PRINCIPLES OF MODELING FOR CYBER PHYSICAL SYSTEMS

Assignment #6

Transition Systems and Linear Temporal Logic

Due Date: 11-15-2018 at 2:00pm (Before the lecture)

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The problems have been adapted from the book - Principles of model checking by Christel Baier and Joost-Pieter Katoen. Remember the following notation:

\Diamond	eventually
	always
\bigcirc	next
U	until
¬	negation
V	or
Λ	and

1 Problem 1

(15 points - 2+2+2+3+3+3)

Consider the following transition system over the set of atomic propositions $\{a, b\}$: Indicate for each of the following LTL formulae the set of states for which these formulae are fulfilled:

- (a) $\bigcirc a$
- (b) $\bigcirc\bigcirc\bigcirc$ a

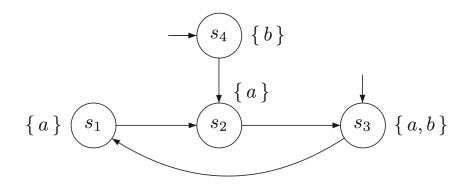


Figure 1.1:

- (c) □ *b*
- (d) □◊ *a*
- (e) \Box ($b \cup a$)
- (f) $\Diamond (a \cup b)$

2 Problem 2

(15 points - 2+2+3+2+3+3)

Consider the transition system *TS* over the set of atomic propositions $AP = \{a, b, c\}$:

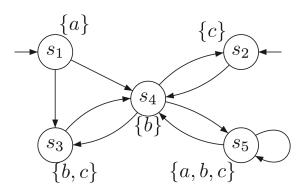


Figure 2.1:

Decide for each of the LTL formulae φ_i below, whether $TS \vDash \varphi_i$ holds. Justify your answers! If $TS \nvDash \varphi_i$, provide a path $\pi \in \text{Paths}(TS)$ such that $\pi \nvDash \varphi_i$

$$\varphi_{1} = \Diamond \Box c$$

$$\varphi_{2} = \Box \Diamond c$$

$$\varphi_{3} = \bigcirc \neg c \rightarrow \bigcirc \bigcirc c$$

$$\varphi_{4} = \Box a$$

$$\varphi_{5} = a \cup \Box (b \vee c)$$

$$\varphi_{6} = (\bigcirc \bigcirc b) \cup (b \vee c)$$

3 PROBLEM 3

(20 points)

Suppose we have two users, *Peter* and *Betsy*, and a single printer device *Printer*. Both users perform several tasks, and every now and then they want to print their results on the *Printer*. Since there is only a single printer, only one user can print a job at a time. Suppose we have the following atomic propositions for *Peter* at our disposal:

- Peter.request ::= indicates that Peter requests usage of the printer;
- *Peter.use* ::= indicates that *Peter* uses the printer;
- Peter.release ::= indicates that Peter releases the printer.

For *Betsy*, similar predicates (propositions) are defined. Specify in LTL the following properties:

- (a) Mutual exclusion, i.e., only one user at a time can use the printer.
- (b) Finite time of usage, i.e., a user can print only for a finite amount of time.
- (c) Absence of individual starvation, i.e., if a user wants to print something, he/she eventually is able to do so.
- (d) Absence of blocking, i.e., a user can always request to use the printer
- (e) Alternating access, i.e., users must strictly alternate in printing.

4 PROBLEM 4

(20 points)

Consider an elevator system that services N > 0 floors numbered 0 through N-1. There is an elevator door at each floor with a call-button and an indicator light that signals whether or not the elevator has been called. For simplicity consider N=4. Present a set of atomic propositions - try to minimize the number of propositions - that are needed to describe the following properties of the elevator system as LTL formulae and give the corresponding LTL formulae:

- (a) The doors are "safe", i.e., a floor door is never open if the elevator is not present at the given floor.
- (b) A requested floor will be served sometime.
- (c) Again and again the elevator returns to floor 0.
- (d) When the top floor is requested, the elevator serves it immediately and does not stop on the way there.

5 PROBLEM 5

(30 points - 10 + 10 + 5 + 5)

Which of the following equivalences (*if and only if*) are correct? Prove the equivalence or provide a counterexample that illustrates that the formula on the left and the formula on the right are not equivalent.

- (a) $\Box \varphi \rightarrow \Diamond \psi \equiv \varphi \cup (\psi \lor \neg \varphi)$
- (b) $\Diamond \Box \varphi \rightarrow \Box \Diamond \psi \equiv \Box (\varphi \cup (\psi \vee \neg \varphi))$
- (c) $\Box\Box(\varphi\cup\neg\psi)\equiv\neg\Diamond(\neg\varphi\wedge\psi)$
- (d) $\Diamond (\varphi \wedge \psi) \equiv \Diamond \varphi \wedge \Diamond \psi$