

## COMPUTATIONAL MATH, SCIENCE AND ENGINEERING DEPARTMENT

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More concurrency

## A long list

Because concurrency is such an important topic there are a number of different types (in C++ and elsewhere).

I'll list them here, but just briefly and you can investigate more on your own if you are interested.



## mutex variety

- timed\_mutex: lock to a timepoint
- recursive\_mutex: allow the **same thread** to recursively call itself
  - for every lock called, an unlock must be called before the mutex is released
  - only the same thread can make the multiple calls. Other threads wait
- recursive\_timed\_mutex: combination of the two above



## more

- `shared_mutex`: two ways to lock:
  - exclusive (using `lock`, `try_lock`). Only 1 thread.
  - shared (using `lock_shared`, `try_lock_shared`). More than 1 thread can use it.
- `shared_timed_mutex`
  - combo of shared and timed.



# semaphores

Semaphore (c++-20) has a count of the number of threads that can share access:

- acquire grants access and decrements the count. When the count hits 0, acquire blocks
- release ups the count, unblocks.
- not tied to one thread. The acquire thread can be different from release.



# lock variety

- `unique_lock`: more general purpose  
`lock_guard`
  - can be constructed without locking
  - can be unlocked
- `scoped_lock`: a `lock_guard` to grab multiple mutexs before moving forward



# barrier

- C++-20, requires g++/clang++ 11 or greater
  - has a count such that all threads reaching the barrier wait until the count reaches 0.



# condition variable

- is a way to have a thread wait and to notify other threads when processing can proceed
  - wait is on a condition (a Boolean function)
  - can notify one or many threads
  - think reader/writer or produce/consumer as natural fits.





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# Deadlock

# Deadlock?

When two (or more) processes are stuck waiting for access to a resource held by a different process.

If true, we have reached an impasse. No deadlocked process can continue unless another process releases a resource.



# Dining Philosophers



- you need two forks to eat
- you can only grab a fork to your direct left or right
- you need to alternate eating and thinking. When thinking (not eating) you must put down both forks

# Four Conditions for Deadlock

- Coffman et. al. 1971
- Necessary conditions for deadlock to exist:
  - **Mutual Exclusion**
    - At least one resource must be held in non-sharable mode
  - **Hold and wait**
    - There exists a process holding a resource, and waiting for another
  - **No preemption**
    - Resources cannot be preempted
  - **Circular wait**
    - There exists a set of processes  $\{P_1, P_2, \dots, P_N\}$ , such that
      - $P_1$  is waiting for  $P_2$ ,  $P_2$  for  $P_3$ , ..., and  $P_N$  for  $P_1$

*All* four conditions must hold for deadlock to occur



# Common causes

- self deadlock
- lack of mutex ordering
- starvation



# selflock example

- how to fix this?



# order example

- how to fix?



# barrier/starvation

- how to fix





Operation	Effect
<i>mutex</i> <i>m</i>	Default constructor; creates an unlocked mutex
<i>m</i> .~ <i>mutex</i> ()	Destroys the mutex (must not be locked)
<i>m</i> .lock()	Locks the mutex (blocks for lock; error if locked and not recursive)
<i>m</i> .try_lock()	Tries to lock the mutex (returns true if lock successful)
<i>m</i> .try_lock_for( <i>dur</i> )	Tries to lock for duration <i>dur</i> (returns true if lock successful)
<i>m</i> .try_lock_until( <i>tp</i> )	Tries to lock until timepoint <i>tp</i> (returns true if lock successful)
<i>m</i> .unlock()	Unlocks the mutex (undefined behavior if not locked)
<i>m</i> .native_handle()	Returns a platform-specific type <code>native_handle_type</code> for nonportable extensions

*Table 18.7. Operations of Mutex Classes, If Available*

Operation	Effect
<i>lock_guard</i> <i>lg</i> ( <i>m</i> )	Creates a lock guard for the mutex <i>m</i> and locks it
<i>lock_guard</i> <i>lg</i> ( <i>m</i> , adopt_lock)	Creates a lock guard for the already locked mutex <i>m</i>
<i>lg</i> .~ <i>lock_guard</i> ()	Unlocks the mutex and destroys the lock guard

Table 18.8. Operations of Class *lock\_guard*