

Decentralized Anonymous Credentials

Christina Garman, Matthew Green, Ian Miers
Johns Hopkins University



JOHNS HOPKINS
UNIVERSITY

Privacy and Identity on the Internet

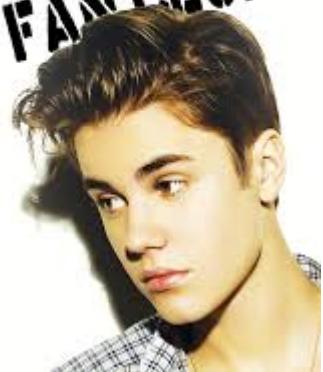
- ▶ Cannot make statements of identity privately
- ▶ But what about identity attributes?



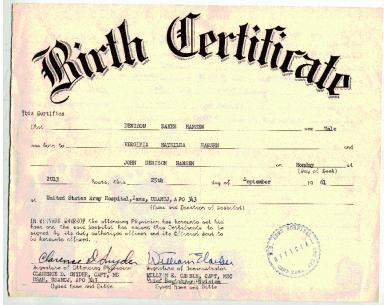
Privacy and Identity on the Internet



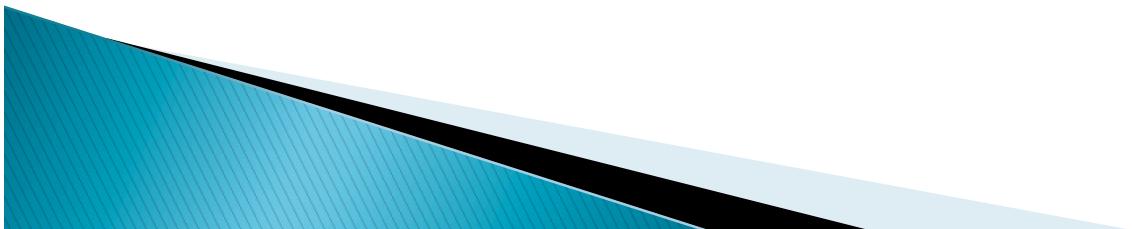
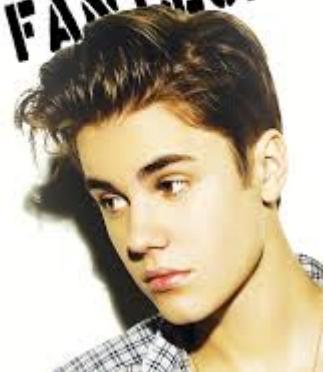
JUSTIN BIEBER
FAN CLUB



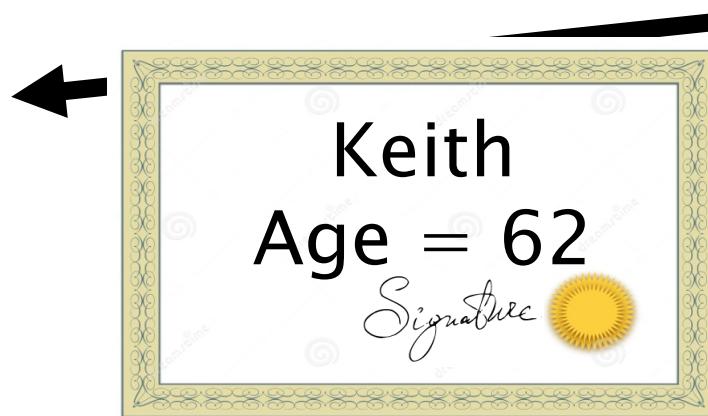
Privacy and Identity on the Internet



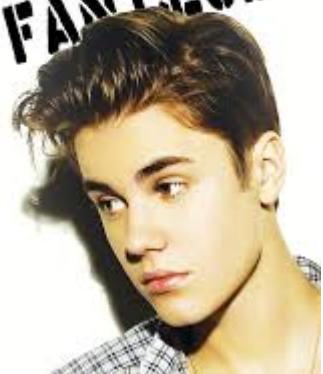
JUSTIN BIEBER
FAN CLUB



Privacy and Identity on the Internet



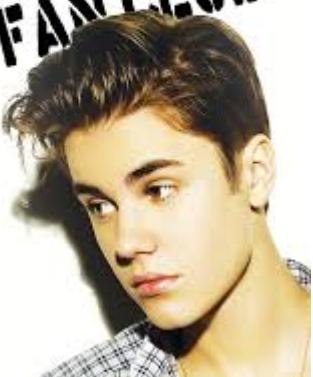
JUSTIN BIEBER
FAN CLUB



Privacy and Identity on the Internet



JUSTIN BIEBER
FAN CLUB

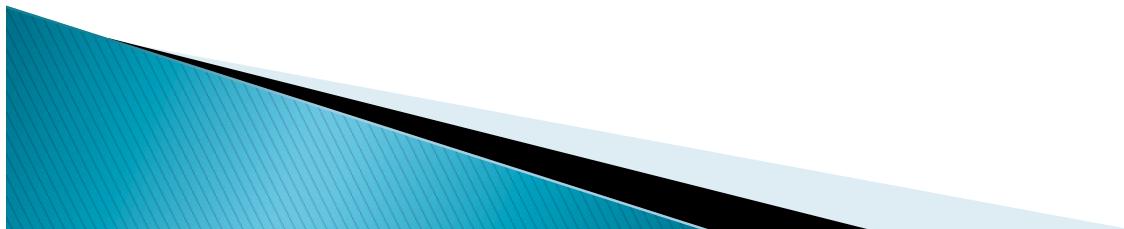
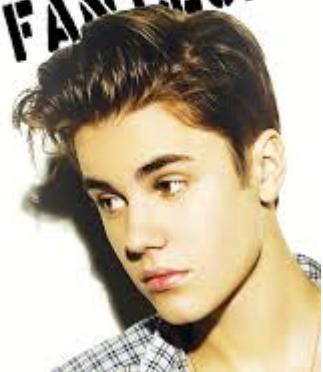


Privacy and Identity on the Internet



“Welcome to
the club!”

JUSTIN BIEBER
FAN CLUB

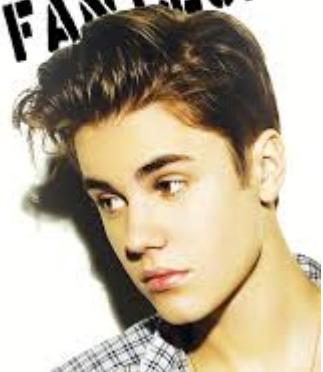


Privacy and Identity on the Internet



“Welcome to
the club!”

JUSTIN BIEBER
FAN CLUB



Privacy and Identity on the Internet



Keith Alexander

earlier today....



Keith liked the Justin Bieber Fan Club.



[Like](#) • [Comment](#) • [Edit](#) • [Delete](#)

Anonymous Credentials

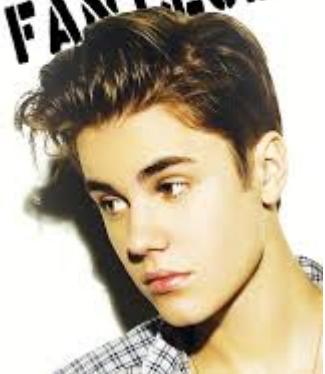
- ▶ Introduced by Chaum [Chaum85] and extended in [Brands00, CL01, CL02, CL03, BCKL08,...]
- ▶ Prove that you have a credential issued by some organization without revealing anything other than that you have the credential
- ▶ Standard techniques use a specialized digital signature



