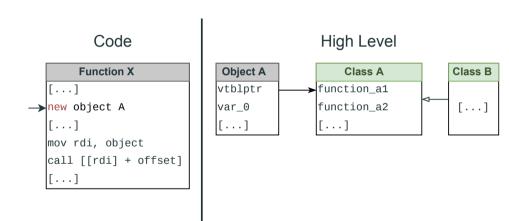
new object A

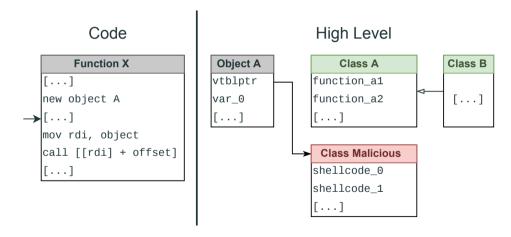
[...]

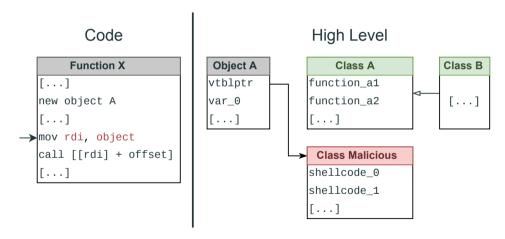
mov rdi, object

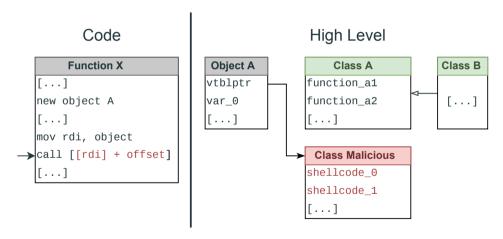
call [[rdi] + offset]

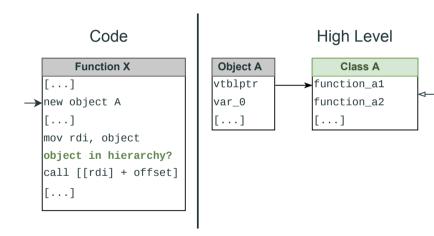
[...]
```





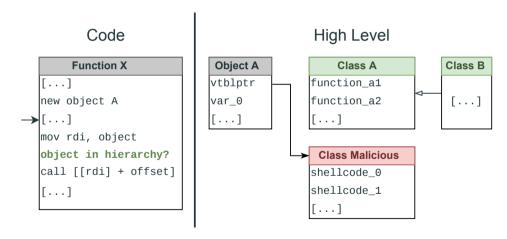


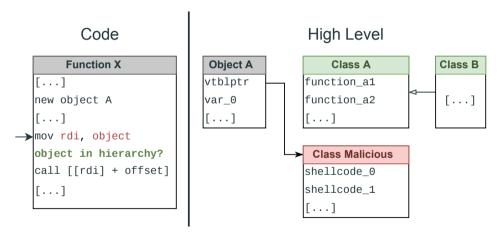


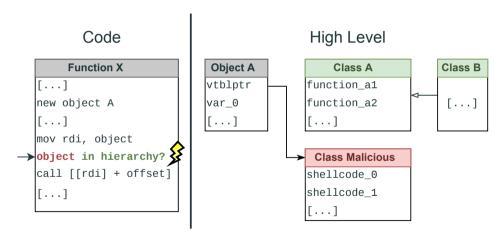


Class B

[...]







Approach

Code

Function X [...] mov rdi, object call [[rdi] + offset] [...]

Source Code object→function_a1()

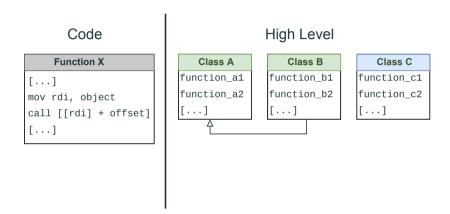
High Level

Class A
function_a1
function_a2
[]

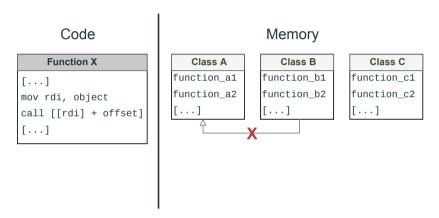


Class C
function_c1
function_c2
[...]

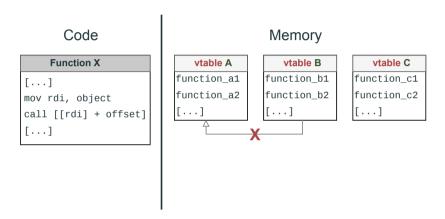
Indirect function calls are prominent C++ artifacts on binary level.



