Notes on specifying user defined types

```
data Exp
= While Exp Exp
 Bool Bool
 If Exp Exp Exp
 Int Int
 Add Exp Exp
  Sub Exp Exp
 | Mul Exp Exp
 Div Exp Exp
 | Leq Exp Exp
  Char Char
 Ceq Exp Exp
 | Pair Exp Exp
 Fst Exp
  Snd Exp
  Cons Exp Exp
  Nil
  Head Exp
  Tail Exp
```

Null Exp

```
data Value data Typ

= IntV Int = TyVar String -- a, b, c
| PairV Addr | TyPair Typ Typ -- (Int . Bool)
| CharV Char | TyFun [Typ] Typ -- Int -> Bool -> Int
| BoolV Bool | TyList Typ -- [Int]
| ConsV Addr | TyCon String -- Bool, Char, etc
```

Recall how we divide the universe of values into types.

Note similarities between PairV and ConsV.

Almost ½ of the language is devoted to Pairs and Lists

Data as Heap Pointers

```
data Value
= IntV Int
| PairV Addr
| CharV Char
| BoolV Bool
| ConsV Addr
| NilV
```

What distinguishes PairV, ConsV, and NilV?

- They have different names
- •They point to consecutive blocks in the heap of different sizes.

Generic Constructors

A constructor with with N arguments, starting at Addr in Heap with name Cname

What operations?

- Construction
 - cons, nil, pair

- Selection
 - head, tail, fst, snd

- Predicate
 - null

Expressions

```
data Exp
                          data Exp
= While Exp Exp
                           = While Exp Exp
 | Bool Bool
                           | Bool SourcePos Bool
 | If Exp Exp Exp
                           | If Exp Exp Exp
 Pair Exp Exp
                           At Exp SourcePos [Exp]
                            | Lambda SourcePos [Vname] Exp
 | Fst Exp
 | Snd Exp
                           | Construction SourcePos Cname [Exp]
 Cons Exp Exp
 Nil
                            | Selection SourcePos Cname Int Exp
 Head Exp
                            | Predicate SourcePos Cname Exp
  Tail Exp
  Null Exp
```

Types

```
data Typ
data Typ
                          = TyVar String
 = TyVar String
                           | TyFun [Typ] Typ
  | TyFun [Typ] Typ
                           TyCon String [Typ]
  | TyPair Typ Typ
  | TyList Typ
                        intT = TyCon "Int" []
  | TyCon String
                        charT = TyCon "Char" []
                        boolT = TyCon "Bool" []
                        stringT = tylist charT
                        typair x y = TyCon "Pair" [x,y]
                        tylist x = TyCon "List" [x]
```

What operations?

Construction

```
- (cons a b), nil, (pair a b)
- (#cons a b), (# nil), (#pair a b)
```

Selection

```
- (head x), (tail x), (fst x), (snd x)
- (!cons 0 x), (!cons 1 x)
- (!pair 0 x) (!pair 1 x)
```

Predicate

```
- (null x), (@not (null x))
- (?nil x), (?cons x)
```

Semantics construction

```
run state (Construction c es) =
   do { (vals, state2) <- interpList vars state es
    ; let count = length es
        (addr, state3) = allocate count vals state2
    ; return(ConV c count addr, state3)}
(#node 3 'x' (#leaf) (#leaf))
                         es
```

Semantics Selection

```
run state (term@(Selection p c n e)) =
  do { (v,state2) <- interpE vars state e
    : case v of
       (ConV d m addr)
          | c==d && n<m
          -> return(access (addr+n) state2,state2)
       (ConV d m addr) | not(c==d) -> error ...
       (ConV d m addr) | not(n<m) -> error ...
       other -> error ("Non Construction in Selection")}
(!pair 0 (@ f 5))
                                 -- this is "fst"
```

Semantics Predicate

```
run state (term@(Predicate p c e)) =
   do { (v,state2) <- interpE vars state e
     ; case v of
       (ConV d m addr)
          | c==d -> return(BoolV True, state2)
       (ConV d m addr) -> return(BoolV False, state2)
       other -> error ("Non construction in Predicate")}
(?Cons (@append x y))
                       6
```

Some samples

```
(global nil [a] (# nil))
(fun head h (x [h]) (!cons 0 x))
(fun tail [a] (x [a]) (!cons 1 x))
(fun fst a (x (a.b)) (!pair 0 x))
(fun isnil Bool (1 [a]) (? nil 1))
(fun list1 [a] (x a) (#cons x nil))
(fun list2 [a] (x a y a)
    (#cons x (#cons y nil)))
(fun list3 [a] (x a y a z a)
      (#cons x (#cons y (#cons z nil))))
(fun snd b (x (a.b)) (!pair 1 x))
(fun fst b (x (a.b)) (!pair 0 x))
```

