

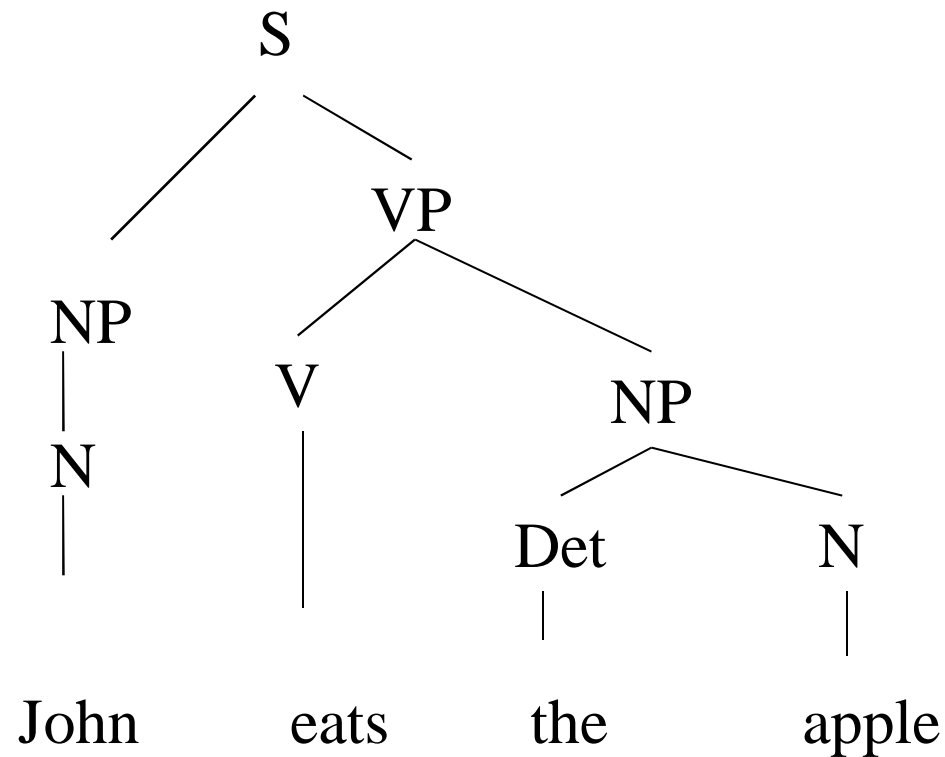
# Grammar development

- Grammar rules
- context free/context sensitive
- Prolog grammar : parser
- Grammar formalism in Prolog:  
Definite Clause Grammars (DCGs)

# Grammar as a production system

*John eats the apple*

Syntax tree:



# Grammar as a production system

Rules:

$$S \rightarrow NP VP$$
$$NP \rightarrow Det N$$
$$NP \rightarrow N$$
$$VP \rightarrow V NP$$
$$VP \rightarrow V$$

# Grammar as a production system

Lexicon:

$N \rightarrow \text{mary}$	$V \rightarrow \text{loves}$	$\text{Det} \rightarrow \text{the}$
$N \rightarrow \text{john}$	$V \rightarrow \text{eats}$	
$N \rightarrow \text{man}$	$V \rightarrow \text{sings}$	
$N \rightarrow \text{apple}$		

# Grammar as a production system

Rules:

$S \rightarrow NP VP$

$NP \rightarrow Det N$

$NP \rightarrow N$

$VP \rightarrow V NP$

$VP \rightarrow V$

$N \rightarrow \text{apple, john, ...}$

$V \rightarrow \text{loves, eats, ...}$

$Det \rightarrow \text{the}$

context free grammar!

# Context free grammar

- Set of *terminal symbols* ( $T$ )
- Set of *non-terminal symbols* ( $NT$ )
- A distinguished NT symbol: *start symbol* ( $S$ )
- Set of (*production-)*rules

of the form:

$$NT \rightarrow \alpha, \text{ with } \alpha \in (NT \cup T)^*$$

$$NT \rightarrow T$$

# Context sensitive grammar

- Set of *terminal symbols* ( $T$ )
- Set of *non-terminal symbols* ( $NT$ )
- A distinguished NT symbol: *start symbol* ( $S$ )
- Set of (*production*) *rules*  
of the form:

$$\alpha A \gamma \rightarrow \alpha \beta \gamma$$

$A$  from  $NT$ ,  
 $\alpha, \gamma, \beta$  from  $(NT \cup T)^*$

# Parsing

- ?- s([john,eats,the,apple]).

yes

- ?- s ([john, the,apple,eats]).

no

- ?- s ([the, eats,john,apple]).

no

- ?- s ([the,apple,eats,john]).

??



