

Realtime Event Summarization from Tweets with Replacements

Lingting Lin, Yunjie Wang, Chen Lin

March 11, 2018

1 Introduction

Accidents, disasters, political rallies...we are eager to gather information about different kinds of live events that happen around us. In the past, we rely on experienced journalists to cover the stories. At now, thanks to the emergence of micro-blogging platforms, we are provided with more instant reports from various information sources. Study shows that the majority (over 85%) of trending topics in microblog sphere are headline news and real-life events [?]. The massive amount of event related tweets highlights the importance of an event summarization system.

To facilitates knowledge management and improves user experiences, the event summary needs to be realtime. On one hand, the response must be realtime, which requires an efficient algorithm to quickly generate a summary. On the other hand, the summary must reflect the current status of the event, which indicates the summary should be updated if necessary.

We have witnessed rapidly increasing popularity of research efforts in event summarization from tweets []. Most previous research works produce the summary in an extractive manner, i.e. they extract most representative tweets based on several criterion, A few researchers also mentioned the realtime factor in event summarizations. However, their summaries are incremental, providing a good coverage, but suffering from poor brevity and inconsistency. In their methods, former summaries can not be deleted or replaced. When there is new information to be covered, the summary gets more verbose. Furthermore, former summaries can be inadequate to deliver up-to-date information.

There are three scenarios when a former summary needs to be deleted or replaced. (1) A better candidate appears for a brief summary. (2) Conflicting information (3) Outdated information. first-hand

In this paper we study event summarization with replacements. Challenges: steaming, minimal storage.

The second challenge

2 Related Work

3 Static Summarization

3.1 Problem Definition

Suppose that we have N tweets to be summarized, within which M tweets are credible and relevant. We are going to select a few representative posts from the tweet universe to form the summary. To model this problem, we use a vector $\tilde{\mathbf{x}} \in R^N$, where each element $\tilde{x}_j \in \{0, 1\}$ is a binary variable. If a tweet i is chosen as the abstractive sentence, the corresponding $\tilde{x}_i = 1$. Otherwise, we set $\tilde{x}_i = 0$. We use another N -dim vector $\tilde{\mathbf{c}} \in R^M$ to describe the loss of choosing each tweet as a candidate. $\tilde{\mathbf{A}} \in R^{M \times N}$ is a similarity matrix, where $\tilde{a}_{i,j}$ is the similarity between a credible and relevant tweet i and a candidate tweet j in the tweet universe. $\mathbf{b} \in R^M$ is a weight vector, where b_i indicates the importance of i being covered in the summary. Our objective is to

$$\min \tilde{\mathbf{c}}^T \tilde{\mathbf{x}} \text{ subject to } \tilde{\mathbf{A}}\tilde{\mathbf{x}} \geq \mathbf{b}, \tilde{\mathbf{x}} \in \{0, 1\}$$

3.2 Methodology Overview

We first transform it to a standard form of bounded linear programming problem by making the following adjustments: $\mathbf{c} = [\tilde{\mathbf{c}}, \mathbf{0}]$, $\mathbf{x} = [\tilde{\mathbf{x}}, \mathbf{z}]^T$, $\mathbf{A} = [\tilde{\mathbf{A}}, -\mathbf{I}]$, where \mathbf{I} is the $M \times M$ identity matrix. Therefore we have

$$\min \mathbf{c}\mathbf{x} \text{ subject to } \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}, \tilde{\mathbf{x}} \leq \mathbf{1} \quad (1)$$

We modify the LP so that there is an easy choice of basic solution. We start by solving

$$\min \mathbf{e}^T \mathbf{s} \text{ subject to } \mathbf{A}\mathbf{x} + \mathbf{I}\mathbf{s} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}, \tilde{\mathbf{x}} \leq \mathbf{1}, \mathbf{s} \geq \mathbf{0} \quad (2)$$

where \mathbf{e} is the vector of all ones, \mathbf{s} are called artificial variables, and $b_i \geq 0$. To solve the above LP relaxation, we adopt the bounded simplex method. In each iterate, only M variables are selected as the basic points. Suppose $\mathbf{A}_{\cdot j}$ is the j th column, and the corresponding columns of basic variables in \mathbf{A} are denoted as $\mathbf{A}_{\mathbf{B}}$, others are denoted as $\mathbf{A}_{\mathbf{N}}$. Non-basic variables are either at upper bound ($j \in U$) or at lower bound ($j \in L$).

Define $\mathbf{c}' = [\mathbf{e}, \mathbf{0}]$, $\mathbf{x}' = [\mathbf{x}, \mathbf{s}]^T$ and $\mathbf{A}' = [\mathbf{A}, \mathbf{I}]$ so that the constraints of the modified LP can be written as $\mathbf{A}'\mathbf{x}' = \mathbf{b}, \mathbf{x}' \geq \mathbf{0}, \tilde{\mathbf{x}} \leq \mathbf{1}$.

Step 1: Let \mathbf{B} be the indices of the artificial variables and set the all non-basic variables to be $x_j = 0$. Then \mathbf{B} is a basis, since the corresponding columns of $\mathbf{A}'_{\mathbf{B}}$ are \mathbf{I} , the identity, the corresponding basic feasible solution is $\mathbf{x} = \mathbf{0}, \mathbf{z} = \mathbf{b}$.