Defining new types

```
(data (Tree a)
   (#tip a)
   (#fork (Tree a) (Tree a)))

{ A type with no arguments }
(data (Color ) (#red) (#blue) (#green))

(data (Result a) (#found a) (#notFound))
```

Example

```
(fun length Int (l [a])
     (local (temp 0)
     (block
        (:= temp 0)
        (while (@not (?nil 1))
            (block
                (:= temp (+ temp 1))
                (:= 1 (@ tail 1))))
        temp)))
```

Abstract Data types

- Data definitions create types that have operations of
 - Construction
 - Selection
 - Predicate
- Other kinds of types are defined by their operations
 - (Env a)
 - lookup ((Env a) -> Int -> (Result a))
 - extend (Int -> a -> (Env a) -> (Env a))
 - empty (Env a)

Example

```
(adt (Env a) [(Char . a)]
  (global empty (Env a) nil)
  (fun extend (Env a) (key Char object a table (Env a))
     (#cons (#pair key object) table))
  (fun lookup (Result a) (tab (Env a) key Char)
     (if (?nil tab) (#notFound)
        (if (= key (@fst (@head tab)))
           (#found (@snd (@head tab)))
          (@lookup (@tail tab) key)))) )
```

Another Example

Modules

- Modules allow breaking a program into separate files
- Track what a file needs from others to compile successfully
- Track what a file might provide to other files
- Control names
- Track types across files.

Sig-Item

 A Sig-Item specifies the type of an item. It says nothing about how it is implemented

- (type (T a b))
- (val x Int)
- (val f (Int -> Bool))
- (data (T x) (#make x Int) (#none Bool))

A signature is a set of Sig-Items

```
(sign Stack
  (type (Stack a))
  (val push (a -> (Stack a)-> (Stack a)))
  (val emptySt (Stack a))
  ( val pop ((Stack a)-> (a . (Stack a)))))
```

Signatures

Appear in programs

```
(sig Stack
  (type (Stack a))
  (val push (a -> (Stack a)-> (Stack a)))
  (val emptySt (Stack a))
  (val pop ((Stack a)-> (a . (Stack a)))))
```

And also in *.sig files

```
(defsig Stack
(sig
(type (Stack a))
(val push (a -> (Stack a)-> (Stack a)))
(val emptySt (Stack a))
(val pop ((Stack a)-> (a . (Stack a)))))
```

Signatures can be read from files

```
(defsig Stack
  (sig
    (type (Stack a))
    (val push (a -> (Stack a)-> (Stack a)))
    (val emptySt (Stack a))
    (val pop ((Stack a)-> (a . (Stack a)))))
(signature Stack "test.sig")
```

SigExp

- A sigExp is a way of creating a set of sig-Items
- There is a syntax for SigExp

```
sigExp :=
  Id
| 'prelude'
| 'everything'
| '(' 'sig' { sigExp } ')'
| '(' hide' sigExp '(' {Id | id } ')' ')'
| '(' 'union' { sigExp } ')'
```

Examples

- prelude
- everything
- (hide prelude (Int Bool nil))
- (sig (val x Int) (data (T) (#a Int) (#b)))
- (union prelude (sig (val x Int) (data (T) (#a Int) (#b))))

Use of SigExp

- A SigExp is used to compute a set of sigItems for three different reasons
- 1. Describe what external functions a file depends on.
 - (module T in sigExp out sigExp)
- 2. Describe what subset of the definitions in a file should be exported
 - (module T in sigExp out sigExp)
- 3. Describe what subset of the exported functions should be imported
 - (import "test.e7" implementing sigExp)
 - (import "test.e7" hiding sigExp)

What needs to be imported

```
(module Small in (sig (val tom Int)) out everything)
(global temp Int 5)
(adt (Stack a) [a]
  (global emptySt (Stack a) (#nil))
  (fun push (Stack a) (x a xs (Stack a)) (#cons x xs))
  (fun pop (a. (Stack a)) (xs (Stack a))
            (#pair (!cons 0 xs) (!cons 1 xs)))
(global www Char 'c')
main
(:= temp (+ tom 1))
```

What should be exported

```
(signature E "envSig.sig")
(module Env2 in prelude out E)
(data (Tree a)
 (#leaf)
 (#node Int a (Tree a) (Tree a)))
Main 0
```

What should be imported

(signature Stack "stack3.sig")

(import "small.e7" implementing Stack)

```
(defsig Stack
(sig (type (Stack a))
        (val push (a -> (Stack a) -> (Stack a)))
        (val emptySt (Stack a))
        (val pop ((Stack a) -> (a.(Stack a))))

))
```