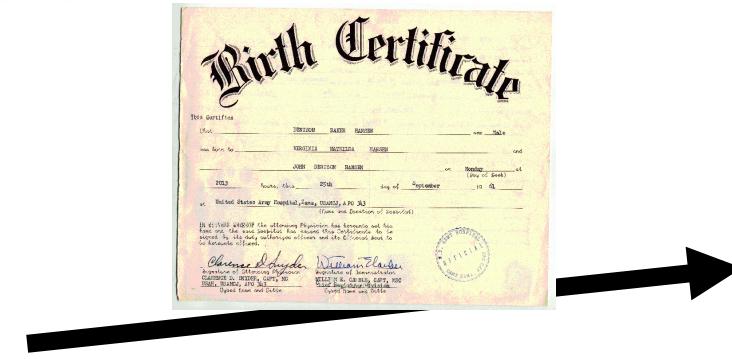
Example of Anonymous Credentials



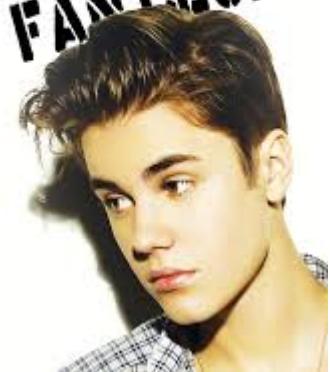
JUSTIN BIEBER
FAN CLUB



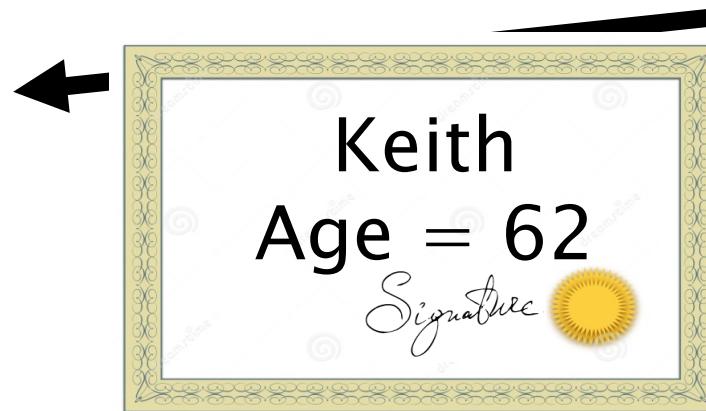
Example of Anonymous Credentials



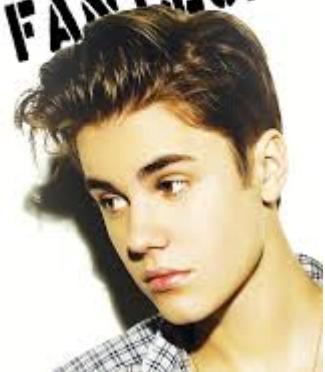
JUSTIN BIEBER
FAN CLUB



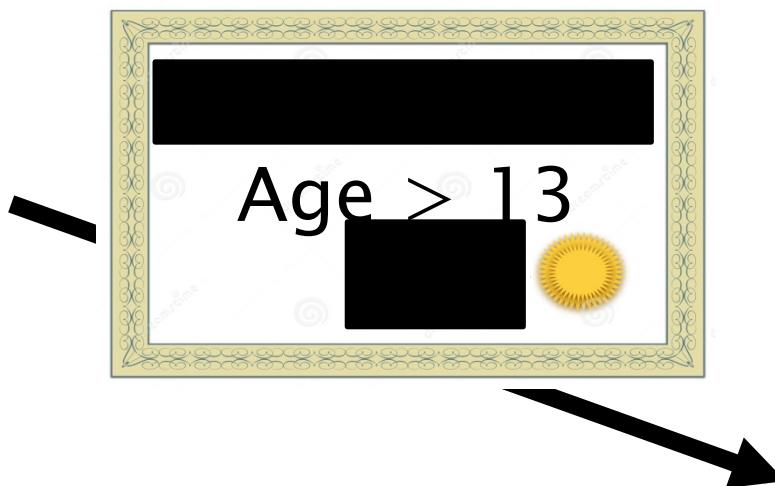
Example of Anonymous Credentials



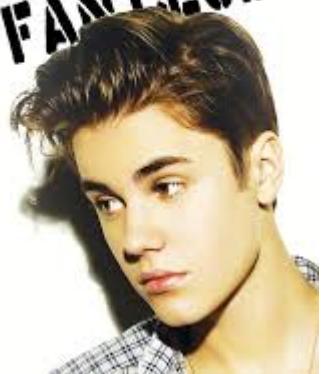
JUSTIN BIEBER
FAN CLUB



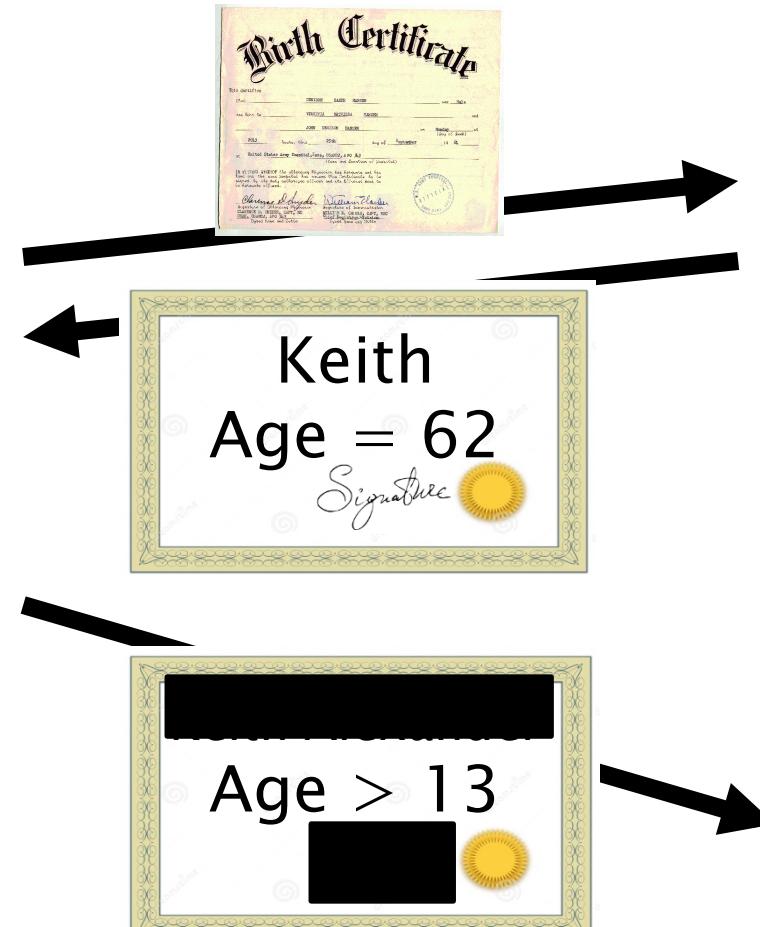
Example of Anonymous Credentials



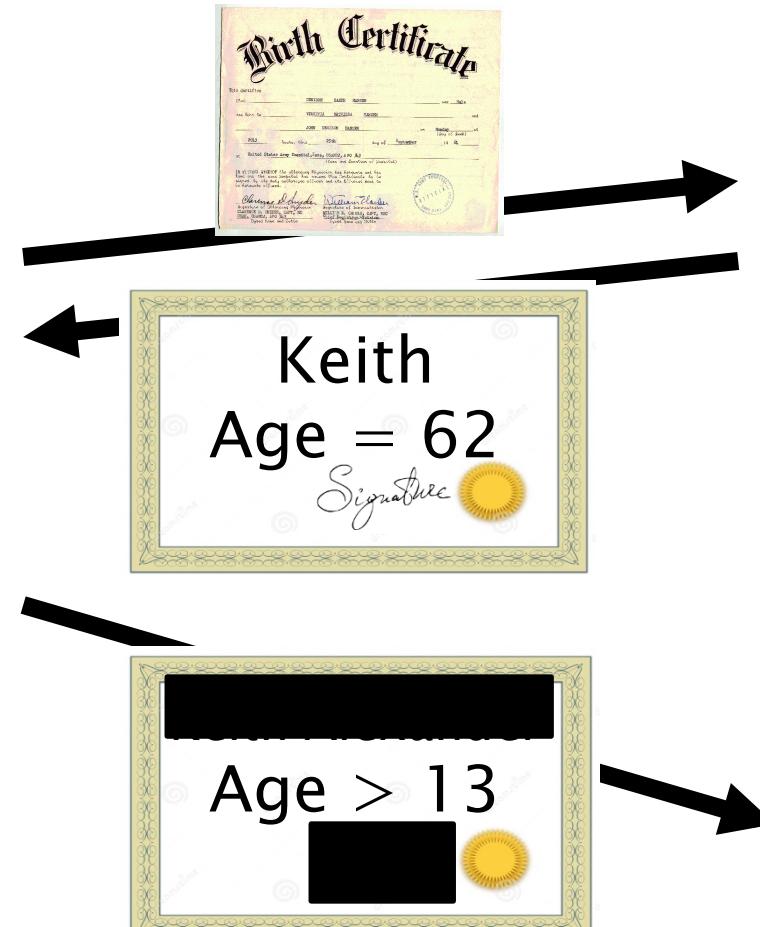
JUSTIN BIEBER
FAN CLUB



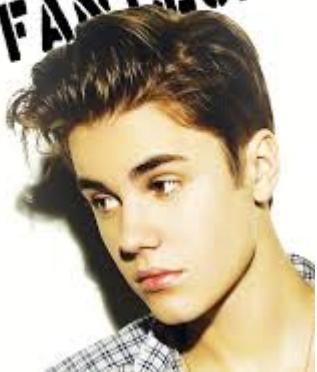
Example of Anonymous Credentials



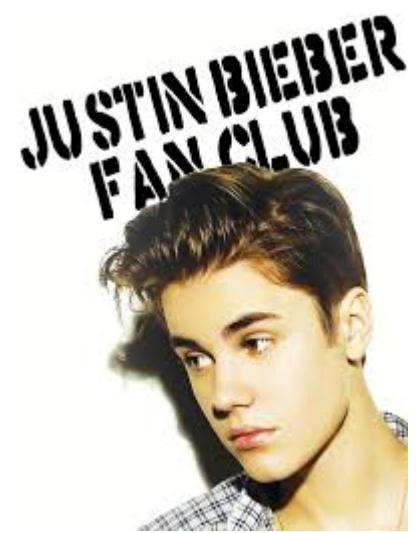
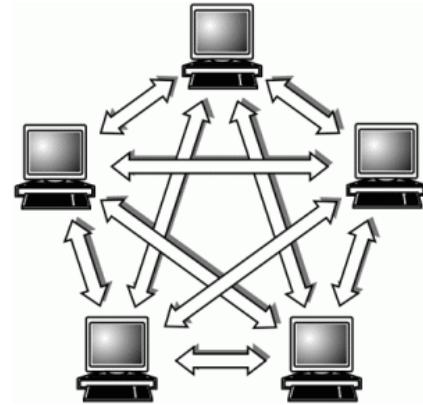
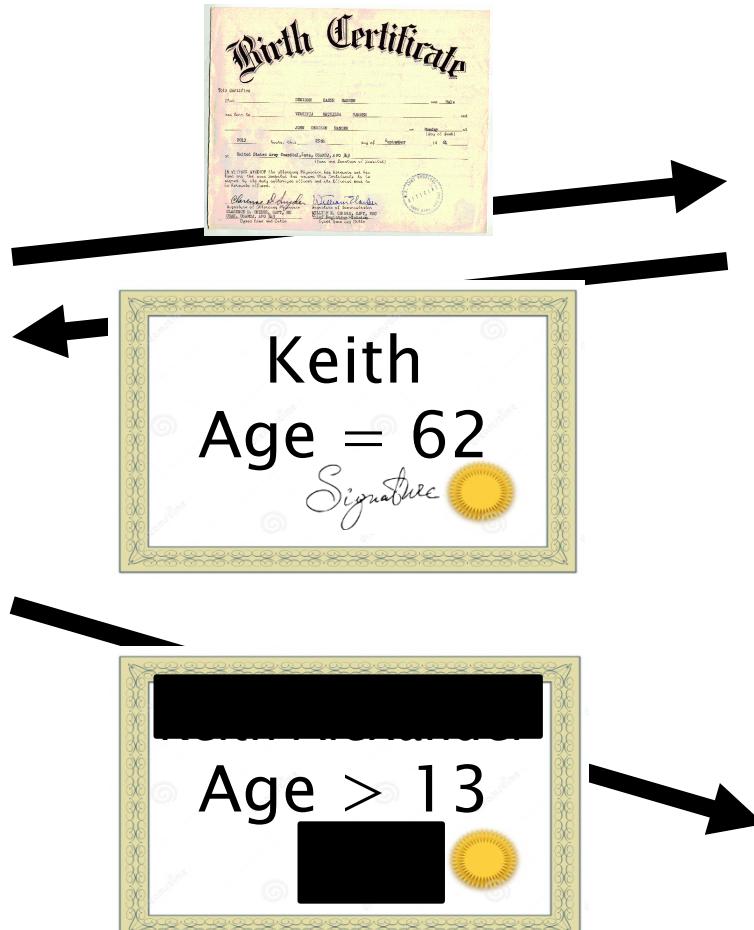
Problems?



JUSTIN BIEBER
FAN CLUB



Solution?



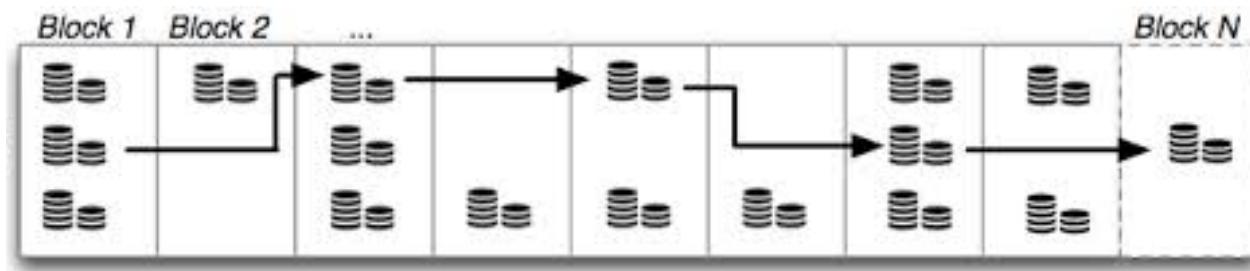
Our Contribution: Decentralized Anonymous Credentials

- ▶ Related to our electronic cash proposal [MGGR13]
 - Zerocoin (decentralized e-cash)
- ▶ **Decentralized anonymous credentials**
 - Decentralized credential issuance
 - Decentralized identity certification
 - Requires:
 - Public append-only ledger
 - Publicly verifiable identity claims



Public Append-Only Ledger

- ▶ Central ledger (audited by users)
- ▶ Broadcast networks
- ▶ Distributed consensus network
 - Bitcoin block chain

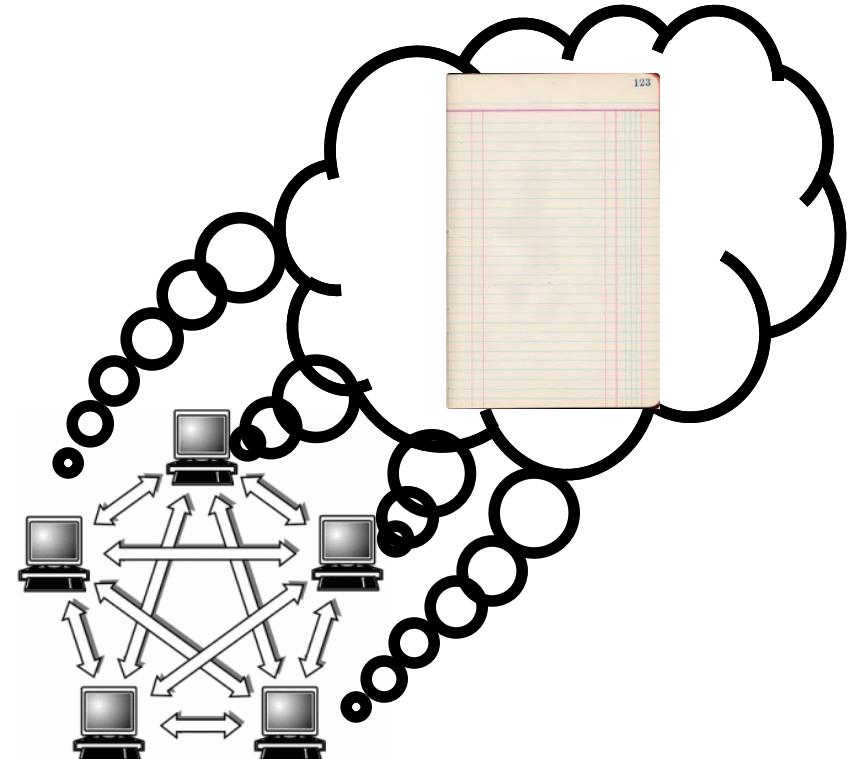


Publicly Verifiable Identity Claims

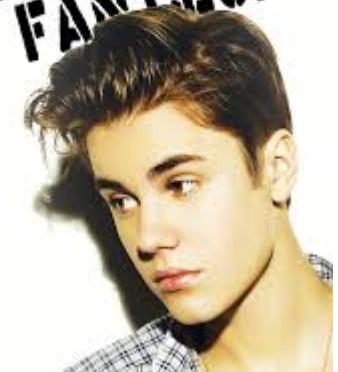
- ▶ Identity assertions are frequently publicly verifiable
- ▶ So why bother with (decentralized) anonymous credentials?
- ▶ Just because an identity assertion is publicly verifiable does not mean we want to link all of the information to every interaction!



