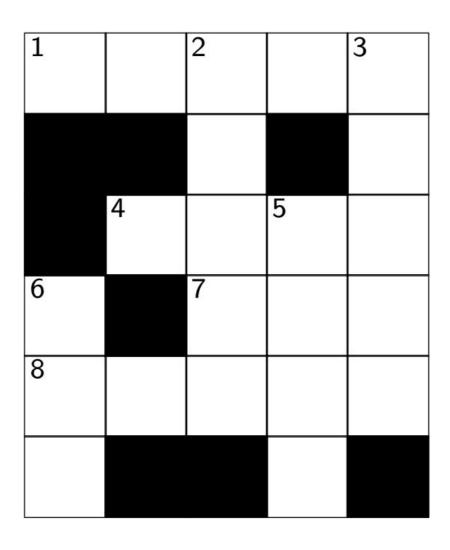
Motivation

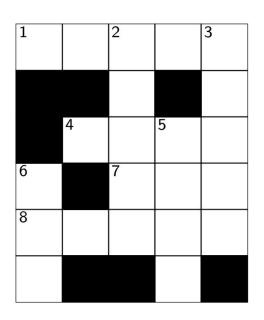
- Constraint Satisfaction Problems (CSPs)
- View logic problems through a systematic framework (allows solutions to be applied to multiple problems and vice versa)
- Apply search algorithms to logic problems (Recursive Backtracking Search)
- Simplify logic problems by eliminating constraint violating solutions (Arc Consistency)

Crossword Puzzle Domain

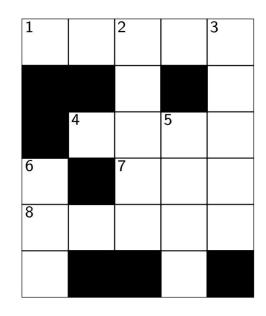


- Words: {AFT, ALE, EEL, HEEL, HIKE, HOSES, KEEL, KNOT, LASER, LEE, LINE, SAILS, SHEET, STEER, TIE}
- Words must start at a position label number and run left-to-right or top-to-bottom only
- Intersection boxes can contain only one letter
- Each word in the list provided can be used only once

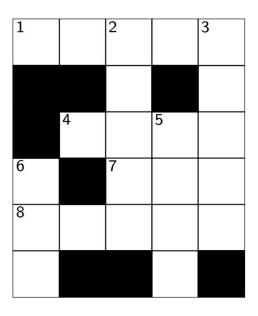
 What are the Variables and their Domains in this CSP?



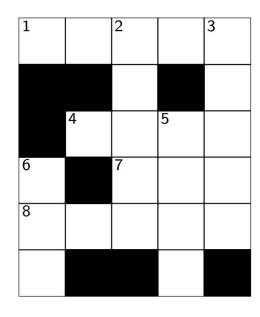
- What are the Variables and their Domains in this CSP?
 - Variables: 1-across, 2-down, 3-down, 4-across, 5-down, 6-down, 7-across and 8-across
 - **Domains:** {AFT, ALE, EEL, HEEL, HIKE, HOSES, KEEL, KNOT, LASER, LEE, LINE, SAILS, SHEET, STEER, TIE} for all variables



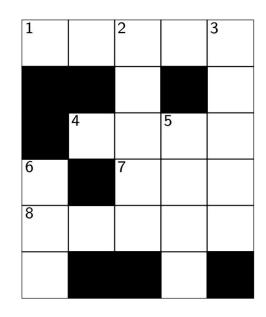
What are the Constraints in this CSP?



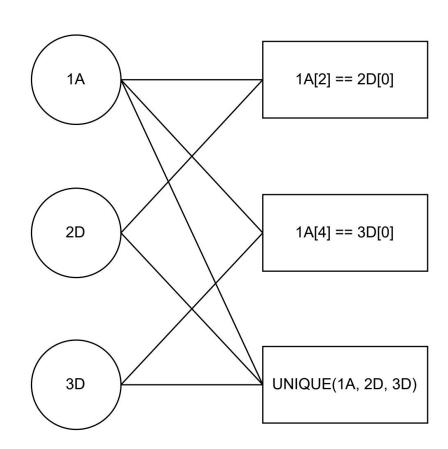
- What are the Constraints in this CSP?
 - Variables can be set to only words of the correct length (these are domain constraints / 1-constraints)
 - Letters used at the intersection of two words must be equal (binary constraints / 2-constraints)
 - Each word is used only once
 - Is this a "global" constraint? Or is it a set of binary constraints, in which each pair of words are linked by a "not-equal-to" predicate?
 - It can be either we usually prefer to use constraints with the fewest number of variables, but this problem is small enough that one global constraint is fine



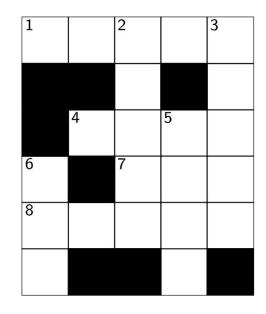
- Start to sketch a constraint graph for the crossword CSP
- Choose one variable, then sketch out the constraints that variable occurs in, and finally include any other variables that are also in scope of the constraints you have already drawn.
- The full constraint graph will be too large to draw easily



- Start to sketch a constraint graph for the crossword CSP
- Use a bipartite graph
- A partial solution is shown, starting from
 1A
- The coordinates refer to the grid in row-column from the top left corner (0-indexed)



- Apply domain consistency to this CSP, and restate any variable domains that change
- Does the structure of the constraint graph change?



