Additional Material for Lecture 6 Type checking

```
data Exp
= While Exp Exp
  Bool Bool
 If Exp Exp Exp
 Int Int
  Add Exp Exp
  Sub Exp Exp
 Mul Exp Exp
  Div Exp Exp
 Leg Exp Exp
  Char Char
 Ceg Exp Exp
 Pair Exp Exp
  Fst Exp
  Snd Exp
  Cons Exp Exp
  Nil
  Head Exp
  Tail Exp
```

Null Exp

```
data Value data Typ

= IntV Int = TyVar String -- a, b, c

| PairV Addr | TyPair Typ Typ -- (Int . Bool)

| CharV Char | TyFun [Typ] Typ -- Int -> Bool -> Int

| BoolV Bool | TyList Typ -- [Int]

| ConsV Addr | TyCon String -- Bool, Char, etc
```

We divide the universe of values into types.

Many expression forms are used to construct values of some type, or to consume values of some type.

Composite types

- Built from type constructors.
- In E6
 - (Int . Bool)
 - [Int]
 - Int -> a -> (Bool . Char)

Static Type Checking

Based on declarations of types

```
(fun append [a] (l [a] m [a] ) { return (l ++ m) }
     (if (@isnil 1) m
         (cons (head 1) (@append (tail 1) m))))
{ generate the list [1,2,...,n] }
(fun gen [Int] (n Int)
     (local (temp nil)
     (block
       (:= temp nil)
       (while (@not (@eq n 0)) (block
                  (:= temp (cons n temp))
                  (:= n (- n 1)))
        temp)))
```

Dynamic Type Checking

- Based on runtime predicates
- Recall language E5

```
exp := var
...

| '(' 'ispair' exp ')'
| '(' 'ischar' exp ')'
| '(' 'ispair' exp ')'
| '(' 'isint' exp ')'
```

Typing judgments 1

 In the context of an environment that maps names to types.

$$\frac{TE(x) = t}{TE \vdash x : t} \text{ (Var)}$$

Typing judgments 2

```
\frac{TE \vdash e_1 : \text{Int} \quad TE \vdash e_2 : \text{Int}}{TE \vdash (+ e_1 e_2) : \text{Int}} \text{ (Add)}
```

```
infer fs vs (term@(Add x y)) =
  do { check fs vs x intT "(+)"
    ; check fs vs y intT "(+)"
    ; return intT}
```

$$\frac{TE \vdash e_1 : \texttt{Bool} \quad TE \vdash e_2 : t \quad TE \vdash e_3 : t}{TE \vdash (\texttt{if} \ e_1 \ e_2 \ e_3) : t} \text{ (If)}$$

```
infer fs vs (term@(If x y z)) =
  do { check fs vs x boolT "if statement test"
     ; t1 <- infer fs vs y
     ; t2 <- infer fs vs z
     ; unify t2 t1 (loc term)
         ["\nWhile infering the term\n
                                           11
         ,show term
         ,"\nThe branches don't match" ]
     ; return t1}
```

Polymorphism

We add to the types, the notion of a type variable.
This can take on any type.

Implementation

The notion of a type variable

Fresh instances

```
(fun hd a (xs [a]) (head x))
                    [a23] -> a23
 (pair (@ hd (@list3 1 2 4))
                     [a25] -> a25
          (@ hd (cons True
                   (cons False nil))))
Each instance of hd gets a
                                       We see
fresh instance of the type
                                       a23 = Int
      `al -> a
                                      a25 = Bool
```

Making a fresh instance

User defined types

data Temp = F Float | C Float

```
boiling (F x) = x >= 212.0
```

boiling
$$(C x) = x >= 100.0$$

Recursive types

```
data Inttree
   = Branch Inttree Inttree
   | Leaf Int
sumleaves (Leaf i) = i
sumleaves (Branch 1 r)
   = sumleaves 1 + sumleaves r
```

Parameterized types

```
data Bintree a
   = Branch (Bintree a) (Bintree a)
    Leaf a
depth (Leaf i) = 0
depth (Branch 1 r)
   = 1 + max (depth 1) (depth r)
```

Enumerations

```
data Day = Mon | Tue |Wed |Thu
|Fri |Sat | Sun
```

```
weekday Sat = False
weekday Sun = False
weekday x = True
```

data Bool = True | False