8.10.1 Galois field

- Field: field is a set of elements and operations of addition and multiplication. The operations must follows rules below
 - Closed: Closure impliles that the sum and product of any two elements in the field are also elements of the field
 - Commutative (ab = ba and a+b = b+a)
 - Associative (a(bc) = (ab)c, and a + (b + c) = (a + b) + c)
 - Distributive law relates multiplication and addition: a(b + c) = ab + ac.
 - Has additive and multiplicative identities (0 and 1) such that a + 0 = a and 1a = a for any element in the field.
 - Elements of a field must have additive and multiplicative inverses. The additive inverse of a is an element b such that a+b = 0 and the multiplicative inverse of a is an element c such that ac = 1.
 - E.g:
 - set of real numbers and addition, multiplication creates field.

8.10.1 Galois field

- Finite field:
 - Denoted by Zp that contains
 - The set of integers {0, 1, .., p-1}
 - Modulo p arithmetic.
 - p is a prime number
- Galois field: GF (p^n) contains
 - p is prime number
 - n is arbitrary positive integer
 - Each element is denoted by polynomial $a1x^0 + a2x^1 + ... + aNx^{N-1}$ where the coefficients at take on values in the set $\{0, 1, ..., p-1\}$.
 - To add two polynomials, for each power of x present in the summands, just add the corresponding coefficients modulo p
 - $a(x) = a1x^0 + a2x^1 + ... + aNx^{N-1}$; $b(x) = b1x^0 + b2x^1 + ... + bNx^{N-1}$
 - $c(x) = a(x)+b(x) = (a1 \oplus b1) + (a2 \oplus b2)x + ... + (aN \oplus bN)x^{N-1}$
 - $ai \bigoplus bi = ai+bi \text{ if } ai+bi < p$ = ai+bi-p if ai+bi >= p
 - Multiplication of two polynomials is done by multilication in modulo x^{N-1} where x^{N-1} is modulo polynomial $a(x) \times b(x)$ modulo $(x^{N-1}) = reminder$ of $((a(x) \times b(x)) / x^{N-1})$

8.10.2 Definition

- Cyclic code uses Galois Field $GF(p^N)$
- Codeword a is considered as polynomials
 - E.g. $a = \{a_1, a_2, ..., a_N\}$ is considered as $a(x) = a_1 x^0 + a_2 x^1 + + a_N x^{N-1}$
- Multiplication is calculated in modulo x^N-1
- Multiple with x is equivalence to right shift its coefficients

$$xa(x) = a_N + a_0 x^1 + a_1 x^2 + \dots + a_{N-1} + a_N (x^N - 1)$$

$$xa(x) \ modulo (x^N - 1) = a_N + a_0 x^1 + a_1 x^2 + \dots + a_{N-1}$$

- Cyclic code is a linear code with the property that any cyclic shift of a code word is also a code word
- A cyclic code has a unique non-zero polynomial of minimal degree
 - This polynomial is called generator polynomial with degree r: $g(x) = g_o + g_1x + ... + g_rx^r$
 - g(x) is the generator polynomial of a cyclic code if and only if it is a factor of (x^N-1)
 - The remainder of division between arbitrary codeword and g(x) = 0
 - If c(x) is codeword then c(x) = m(x) g(x)

8.10.2. Definition(Cont.)

Generator matrix:

$$G = \begin{bmatrix} g_0 & g_1 & \dots & g_r & 0 & 0 & \dots & 0 \\ 0 & g_0 & \dots & g_{r-1} & g_r & 0 & \dots & 0 \\ & & \vdots & & & & \\ 0 & \dots & 0 & g_0 & g_1 & \dots & g_r & 0 \\ 0 & \dots & 0 & 0 & g_0 & \dots & g_{r-1} & g_r \end{bmatrix}$$

• G is a cyclic matrix (each row is obtained by shifting the previous row one column to the right).

8.10.2. Definition(Cont.)

• Since g(x) is the factor of (x^N-1) , that

$$(x^N-1)=g(x)\;h(x)$$

Where h(x) is called check parity matrix

- If c(x) is codeword then c(x) h(x) = m(x) g(x) h(x) modulo $(x^N-1) = 0$
- $h(x) = h_0 + h_0 x + ... + h_k x^{k-1}$
- Check parity matrix H: is a cyclic matrix (each row is obtained by shifting the previous row one column to the right).
 - First row is h(x)

8.10.3. Encoding and decoding

- Encoding process is multiple generator polynomial g(x) with carrying information (message) polynomial m(x)
 - c(x) = m(x) g(x)
- Decoding process:
 - Syndrome S is remainder of division between received polynomial r(x) and g(x)
 - $S = r(x) \mod g(x) \mod (x^N-1)$
 - If $S = 0 \rightarrow codeword$
 - If $S \neq 0 \rightarrow S = e(x) \mod g(x) \mod (x^N-1)$
 - Can find error polynomial e(x) from S