Overview



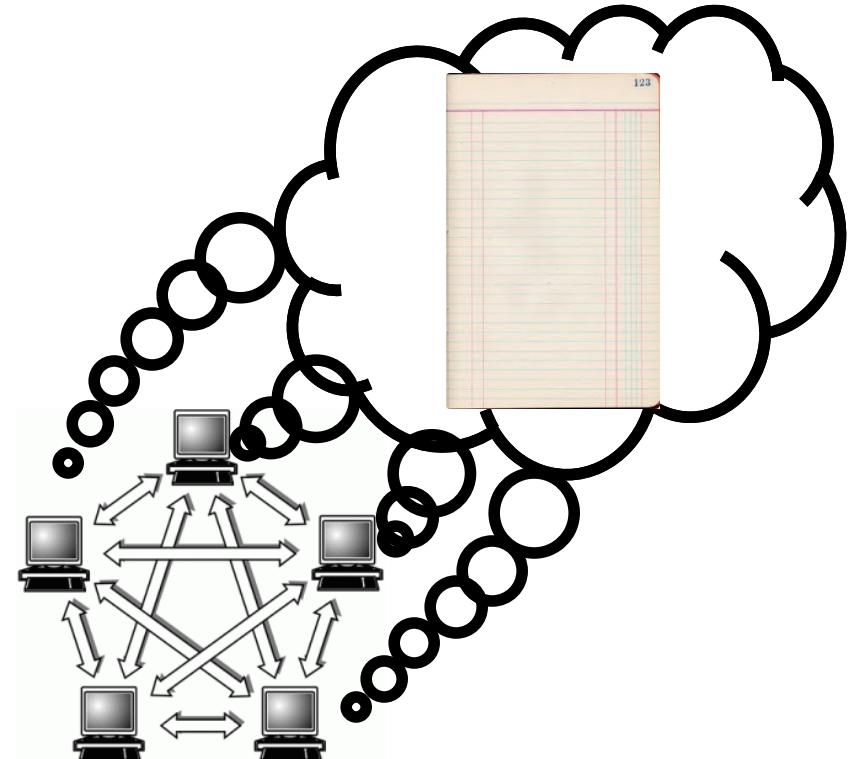
JUSTIN BIEBER
FAN CLUB



Overview



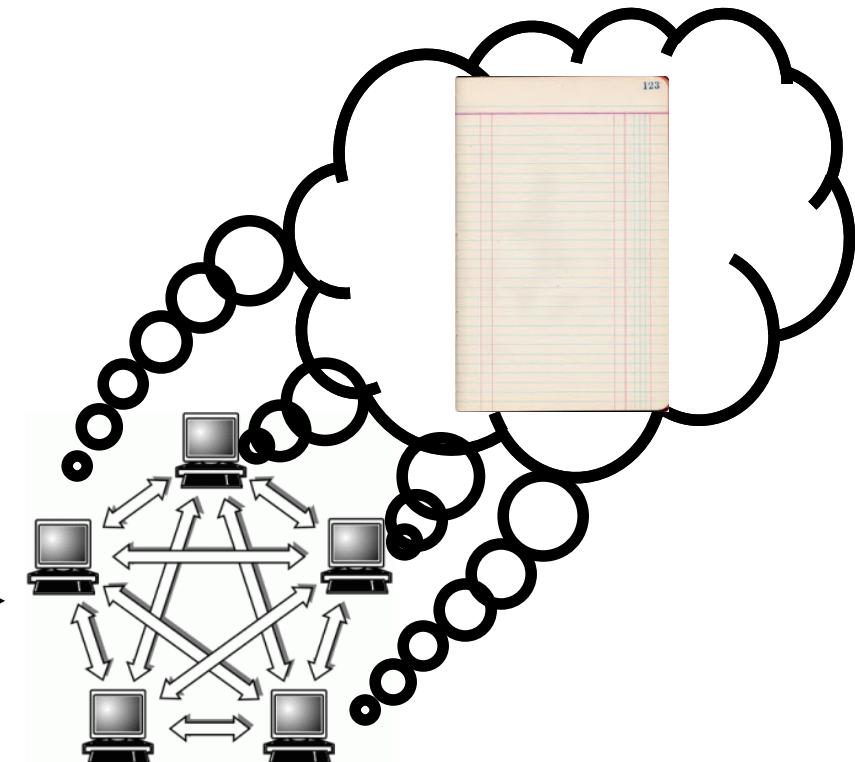
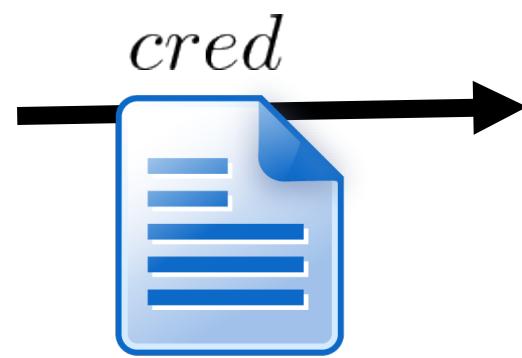
cred



