

Realtime Event Summarization

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3.1 Problem Definition

Suppose that we have N tweets to be summarized, within which M tweets are credible and relevant. We are going to select a few representative posts from the tweet universe to form the summary. To model this problem, we use a vector $\tilde{\mathbf{x}} \in R^N$, where each element $\tilde{x}_j \in \{0, 1\}$ is a binary variable. If a tweet i is chosen as the abstractive sentence, the corresponding $\tilde{x}_i = 1$. Otherwise, we set $\tilde{x}_i = 0$. We use another N -dim vector $\tilde{\mathbf{c}} \in R^N$ to describe the loss of choosing each tweet as a candidate. $\tilde{\mathbf{A}} \in R^{M \times N}$ is a similarity matrix, where $\tilde{a}_{i,j}$ is the similarity between a credible and relevant tweet i and a candidate tweet j in the tweet universe. $\mathbf{b} \in R^M$ is a weight vector, where b_i indicates the importance of i being covered in the summary. Our objective is to

$$\min \tilde{\mathbf{c}}^T \tilde{\mathbf{x}} \text{ subject to } \tilde{\mathbf{A}} \tilde{\mathbf{x}} \geq \mathbf{b}, \tilde{\mathbf{x}} \in \{0, 1\}$$

3.2 Methodology Overview

We first transform it to a standard form of bounded linear programming problem by making the following adjustments: $\mathbf{c} = [\tilde{\mathbf{c}}, \mathbf{0}]$, $\mathbf{x} = [\tilde{\mathbf{x}}, \mathbf{z}]^T$, $\mathbf{A} = [\tilde{\mathbf{A}}, -\mathbf{I}]$, where \mathbf{I} is the $M \times M$ identity matrix. Therefore we have

$$\min \mathbf{c} \mathbf{x} \text{ subject to } \mathbf{A} \mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}, \tilde{\mathbf{x}} \leq \mathbf{1} \quad (1)$$

We modify the LP so that there is an easy choice of basic solution. We start by solving

$$\min \mathbf{e}^T \mathbf{s} \text{ subject to } \mathbf{A} \mathbf{x} + \mathbf{I} \mathbf{s} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}, \tilde{\mathbf{x}} \leq \mathbf{1}, \mathbf{s} \geq \mathbf{0} \quad (2)$$

where \mathbf{e} is the vector of all ones, \mathbf{s} are called artificial variables, and $b_i \geq 0$. To solve the above LP relaxation, we adopt the bounded simplex method. In

each iterate, only M variables are selected as the basic points. Suppose $\mathbf{A}_{.j}$ is the j th column, and the corresponding columns of basic variables in \mathbf{A} are denoted as $\mathbf{A}_{\mathbf{B}}$, others are denoted as $\mathbf{A}_{\mathbf{N}}$. Non-basic variables are either at upper bound ($j \in U$) or at lower bound ($j \in L$).

Define $\mathbf{c} = [\mathbf{e}, \mathbf{0}]$, $\mathbf{x}' = [\mathbf{x}, \mathbf{s}]^T$ and $\mathbf{A}' = [\mathbf{A}, \mathbf{I}]$ so that the constraints of the modified LP can be written as $\mathbf{A}' \mathbf{x}' = \mathbf{b}$, $\mathbf{x}' \geq \mathbf{0}$, $\tilde{\mathbf{x}} \leq \mathbf{1}$.

Step 1: Solve $\mathbf{x}'_{\mathbf{B}} = \mathbf{A}'_{\mathbf{B}}^{-1}(\mathbf{b} - \mathbf{u})$, where $\mathbf{u} = \sum_{j \in U} \mathbf{A}'_{.j}$. Let \mathbf{B} be the indices of the artificial variables and set the all non-basic variables to be $x_j = 0$. Then \mathbf{B} is a basis, since the corresponding columns of $\mathbf{A}'_{\mathbf{B}}$ are \mathbf{I} , the identity, the corresponding basic feasible solution is $\mathbf{x} = \mathbf{0}$, $\mathbf{z} = \mathbf{b}$.

