# Reverse Words in a String

```
Given an input string, reverse the string word by word.
```

```
For example,

Given s = "the sky is blue",

return "blue is sky the".

public class Solution {
    public shiring reverseWords(String s) {
        StringsUffer revStr = new StringsUffer():
        if (s == null || s.length() == 0)
        {
            return ";
        }
        Strings[] str = s.split(");
        intn = strlength;
        for (int1 = n - 1; 1>= 0; -)
        {
            (if (stri), equals("))
            {
                  revStr.append(strii), append(");
            }
        }
        return revStr.length() == 0 ? ": revStr.substring(0, revStr.length() - 1);
    }
```

# **Evaluate Reverse Polish Notation**

Evaluate the value of an arithmetic expression in Reverse Polish Notation.

Valid operators are +, -, \*, /. Each operand may be an integer or another expression.

Some examples:

# Sort List

Sort a linked list in  $O(n \log n)$  time using constant space complexity.

```
cur.next = hi;
hi = hi.next;
        }
cur = cur.next;
     if (lo != null)
        cur.next = lo;
     }
else if (hi != null)
{
      }
return head.next;
 Reorder List
Given a singly linked list L: L0 \rightarrow L1 \rightarrow ... \rightarrow Ln-1 \rightarrow Ln,
reorder it to: L0 \rightarrow Ln \rightarrow L1 \rightarrow Ln-1 \rightarrow L2 \rightarrow Ln-2 \rightarrow ...
 You must do this in-place without altering the nodes' values.
For example,
Given \{1,2,3,4\}, reorder it to \{1,4,2,3\}.

    break the linked list in the middle into two list;
    reverse the second linked list
    merge the two list by inserting the second one after the first one.

fast = fast.next.next;
slow = slow.next;
      }
ListNode second = slow.next;
slow.next = null;
      second = reverseOrder(second);
ListNode curr = head;
ListNode curr2 = second;
while (curr2!= null)
        ListNode temp1 = curr.next;
ListNode temp2 = curr2.next;
        curr.next = curr2;
curr = temp1;
curr2.next = curr;
        curr2 = temp2;
    public ListNode reverseOrder(ListNode list)
      if (list == null || list.next == null)
     return list;
     ListNode currNode, nextNode;
currNode = list;
nextNode = list.next;
list.next = null;
while(nextNode!= null)
        ListNode loopNode = nextNode.next;
nextNode.next = currNode;
currNode = nextNode;
nextNode = loopNode;
      list = currNode;
  return list;
 Linked List Cycle II
 Given a linked list, return the node where the cycle begins. If there is no cycle, return null.
 Follow up:
Can you solve it without using extra space?
 循环中奇数和偶数个结点;
l.使用两个指针slow,fast。两个指针都从表头开始走,slow每次走一步,fast每次走两步,如果fast遇到null,则说明没有环,返回false;如果slow==fast,说明有环
2.第一次相遇后,让slow,fast继续走.slow从head开始走,slow,fast各走一步,再次相遇的就是循环起始结点
if (fast == null || fast.next == null) {
          return null;
```

```
if (fast == slow)
         slow = slow.next;
fast = fast.next;
 Linked List Cycle
 Given a linked list, determine if it has a cycle in it.
 Can you solve it without using extra space?
return false;
      ListNode slow = head;
ListNode fast = head;
ListNode fast = head;
while (true)
{
if (fast == null || fast.next == null)
{
return false;
           return false;
         }
slow = slow.next;
fast = fast.next.next;
if (slow == fast)
{
            return true;
 Binary Tree Preorder Traversal
 Given a binary tree, return the preorder traversal of its nodes' values.
 Given binary tree {1,#,2,3},
 return [1,2,3].
 Recursive solution
/**
    * Definition for binary tree
    * public class TreeNode {
    * int vol;
    * TreeNode left;
    * TreeNode light;
    * TreeNode(int x) { val = x; }
    *}
 */
public class Solution {
   public List<nteger> preorderTraversal(TreeNode root) {
      List<nteger> list = new ArrayList<Integer> (|;
      if (root == null)
         return Collections.EMPTY_LIST;
      }
list.add{root.val);
list.addAll(preorderTraversal(root.left));
list.addAll(preorderTraversal(root.right));
...adAll(
return list;
}
 Binary Tree Postorder Traversal
 Given a binary tree, return the postorder traversal of its nodes' values.
 Given binary tree {1,#,2,3},
 return [3,2,1].
 Note: Recursive solution is trivial, could you do it iteratively?
```

