

1.1. Conceptual data model Naturual expressions without constraints imposed by DBMS E-R model Expressed by E-R diagram Data model Logical model Independent of DBMS DBMS dependent

1.2. Logical Data Model

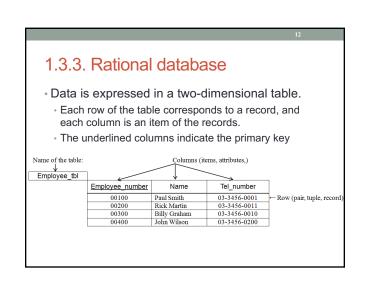
- 3 types
- · relational model,
- network model,
- · and hierarchical model

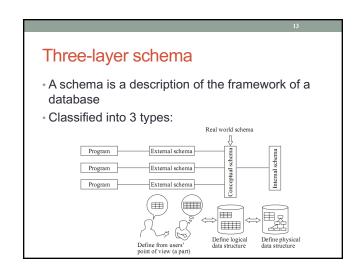
1.3. Physical Data Model

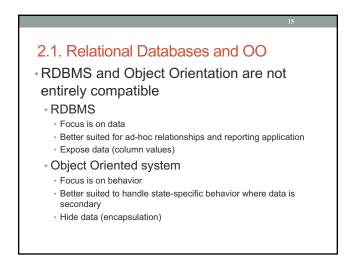
- Logical data models, when they are implemented, become physical data models:
- · relational databases,
- network databases,
- · or hierarchical databases

1.3.2. Network Database Parent records and child records do not have 1-to-n (1:n) correspondences; rather, they are in many-to-many (m:n) correspondence Sometimes called CODASYL database Baseball club Swimming club Track & field club Parent Billy John Susie Bobby Jerry Tommy Nancy Child

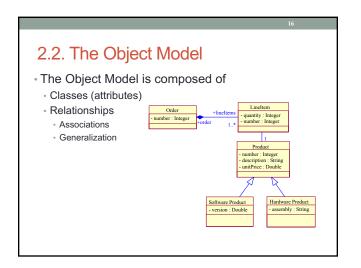
1.3.1. Hierarchical Database (Tree-Structure Database) Divides records into parents and children and shows the relationship with a hierarchical structure 1-to-many (1:n) correspondences between parent records and child records Baseball club Swimming club Track & field club Parent Billy John Susie Bobby Jerry Tommy Nancy Child

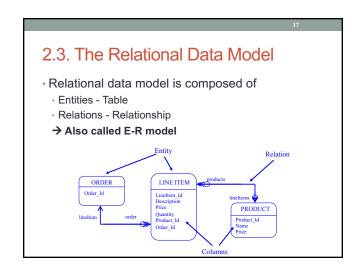






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2.3.2. Relations/Relationships

Relations between entities or relationship between tables

Multiplicity/Cardinality

One-to-one (1:1)

One-to-many (1:m)

Many-to-one (m:1)

Many-to-many (m:n)

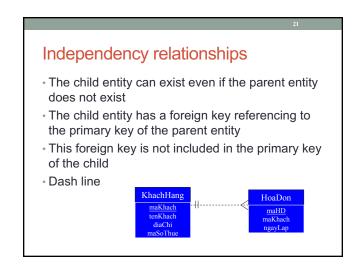
(Normally, many-to-many relation is devided to one-to-many and many-to-one relations)

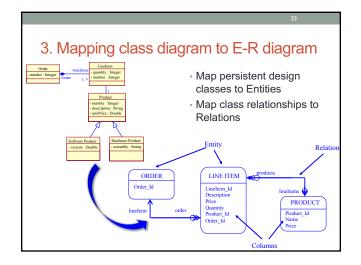
2.3.1. Entities/Tables • Entities is mapped to table when design physical database • Including • Columns: Attributes • Rows: Concrete values of attributes Columns Colu

Dependency relationships

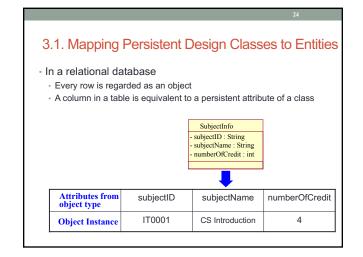
- The child entity can exist only when the parent entity exists
- The child entity has a foreign key referencing to the primary key of the parent entity
- This foreign key is included in the primary key of the child
- Solid line

