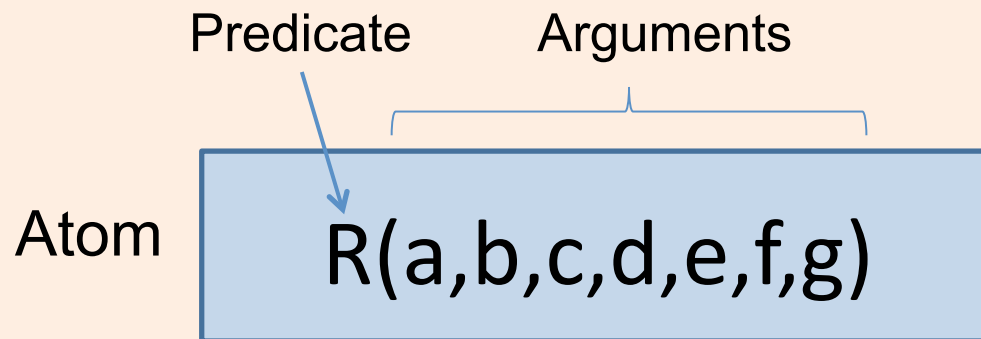


# ICS 321 Data Storage & Retrieval

## Algebraic and Logical Query Languages (ii)

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# Datalog : Database Logic



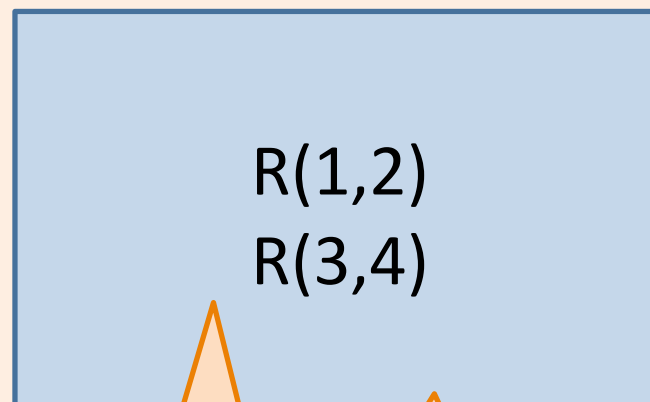
- A (relational) **atom**
  - Consists of a **predicate** and a list of **arguments**
  - Arguments can be **constants** or **variables**
  - Takes on Boolean value (true or false)
- A relation  $R$  can be represented as a predicate  $R$ 
  - A tuple  $\langle a,b,c,d,e,f,g \rangle$  is in  $R$  iff the atom  $R(a,b,c,d,e,f,g)$  is true.

# Example: tables in datalog

**R**

A	B
1	2
3	4

**Datalog**



True by default.