/\*\*
 \* Definition for binary tree
 \* public class TreeNode {
 int val;

```
* TreeNode left;
* TreeNode injth:
* TreeNode injth:
* TreeNode (injt) { val = x; }
*/
public class Solution {
    public List-integer> postorderTraversol([TreeNode root) {
        List-integer> list = new ArrayList-integer>();
        if (root == null)
        {
            return Collections.EMPTY_LIST;
            ist.addAll(postorderTraversol(root.left));
            ist.addAll(postorderTraversol(root.right));
            ist.addAll(postorderTraversol(root.right));
            return list;
}
```

#### Word Break

Given a string s and a dictionary of words dict, determine if s can be segmented into a space-separated sequence of one or more dictionary words.

### Word Break II

Given a string s and a dictionary of words dict, add spaces in s to construct a sentence where each word is a valid dictionary word.

Return all such possible sentences.

```
For example, given
s = "catsanddog".
dict = ["cat", "cats", "and", "sand", "dog"].
A solution is ["cats and dog", "cat sand dog"].
Submission Result: Time Limit Exceeded
```

```
public class Solution {
    public last-String> wordBreak(String s, Set-String> dict) {
        List-String> wist = new ArrayList-String>();

        if (s == null || dict.size() == 0)
        {
            return list;
        }
        helpbuilder(s, dict, 0, "", list);
        return list;
    }
    private void helpBuilder(String s, Set-String> dict, int start, String litem, List-String> list){
        if (start > slength()){
            inst add(lemn);
            return
        inst add(lemn);
        return
        if (dict.contains(curnStr.toString())){
            curnStr.append(s.charAf(i));
        if (dict.contains(curnStr.toString())){
            String append(s.charAf(i));
        if (dict.contains(curnStr.toString())) ? ? (item + "" + curnStr.toString()) : curnStr.toString();
        helpBuilder(s, dict, i + 1, newItem, list);
    }
}
```

# Single Number

Given an array of integers, every element appears twice except for one. Find that single one.

#### Note:

Your algorithm should have a linear runtime complexity. Could you implement it without using extra memory?

```
int cnt = 0;
int x = cnt+;
cnt = 0;
int y = +cnt;
System.out.println(x + " " + y); //0 1//
int cnt = 0;
```

```
public class Solution {
  public int singleNumber(int[] A) {
    if (A.length <= 1 ) {
      return A[0];
    }</pre>
      int len = A.length;
HashMap<Integer, Integer> map = new HashMap<Integer, Integer>();
     for (int i = 0; i < len; i++) {
    if (!map.containskey(A[i])) {
        int count = 0;
        map.put(A[i], ++count);
    }
        } else { map.put(A[i], map.get(A[i]) + 1);
     int results = 0;
for (Integer j : map.keySet()) {
    if (map.get(j) == 1) {
        results = j;
    }
return results;
 Candy
 There are N children standing in a line. Each child is assigned a rating value.
 You are giving candies to these children subjected to the following requirements:

    Each child must have at least one candy.

    Children with a higher rating get more candies than their neighbors.

What is the minimum candies you must give?
通常是要求的变量跟左右元素有关系的题目: 两边扫描的方法
一些例子:
12333 =》 8 因为candy数可以是
12311
12323=》9 因为candy数可以是
12312
 思路:
 1
        d[i] 是给第i个小孩最少几块糖 rank[i] > rank[i-1],必须比前一个多给一块,d[i] = d[i-1] + 1 rank[i] <= rank[i-1],两个排名一样,第二个就给一块就行了,d[i] = 1
 基本思路就是进行两次扫描,一次从左往右,一次从右往左。第一次扫描的时候维护对于每一个小孩左边所需要是少的糖果数量。存入数组对应元素中,第二次扫描的时候维护右边所需的最少糖果数,并且比较将左边和右边大的糖果数量
存入结果数组对应元素中。这样两遍扫描之后就可以得到每一个所需要的最最少糖果量,从而累加得出结果。方法只需要两次扫描,所以时间复杂度是O(2*n)=O(n)。空间上需要一个长度为n的数组,复杂度是O(n)。
public class Solution {
   public int candy(int[] ratings) {
    if (ratings.length == 0) {
        return 0;
    }
      int[] num = new int[ratings.length];
num[0] = 1;
     for (int i = 1; i < ratings.length; i++) {
    if (ratings[i] > ratings[i - 1]) {
        num[i] = num[i - 1] + 1;
    }
       }
else if (ratings[i] < ratings[i - 1]) {
num[i] = 1;
     int res = num[ratings.length - 1];
for (int i = ratings.length - 2; i >= 0; i-) {
   int curr = 1;
   if (ratings[i] > ratings[i + 1]) {
      curr = num[i + 1] + 1;
        }
res += Math.max(num[i], curr);
num[i] = curr;
```