Step 2: Pricing. $\mathbf{y} = \mathbf{B}^{T-1} \mathbf{e}_{\mathbf{B}}$

Step 3: Compute $\bar{c}_j = c_j - \mathbf{y}^T \mathbf{A}'_{\cdot j}$. If $x_j = 0, \bar{c}_j > 0$ and $x_j = 1, \bar{c}_j < 0$, then the problem is solved. Else pick q the most contradictive variable, i.e. the most negative \bar{c}_q for $x_q = 0$ as the entering index (step 4 in Algorithm 1).

Step 4: Choose the outing index(step 5 to step 27 in Algorithm 1) and update the basis and the value of variables.

Step 5: If all artificial variables are non-basic or some artificial variables are in the basis but all $x_{z_i} = 0$ then go to Step 6,otherwise return to Step 2.

Step 6: If some artificial variables are in the basis,remove them from basis. Change $\mathbf{c}' = [\tilde{\mathbf{c}}, \mathbf{0}]$, repeat Step 2 to Step 4.

4 Update Summarization

$$\min c_0 x_0 + c_1 x_1 s.t. \begin{bmatrix} A & D \\ D^T & B \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \end{bmatrix} \geq \begin{bmatrix} b_0 \\ b_1 \end{bmatrix} \quad (3)$$

$$D^T x_0 \geq b_1$$

Input: $\mathbf{c} = [\mathbf{e}, 0]$, $\mathbf{x}' = [\mathbf{x}, \mathbf{s}]^T$ and $\mathbf{A}' = [\mathbf{A}, \mathbf{I}]$
Output: $\mathbf{x}'_{\mathbf{B}} = \mathbf{A}'_{\mathbf{B}}^{-1}(\mathbf{b} - \mathbf{u})$, where $\mathbf{u} = \sum_{j \in U} \mathbf{A}'_{\cdot,j}$

```

1 Pricing  $\mathbf{y} = \mathbf{B}^{T^{-1}} \mathbf{e}_{\mathbf{B}}$ ;
2 Compute  $\bar{c}_j = c_j - \mathbf{y}^T \mathbf{A}'_{\cdot,j}$ ;
3 while  $\exists x_j = 0, \bar{c}_j < 0$  or  $x_j = 1, \bar{c}_j < 0$  do
4    $q = \arg \max_j \{\bar{c}_j \forall j \in \{\mathbf{z}_{\mathbf{N}}, \mathbf{s}_{\mathbf{N}}, \mathbf{x}_{\mathbf{L}}\}, -\bar{c}_j \forall j \in \mathbf{x}_{\mathbf{U}}\}$ ;
5    $d = \mathbf{B}^{-1} \mathbf{A}'_{\cdot,q}$ ;
6   if  $q \in \{\mathbf{z}_{\mathbf{N}}, \mathbf{s}_{\mathbf{N}}\}$  then
7      $x_q^{new} = \min_i \left\{ \left\{ \frac{x_{\mathbf{B}_i}}{d_i}, \frac{s_{\mathbf{B}_i}}{d_i}, \frac{z_{\mathbf{B}_i}}{d_i} \right\} \forall d_i > 0, \left\{ \frac{x_{\mathbf{B}_i} - 1}{d_i} \right\} \forall d_i < 0 \right\}$ ;
8      $\mathbf{x}'_{\mathbf{B}^{old}} = \mathbf{x}'_{\mathbf{B}^{old}} - d x_q^{new}$ ;
9      $p = \arg \min_i$ ;
10     $\mathbf{B}^{new} \leftarrow \mathbf{B}^{old} - \{p\} \cup \{q\}$ ;
11  end
12  if  $q \in L$  then
13     $x_q^{new} = \min_i \left\{ \left\{ \frac{x_{\mathbf{B}_i}}{d_i}, \frac{s_{\mathbf{B}_i}}{d_i}, \frac{z_{\mathbf{B}_i}}{d_i} \right\} \forall d_i > 0, \left\{ \frac{x_{\mathbf{B}_i} - 1}{d_i} \right\} \forall d_i < 0, 1 \right\}$ ;
14     $\mathbf{x}'_{\mathbf{B}^{old}} = \mathbf{x}'_{\mathbf{B}^{old}} - d x_q^{new}$ ;
15    if  $x_q^{new} \neq 1$  then
16       $p = \arg \min_i$ ;
17       $\mathbf{B}^{new} \leftarrow \mathbf{B}^{old} - \{p\} \cup \{q\}$ ;
18    end
19  end
20  if  $q \in U$  then
21     $x_q^{new} = 1 - \min_i \left\{ \left\{ \frac{x_{\mathbf{B}_i}}{-d_i}, \frac{s_{\mathbf{B}_i}}{-d_i}, \frac{z_{\mathbf{B}_i}}{-d_i} \right\} \forall d_i < 0, \left\{ \frac{1 - x_{\mathbf{B}_i}}{d_i} \right\} \forall d_i > 0, 1 \right\}$ ;
22     $\mathbf{x}'_{\mathbf{B}^{old}} = \mathbf{x}'_{\mathbf{B}^{old}} + d(1 - x_q^{new})$ ;
23    if  $x_q^{new} \neq 0$  then
24       $p = \arg \min_i$ ;
25       $\mathbf{B}^{new} \leftarrow \mathbf{B}^{old} - \{p\} \cup \{q\}$ ;
26    end
27  end
28 end

```

Algorithm 1: the bounded simplex method