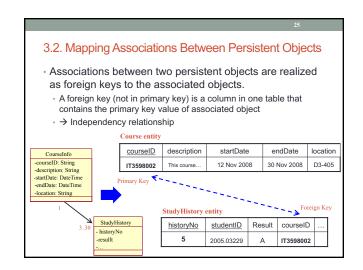


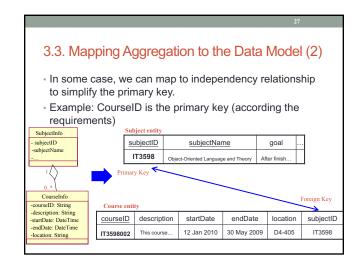


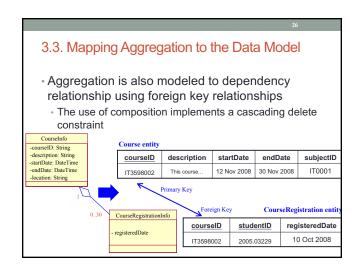


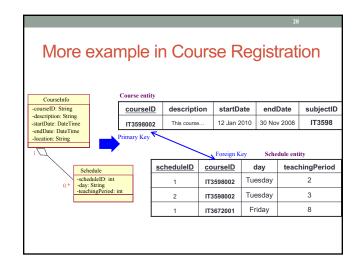
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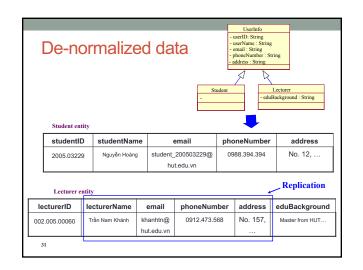


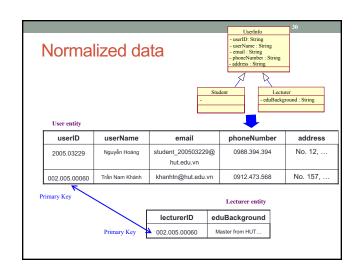


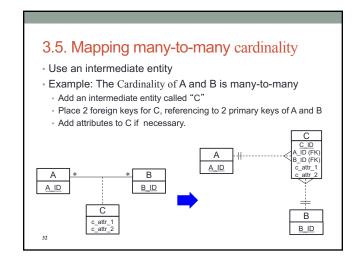


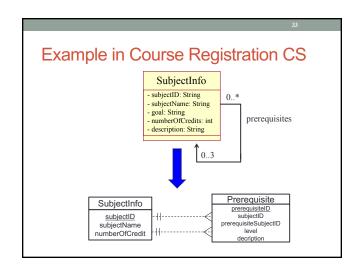
3.4. Modeling Inheritance in the Data Model

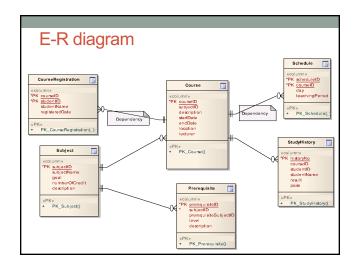
- A Data Model does not support modeling inheritance in a direct way
- Two options:
- Use separate tables (normalized data)
- Duplicate all inherited associations and attributes (denormalized data)











Content

- 1. Data models
- 2. Object model and Rational Data Model
- 3. Mapping class diagram to E-R diagram
- 4. Normalization

4.1. Overview of Normalization

- Normalization: the process of steps that will identify, for elimination, redundancies in a database design.
- Purpose of Normalization: to improve
- storage efficiency
- data integrity
- and scalability

4.1. Overview of Normalization (2)

- In relational model, methods exist for quantifying how efficient a database is.
- These classifications are called normal forms (or NF), and there are algorithms for converting a given database between them.
- Normalization generally involves splitting existing tables into multiple ones, which must be re-joined or linked each time a query is issued

4.3. Normal Forms

- Edgar F. Codd originally established three normal forms: 1NF, 2NF and 3NF.
- There are now others that are generally accepted, but 3NF is widely considered to be sufficient for most applications.
- Most tables when reaching 3NF are also in BCNF (Boyce-Codd Normal Form).

4.2. History



 Edgar F. Codd first proposed the process of normalization and what came to be known as the 1st normal form in his paper A Relational Model of Data for Large Shared Data Banks Codd stated:

"There is, in fact, a very simple elimination procedure which we shall call normalization. Through decomposition nonsimple domains are replaced by 'domains whose elements are atomic (nondecomposable) values".



Functionally determines

• In a table, a set of columns X, functionally determines another column Y...