# Derivations

Rules:

1.  $S \rightarrow NP VP$
2.  $NP \rightarrow Det N$
3.  $NP \rightarrow N$
4.  $VP \rightarrow V NP$
5.  $VP \rightarrow V$

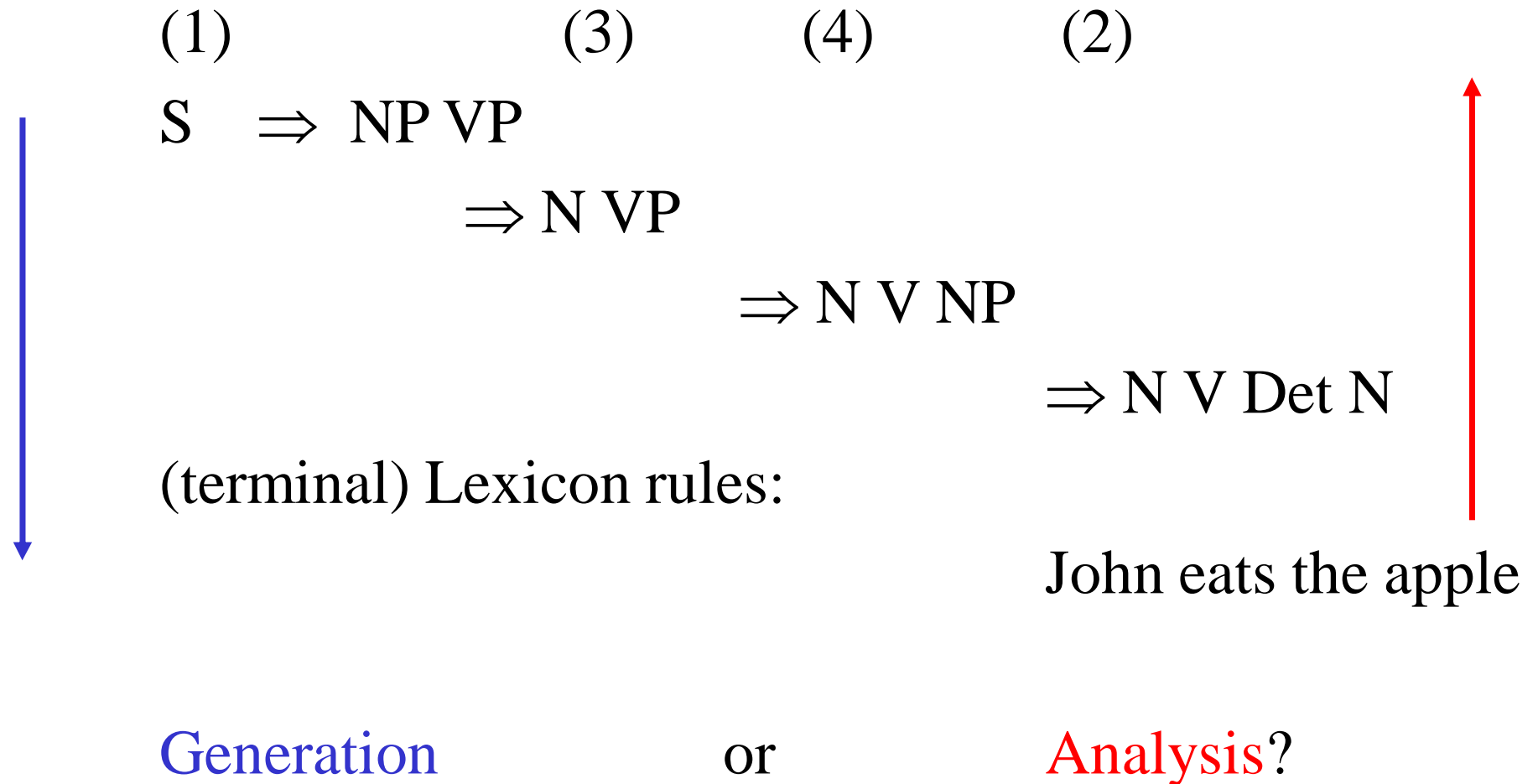
# Derivations

(1)	(3)	(4)	(2)
S	$\Rightarrow$ NP VP		
	$\Rightarrow$ N VP		
		$\Rightarrow$ N V NP	
			$\Rightarrow$ N V Det N

(terminal) lexicon rules:

John eats the apple

# Derivations: direction ?



# Generation

- $s(S) \text{ :- np(NP), vp(VP), append(NP,VP,S).$
- $np(NP) \text{ :- det(Det), n(N), append(Det,N,NP).$
- $np(NP) \text{ :- n(N), NP=N.}$
- $vp(VP) \text{ :- v(V), np(NP), append(V,NP,VP).$
- $vp(VP) \text{ :- v(V), VP=V.}$

# Lexicon

- n([mary]).
- n([apple]).
- n([john]).
- n([man]).
- v([loves]).
- v([eats]).
- v([sings]).
- det([the]).

