Reverse Words in a String

```
Given an input string, reverse the string word by word.
```

```
For example,

Given s = "the sky is blue",

return "blue is sky the".

public class Solution {
    public String reverseWords(String s) {
        StringBuffer revStr = new StringBuffer();

        if (s == null || slength() == 0)
        {
            return ";
        }
        StringBl str = s.split(");
        intn = strlength;
        for (int1 = n - 1; 1>= 0; -)
        {
            if (istri) | equals("))
            revStr.append(strii)|.append(");
        }
        return revStr.length() == 0 ? ": revStr.substring(0, revStr.length() - 1);
}
```

Evaluate Reverse Polish Notation

Evaluate the value of an arithmetic expression in Reverse Polish Notation.

Valid operators are +, -, *, /. Each operand may be an integer or another expression.

Some examples:

Sort List

Sort a linked list in $O(n \log n)$ time using constant space complexity.

```
cur.next = hi;
hi = hi.next;
        }
cur = cur.next;
     if (lo != null)
        cur.next = lo;
     }
else if (hi != null)
{
      }
return head.next;
 Reorder List
Given a singly linked list L: L0 \rightarrow L1 \rightarrow ... \rightarrow Ln-1 \rightarrow Ln,
reorder it to: L0 \rightarrow Ln \rightarrow L1 \rightarrow Ln-1 \rightarrow L2 \rightarrow Ln-2 \rightarrow ...
 You must do this in-place without altering the nodes' values.
For example,
Given \{1,2,3,4\}, reorder it to \{1,4,2,3\}.

    break the linked list in the middle into two list;
    reverse the second linked list
    merge the two list by inserting the second one after the first one.
fast = fast.next.next;
slow = slow.next;
      }
ListNode second = slow.next;
slow.next = null;
      second = reverseOrder(second);
ListNode curr = head;
ListNode curr2 = second;
while (curr2!= null)
        ListNode temp1 = curr.next;
ListNode temp2 = curr2.next;
        curr.next = curr2;
curr = temp1;
curr2.next = curr;
        curr2 = temp2;
    public ListNode reverseOrder(ListNode list)
      if (list == null || list.next == null)
     return list;
     ListNode currNode, nextNode;
currNode = list;
nextNode = list.next;
list.next = null;
while(nextNode!= null)
        ListNode loopNode = nextNode.next;
nextNode.next = currNode;
currNode = nextNode;
nextNode = loopNode;
      list = currNode;
  return list;
 Linked List Cycle II
 Given a linked list, return the node where the cycle begins. If there is no cycle, return null.
 Follow up:
Can you solve it without using extra space?
 循环中奇数和偶数个结点;
l.使用两个指针slow,fast。两个指针都从表头开始走,slow每次走一步,fast每次走两步,如果fast遇到null,则说明没有环,返回false;如果slow==fast,说明有环
2.第一次相遇后,让slow,fast继续走.slow从head开始走,slow,fast各走一步,再次相遇的就是循环起始结点
if (fast == null || fast.next == null) {
          return null;
```

```
if (fast == slow)
         slow = slow.next;
fast = fast.next;
 Linked List Cycle
 Given a linked list, determine if it has a cycle in it.
 Can you solve it without using extra space?
return false;
      ListNode slow = head;
ListNode fast = head;
ListNode fast = head;
while (true)
{
if (fast == null || fast.next == null)
{
return false;
           return false;
         }
slow = slow.next;
fast = fast.next.next;
if (slow == fast)
{
            return true;
 Binary Tree Preorder Traversal
 Given a binary tree, return the preorder traversal of its nodes' values.
 Given binary tree {1,#,2,3},
 return [1,2,3].
 Recursive solution
/**
    * Definition for binary tree
    * public class TreeNode {
    * int vol;
    * TreeNode left;
    * TreeNode left;
    * TreeNode (ight;
    * TreeNode(int x) { val = x; }
    *}
 */
public class Solution {
   public List<nteger> preorderTraversal(TreeNode root) {
      List<nteger> list = new ArrayList<Integer> (|;
      if (root == null)
         return Collections.EMPTY_LIST;
      }
list.add{root.val);
list.addAll(preorderTraversal(root.left));
list.addAll(preorderTraversal(root.right));
...adAll(
return list;
}
 Binary Tree Postorder Traversal
 Given a binary tree, return the postorder traversal of its nodes' values.
 Given binary tree {1,#,2,3},
 return [3,2,1].
 Note: Recursive solution is trivial, could you do it iteratively?
```

