

8.10.1 Galois field

- Field: field is a set of elements and operations of addition and multiplication. The operations must follow rules below
 - Closed: Closure implies that the sum and product of any two elements in the field are also elements of the field
 - Commutative ($ab = ba$ and $a+b = b+a$)
 - Associative ($a(bc) = (ab)c$, and $a + (b + c) = (a + b) + c$)
 - Distributive law relates multiplication and addition: $a(b + c) = ab + ac$.
 - Has additive and multiplicative identities (0 and 1) such that $a + 0 = a$ and $1a = a$ for any element in the field.
 - Elements of a field must have additive and multiplicative inverses. The additive inverse of a is an element b such that $a+b = 0$ and the multiplicative inverse of a is an element c such that $ac = 1$.
 - E.g:
 - set of real numbers and addition, multiplication creates field.

8.10.1 Galois field

- Finite field:
 - Denoted by \mathbb{Z}_p that contains
 - The set of integers $\{0, 1, \dots, p-1\}$
 - Modulo p arithmetic.
 - p is a prime number
- Galois field: $\text{GF}(p^n)$ contains
 - p is prime number
 - n is arbitrary positive integer
 - Each element is denoted by polynomial $a_1x^0 + a_2x^1 + \dots + a_Nx^{N-1}$ where the coefficients a_i take on values in the set $\{0, 1, \dots, p-1\}$.
 - To add two polynomials, for each power of x present in the summands, just add the corresponding coefficients modulo p
 - $a(x) = a_1x^0 + a_2x^1 + \dots + a_Nx^{N-1}$; $b(x) = b_1x^0 + b_2x^1 + \dots + b_Nx^{N-1}$
 - $c(x) = a(x) + b(x) = (a_1 \oplus b_1) + (a_2 \oplus b_2)x + \dots + (a_N \oplus b_N)x^{N-1}$
 - $a_i \oplus b_i = a_i + b_i$ if $a_i + b_i < p$
 $= a_i + b_i - p$ if $a_i + b_i \geq p$
 - Multiplication of two polynomials is done by multiplication in modulo x^{N-1} where x^{N-1} is modulo polynomial $a(x) \times b(x) \text{ modulo } (x^{N-1}) = \text{remainder of } ((a(x) \times b(x)) / x^{N-1})$

8.10.2 Definition

- *Cyclic code uses Galois Field $GF(p^N)$*
- *Codeword a is considered as polynomials*
 - *E.g. $a = \{a_1, a_2, \dots, a_N\}$ is considered as $a(x) = a_1x^0 + a_2x^1 + \dots + a_Nx^{N-1}$*
- *Multiplication is calculated in modulo $x^N - 1$*
- *Multiple with x is equivalence to right shift its coefficients*
$$xa(x) = a_N + a_0x^1 + a_1x^2 + \dots + a_{N-1} + a_N(x^N - 1)$$
$$xa(x) \text{ modulo } (x^N - 1) = a_N + a_0x^1 + a_1x^2 + \dots + a_{N-1}$$
- *Cyclic code is a linear code with the property that any cyclic shift of a code word is also a code word*
- *A cyclic code has a unique non-zero polynomial of minimal degree*
 - *This polynomial is called generator polynomial with degree r :*
$$g(x) = g_0 + g_1x + \dots + g_rx^r$$
 - *$g(x)$ is the generator polynomial of a cyclic code if and only if it is a factor of $(x^N - 1)$*
 - *The remainder of division between arbitrary codeword and $g(x) = 0$*
 - *If $c(x)$ is codeword then $c(x) = m(x)g(x)$*

8.10.2. Definition(Cont.)

- Generator matrix:

$$G = \begin{bmatrix} g_0 & g_1 & \dots & g_r & 0 & 0 & \dots & 0 \\ 0 & g_0 & \dots & g_{r-1} & g_r & 0 & \dots & 0 \\ & & & \vdots & & & & \\ 0 & \dots & 0 & g_0 & g_1 & \dots & g_r & 0 \\ 0 & \dots & 0 & 0 & g_0 & \dots & g_{r-1} & g_r \end{bmatrix}$$

- G is a cyclic matrix (each row is obtained by shifting the previous row one column to the right).

8.10.2. Definition(Cont.)

- Since $g(x)$ is the factor of (x^N-1) , that

$$(x^N-1) = g(x) h(x)$$

Where $h(x)$ is called check parity matrix

- If $c(x)$ is codeword then $c(x) h(x) = m(x) g(x) h(x)$ modulo $(x^N-1) = 0$
- $h(x) = h_0 + h_1x + \dots + h_kx^{k-1}$
- Check parity matrix H: is a cyclic matrix (each row is obtained by shifting the previous row one column to the right).
 - First row is $h(x)$

8.10.3.Encoding and decoding

- Encoding process is multiple generator polynomial $g(x)$ with carrying information (message) polynomial $m(x)$
 - $c(x) = m(x) g(x)$
- Decoding process:
 - Syndrome S is remainder of division between received polynomial $r(x)$ and $g(x)$
 - $S = r(x) \bmod g(x) \text{ modulo } (x^N - 1)$
 - *If $S = 0 \rightarrow \text{codeword}$*
 - *If $S \neq 0 \rightarrow S = e(x) \bmod g(x) \text{ modulo } (x^N - 1)$*
 - *Can find error polynomial $e(x)$ from S*

8.10.3.Encoding and decoding (Cont.)