# Generation

- ?- s(S)
- $X = [\text{the,mary,loves,the,mary}] ;$
- $X = [\text{the,mary,loves,the,apple}] ;$
- ...
- $X = [\text{john,eats,the,apple}] ;$
- ...

# Parsing?

?- s([john,eats,the,apple]).

- $s(S) \text{ :- np(NP), vp(VP), append(NP,VP,S).}$
- $np(NP) \text{ :- det(Det), n(N), append(Det,N,NP).}$
- $np(NP) \text{ :- n(N), NP=N.}$
- $vp(VP) \text{ :- v(V), np(NP), append(V,NP,VP).}$
- $vp(VP) \text{ :- v(V), VP=V.}$

# Parsing

?- s([john,eats,the,apple]).

- $s(S) \text{ :- append(NP,VP,S), np(NP), vp(VP).}$
- $np(NP) \text{ :- append(Det,N,NP), det(Det), n(N).}$
- $np(NP) \text{ :- NP=N, n(N).}$
- $vp(VP) \text{ :- append(V,NP,VP), v(V), np(NP).}$
- $vp(VP) \text{ :- VP=V, v(V).}$



# Is parsing efficient ?

?- s([john,eats,the,apple]).

- $s(S) \text{ :- append}(NP,VP,S), np(NP), vp(VP).$
- $np(NP) \text{ :- append}(\text{Det},N,NP), \text{det}(\text{Det}), n(N).$
- $np(NP) \text{ :- NP=N}, n(N).$
- $vp(VP) \text{ :- append}(V,NP,VP), v(V), np(NP).$
- $vp(VP) \text{ :- VP=V}, v(V).$

trace !!

## Is parsing efficient ?

No!

Use two lists instead:

- Input
- Part of the input that remains after cutting off the part recognized by the rule:

**Difference lists !**

```
s([john,eats,the,apple],[]) :-  
    np([john,eats,the,apple],[eats,the,apple]),  
    vp([eats,the,apple,[]]).
```

# More efficient parsing

Schematically: ..

- $s(\text{NPVP}, []) \quad :- \quad \text{np}(\text{NPVP}, \text{VP}), \text{vp}(\text{VP}, []).$
- $\text{np}(\text{NPVP}, \text{VP}) \quad :- \quad \text{det}(\text{NPVP}, \text{NVP}), \text{n}(\text{NVP}, \text{VP}).$
- $\text{np}(\text{NVP}, \text{VP}) \quad :- \quad \text{n}(\text{NVP}, \text{VP}).$
- $\text{vp}(\text{VP}, []) \quad :- \quad \text{v}(\text{VP}, \text{NP}), \text{np}(\text{NP}, []).$
- $\text{vp}(\text{VP}, []) \quad :- \quad \text{v}(\text{VP}, []).$

# More efficient parsing

Schematically, ..more generally!

- $s(SR, R) \quad :- \quad np(SR, VPR), \quad vp(VPR, R).$
- $np(DetNPR, R) \quad :- \quad det(DetNPR, NR), \quad n(NR, R).$
- $np(NR, R) \quad :- \quad n(NR, R).$
- $vp(VNPR, R) \quad :- \quad v(VNPR, NPR), \quad np(NPR, R).$
- $vp(VR) \quad :- \quad v(VR, R).$

# More efficient parsing

## Lexicon

- `n([mary|R],R).`
- `n([apple|R],R).`
- `n([john|R],R).`
- ...
- `v([loves|R],R).`
- `v([eats|R],R).`
- ...
- `det([the|R],R).`
- ...

# More efficient parsing **and** Generation

- ?- s([john,eats,the,apple],[]).

yes

- ?- s(X,[]).

X = [the,mary,loves,the,mary] ;

...

- ?- np([the,apple],[]).

yes

- ...

# A friendlier type of grammar rule writing:

- Definite clause grammars
- A formalism implemented in prolog

## A DCG notation for our grammar:

Rules: omit lists (they are added in compilation to Prolog)

s --> np, vp.

np --> det, n.

np --> n.

vp --> v, np.

vp --> v.



## A DCG notation for our grammar :

Lexicon: name word only (rest is added automatically)

n --> [mary].

n --> [apple].

...

v --> [loves].

v --> [eats].

...

det --> [the].

..

A DCG notation for our grammar :

DCG rule operator is recognized: therefore simply *consult*  
DCG

*listing* shows that internally difference lists are used.

# Grammars/DCGs

Question: Are derivations unique?

- In our grammar?
- In an extended grammar, with:
- $PP \rightarrow P\ NP$
- $NP \rightarrow NP\ PP$
- $VP \rightarrow VP\ PP$
- $N \rightarrow \{\text{worm, knife}\}$
- $P \rightarrow \{\text{with, in}\}$

cf: *John eats the apple with the worm with the knife.*

# Grammars/DCGs

How can the derivations/derivation trees be made visible ?

What about constraints like Subj-Verb agreement?