/**
 * Definition for binary tree
 * public class TreeNode {
 int val;

```
* TreeNode left;
* TreeNode injth:
* TreeNode injth:
* TreeNode (injt) { val = x; }
*/
public class Solution {
    public List-integer> postorderTraversol([TreeNode root) {
        List-integer> list = new ArrayList-integer>();
        if (root == null)
        {
            return Collections.EMPTY_LIST;
            ist.addAll(postorderTraversol(root.left));
            list.addAll(postorderTraversol(root.right));
            list.add(postorderTraversol(root.right));
            return list;
}
```

Word Break

Given a string s and a dictionary of words dict, determine if s can be segmented into a space-separated sequence of one or more dictionary words.

Word Break II

Given a string s and a dictionary of words dict, add spaces in s to construct a sentence where each word is a valid dictionary word.

Return all such possible sentences.

```
For example, given
s = "catsanddog".
dict = ["cat", "cats", "and", "sand", "dog"].
A solution is ["cats and dog", "cat sand dog"].
Submission Result: Time Limit Exceeded
```

```
public class Solution {
    public last-String> wordBreak(String s, Set-String> dict) {
        List-String> wordBreak(String s, Set-String> dict) {
        List-String> list = new ArrayList-String>();
        if (s == null || dict.size() == 0)
        {
            return list,
        }
        helpBuilder(s, dict, 0, **, list);
        return list;
    }
    private void helpBuilder(String s, Set-String> dict, int start, String item, List-String> list){
        if (start > s.length()){
            list.ad(ptem);
            return
        ist.and(ptem);
        return
        ist.and(ptem);
        return
        if (dict.contains(currStr.toString()));
        if (dict.contains(currStr.toString()));
        String newtlem = item.neght() > 0 ? (item + ** + currStr.toString()) : currStr.toString();
        helpBuilder(s, dict, i + 1, newtlem, list);
    }
}
```

Single Number

Given an array of integers, every element appears twice except for one. Find that single one.

Note:

Your algorithm should have a linear runtime complexity. Could you implement it without using extra memory?

```
int cnt = 0;
int x = cnt+;
cnt = 0;
int y = +cnt;
System.out.println(x + " " + y); //0 1//
int cnt = 0;
```

```
public class Solution {
  public int singleNumber(int[] A) {
    if (A.length <= 1 ) {
      return A[0];
    }</pre>
      int len = A.length;
HashMap<Integer, Integer> map = new HashMap<Integer, Integer>();
     for (int i = 0; i < len; i++) {
    if (!map.containskey(A[i])) {
        int count = 0;
        map.put(A[i], ++count);
    }
        } else { map.put(A[i], map.get(A[i]) + 1);
     int results = 0;
for (Integer j : map.keySet()) {
    if (map.get(j) == 1) {
        results = j;
    }
return results;
 Candy
 There are N children standing in a line. Each child is assigned a rating value.
 You are giving candies to these children subjected to the following requirements:

    Each child must have at least one candy.

    Children with a higher rating get more candies than their neighbors.

What is the minimum candies you must give?
通常是要求的变量跟左右元素有关系的题目: 两边扫描的方法
一些例子:
12333 =》 8 因为candy数可以是
12311
12323=》9 因为candy数可以是
12312
 思路:
 1
        d[i] 是给第i个小孩最少几块糖 rank[i] > rank[i-1],必须比前一个多给一块,d[i] = d[i-1] + 1 rank[i] <= rank[i-1],两个排名一样,第二个就给一块就行了,d[i] = 1
 基本思路就是进行两次扫描,一次从左往右,一次从右往左。第一次扫描的时候维护对于每一个小孩左边所需要是少的糖果数量。存入数组对应元素中,第二次扫描的时候维护右边所需的最少糖果数,并且比较将左边和右边大的糖果数量
存入结果数组对应元素中。这样两遍扫描之后就可以得到每一个所需要的最最少糖果量,从而累加得出结果。方法只需要两次扫描,所以时间复杂度是O(2*n)=O(n)。空间上需要一个长度为n的数组,复杂度是O(n)。
public class Solution {
   public int candy(int[] ratings) {
    if (ratings.length == 0) {
        return 0;
    }
      int[] num = new int[ratings.length];
num[0] = 1;
     for (int i = 1; i < ratings.length; i++) {
    if (ratings[i] > ratings[i - 1]) {
        num[i] = num[i - 1] + 1;
    }
       }
else if (ratings[i] < ratings[i - 1]) {
num[i] = 1;
     int res = num[ratings.length - 1];
for (int i = ratings.length - 2; i >= 0; i-) {
    int curr = 1;
    if (ratings[i] > ratings[i + 1]) {
        curr = num[i + 1] + 1;
        }
res += Math.max(num[i], curr);
num[i] = curr;
```