#### Gas Station

return res; }

There are N gas stations along a circular route, where the amount of gas at station i is gas[i].

You have a car with an unlimited gas tank and it costs cost[i] of gas to travel from station i to its next station (i+1). You begin the journey with an empty tank at one of the gas stations.

Return the starting gas station's index if you can travel around the circuit once, otherwise return -1.

#### Note:

The solution is guaranteed to be unique.

```
 \begin{aligned} & \text{public class Solution} \, \{ \\ & \text{public int canCompleteCircuit(int] gas, int[] cost)} \, \{ \\ & \text{if (gas.length} = 0 \, || \, \text{cost.length} \, |= 0 \, || \, \text{gas.length} \, |= \, \text{cost.length}) \, \{ \\ & \text{return} - 1; \\ & \} \\ & \text{int sum} = 0; \\ & \text{int total} = 0; \\ & \text{int total} = 0; \\ & \text{int start} = 0; \\ & \text{for (int i} = 0; i < \, \text{gas.length}; i++) \, \{ \\ & \text{sum} + 2 \, \text{gas[i]} - \, \text{cost[i]}; \\ & \text{total} + 2 \, \text{sum}; \\ & \text{if (sum} < 0) \, \{ \end{aligned}
```

```
start = i + 1;
sum = 0;
}
      if (total < 0) {
return -1;
return start;
 Palindrome Partitioning
 Given a string s, partition s such that every substring of the partition is a palindrome.
 Return all possible palindrome partitioning of s.
 For example, given s = "aab",
 Return
  Ε
       ["aa","b"],
["a","a","b"]
  recursive solution, two pointers; a start pointer that increases recursively, when the end index pointer is fixed(a,ab -> a,a,b);
  an end index pointer increases iteratively, when the start pointer is fixed (a->aa). \mbox{ds}
  \begin{array}{ll} \text{public class Solution (} \\ \text{public List-List-String>> partition(String s) {} \\ \text{intn = slength();} \\ \text{List tol = new ArrayList();} \\ \text{List Soling> head = new ArrayList-String>();} \\ \text{if $(s=-nul)$} & | n==0 \} \\ \text{} \\ \text{} \end{array} 
      getPal(s, tol, head, 0);
return tol;
    private void getPal(String s, List tol, List-String> head, int start) {
    int n = slength();
    if (n = start) {
        tol.ada(new ArrayList-String>(head));
        } else {
       for (int i = start; i < n; i++) {
String sub = s.substring(start, i+1);
String reverseSub = new StringBuffer(sub).reverse().toString();
        if (sub.equals(reverseSub) && !sub.equals("")) {
    head.add(sub);
    getPal(s, tol, head, i + 1);
    head.remove(head.size() - 1);
}
 Palindrome Partitioning II
 Given a string s, partition s such that every substring of the partition is a palindrome.
 Return the minimum cuts needed for a palindrome partitioning of s.
 For example, given s = "aab".
 Return 1 since the palindrome partitioning ["aa", "b"] could be produced using 1 cut.
  Submission Result: Time Limit Exceeded
```

#### Letter Combinations of a Phone Number

Given a digit string, return all possible letter combinations that the number could represent.