At high level, a class may *inherit* from another.



Inheritance relationships are not readily available in the binary.



Classes are (loosely) represented by *vtables* instead.

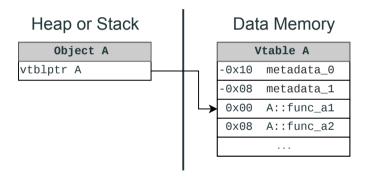
- MARX **statically** recovers the class hierarchy from an x86-64 binary (Itanium ABI).
 - · Represents class hierarchy as set of vtables.
 - · Ties set of vtables to indirect function call.

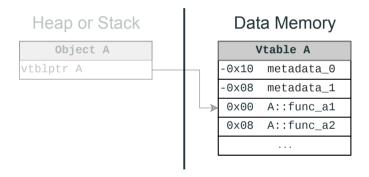
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- Analysis steps:
 - 1. Identify vtables,
 - 2. inspect their usages,

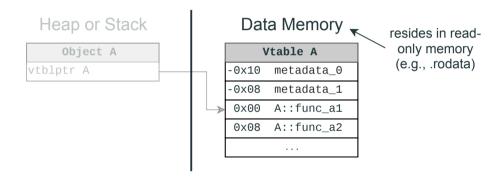
- 3. refine results using heuristics,
- 4. merge results across modules.

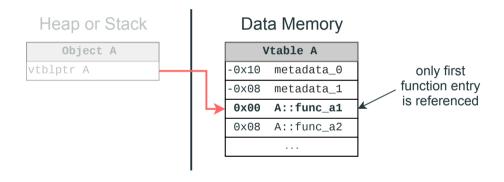
1. Vtable Candidates

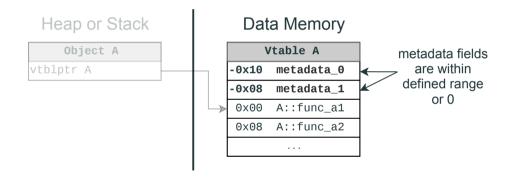
Approach



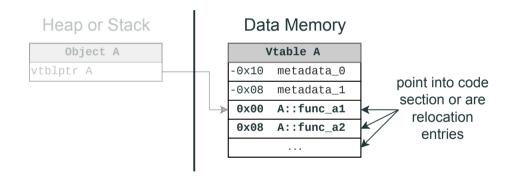






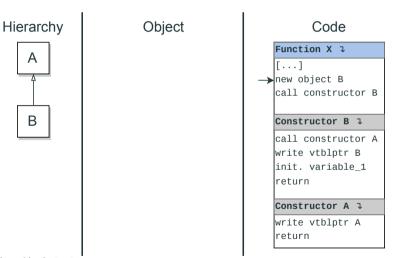


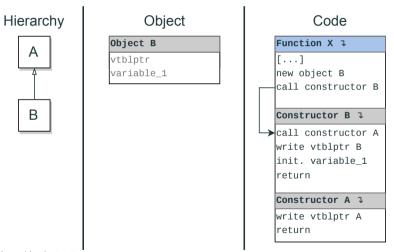
Approach - Vtable Identification

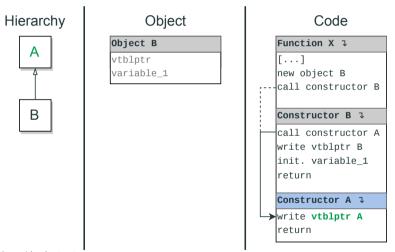


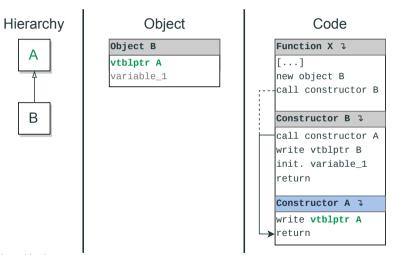
Approach

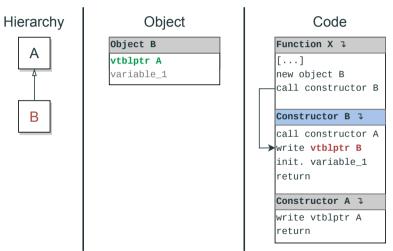
2. Vtable Usages

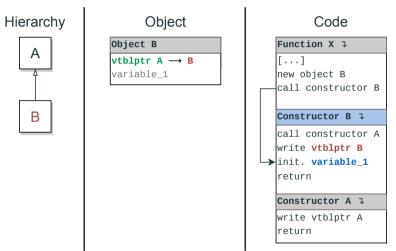




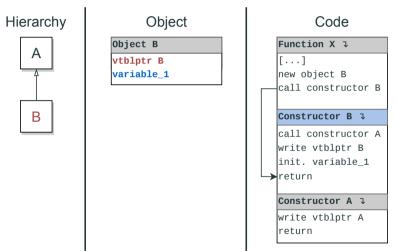








APPROACH - Overwrite Analysis



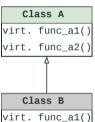
APPROACH - Path Generation

- In a function, build paths visiting
 - 1. indirect calls,
 - 2. direct calls to new, and
 - 3. instructions operating on vtables.