Gas Station

return res; }

There are N gas stations along a circular route, where the amount of gas at station i is gas[i].

You have a car with an unlimited gas tank and it costs cost[i] of gas to travel from station i to its next station (i+1). You begin the journey with an empty tank at one of the gas stations.

Return the starting gas station's index if you can travel around the circuit once, otherwise return -1.

Note:

The solution is guaranteed to be unique.

```
 \begin{aligned} & \text{public class Solution} \, \{ \\ & \text{public int canCompleteCircuit(int] gas, int[] cost)} \, \{ \\ & \text{if (gas.length} = 0 \, || \, \text{cost.length} \, |= 0 \, || \, \text{gas.length} \, |= \, \text{cost.length}) \, \{ \\ & \text{return} - 1; \\ & \} \\ & \text{int sum} = 0; \\ & \text{int total} = 0; \\ & \text{int total} = 0; \\ & \text{int start} = 0; \\ & \text{for (int i} = 0; i < \, \text{gas.length}; i++) \, \{ \\ & \text{sum} + 2 \, \text{gas[i]} - \, \text{cost[i]}; \\ & \text{total} + 2 \, \text{sum}; \\ & \text{if (sum} < 0) \, \{ \end{aligned}
```

```
start = i + 1;
sum = 0;
}
      if (total < 0) {
return -1;
return start;
 Palindrome Partitioning
 Given a string s, partition s such that every substring of the partition is a palindrome.
 Return all possible palindrome partitioning of s.
 For example, given s = "aab",
 Return
  Ε
       ["aa","b"],
["a","a","b"]
  recursive solution, two pointers; a start pointer that increases recursively, when the end index pointer is fixed(a,ab -> a,a,b);
  an end index pointer increases iteratively, when the start pointer is fixed (a->aa). \mbox{ds}
  \begin{array}{ll} \text{public class Solution (} \\ \text{public List-List-String>> partition(String s) {} \\ \text{intn = slength();} \\ \text{List tol = new ArrayList();} \\ \text{List Soling> head = new ArrayList-String>();} \\ \text{if $(s=-nul)$} & | n==0 \} \\ \text{} \\ \text{} \end{array} 
      getPal(s, tol, head, 0);
return tol;
    private void getPal(String s, List tol, List-String> head, int start) {
    int n = slength();
    if (n = start) {
        tol.ada(new ArrayList-String>(head));
        } else {
       for (int i = start; i < n; i++) {
String sub = s.substring(start, i+1);
String reverseSub = new StringBuffer(sub).reverse().toString();
        if (sub.equals(reverseSub) && !sub.equals("")) {
    head.add(sub);
    getPal(s, tol, head, i + 1);
    head.remove(head.size() - 1);
}
 Palindrome Partitioning II
 Given a string s, partition s such that every substring of the partition is a palindrome.
 Return the minimum cuts needed for a palindrome partitioning of s.
 For example, given s = "aab".
 Return 1 since the palindrome partitioning ["aa", "b"] could be produced using 1 cut.
  Submission Result: Time Limit Exceeded
```

Letter Combinations of a Phone Number

Given a digit string, return all possible letter combinations that the number could represent.

A mapping of digit to letters (just like on the telephone buttons) is given below.

```
Input:Digit string "23"
Output: ["ad", "ae", "af", "bd", "be", "bf", "cd", "ce", "cf"].
```

Note:

Although the above answer is in lexicographical order, your answer could be in any order you want.

DFS: add(ad), add(ae), add(af), add(bd).....

```
public class Solution {
    public List-String= letter/Combinations(String digits) {
        String[I map = new String[10];
        map(0] = ",
        map(1) = ",
        map(2) = ",
        map(3) = "def;
        map(4) = ",
        map(5) = ",
        map(6) = ",
        map(7) = ",
        map(8) = ",
```