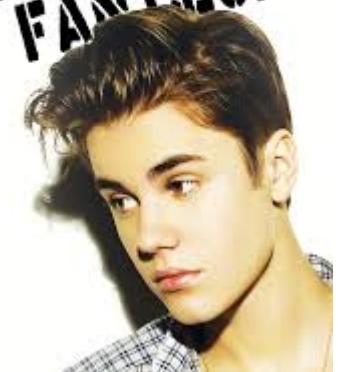
JUSTIN BIEBER
FAN CLUB



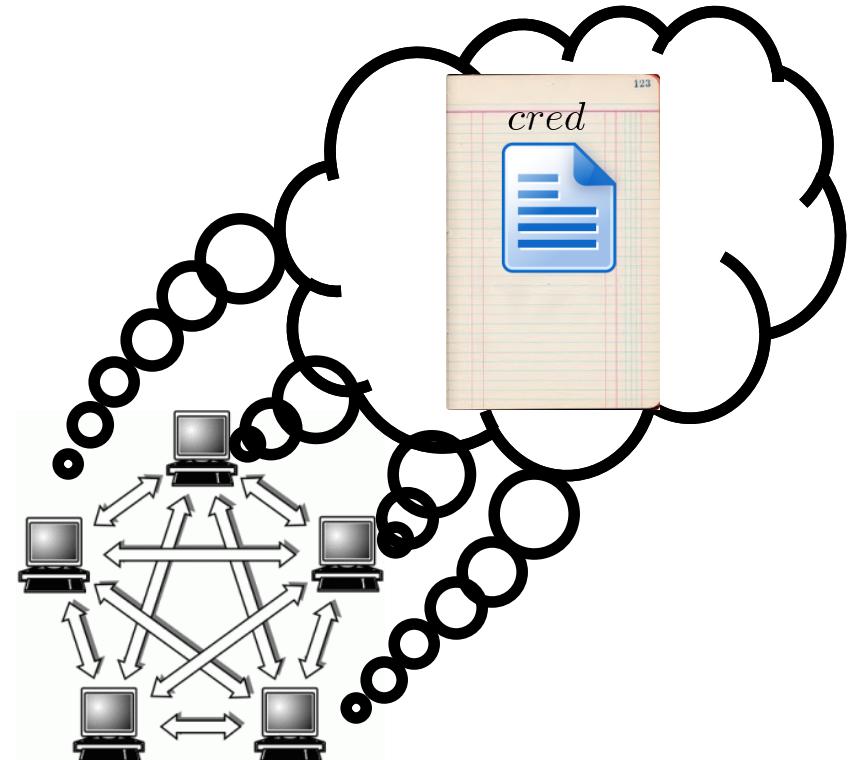
Overview



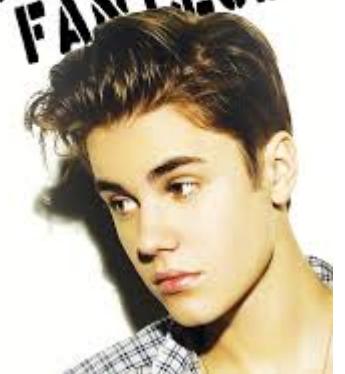
JUSTIN BIEBER
FAN CLUB



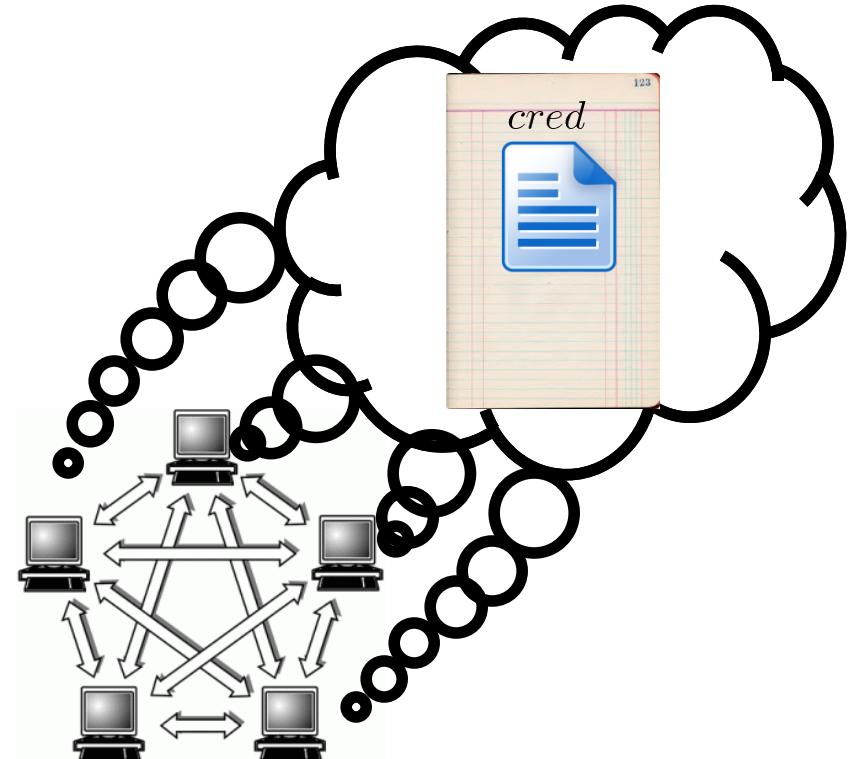
Overview



JUSTIN BIEBER
FAN CLUB

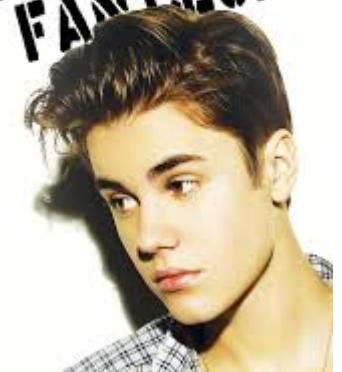


Overview

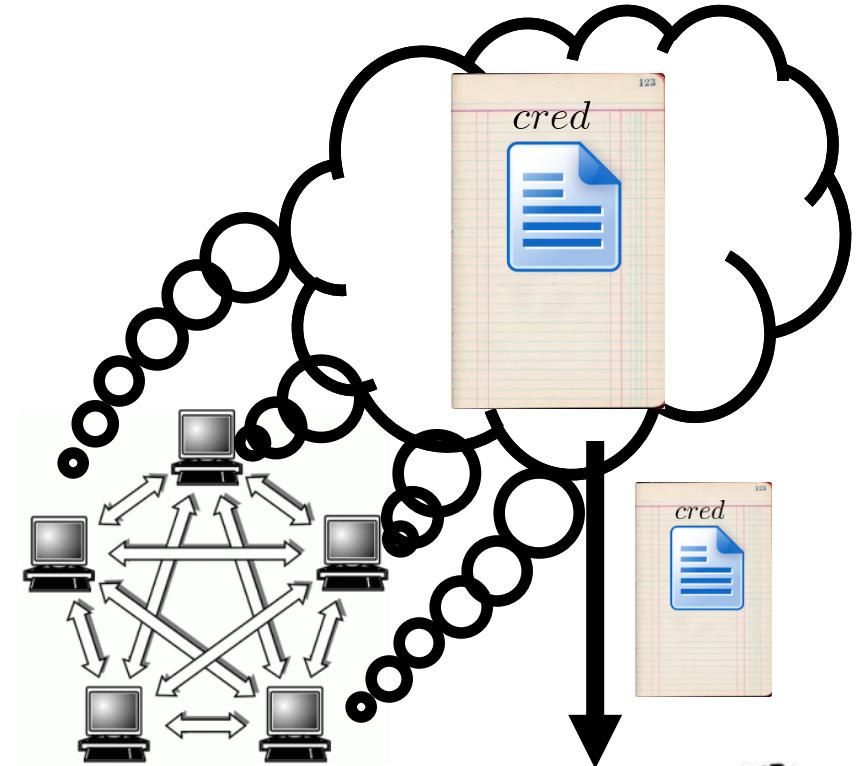


“A credential
on the ledger
says age > 13.”

JUSTIN BIEBER
FAN CLUB



Overview



“A credential
on the ledger
says age > 13.”

JUSTIN BIEBER
FAN CLUB



Cryptographic Building Blocks

- ▶ Commitments
- ▶ Zero-knowledge proofs
- ▶ Accumulators



Commitments

- ▶ Allow you to commit to and later reveal a value
- ▶ Binding: value cannot be tampered with
- ▶ Hiding: value cannot be read until revealed
- ▶ We use Pedersen commitments

$$C = g^x h^r \bmod q$$



Zero-knowledge Proofs

- ▶ Zero-knowledge [Goldwasser, Micali 1980s, and beyond]
- ▶ Prove a statement without revealing any other information
- ▶ Specific variant: non-interactive proof of knowledge
- ▶ Here we prove we know:
 1. The opening for a credential
 2. That the credential is in the ledger



An inefficient approach...

- ▶ Inefficient proof

- Identify all valid credentials in the ledger (call them $C\downarrow 1, \dots, C\downarrow N$)
- Prove that you know the opening of a credential C and $C = C\downarrow 1 \vee C = C\downarrow 2 \vee \dots \vee C = C\downarrow N$
- This “OR” proof is $O(N)$



Cryptographic Accumulators

- ▶ Allow constant size set membership proofs
- ▶ Strong RSA accumulator originally due to Benaloh and de Mare
- ▶ Efficient proof for accumulation of primes proposed by Camenisch and Lysyanskaya '01

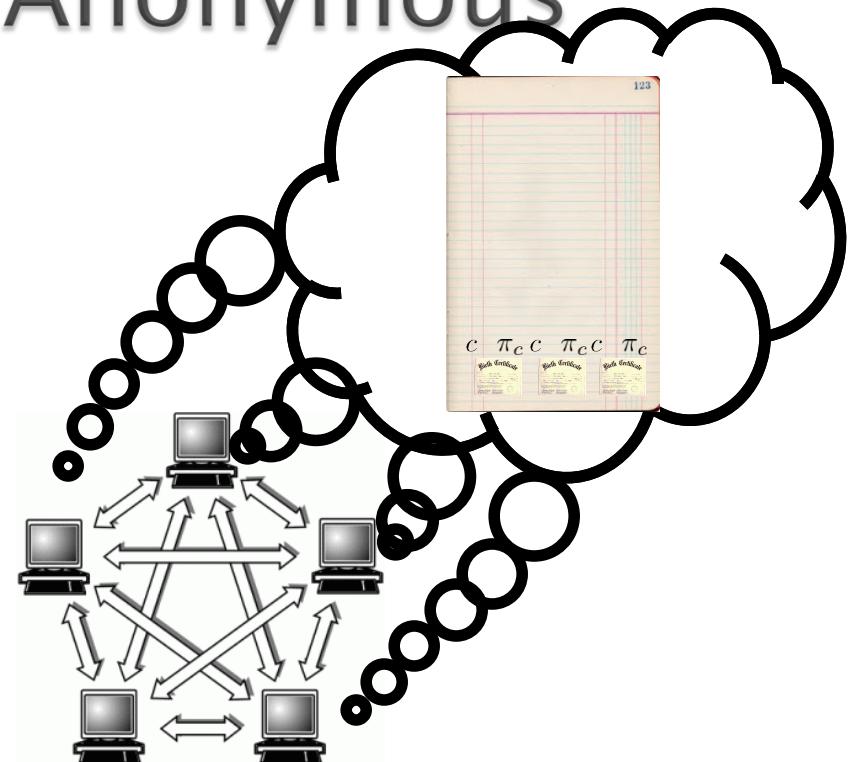
$$N = p \cdot q, u \in QR_N (u \neq 1)$$

$$A = u^{C_1 \cdot C_2 \cdot \dots \cdot C_n} \bmod N$$

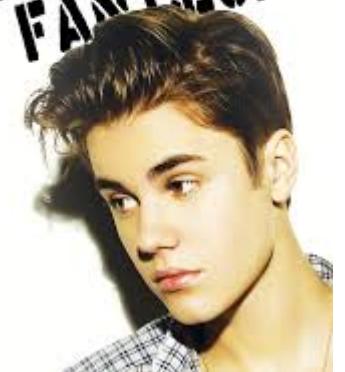
$$w_i = u^{C_1 \cdot C_2 \cdot \dots \cdot C_{i-1} \cdot C_{i+1} \cdot \dots \cdot C_n} \bmod N$$



Basic Decentralized Anonymous Credentials



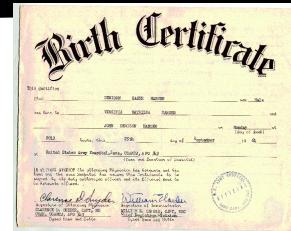
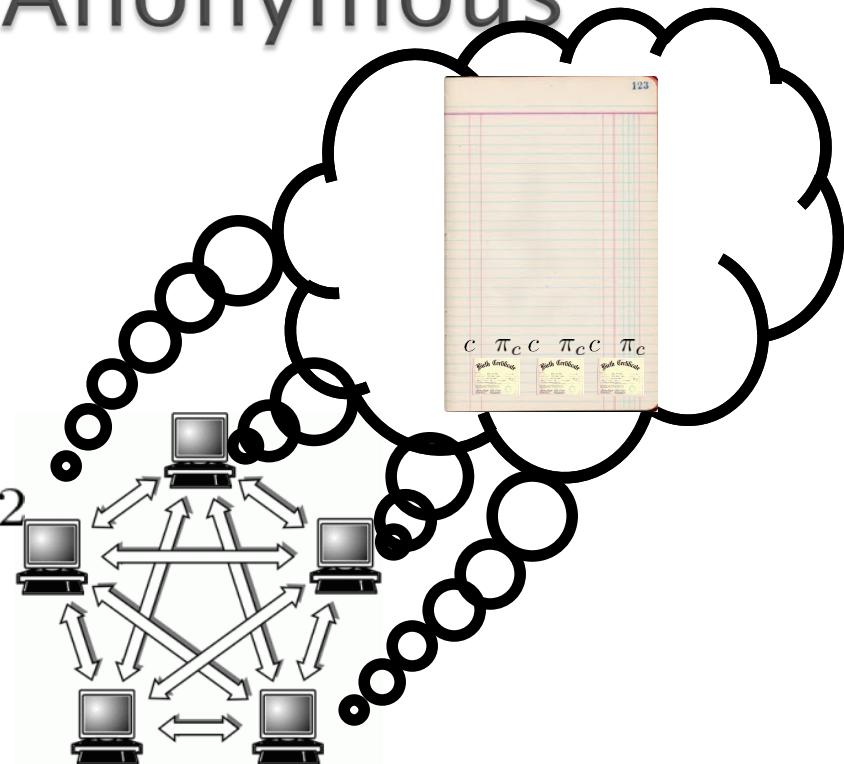
JUSTIN BIEBER
FAN CLUB



Basic Decentralized Anonymous Credentials



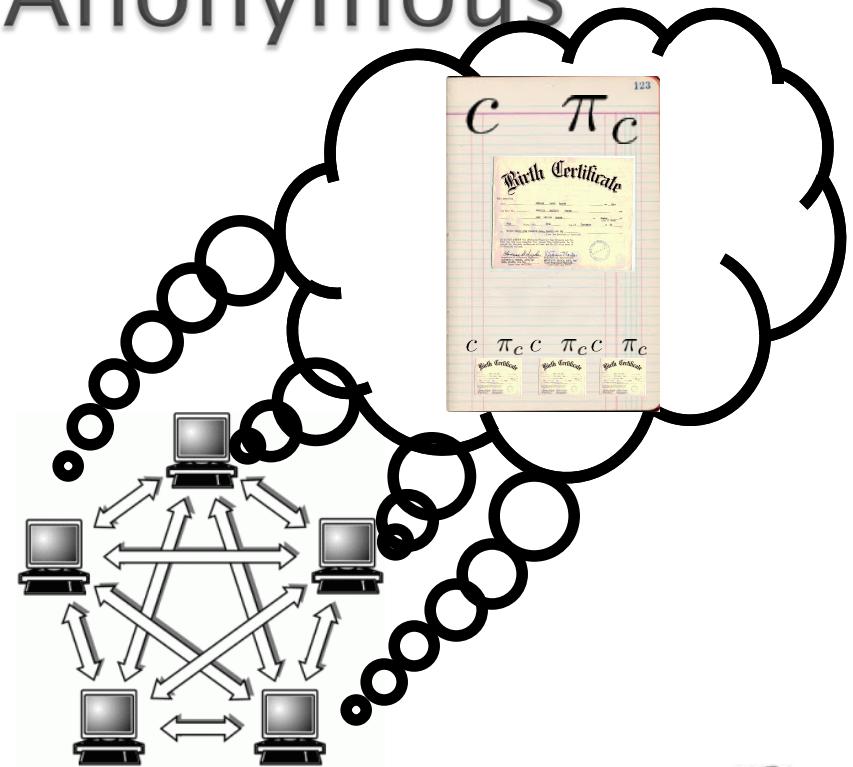
$$c = g_0^r g_1^{sk} g_2^{\text{age}=62}$$