- If generator matrix G is transformed into canonical form, codeword is in systematic form
 - $c(x) = m(x) + d(x) x^k$
Where $d(x)$ is a polynomial has degree of $n-k-1$
- Since $c(x) \bmod g(x) \bmod (x^N - 1) = 0$,
then $d(x) = m(x) x^{N-k} \bmod g(x) \bmod (x^N - 1)$

8.10.4. Cyclic Redundancy Check Codes

- Is cyclic systematic code
- Used for send or store the information
- Codeword $c(x) = m(x) - \text{crc}$
 - $\text{crc} = m(x) \bmod g(x) \bmod (x^N - 1)$
- Decoding
 - Let $r(x) = m'(x) - \text{crc}'$ where $m'(x) = m(x) + e_1(x)$; $\text{crc}' = \text{crc} + e_2(x)$
 - $e_1(x)$ first L symbol of $e(x)$
 - $e_2(x)$ remaining N-L symbols of $e(x)$
 - $S = m'(x) \bmod g(x) \bmod (x^N - 1) - \text{crc}'$
 - $S = 0 \rightarrow$ no error
 - $S \neq 0 \rightarrow S = e_1(x) \bmod g(x) \bmod (x^N - 1) - e_2(x)$
 - Calculate $e_1(x), e_2(x)$ from S

Example

- Let

$$M(x) = a_m \cdot x^m + a_{m-1} \cdot x^{m-1} + a_{m-2} \cdot x^{m-2} + \dots + a_1 \cdot x^1 + a_0$$

$$G(x) = a_n \cdot x^n + a_{n-1} \cdot x^{n-1} + a_{n-2} \cdot x^{n-2} + \dots + a_1 \cdot x^1 + a_0$$

- Codeword: $T(x) = a_n \cdot x^n \cdot M(x)$

→ left shift n bit of M(x)

- CRC string R(X):

$$R(x) = T(x) \% G(x)$$

Example

- CRC-4 $\rightarrow G(X): x^4 + x + 1$ (10011)
- Input: $M(X) = x^7 + x^5 + x$ (1010_0010)
- To extend to $T(X)$: $M(X)$ with 4 zeros shifted 4 positions:
 $\rightarrow 1010_0010_0000$

Calculate codeword

The diagram illustrates the calculation of a codeword using XOR operations. It shows a data word, a generator polynomial, and a series of XOR steps to produce a remainder.

Data Word: 1 0 1 0 0 0 1 0 0 0 0 0

Generator Polynomial: 1 0 0 1 1

XOR Steps:

```
XOR
1 0 1 0 0 0 1 0 0 0 0 0
1 0 0 1 1
-----
0 0 1 1 1 0
0 0 0 0 0
-----
0 1 1 1 0 1
1 0 0 1 1
-----
0 1 1 1 0 0
1 0 0 1 1
-----
0 1 1 1 1 0
1 0 0 1 1
-----
0 1 1 0 1 0
1 0 0 1 1
-----
0 1 0 0 1 0
1 0 0 1 1
-----
0 0 0 0 1 0
0 0 0 0 0
-----
0 0 0 1 0
```

Final Codeword: 1 0 1 0 0 0 1 0 0 0 1 0

Test: errors

XOR

1 0 1 0 0	1	1 0 0 0 1 0	1 0 0 1 1
1 0 0 1 1			1 0 1 1 1 0 1 0
<hr/>			
0 0 1 1 1	1		
0 0 0 0 0			
<hr/>			
0 1 1 1 1 1			
1 0 0 1 1			
<hr/>			
0 1 1 0 0 0			
1 0 0 1 1			
<hr/>			
0 1 0 1 1 0			
1 0 0 1 1			
<hr/>			
0 0 1 0 1 0			
0 0 0 0 0			
<hr/>			
0 1 0 1 0 1			
1 0 0 1 1			
<hr/>			
0 0 1 1 0 0			
0 0 0 0 0			
<hr/>			
0	1 1 0 0		

XOR

1 0 1 0 0	1	1 0 0 1 1 0	1 0 0 1 1
1 0 0 1 1			1 0 1 1 1 0 1 0
<hr/>			
0 0 1 1 1	1		
0 0 0 0 0			
<hr/>			
0 1 1 1 1 1			
1 0 0 1 1			
<hr/>			
0 1 1 0 0 0			
1 0 0 1 1			
<hr/>			
0 1 0 1 1 0			
1 0 0 1 1			
<hr/>			
0 0 1 0 1	1		
0 0 0 0 0			
<hr/>			
0 1 0 1 1 1			
1 0 0 1 1			
<hr/>			
0 0 1 0 0 0			
0 0 0 0 0			
<hr/>			
0	1 0 0 0		

CRC does not work

