

Sec 2.2 Operations on Sets

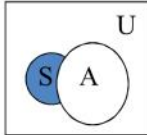
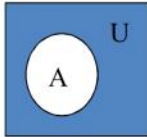
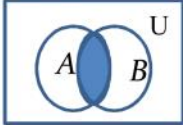
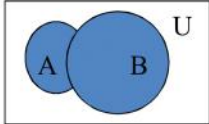
June 12, 2016 10:17 AM

Section 2.2 Operations on Sets

Comp 232
Robert Mearns

1. Define the operations on a set:

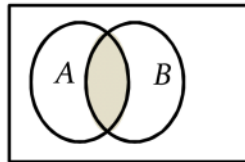
Final result is shaded area below

| Term | Symbol | Set Builder expression | Venn Diagram |
|---------------------------------|------------------------------------|--|---|
| Compliment of A compare to S | $S - A$ | $\{x \mid (x \in S) \wedge (x \notin A)\}$ |  |
| Compliment of A compare to U | $U - A$ denoted as \bar{A} | $\{x \mid (x \in U) \wedge (x \notin A)\}$ $\equiv \{x \mid x \in \bar{A}\}$ $\equiv \{x \mid \neg(x \in A)\}$ |  |
| Intersection | $A \cap B$ | $\{x \mid (x \in A) \wedge (x \in B)\}$ |  |
| Union | $A \cup B$ | $\{x \mid (x \in A) \vee (x \in B)\}$ |  |

Ex 1: let $U = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, $A = \{4, 5, 6, 7\}$, $B = \{0, 5, 6\}$, $C = \{8, 9\}$

$$\begin{aligned} A - B &= \{4, 7\} \\ \bar{A} &= \{0, 1, 2, 3, 8, 9\} \\ A \cap B &= \{5, 6\} \\ B \cap C &= \{\} = \emptyset \\ B \cup C &= \{0, 5, 6, 8, 9\} \end{aligned}$$

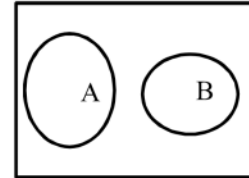
Ex 2: The Cardinality of Two Sets



$$|A \cup B| \neq |A| + |B| \quad \text{Why?}$$

We are counting elements in the Intersection twice.

$$\begin{aligned} |A \cup B| &= |A| + |B| - |A \cap B| \\ \text{or } n(A \cup B) &= n(A) + n(B) - n(A \cap B) \end{aligned}$$



$$|A \cap B| = 0$$

A, B are called Disjoint sets

Ex 3: Let A be the set of students in the class who have taken Calculus, B be the set of students in class who have taken Matrices. $n(A) = 15$, $n(B) = 13$, $n(A \cap B) = 5$

$$n(A \cup B) = 15 + 13 - 5 = 23$$

$n(A \cup B)$ is the count of the number of student who have taken Calculus or Matrices or both.

2. Equivalence between the three systems: Logic symbols, Boolean Algebra symbols, Set symbols

a) Recall from the Table of Set operations that:

Compliment involves Not: $\bar{A} = \{x \mid \neg(x \in A)\} = \{x \mid (x \notin A)\}$

Intersection involves And: $A \cap B = \{x \mid (x \in A) \wedge (x \in B)\}$

Union involves OR: $A \cup B = \{x \mid (x \in A) \vee (x \in B)\}$

b) This leads to the following equivalences:

| Word | Set notation | Logic notation | Boolean Algebra Notation |
|------|--------------|----------------|--------------------------|
| Not | \bar{A} | $\neg A$ | \bar{A} |
| And | $A \cap B$ | $A \wedge B$ | $A * B = AB$ |
| Or | $A \cup B$ | $A \vee B$ | $A + B$ |

c) This leads to the equivalence of the logic properties previously developed. The properties table is summarized on the next slide.

It needs to be completed (see p 27 & 130 of Rosen Text 7th ed.)

d) Summary of Equivalent Notations:

