

## A. type check

### A.1 Syntax

$e, \tau ::=$		Expressions
	$x$	Variable
	$\star$	Type of Types
	$e_1 e_2$	Application
	$\lambda x. e$	Abstraction
	$\lambda x : \sigma. e$	Abstraction with annotation
	$\text{cast}^\uparrow e$	Cast Up
	$\text{cast}_\downarrow e$	Cast Down
	$\Pi(x : \tau). \tau_2$	Explicit Pi type
	$\text{let } x = e_1 \text{ in } e_2$	Let binding
	$e :: \sigma$	Annotation
$\sigma ::=$		Polytype
	$\forall x : \tau. \rho$	Implicit Pi Type
$\rho ::=$		Rhotype
	$\tau$	Mono
	$\Pi(x : \sigma_1). \sigma_2$	FunPoly

$$\begin{array}{c}
 \frac{\Gamma \vdash e : \sigma}{\Gamma \vdash (e :: \sigma) : \sigma} \text{ANN} \\
 \boxed{\Gamma \vdash \sigma : \star} \\
 \frac{\Gamma \vdash \tau : \star \quad \Gamma, x : \tau \vdash \rho : \star}{\Gamma \vdash (\forall x : \tau. \rho) : \star} \text{IMPLICITPI} \\
 \boxed{\Gamma \vdash \rho : \star} \\
 \frac{\Gamma \vdash \sigma_1 : \star \quad \Gamma, x : \sigma_1 \vdash \sigma_2 : \star}{\Gamma \vdash (\Pi x : \sigma_1. \sigma_2) : \star} \text{FUNPOLY}
 \end{array}$$

### A.2 non syntax-directed Typing

$$\boxed{\Gamma \vdash e : \sigma}$$

$$\Gamma \vdash \star : \star \text{ AX}$$

$$\frac{x : \sigma \in \Gamma}{\Gamma \vdash x : \sigma} \text{VAR}$$

$$\frac{\Gamma \vdash e_1 : (\Pi x : \sigma_1. \sigma_2) \quad \Gamma \vdash e_2 : \sigma_1}{\Gamma \vdash e_1 e_2 : \sigma_2[x \mapsto e_2]} \text{APP}$$

$$\frac{\Gamma, x : \tau \vdash e : \sigma}{\Gamma \vdash (\lambda x. e) : (\Pi x : \tau. \sigma)} \text{LAM}$$

$$\frac{\Gamma, x : \sigma \vdash e : \sigma_2}{\Gamma \vdash (\lambda x : \sigma. e) : (\Pi x : \sigma. \sigma_2)} \text{LAMANN}$$

$$\frac{\Gamma \vdash \tau_1 : \star \quad \Gamma, x : \tau_1 \vdash \tau_2 : \star}{\Gamma \vdash (\Pi x : \tau_1. \tau_2) : \star} \text{EXPLICITPI}$$

$$\frac{\Gamma \vdash e : \tau_2 \quad \Gamma \vdash \tau_1 : \star \quad \tau_1 \longrightarrow \tau_2}{\Gamma \vdash (\text{cast}^\uparrow e) : \tau_1} \text{CASTUP}$$

$$\frac{\Gamma \vdash e : \tau_1 \quad \Gamma \vdash \tau_2 : \star \quad \tau_1 \longrightarrow \tau_2}{\Gamma \vdash (\text{cast}_\downarrow e) : \tau_2} \text{CASTDOWN}$$

$$\frac{\Gamma \vdash e_1 : \sigma \quad \Gamma, x : \sigma \vdash e_2 : \tau}{\Gamma \vdash (\text{let } x = e_1 \text{ in } e_2) : \tau} \text{LET}$$

$$\frac{\Gamma \vdash e : \forall(\alpha_i : \tau_i). \sigma \quad \tau_{\beta i} : \tau_i}{\Gamma \vdash e : \sigma[\alpha_i \mapsto \tau_{\beta i}]} \text{INST}$$

$$\frac{\Gamma \vdash e : \sigma \quad \alpha_i \notin \text{free}(\Gamma) \quad \tau_i : \star \quad \alpha_i : \tau_i \vdash \sigma : \star}{\Gamma \vdash e : \forall(\alpha_i : \tau_i). \sigma} \text{GEN}$$