Step 2: Pricing. $\mathbf{y} = \mathbf{B}^{T-1} \mathbf{e}_{\mathbf{B}}$

Step 3: Compute $\bar{c}_j = c_j - \mathbf{y}^T \mathbf{A}'_{.j}$. If $x_j = 0$, $\bar{c}_j > 0$ and $x_j = 1$, $\bar{c}_j < 0$, then the problem is solved. Else pick q the most contradictive variable, i.e. the most negative \bar{c}_q for $x_q = 0$ as the entering index.

Step 4: $d = \mathbf{B}^{-1} \mathbf{A}'_{.q}$.

Step 5: There are three subcases as follows:

If $q \in \{\mathbf{z}_{\mathbf{N}}, \mathbf{s}_{\mathbf{N}}\}$, calculate $\min_i \left\{ \left\{ \frac{x_{\mathbf{B}_i}}{d_i}, \frac{s_{\mathbf{B}_i}}{d_i}, \frac{z_{\mathbf{B}_i}}{d_i} \right\} \forall d_i > 0, \left\{ \frac{x_{\mathbf{B}_i}-1}{d_i} \right\} \forall d_i < 0 \right\}$.

Let $p = i$, $x_q^{new} = \min$, $\mathbf{x}'_{\mathbf{B}^{old}} = \mathbf{x}'_{\mathbf{B}^{old}} - dx_q^{new}$, $\mathbf{B}^{new} \leftarrow \mathbf{B}^{old} - \{p\} \cup \{q\}$.

If $q \in L$, calculate $\min_i \left\{ \left\{ \frac{x_{\mathbf{B}_i}}{d_i}, \frac{s_{\mathbf{B}_i}}{d_i}, \frac{z_{\mathbf{B}_i}}{d_i} \right\} \forall d_i > 0, \left\{ \frac{x_{\mathbf{B}_i}-1}{d_i} \right\} \forall d_i < 0, 1 \right\}$, $x_q^{new} = \min$, $\mathbf{x}'_{\mathbf{B}^{old}} = \mathbf{x}'_{\mathbf{B}^{old}} - dx_q^{new}$. If $x_q^{new} \neq 1$, let $p = i$, $\mathbf{B}^{new} \leftarrow \mathbf{B}^{old} - \{p\} \cup \{q\}$.

If $q \in U$, calculate $\min_i \left\{ \left\{ \frac{x_{\mathbf{B}_i}}{-d_i}, \frac{s_{\mathbf{B}_i}}{-d_i}, \frac{z_{\mathbf{B}_i}}{-d_i} \right\} \forall d_i < 0, \left\{ \frac{1-x_{\mathbf{B}_i}}{d_i} \right\} \forall d_i > 0, 1 \right\}$, $x_q^{new} = 1 - \min$, $\mathbf{x}'_{\mathbf{B}^{old}} = \mathbf{x}'_{\mathbf{B}^{old}} + d(1 - x_q^{new})$. If $x_q^{new} \neq 0$, let $p = i$, $\mathbf{B}^{new} \leftarrow \mathbf{B}^{old} - \{p\} \cup \{q\}$.

Step 6: If all artificial variables are non-basic or some artificial variables are in the basis but all $x_{z_i} = 0$ then go to Step 7, otherwise return to Step 2.

Step 7: If some artificial variables are in the basis, remove them from basis. Change $\mathbf{c} = [\tilde{\mathbf{c}}, \mathbf{0}]$, repeat Step 2 to Step 5.

4 Update Summarization

$$\min c_0 x_0 + c_1 x_1 \text{ s.t. } \begin{bmatrix} A & D \\ D^T & B \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \end{bmatrix} \geq \begin{bmatrix} b_0 \\ b_1 \end{bmatrix} \quad (3)$$

$$D^T x_0 \geq b_1$$