 $X \rightarrow Y$

... if and only if each X value is associated with at most one Y value in a table.

• i.e. if you know X then there is only **one** possibility for Y.

Normal forms so Far...

First normal form

All data values are atomic, and so everything fits into a mathematical relation.

Third normal form

As 1NF plus no non-primary-key attribute is partially dependant on the primary key

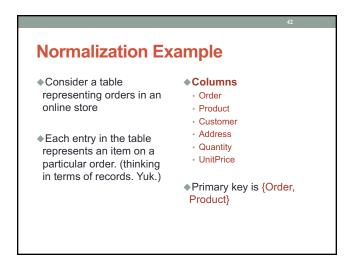
Functional Dependencies

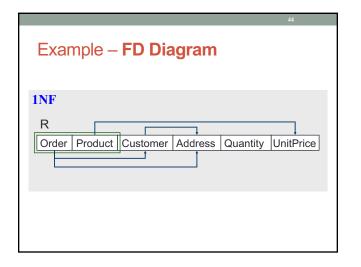
■ Each order is for a single customer {Order} → {Customer}

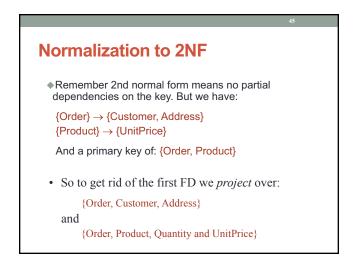
■ Each customer has a single address {Customer} → {Address}

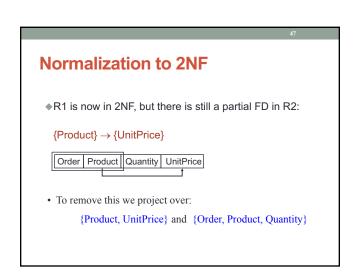
■ Each product has a single price {Product} → {UnitPrice}

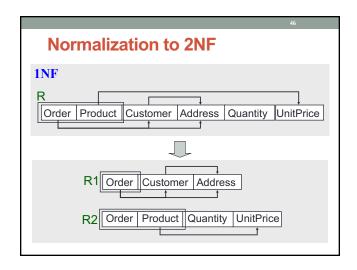
■ FD's 1 and 2 are transitive {Order} → {Address}

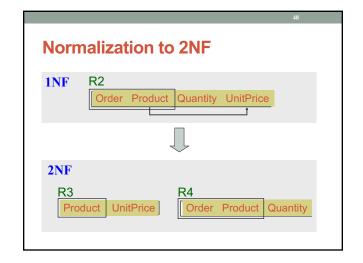


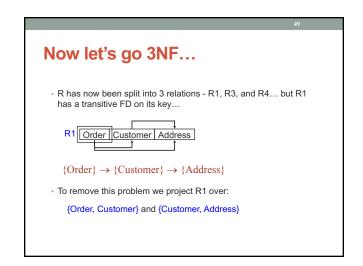


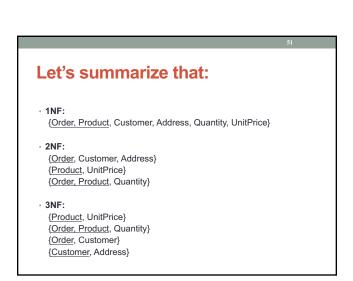


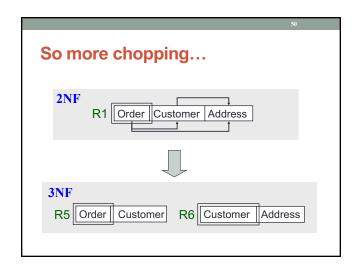


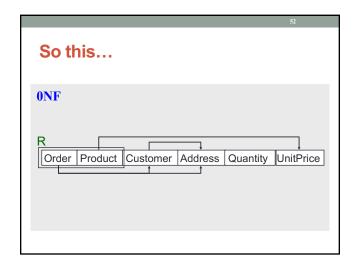




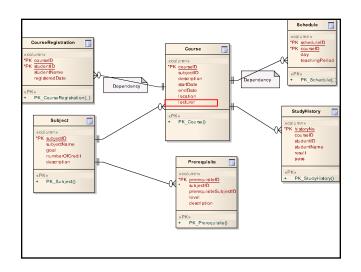












"Register for course" use case

- Make the E-R diagram from the previous step for "Register for course" use case to become:
- The first normal form
- The second normal form
- · The third normal form