 π_c 

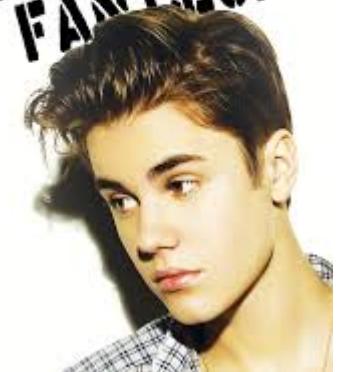
JUSTIN BIEBER
FAN CLUB



Basic Decentralized Anonymous Credentials



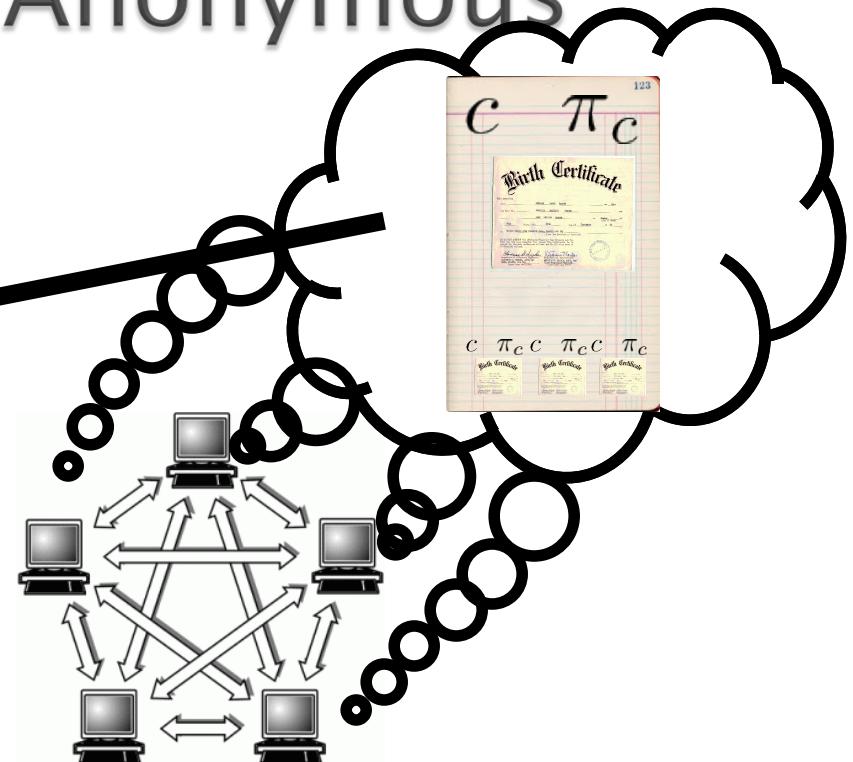
JUSTIN BIEBER
FAN CLUB



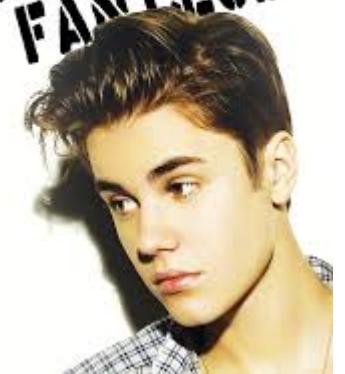
Basic Decentralized Anonymous Credentials



c_1, c_2, \dots, c_n

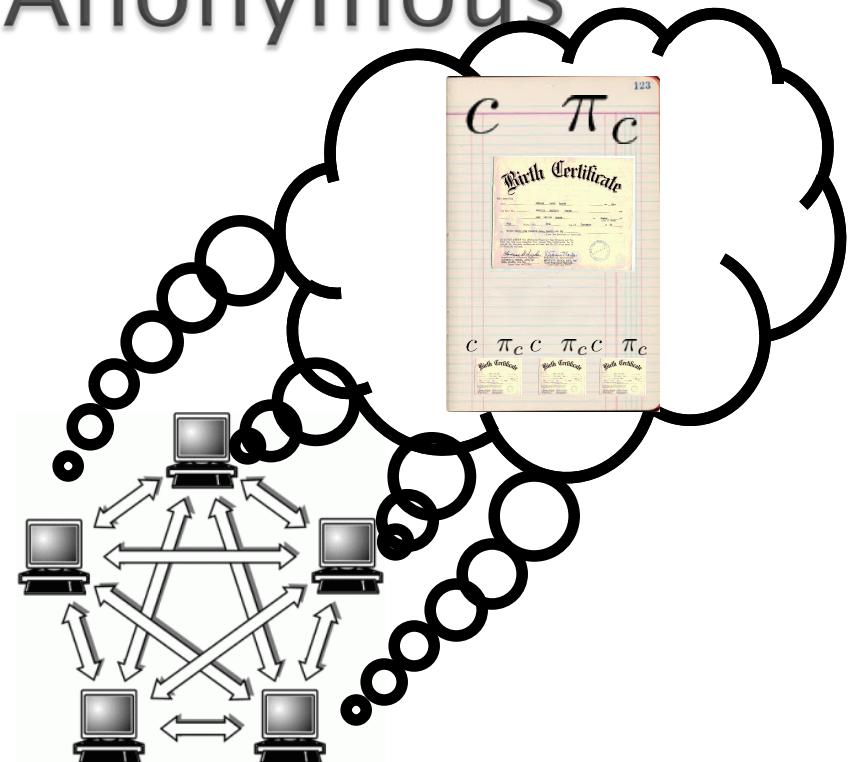


JUSTIN BIEBER
FAN CLUB

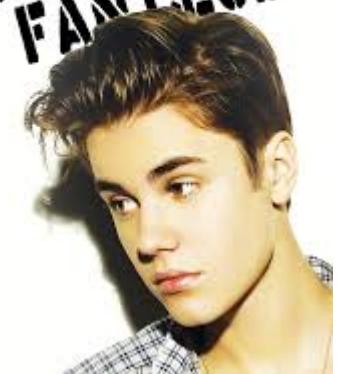


Basic Decentralized Anonymous Credentials

$$A = u^{c_1 \cdot c_2 \cdot \dots \cdot c_n} \bmod N$$

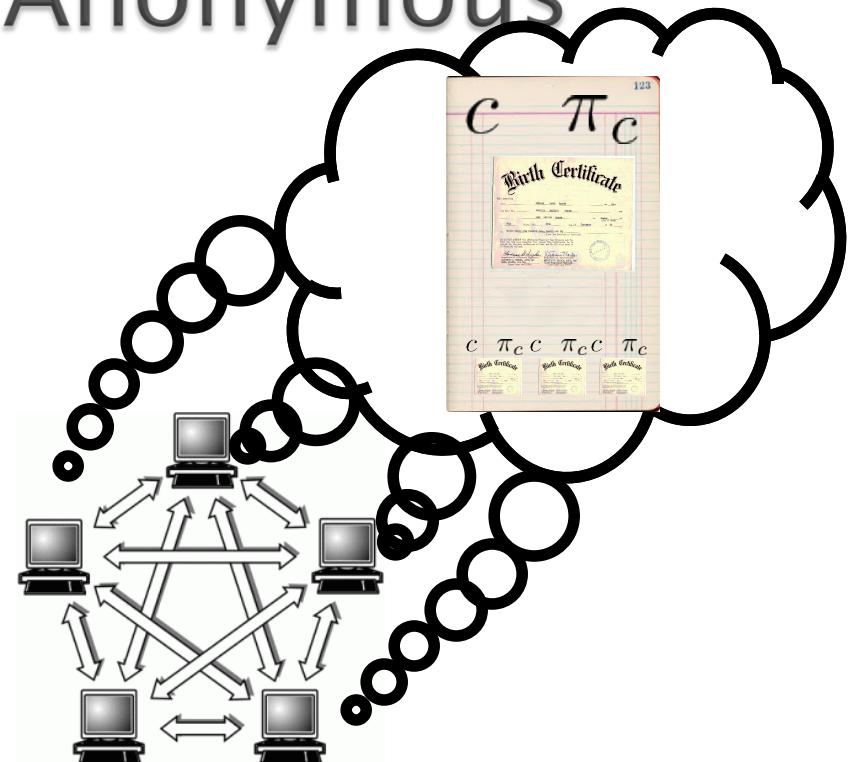
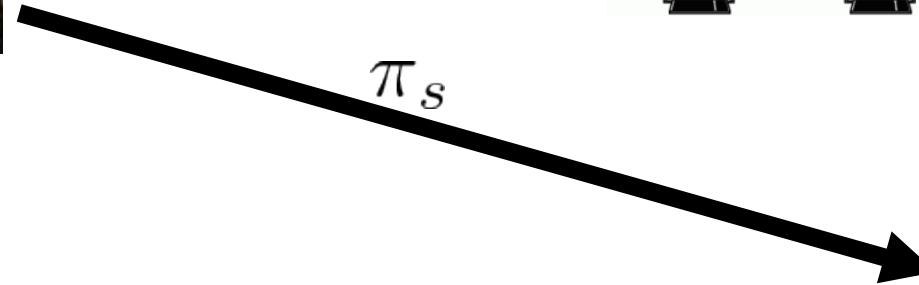