- Apply domain consistency to this CSP, and restate any variable domains that change
 - Domains are reduced to those with answers with word lengths that are consistent with the number of blank spaces, so all variable domains are changed
 - Domains after applying node consistency:

```
dom[1A] = {HOSES, LASER, SAILS, SHEET, STEER}
dom[2D] = {HOSES, LASER, SAILS, SHEET, STEER}
dom[3D] = {HOSES, LASER, SAILS, SHEET, STEER}
dom[3D] = {HOSES, LASER, SAILS, SHEET, STEER}
dom[4A] = {HEEL, HIKE, KEEL, KNOT, LINE}
dom[8A] = {HOSES, LASER, SAILS, SHEET, STEER}
```

- Does the structure of the constraint graph change?
 - No, as domain constraints (which apply to only 1 variable) are not represented in the constraint graph

Backtracking Search

- Apply backtracking search to the domain-consistent constraint graph
- Record the number of nodes expanded in the search procedure
- You can trace the algorithm manually, e.g. by sketching a search tree, or develop code to answer this question (reusing some of your tree search code from previous weeks)

```
function Backtracking-Search(csp) returns solution/failure return Recursive-Backtracking(\{\}, csp) function Recursive-Backtracking(assignment, csp) returns soln/failure if assignment is complete then return assignment var \leftarrow Select-Unassigned-Variable(Variables[csp], assignment, csp) for each value in Order-Domain-Values(var, assignment, csp) do if value is consistent with assignment given Constraints[csp] then add \{var = value\} to assignment result \leftarrow Recursive-Backtracking(assignment, csp) if result \neq failure then return result remove \{var = value\} from assignment return failure
```

Backtracking Search

• As a demonstration, we can manually trace Backtracking-Search:

- 1. Select 1-across, as the first unassigned variable.
- Check if the first value in the domain of 1-across conflicts with the existing (empty) assignment; it doesn't, so expand a node for 1-across=HOSES.
- 3. Now select 2-down, as the next unassigned variable; this is using the recursive function call.
- 4. Check if the first value in the domain of 2-down, HOSES, conflicts with the existing assignment it does! Throw it away.
- 5. The second value, LASER also conflicts with the existing assignment, so throw it away. At this point we have iterated over the inner if statement twice.
- 6. The third value in the domain, 2-down=SAILS is consistent with 1-across=HOSES, so expand a child node of 1-across=HOSES given by 2-down=SAILS
- 7. Select 3-down as the next unassigned variable.
- 8. . . . and so on.

Backtracking Search

• Eventually, you will find the solution as:

Н	0	S	Е	S
		Α		Т
	Н	Ī	K	Е
Α		L	Е	Е
L	А	S	Е	R
Е			L	

- Apply arc-consistency to this CSP (manually or in code)
- Record the number of arc-consistency check operations that are performed
- What is the outcome of applying arc consistency?

```
function AC-3(csp) returns false if an inconsistency is found and true otherwise
  inputs: csp, a binary CSP with components (X, D, C)
  local variables: queue, a queue of arcs, initially all the arcs in csp
  while queue is not empty do
     (X_i, X_j) \leftarrow \text{REMOVE-FIRST}(queue)
    if REVISE(csp, X_i, X_i) then
       if size of D_i = 0 then return false
       for each X_k in X_i. NEIGHBORS - \{X_i\} do
          add (X_k, X_i) to queue
  return true
function REVISE(csp, X_i, X_j) returns true iff we revise the domain of X_i
  revised \leftarrow false
  for each x in D_i do
    if no value y in D_i allows (x,y) to satisfy the constraint between X_i and X_j then
       delete x from D_i
       revised \leftarrow true
  return revised
```

• The first few operations of AC-3:

- a) 1-across: {HOSES, LASER, SAILS, SHEET, STEER}
- b) Check for a value in dom(2-down) consistent with each element of dom(1-across) 2-down: {HOSES, LASER, SAILS, SHEET, STEER}
- c) Remove SAILS, SHEET, STEER from dom(1-across), because there is no value in dom(2-down) that is consistent with the value assignments.

 Set 1-across = {HOSES, LASER}
- d) Next, check for a value in dom(3-down) consistent with each element of dom(1-across) 3-down: {HOSES, LASER, SAILS, SHEET, STEER}
- e) Remove LASER from dom(1-across)
- f) 1-across must be HOSES, as all other values are not arc consistent.
- g) Now move on to the next variable (e.g. 2-down). . .

• If needed, apply backtracking search to the arc-consistent CSP

 Compare the number of search expansion and/or consistency check operations of backtracking search and (arc-consistency + backtracking search)

- If needed, apply backtracking search to the arc-consistent CSP
 - If you apply the arc-consistency algorithm correctly, you will find that it solves the cross-word puzzle, and no further search is required
 - Note that applying arc consistency will not always lead to a unique solution, so search on the reduced may still be required for some problems
- Compare the number of search expansion and/or consistency check operations of backtracking search and (arc-consistency + backtracking search)

- If needed, apply backtracking search to the arc-consistent CSP
 - If you apply the arc-consistency algorithm correctly, you will find that it solves the cross-word puzzle, and no further search is required
 - Note that applying arc consistency will not always lead to a unique solution, so search on the reduced may still be required for some problems
- Compare the number of search expansion and/or consistency check operations of backtracking search and (arc-consistency + backtracking search)
 - Even allowing for different variations on counting the number of operations, AC-3 solves the problem in many fewer operations than backtracking search