A mapping of digit to letters (just like on the telephone buttons) is given below.

```
Input:Digit string "23"
Output: ["ad", "ae", "af", "bd", "be", "bf", "cd", "ce", "cf"].
```

Although the above answer is in lexicographical order, your answer could be in any order you want.

DFS: add(ad), add(ae), add(af), add(bd).....

```
public class Solution {
    public List<String> letterCombinations(String digits) {
                           String| map = new String[10]:
map[0] = ";
map[0] = ";
map[2] = "abc";
map[2] = "abc";
map[3] = "def;
map[4] = "gh";
map[6] = "mon";
map[6] = "mon";
map[6] = "mon";
map[6] = "wwyz";
                              int n = digits.length();
List<String> comb = new ArrayList<String>();
char[] trace = new char[n];
                              if (digits == null || n == 0) {
  comb.add(new String(trace));
  return comb;
                                 LC(digits, 0, 0, comb, trace, map); return comb;
         private void LC(String digits, int L int d, List<String> comb, char[] trace. String[] map] {
    int n = digits.length();
    if (m)
    condad(new String(trace));
    else {
        int ind = integer, parseInt(digits, substring(I, I + 1));
        if (ind > I);
        if (map | Ind);
        ind private | (some substring(I, I + 1));
        ind (ind | map | Ind);
        ind (ind | map | Ind);

                                                   } else {
LC(digits, I + 1, d, comb, trace, map);
```

# String to Integer (atoi)

Implement atoi to convert a string to an integer.

Hint: Carefully consider all possible input cases. If you want a challenge, please do not see below and ask yourself what are the possible input cases.

Notes: It is intended for this problem to be specified vaguely (ie, no given input specs). You are responsible to gather all the input requirements up front.

#### Requirements for atoi:

The function first discards as many whitespace characters as necessary until the first non-whitespace character is found. Then, starting from this character, takes an optional initial plus or minus sign followed by as many numerical digits as possible, and interprets them as a numerical value

The string can contain additional characters after those that form the integral number, which are ignored and have no effect on the behavior of this function.

If the first sequence of non-whitespace characters in str is not a valid integral number, or if no such sequence exists because either str is empty or it contains only whitespace characters, no conversion is

If no valid conversion could be performed, a zero value is returned. If the correct value is out of the range of representable values, INT\_MAX (2147483647) or INT\_MIN (-2147483648) is returned.

1. get rid of white space using trim():

2. '+' '-' look at first char of the string, if it is '-' then return neg num; '+-' if the second char is not a digit, then return 0.

3. not valid digit char, ('0'  $\sim$  '9'), then not add. num = num\*10 + charToInteger;

a. int add = str.charAt(i) - '0';

b. int add = Character.getNumericValue(str.charAt(i));

c. int add =Character.digit(str.charAt(i), 10);