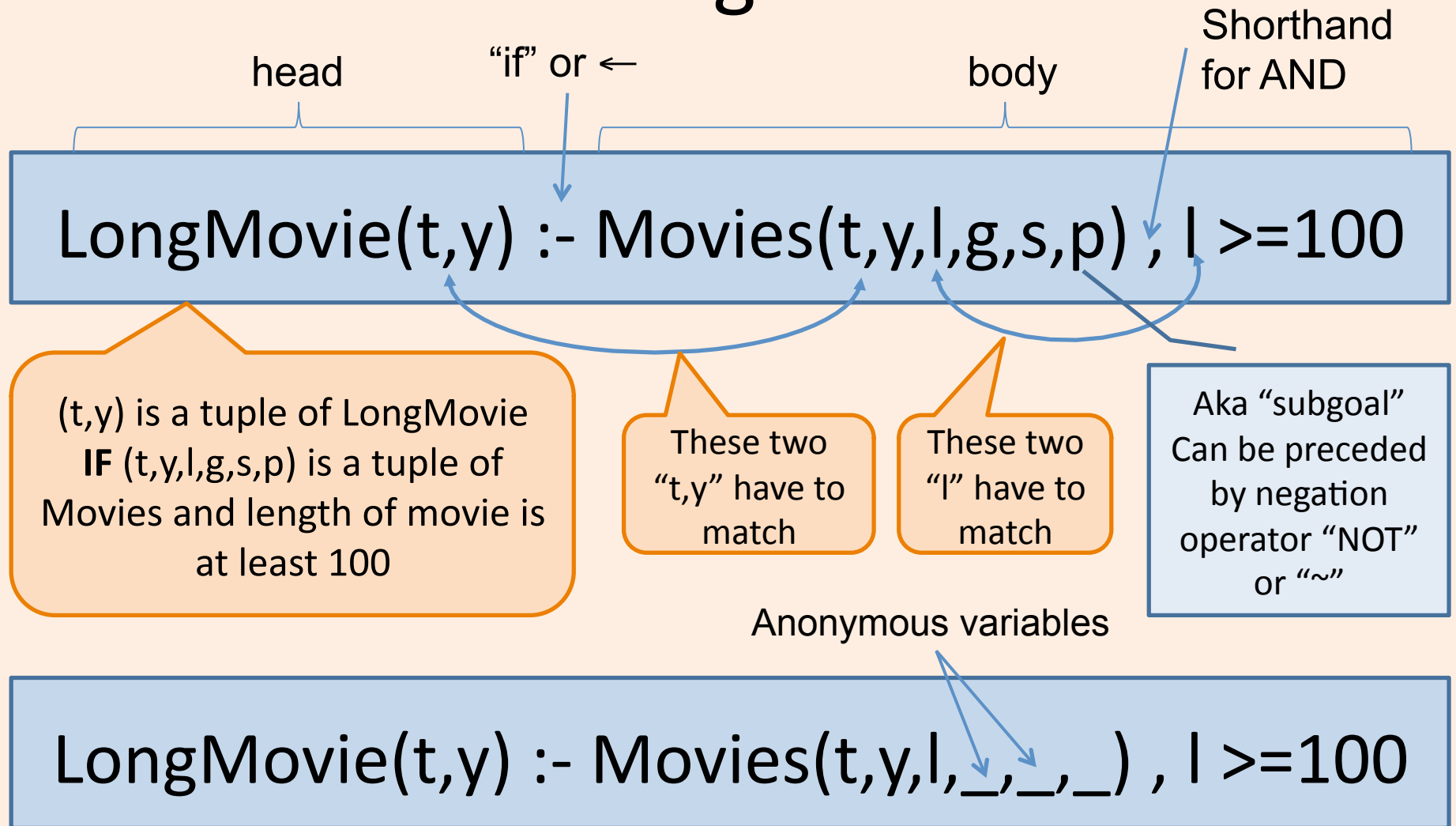
R(1,4) would  
be false

# Arithmetic Atoms

$$x < y$$
$$x+1 \geq y+4*z$$

Can contain  
both constants  
and variables.

# Datalog Rules



# Safety Condition for Datalog Rules

Every **variable** that appears anywhere in the rule **must** appear in some **nonnegated, relational subgoal** of the body

- Without the safety condition, rules may be underspecified, resulting in an infinite relation (not allowed).
- Examples
  - $\text{LongMovie}(t,y) \text{ :- Movies}(t,y,l,\_,\_,\_) , l \geq 100$
  - $P(x,y) \text{ :- } Q(x,z), \text{ NOT } R(w,x,z), x < y$

# Alternative Interpretation: Consistency

## Datalog

Q(1,2)  
Q(1,3)  
R(2,3)  
R(3,1)  
 $P(x,y) :- Q(x,z), R(z,y), \text{NOT } Q(x,y)$

- For each consistent assignment of nonnegated, relational subgoal,
- Check the negated, relational subgoals and the arithmetic subgoals for consistency

Q(x,z)	R(z,y)	Consistent?	NOT Q(x,y)	Head
(1,2)	(2,3)	Yes	false	
(1,2)	(3,1)	No, z=2,3		
(1,3)	(2,3)	No, z=2,3		
(1,3)	(3,1)	Yes	true	P(1,1)

# Intensional vs Extensional

## Datalog

Q(1,2)

Q(1,3)

R(2,3)

R(3,1)

P(x,y) :- Q(x,z), R(z,y), NOT Q(x,y)

extensional

intensional

- **Extensional** predicates – relations stored in a database
- **Intensional** predicates – computed by applying one or more datalog rules



# What about bag semantics ?

- Datalog still works if there are no negated, relational subgoals.
- Treat duplicates like non-duplicates

## Datalog

R(1,2)  
R(1,2)  
S(2,3)  
S(4,5)  
S(4,5)  
H(x,z) :- R(x,y), S(y,z)

R(x,y)	S(y,z)	Consistent?	Head
(1,2)	(2,3)	Yes	H(1,3)
(1,2)	(4,5)	No, y=2,4	
(1,2)	(4,5)	No, y=2,4	
...	...	...	...

# Example 1

## Datalog

Answer(x,y) :- A(x,y)

Answer(x,y) :- B(x,y)

# Example 2

## Datalog

```
Answer(x,y) :- A(x,y), B(x,y)
```

# Example 3

## Datalog

```
Answer(x,y) :- A(x,y), NOT B(x,y)
```

# Example 4

## Datalog

```
Answer(x,y) :- A(x,y), x > 10, y = 200
```

# Example 5

## Datalog

```
Answer(x) :- A(x,y)
```

# Example 6

## Datalog

$\text{Answer}(w,x,y,z) \text{ :- } A(w,x), B(y,z)$

# Example 7

## Datalog

$\text{Answer}(w,x,y) \text{ :- } A(w,x), B(x,y)$



# Example 8

## Datalog

$\text{Answer}(w,x,z) \text{ :- } A(w,x), B(y,z), x > y$

# Example 9

## Datalog

```
Path(x,y) :- Edge(x,y)  
Path(x,z) :- Edge(x,y), Edge(y,z)
```