Adapted from Ex 8.4.2. In this problem, we consider indexes for the relation

   Ships(name, class, launched)

from our running battleships exercise. Assume:

  i. `name` is the key.

  ii. The relation `Ships` is stored over 50 pages.

  iii. The relation is clustered on `class` so we expect that only one disk access is needed to find the ships of a given class.

  iv. On average, there are 5 ships of a class, and 25 ships launched in any given year.

  v. With probability p_1 the operation on this relation is a query of the form

      SELECT * FROM Ships WHERE name $= n$.

  vi. With probability p_2 the operation on this relation is a query of the form

      SELECT * FROM Ships WHERE class $= c$.

  vii. With probability p_3 the operation on this relation is a query of the form

      SELECT * FROM Ships WHERE launched $= y$.

  viii. With probability $1 - p_1 - p_2 - p_3$ the operation on this relation is an insertion of a new tuple into Ships.

You can also make assumptions about accessing indexes and finding empty space for insertions that were made in Example 8.14.

(a) If you can only create one index, how would you decide what index to create ?

(b) If you can create any number of indexes, what are the possible index combinations ?

(c) Consider the creation of indexes on `name`, `class`, and `launched`. For each combination of indexes, estimate the average cost of an operation. As a function of p_1, p_2, and p_3, what is the best choice of indexes ?