4. overflow: inter > (Integer.MAX\_VALUE - add)/10

Integer.MAX VALUE = 2147483647

```
Integer.MINVALUE = -2147483648
public class Solution {
  public int atoi(String str) {
      int inter = 0;
if (str == null || str.length() == 0) {
    return 0;
     for (int i = k; i < n; i++) {
    char st = str.charAt(i);
    if (st < '0' || st > '9') break;
         int stToInt = (st - '0');
if (inter > (Integer.MAX_VALUE - stToInt)/10 ) {
  overflow = true;
  break;
          inter = inter * 10 + stToInt;
```

```
} if (overflow) {
   if (sign == 1) {
      return Integer.MAX_VALUE;
   } else {
      return Integer.MIN_VALUE;
   }
 Valid Parentheses
 Given a string containing just the characters '(', ')', '{', '}', '[' and ']', determine if the input string is valid.
 The brackets must close in the correct order, "()" and "()[]{}" are all valid but "(]" and "([)]" are not.
  stack push in '(' and pop out '(' and check if that match with the current char ')', etc.
 public class Solution {
   public boolean is Valid(String s) {
      if (s == null || s.length() <= 1) {
        return false;
   }
       int n = slength();

Stack<Character> comp = new Stack<Character>();

for (int i = 0; t = n; i++) {

    char st = s.char.kt();

    if (st == "(1 | st == "(") {

        comp.push(st);

    }
          }
if (!comp.isEmpty()) {
    if (st == '|') {
        char t = comp.pop();
        if (t!= '|') {
            return false;
        }
}
             }
} else return false;
        }
if (!comp.isEmpty()) return false;
,...omp.isE
return true;
}
```

#### **Generate Parentheses**

Given n pairs of parentheses, write a function to generate all combinations of well-formed parentheses.

For example, given n = 3, a solution set is:

```
"(((()))", "(()())", "(()())", "()(())", "()(())"
DFS, recursion, hash set to store unique string
add left parentheses first, left and right reaches n then add to hash set, if left reaches n then add right parentheses.
public class Solution {
  public List<String> generateParenthesis(int n) {
    List<String> list = new ArrayList<String>();
    if (n == 0);
        return list;
    }
}
     }
HashSet<String> hash = new HashSet<String>();
dfs(n, 0, 0, hash, "', list);
     return list;
  }
private void dfs(int n, int left, int right, HashSet<String> hash, String s, List<String> list) {
if (left < right) return;
     if (left == n && right == n) {
    if (lhash.contains(s))
    { hash.add(s);
        list.add(s);
    }
    dfs(n, left + 1, right, hash, s + "(", list);
dfs(n, left, right + 1, hash, s + ")", list);
Add Binary
Given two binary strings, return their sum (also a binary string).
For example,
a = "11"
b = "1"
Return "100".
 11
+ 1
 100
```

math: 个位数是 (a+b+carry) %2 and carry=(a+b+carry)/2;last step, if carry is 1, add one additional 1 to front