| Logic symbols | Set symbols | Boolean Algebra | |
|--|---|---|------------------|
| $\neg (A \wedge B) \equiv \neg A \vee \neg B$ $\neg (A \vee B) \equiv \neg A \wedge \neg B$ | $\overline{(A \cap B)} =$ $\overline{(A \cup B)} =$ | $\overline{p} \overline{q} =$ $\overline{p + q} =$ | DeMorgan's Rules |
| $A \wedge B \equiv$ $A \vee B \equiv$ | $A \cap B =$ $A \cup B =$ | $p q = q p$ $p + q = q + p$ | Commutative |
| $C \wedge (A \wedge B) \equiv$ $C \vee (A \vee B) \equiv$ | $C \cap (A \cap B) =$ $C \cup (A \cup B) =$ | $c (p q) = (c p) q$ $c + (p + q) = (c + p) + q$ | Associative |
| $A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C)$ $A \vee (B \wedge C) \equiv (A \vee B) \wedge (A \vee C)$ | $A \cap (B \cup C) =$ $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$ | $p (q + r) =$ $p + (q r) = (p + q)(p + r)$ | Distributive |
| $A \wedge A \equiv$ $A \vee A \equiv$ | $A \cap A =$ $A \cup A =$ | $p p = p$ $p + p = p$ | Idempotent |
| $A \wedge T \equiv$ $A \vee F \equiv$ | $A \cap U =$ $A \cup \Phi =$ | $p * 1 = p$ $p + 0 = p$ | Identities |
| $A \wedge F \equiv$ $A \vee T \equiv T$ | $A \cap \Phi =$ $A \cup U = U$ | $p * 0 =$ $p + 1 = 1$ | Domination |
| $A \wedge \neg A \equiv F$ $A \vee \neg A \equiv T$ | $A \cap \bar{A} = \emptyset$ $A \cup \bar{A} = U$ | $p \bar{p} = 0$ $p + \bar{p} = 1$ | Negation |

3. Verification of set expressions

a) Notation: Uppercase letters and $\{ \}$ brackets are read as “set”:

Ex: $\{ x \mid x \in (A \cap B) \}$ is “set of x such that x is contained in the intersection of sets A, B ”

b) Summary of equivalent statements:

| | |
|--|--------------------------------------|
| $\{x \mid x \in \bar{A}\} \equiv \{x \mid x \notin A\}$ | Def. of Compliment of a set (Not) |
| $\{x \mid x \notin A\} \equiv \{x \mid \neg(x \in A)\}$ | Def. of \notin |
| $\{x \mid x \in (A \cap B)\} \equiv \{x \mid x \in A \wedge x \in B\}$ | Def. of Intersection of sets (And) |
| $\{x \mid x \in (A \cup B)\} \equiv \{x \mid x \in A \vee x \in B\}$ | Def. of Union of sets (Or) |

c) Method 1 Use: LHS set = RHS set iff [LHS set \subseteq RHS set and RHS set \subseteq LHS]

Ex: Show that: $A - B = A \cap \bar{B}$

Step 1 Consider $x \in A - B$

$$\rightarrow x \in A \wedge x \notin B$$

$$\rightarrow x \in A \wedge x \in \bar{B}$$

$$\rightarrow x \in A \cap \bar{B}$$

$$\rightarrow A - B \subseteq A \cap \bar{B}$$

Step 2 Consider $x \in A \cap \bar{B}$

In a similar way it can be shown $x \in A - B$

$$A \cap \bar{B} \subseteq A - B$$

$$[A \cap \bar{B} \subseteq A - B \text{ and } A - B \subseteq A \cap \bar{B}] \rightarrow A - B = A \cap \bar{B}$$

Def A-B

Def compliment

Def \cap

Def subset \subseteq

Def of equal sets

d) Method 2 Use Set-Builder notation without assuming set identities

Ex: Show that: $\overline{\overline{A} \cup B} = A \cap \overline{B}$

$$\overline{\overline{A} \cup B} = \{x \mid x \notin (\overline{A} \cup B)\}$$

$$= \{x \mid \neg[x \in (\overline{A} \cup B)]\}$$

Def. of \notin

$$= \{x \mid \neg[x \in \overline{A}] \vee \neg(x \in B)\}$$

Def. of Union

$$= \{x \mid \neg[(x \notin A) \vee (x \in B)]\}$$

Def of Compliment

$$= \{x \mid \neg[\neg(x \in A) \vee (x \in B)]\}$$

Def of \neg

$$= \{x \mid [(x \in A) \wedge \neg(x \in B)]\}$$

De Morgan

$$= \{x \mid (x \in A) \wedge \neg(x \in B)\}$$

Def. of \notin

$$= x \in A \wedge x \in \overline{B}$$

Def of Compliment

$$= A \cap \overline{B}$$

Def. of Intersection

QED

4. Applications: Representing sets with Boolean strings

Ex 1: Consider data kept by a provincial government concerning 5 cities and their total property evaluation in millions of dollars, number of properties and the tax rate per 100\$ evaluation.

Universal set is $U = \{A, B, C, D, E\}$.