JUSTIN BIEBER
FAN CLUB

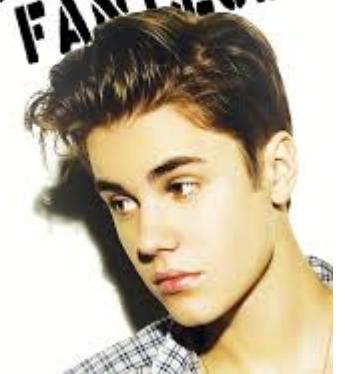


Basic Decentralized Anonymous Credentials

$$A = u^{c_1 \cdot c_2 \cdots c_n} \bmod N$$

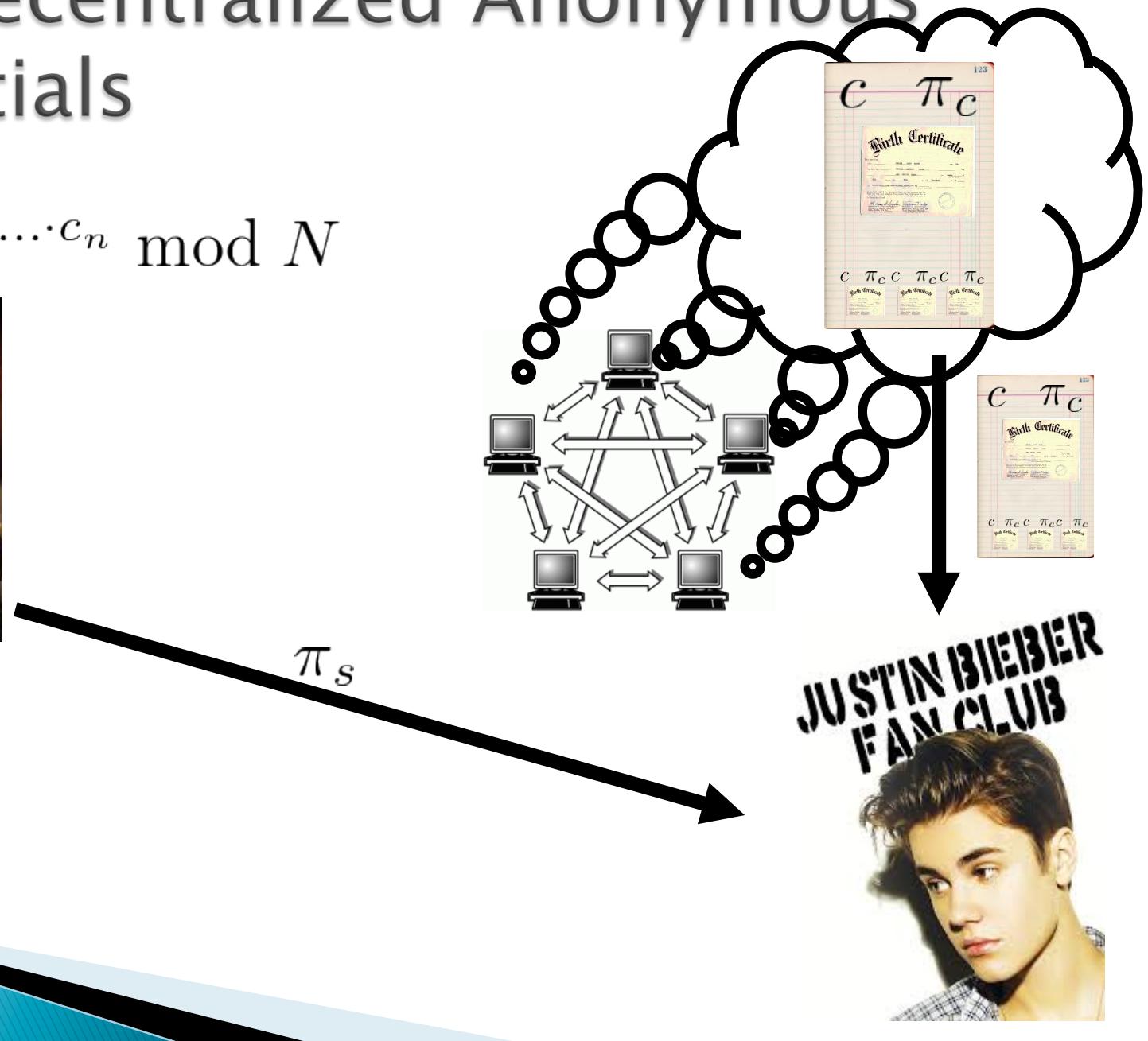


JUSTIN BIEBER
FAN CLUB



Basic Decentralized Anonymous Credentials

$$A = u^{c_1 \cdot c_2 \cdots c_n} \bmod N$$

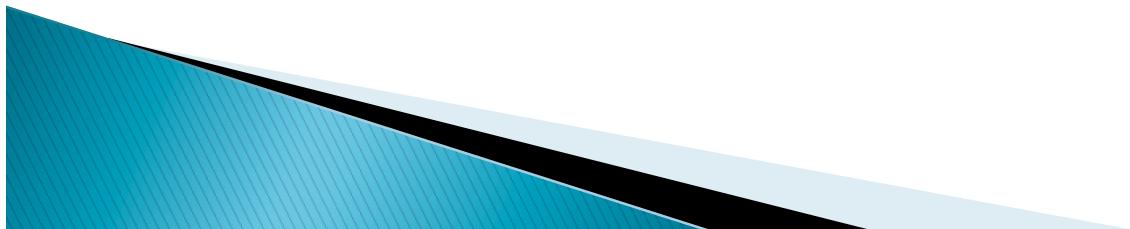
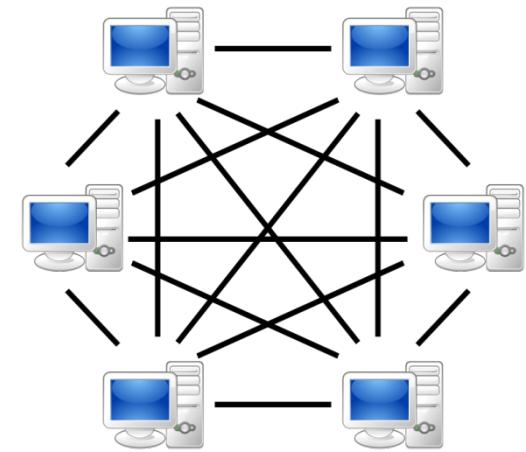


Applications

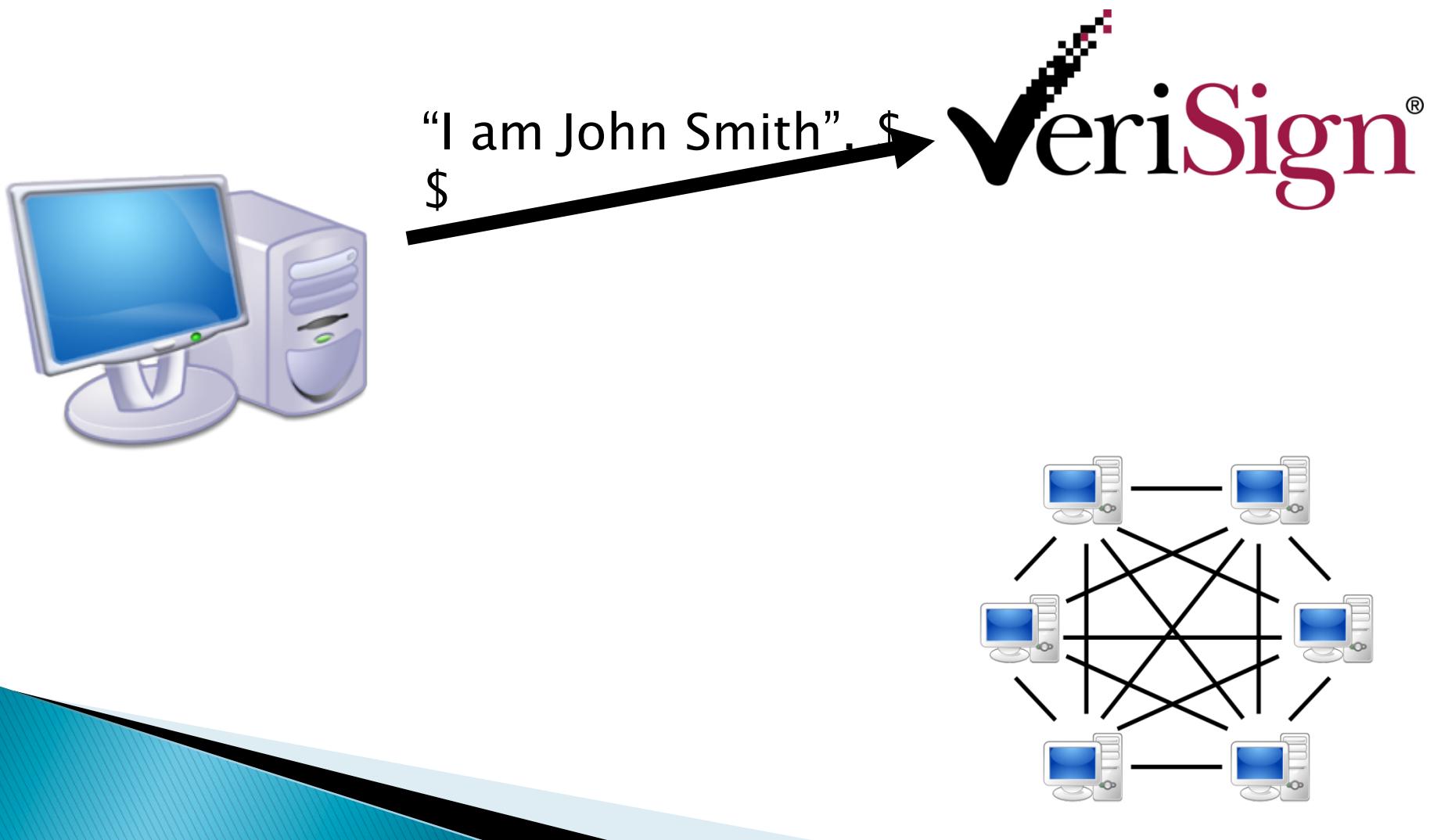
- ▶ Anonymous resource management in ad hoc networks
- ▶ Decentralized Direct Anonymous Attestation
- ▶ Auditable credentials
- ▶ Mitigating Sybil attacks in ad hoc networks