- Resolve indirect function calls for known vtables in current context.
- Follows calls with a **call depth** of 2.

Approach

3. Vtable Heuristics

Hierarchy



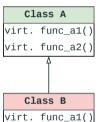
virt. func_b1()

Data Memory

Vtable A			
-0×10	0		
-0x08	0		
0×00	A::func_a1		
0x08	A::func_a2		

Vtable B		
-0×10	0	
-0x08	0	
0×00	B::func_a1	
0x08	A::func_a2	
0x10	B::func b1	

Hierarchy

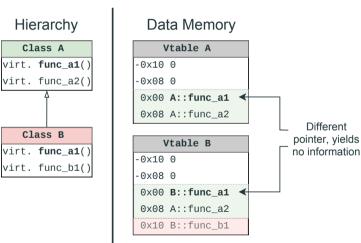


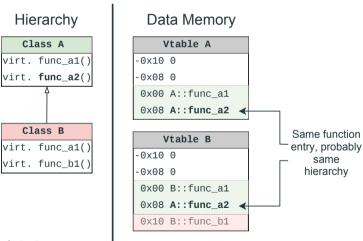
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-0x08	Θ		
0×00	B::func_a1		
0x08	A::func_a2		
0×10	B::func_b1		





4. Inter-Modular Dependencies

Approach



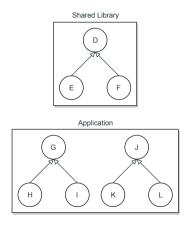
APPROACH – Inter-Modular Dependencies

Applications may consist of several modules.

Shared Library	
Application	

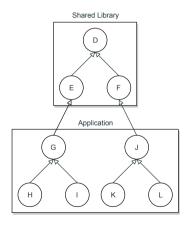
APPROACH - Inter-Modular Dependencies

Independent analysis yields, e.g., three distinct hierarchies.



APPROACH - Inter-Modular Dependencies

Unless we merge results across modules, we *underestimate* the hierarchy.





Evaluation

Application	No. of Hierarchies	Exactly Reconstructed	Time (h)
MySQL Server	78	69 (88%)	11:36
MongoDB	158	137 (87%)	1:08
Node.js	59	55 (93%)	0:33

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48 min analysis time on average for 5 real-world applications.

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Application	Overestimated	Underestimated	Not Found
MySQL Server	1 (1%)	7 (9%)	1 (1%)
MongoDB	0	8 (5%)	13 (8%)
Node.js	2 (3%)	2 (3%)	0

48 min analysis time on average for 5 real-world applications.

EVALUATION – Binary Hardening

• Use MARX' results to harden legacy binaries.

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- Implemented and tested two applications:

Lack of precision of analysis?

EVALUATION – Precision

- Type-safe Object Reuse allows imprecision, lowered security.
- Control-Flow Integrity may break applications.

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 - Enrich static results with dynamic analysis (unit tests).
 - Fall-back to computationally intensive **slow path** (*PathArmor*).

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• Please refer to the paper for further details.

Conclusion

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- Approach is **compiler-dependent**, optimizations may break assumptions.
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- Analysis is prone to **path explosion** (inter/intra-procedural path generation).
- Applications building upon results have to tolerate imprecision.

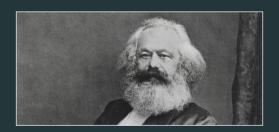
CONCLUSION

- MARX succeeds in **recovering the majority** of class hierarchies.
 - · Large, real-world applications.
 - Promising results.

- · Results are applicable to **security-related** use cases on the **binary level**.
- Our C++ open-source implementation based on libVEX is available at

https://github.com/RUB-SysSec/Marx.

Thank you for your attention.



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