Sets are represented by bit strings: (i) Bit string 11111 rep. U (ii) Bit string 11001 rep.

| | String | City A | City B | City C | City D | City E |
|---|--------|--------|--------|--------|--------|--------|
| $1\text{m\$} < \text{eval} \leq 5\text{m\$}$ | B_1 | 1 | 0 | 0 | 0 | 0 |
| $5\text{m\$} < \text{eval} \leq 15\text{m\$}$ | B_2 | 0 | 0 | 1 | 0 | 0 |
| $\text{eval} > 15\text{m\$}$ | B_3 | 0 | 1 | 1 | 0 | 1 |
| $0 < \# \text{ properties} \leq 300$ | B_4 | 0 | 0 | 1 | 1 | 0 |
| $300 < \# \text{ prop} \leq 1000$ | B_5 | 1 | 1 | 0 | 0 | 1 |
| $1\$ < \text{tax rate} \leq 1.30\$$ | B_6 | 0 | 1 | 1 | 0 | 0 |
| $1.30\$ < \text{tax rate}$ | B_7 | 1 | 0 | 0 | 1 | 1 |

- a) Form set S_a : designates cities with $(1\text{m\$} < \text{eval} \leq 5\text{m\$})$ and $(300 < \# \text{ prop} \leq 1000)$ and $(\text{tax rate} > 1.30\$)$

We need bit string $B_a = \text{bit-wise And with } B_1, B_5, B_7 = 10000 \rightarrow s_a = \{A\}$

| | String | City A | City B | City C | City D | City E |
|---|--------|--------|--------|--------|--------|--------|
| $1\text{m\$} < \text{eval} \leq 5\text{m\$}$ | B_1 | 1 | 0 | 0 | 0 | 0 |
| $5\text{m\$} < \text{eval} \leq 15\text{m\$}$ | B_2 | 0 | 0 | 1 | 0 | 0 |
| $\text{eval} > 15\text{m\$}$ | B_3 | 0 | 1 | 1 | 0 | 1 |
| $0 < \# \text{ properties} \leq 300$ | B_4 | 0 | 0 | 1 | 1 | 0 |
| $300 < \# \text{ prop} \leq 1000$ | B_5 | 1 | 1 | 0 | 0 | 1 |
| $1\$ < \text{tax rate} \leq 1.30\$$ | B_6 | 0 | 1 | 1 | 0 | 0 |
| $1.30\$ < \text{tax rate}$ | B_7 | 1 | 0 | 0 | 1 | 1 |

b) Form the set S_b that designates cities with $5\text{m\$} < \text{eval} \leq 15\text{m\$}$ or tax rate $> 1.30\$$

We need bit string $B_b =$ bit-wise Or with B_2 or $B_7 = 1011 \rightarrow s_b = \{A, C, D, E\}$

c) Form the set S_c that designates cities with missing total evaluation data

We need bit string $B_c = \text{Not } B_1 \text{ and Not } B_2 \text{ and Not } B_3 = \text{Not}(B_1 \text{ or } B_2 \text{ or } B_3)$
 $\text{Not}(\text{bit-wise Or with } B_1, B_2, B_3) = \overline{11101} = 00010 \rightarrow s_c = \{D\}$

Ex 2: Adapt the above procedure to find the set S_d that designates cities that have more than one total evaluation

$$B_d = (B_1 \wedge B_2) \vee (B_2 \wedge B_3) \vee (B_3 \wedge B_1) = 00100 \rightarrow s_d = \{C\}$$

Ex 3: Consider data collected concerning the population in various postal codes designated by the Boolean strings below:

| First Letter of Postal code | A | B | G | H | I | J | M | T |
|---|---|---|---|---|---|---|---|---|
| U designates the Universal set $U =$ | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| $B_1 =$ B_1 designates postal codes where population ≤ 1 million | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 |
| $B_2 =$ B_2 designates postal codes where 1 million $<$ pop. ≤ 2 million | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| $B_3 =$ B_3 designates postal codes where 2 million $<$ pop. ≤ 3 million | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 1 |
| $B_4 =$ B_4 designates postal codes where the majority of population speak French | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 0 |

Form a Boolean string B_i for each condition below and using the results write the final set S_i :

- Universal set
- Set where (1 million $<$ pop. ≤ 3 million)
- (Majority speak French) and (2 million $<$ pop. ≤ 3 million)
- A set that designates if there is some postal code where pop. has been omitted (data error)
- A set that designates if a postal code exists in more than one pop. size (data error)

Answer:

- $B_a = U, S_a = \{A, B, G, H, I, J, M, T\}$
- $B_b = B_2 \vee B_3 = 1010\ 1111, S_b = \{A, G, I, J, M, T\}$
- $B_c = B_3 \vee B_4 = 0010\ 0000, S_c = \{G\}$
- $B_d = \neg(B_1 \vee B_2 \vee B_3) = 1111\ 1111 = 0000\ 0000, S_d = \emptyset$
- $B_e = (B_1 \wedge B_2) \vee (B_2 \wedge B_3) \vee (B_1 \wedge B_3) = 0000\ 1000, S_e = \{I\}$

17