0+1=1; 1+0=1;0+0=0;1+1=10 (last is 0 and add 1(carry) to front)

```
public class Solution {
   public String addBinary(String a, String b) {
              \begin{array}{l} \text{if } (a == \text{null \&\& b} == \text{null) return } \text{'''}; \\ \text{if } (a == \text{null \&\& b.length(l)} > 0) \text{ return b}; \\ \text{if } (b == \text{null \&\& a.length(l)} > 0) \text{ return a}; \\ \end{array}
               int carry = 0;
StringBuffer str = new StringBuffer();
              \label{eq:continuity} \begin{cases} \text{for (int i = a.length() - 1, i = b.grgh() - 1; i >= 0 \text{ | } i >= 0; i -, j -) \{ \\ \text{int } 02 = i >= 0? \text{ a.charAt(i) - 0°; 0;} \\ \text{int } 02 = i >= 0? \text{ b.charAt(i) - 0°; 0;} \\ \text{st. inserf(). } (\text{cot})[(02 + b2 + \text{carry}) \%; 2 + \text{V}) \}; \\ \text{carry = } (a2 + b2 + \text{carry})/2; \\ \end{cases} 
             if (carry == 1) {
    str.insert(0, '1');
               return str.toString();
   Restore IP Addresses
  Given a string containing only digits, restore it by returning all possible valid IP address combinations.
 Given "25525511135",
  return ["255.255.11.135", "255.255.111.35"]. (Order does not matter)
  brute-force algorithm,
   valid IP address: each substring length < 4; no '00', i,e, if length>1 then the first char cannot be 0; each value<255
  public class Solution {
   public List-String> restorelp.Addresses(String s) {
      List-String> list = new ArrayList-String>();
      // s = s.trin();
      int n = s.ength();
      if (s == null || (n > 12)){
            return list;
      }
            \begin{array}{l} \text{feture so.} \\ \text{Int } i = 0, \\ \text{for } [nt]_i = i + 1; j < n - 2; j + j + 1 \\ \text{for } [nt]_i = i + 1; j < n - 1; k + j + j \\ \text{for } [nt]_i = k + 1; j < n; 2 + j + j \\ \text{Strip}_i = 0 = s.substring[i, j], b = s.substring[j, k], c = s.substring[k, z], d = s.substring[z]; \\ \text{If } [siB^2(c, b, c, c, d)] \\ \text{if } st.add(a + "." + b + "." + c + "." + d]; \\ \end{array} 
       | Private boolean BIP(String a, String b, String c), String d) {
| If (a.length() > 4 | b.length() > 4 | c.length() > 4 | d.length() > 4 |
| return folds:
| If (a.length() > 1 && c.charAt(0) == '0') | | (c.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') | | (b.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
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| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.length() > 1 && c.charAt(0) == '0') |
| If (a.len
               if (a0 > 255 || b0 > 255 || c0 > 255 || d0 > 255)
return false;
return true;
   Valid Palindrome
  Given a string, determine if it is a palindrome, considering only alphanumeric characters and ignoring cases
   "A man, a plan, a canal: Panama" is a palindrome.
  "race a car" is not a palindrome.
  Have you consider that the string might be empty? This is a good question to ask during an interview.
   For the purpose of this problem, we define empty string as valid palindrome.
   test cases: ab@a, a, .a, "
    \begin{array}{ll} \text{public class Solution \{} \\ \text{public boolean isPalindrome(String s) \{} \\ \text{s = s.replaceAli("\\s'.")");} \\ \text{s = s.folowerCase();} \\ \text{if $\{s == null \mid | s.length() <= 1\}$ return true;} \\ \end{array} 
             \begin{array}{l} r = s.charAt(j); \\ if \left(r < '0' \mid \mid r > 'z' \mid \mid (r < 'a' && r > '9')\right) \{ \end{array}
```

# Remove Duplicates from Sorted Array

return true;

Given a sorted array, remove the duplicates in place such that each element appear only once and return the new length.

Do not allocate extra space for another array, you must do this in place with constant memory.

```
For example,
```

```
Given input array A = [1,1,2],
```

Your function should return length = 2, and A is now [1,2].

solution: remove duplicates in the sorted array and return the length;

test case: null, 1, [1, 2], [1,1,2]. when duplicates occurs, take advantage of the sorted array, increase pointer i(i=0) until no duplicates, and assign A[i](j=0) as A[i], increase j, increase i.

```
public class Solution {
    public intermoveDuplicates(int[] A) {
        int = A.length;
        if (A = null || n = 0)
            return 0;
        int = 0;
        int | -0;
        int |
```

# Remove Duplicates from Sorted Array II

Follow up for "Remove Duplicates":

What if duplicates are allowed at most twice?

For example,

Given sorted array A = [1,1,1,2,2,3],

Your function should return length = 5, and A is now [1,1,2,2,3].

solution: take the first two and add to the 'new' array, skip the redundant ones.

```
 \begin{aligned} & \text{public class Solution } \{ & \text{public int remove Duplicates [int]] A) } \{ & \text{if } \{A = \text{null} \ \| A \text{longth} = = 0 \} \text{ return 0}; \\ & \text{int } n = A \text{length}; \\ & \text{int } n = A \text{length}; \\ & \text{for } [\text{int}] = 0, \\ & \text{for } [\text{int}] = 0, \\ & \text{for } [\text{int}] = A \text{[i]} = A \text{[i]} + A \text{[i]}; \\ & \text{while } [\text{i+2} < \text{n \&\& A[i]} == A \text{[i+1] \&\& A[i]} == A \text{[i+2]} \} \\ & \text{i++}; \\ & \text{} \} \\ & \text{} \\ & \text{return}; \\ & \text{} \} \end{aligned}
```