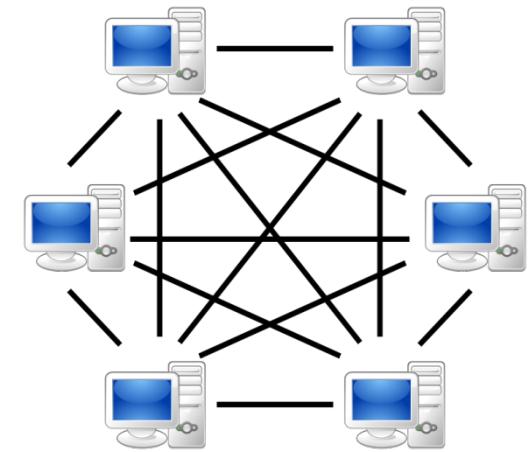
Protecting Against Sybil Attacks



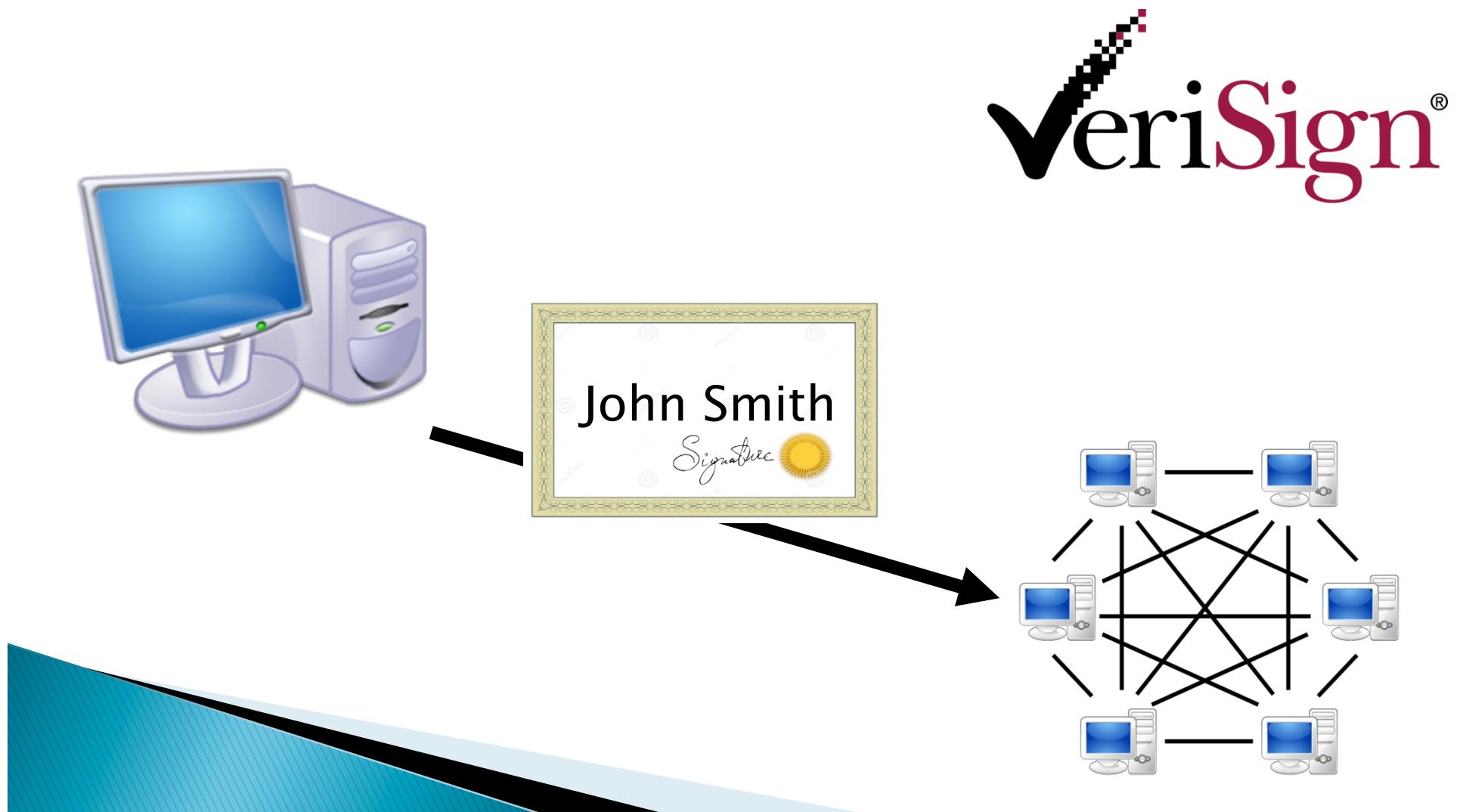
Protecting Against Sybil Attacks



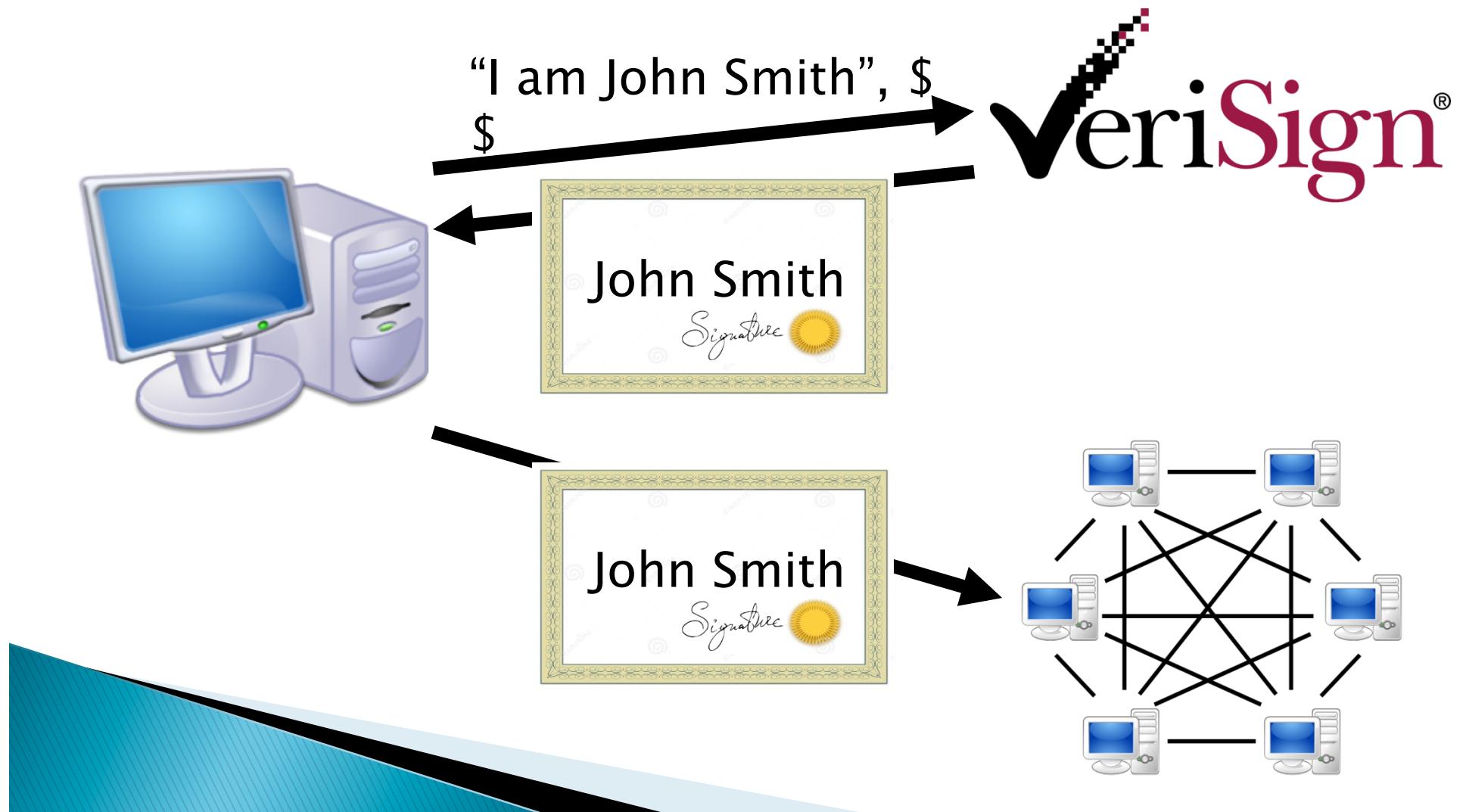
Protecting Against Sybil Attacks



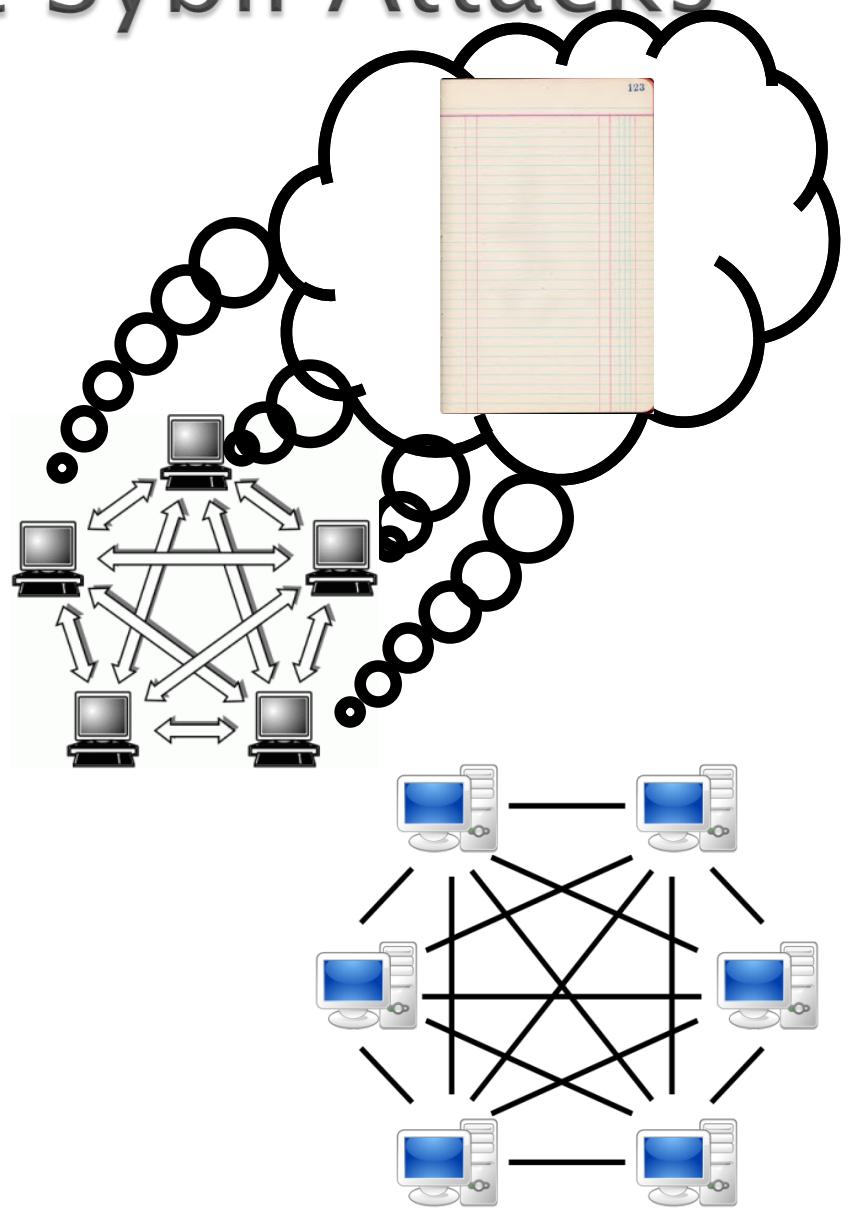
Protecting Against Sybil Attacks



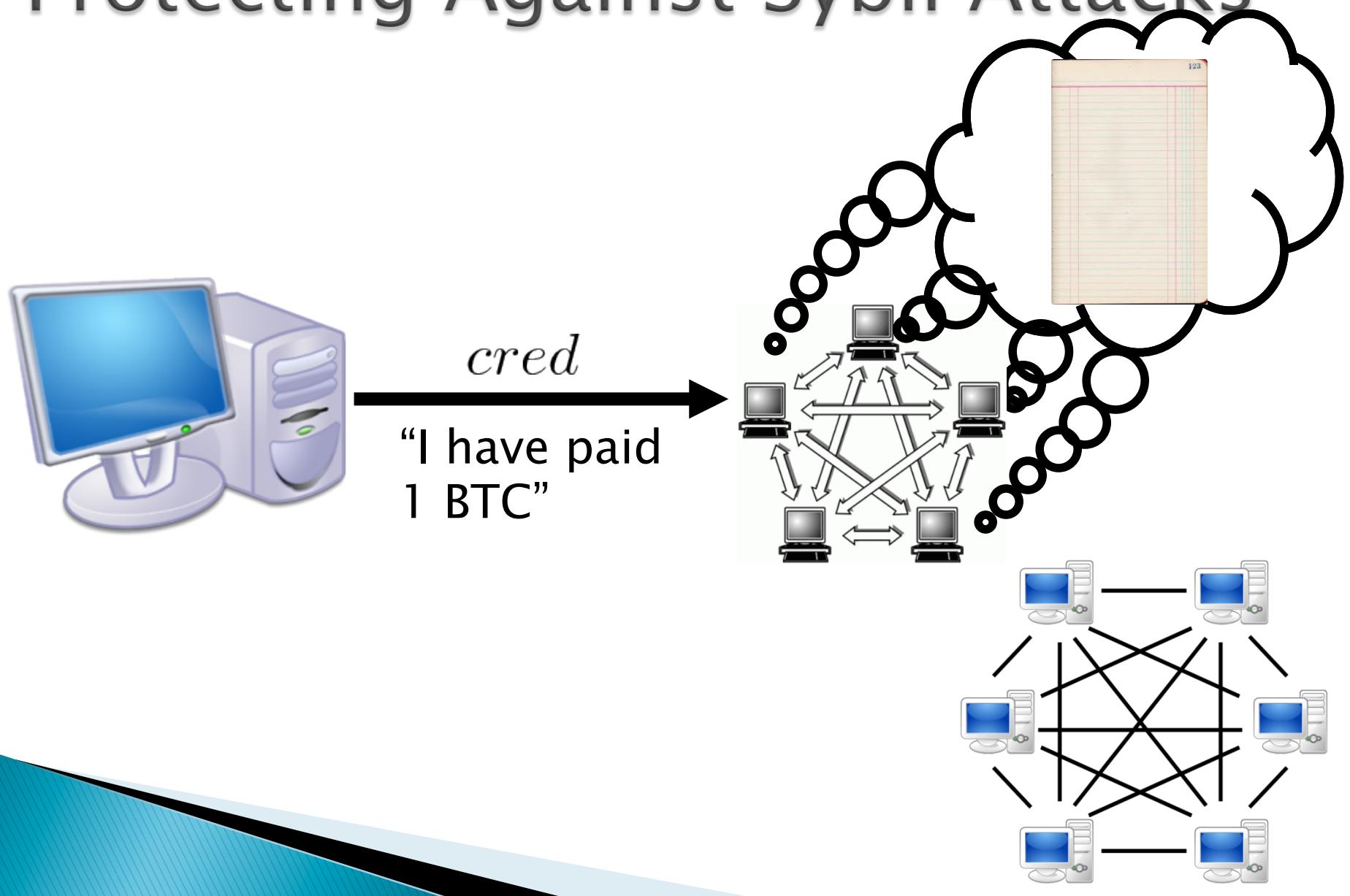
Protecting Against Sybil Attacks



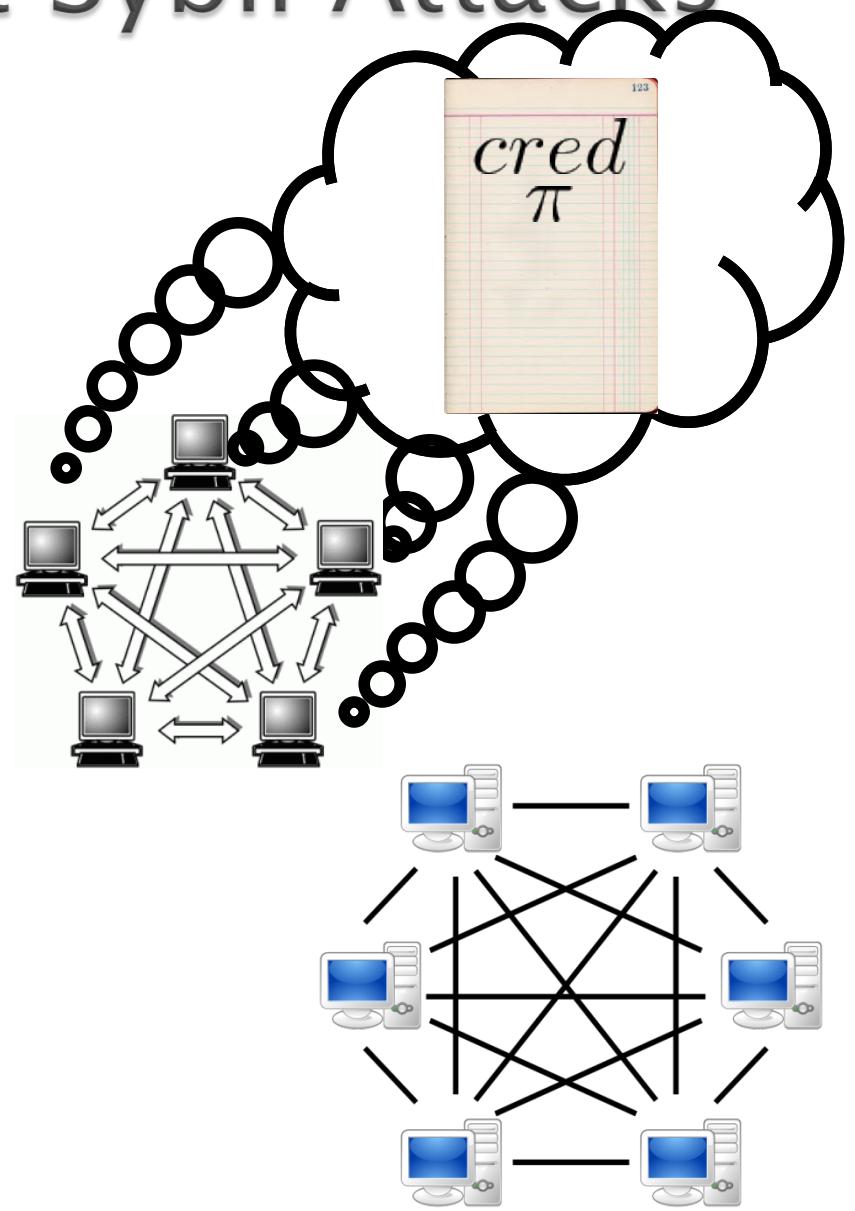
Protecting Against Sybil Attacks



Protecting Against Sybil Attacks



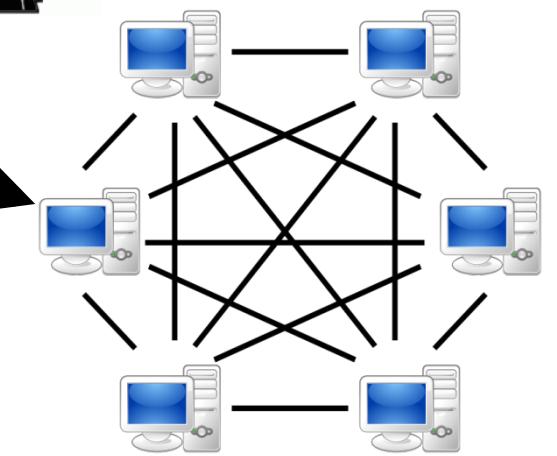
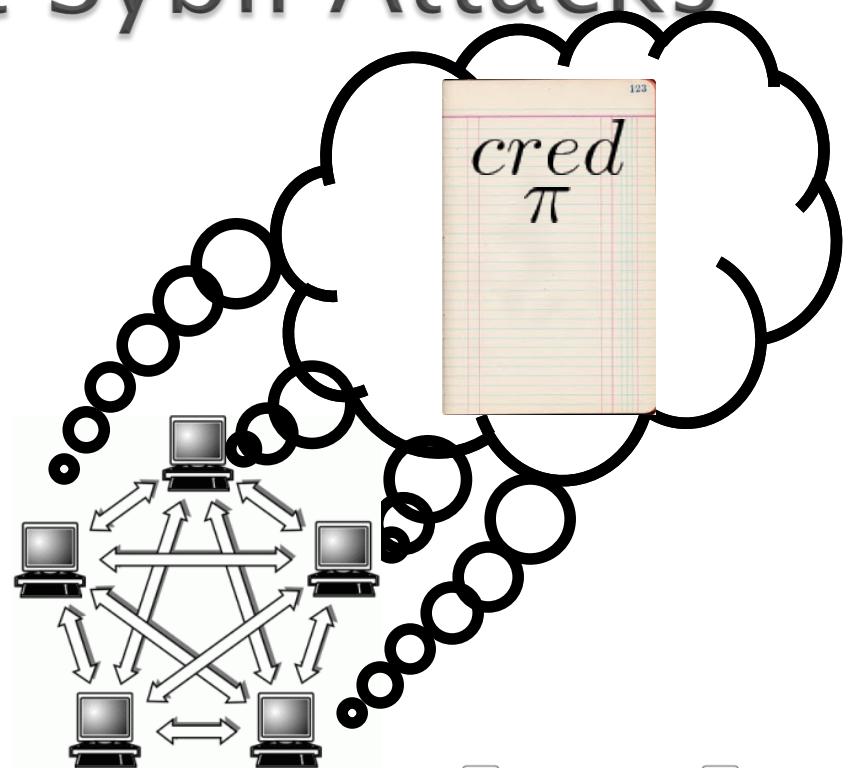
Protecting Against Sybil Attacks



Protecting Against Sybil Attacks



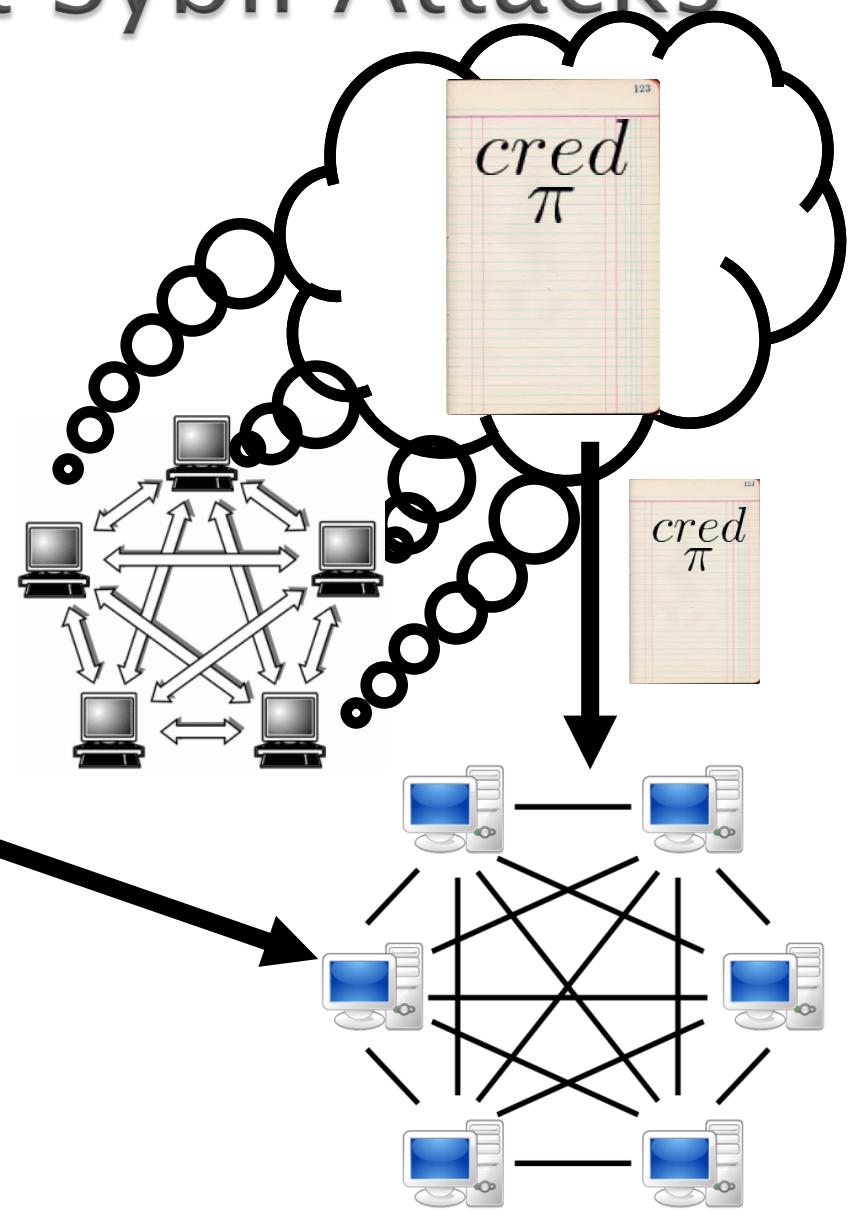
“I am not
a Sybil”



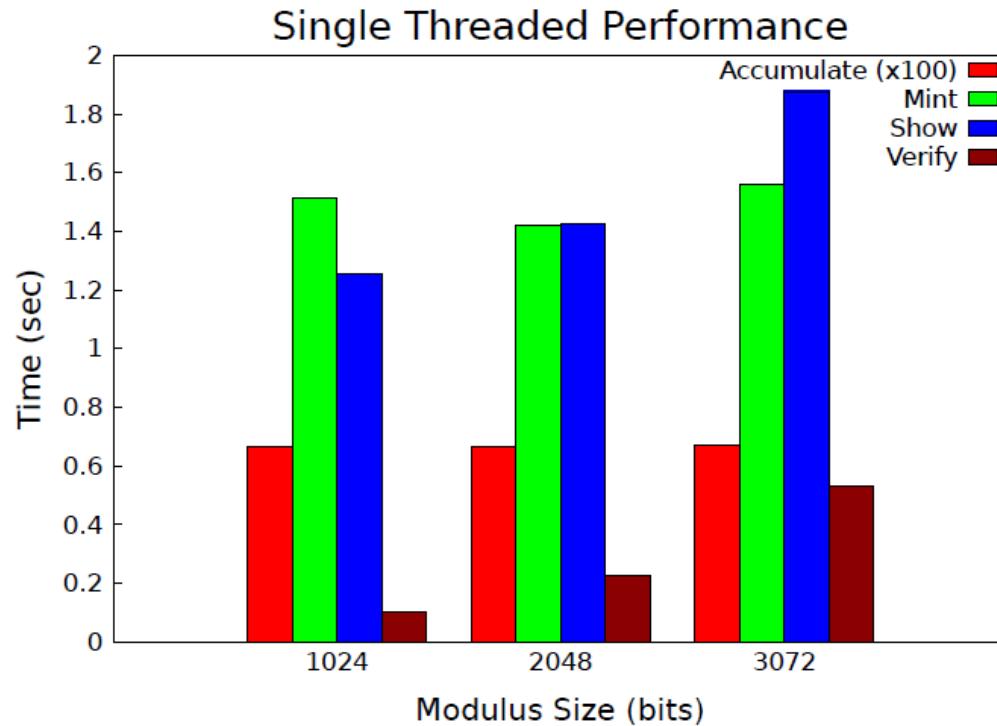
Protecting Against Sybil Attacks



“I am not
a Sybil”



Performance



- ▶ Basic scheme implemented as stand-alone library
 - Proofs 50 KB

Future Work

- ▶ Better, smaller “proofs” of knowledge:
- ▶ Succinct Non-Interactive ARguments of Knowledge (zkSNARKs) [PHGR13, BCGTV13]
 - 288 byte proof for arbitrary-sized arithmetic circuits
 - 8 ms verification time
- ▶ Additional applications?



Questions?



JOHNS HOPKINS
UNIVERSITY

Potential Alternatives

- ▶ Threshold cryptography
 - High setup cost for large number of parties
 - Difficult for parties to come and go
- ▶ Ring signatures [RST01]
 - Grow linearly with the number of participating signers
 - Expensive to generate



Non Publicly Verifiable Credentials

- ▶ Credential transform service
- ▶ Allows user to transform a credential to an anonymous credential without additional trust assumption
- ▶ Works for any statement that an authority can certify



Proof of Work for Sybil Attacks

- ▶ Proof of resource expenditure instead of payment
- ▶ Cannot reuse proof of work with different peers
 - Not anonymous
 - Clonable
- ▶ Do not want to have to do a proof of work with each peer in the system
- ▶ Instead do one proof of work per k interactions



Resource Management

- ▶ Publicly verifiable proofs of resources
- ▶ File storage, bandwidth, etc.
- ▶ Do not want to link resources provided to resources consumed
 - Files uploaded vs. files downloaded