8.10.3. Encoding and decoding (Cont.)

- If generator matrix G is transformed into canonical form, codeword is in systematic form
 - $c(x) = m(x) + d(x) x^k$ Where d(x) is a polynomial has degree of n-k-1
- Since c(x) mod g(x) modulo(x^N-1) = 0, then $d(x) = m(x) x^{N-k} \mod g(x) \mod (x^N-1)$

8.10.4. Cyclic Redundancy Check Codes

- Is cyclic systematic code
- Used for send or store the information
- Codeword c(x) = m(x) crc
 - crc = m(x) mod g(x) modulo $(x^N 1)$
- Decoding
 - Let r(x) = m'(x) crc' where $m'(x) = m(x) + e_1(x)$; $crc' = crc + e_2(x)$
 - $e_1(x)$ first L symbol of e(x)
 - $e_2(x)$ remaining N-L symbols of e(x)
 - S= m'(x) mod g(x) modulo $(x^N 1) \text{crc}'$
 - S= 0 \rightarrow no error
 - $S \neq 0 \rightarrow S = e_1(x) \mod g(x) \mod (x^N 1) e_2(x)$
 - Calculate $e_1(x)$, $e_2(x)$ from S

Example

• Let

$$M(x) = a_m \cdot x^m + a_{m-1} \cdot x^{m-1} + a_{m-2} \cdot x^{m-2} + \dots + a_1 \cdot x^1 + a_0$$

$$G(x) = a_n \cdot x^n + a_{n-1} \cdot x^{n-1} + a_{n-2} \cdot x^{n-2} + \dots + a_1 \cdot x^1 + a_0$$

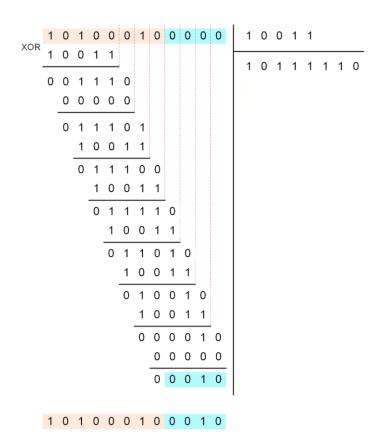
- Codeword: $T(x) = a_n.x^n.M(x)$
 - \rightarrow left shift n bit of M(x)
- CRC string R(X):

$$R(x) = T(x) \% G(x)$$

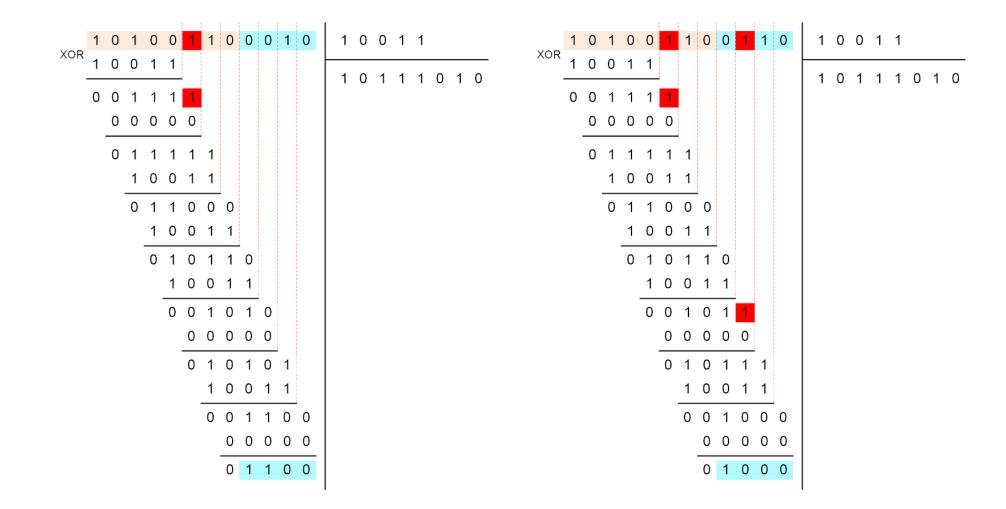
Example

- CRC-4 \rightarrow G(X): x^4 + x + 1 (10011)
- Input: $M(X) = x^7 + x^5 + x (1010_0010)$
- To extent to T(X): M(X) wi Ifted 4 positions: \rightarrow 1010 0010 0000

Calculate codeword



Test: errors



CRC does not work

