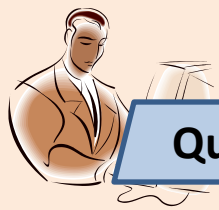


ICS 321 Fall 2013

Overview of Query Processing

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Query

Parse Query

Enumerate
Plans

Estimate
Cost

Choose
Best Plan

Evaluate
Query Plan

Result

SELECT * FROM Reserves WHERE sid=101

$\sigma_{\text{Sid}=101}$

Reserves

B

fetch

Reserves

IDXSCAN
(sid=101)

Index(sid)

SCAN (sid=101)

Reserves

32.0

25.0

Pick B

Optimizer



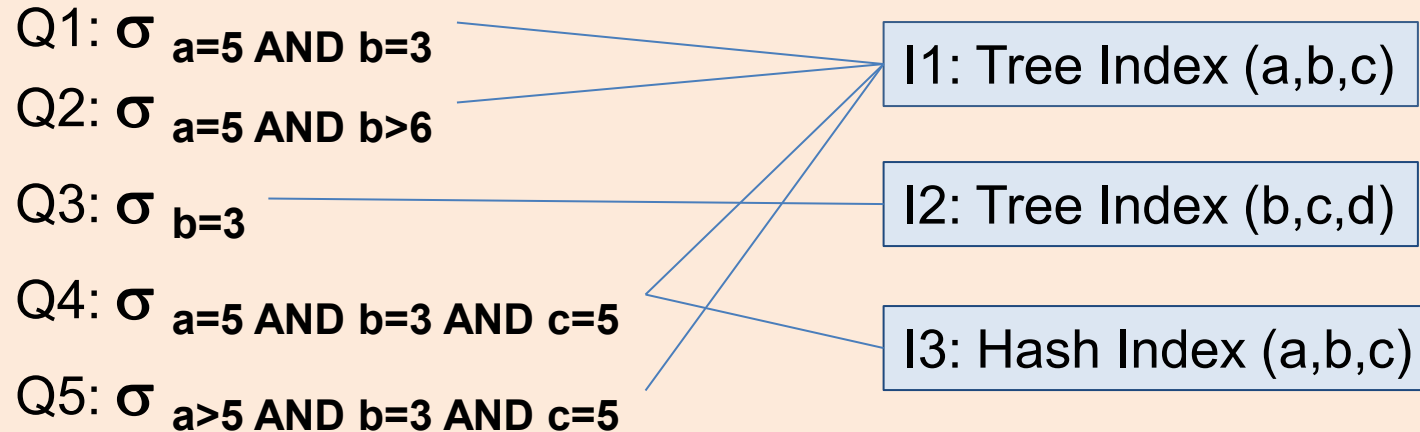
Query Processing

- **Query Execution Plan (QEP):** tree of database operators.
 - At high-level, relational algebra operators are used
 - At low-level, RA operators with particular implementation algorithm.
- **Plan enumeration:** find equivalent plans
 - Different QEPs that return the same results
 - Query rewriting : transformation of one QEP to another equivalent QEP.
- **Cost estimation:** a mapping of a QEP to a cost
 - **Cost Model:** a model of what counts in the cost estimate. Eg. Disk accesses, CPU cost ...
- **Query Optimizer:**
 - Explores the space of equivalent plan for a query
 - Chooses the best plan according to a cost model

Access Paths

- An **access path** is a method of retrieving tuples.
Eg. Given a query with a selection condition:
 - File or table scan
 - Index scan
- **Index matching problem:** given a selection condition, which indexes can be used for the selection, i.e., matches the selection ?
 - Selection condition normalized to conjunctive normal form (CNF), where each term is a *conjunct*
 - Eg. (day<8/9/94 **AND** rname='Paul') **OR** bid=5 **OR** sid=3
 - **CNF:** (day<8/9/94 **OR** bid=5 **OR** sid=3) **AND** (rname='Paul' **OR** bid=5 **OR** sid=3)

Index Matching



- A tree index matches a selection condition if the selection condition is a prefix of the index search key.
- A hash index matches a selection condition if the selection condition has a term *attribute=value* for every attribute in the index search key

One Approach to Selections

1. Find the *most selective access path*, retrieve tuples using it
2. Apply remaining terms in selection not matched by the chosen access path


- The **selectivity** of an access path is the size of the result set (in terms of tuples or pages).
 - Sometimes selectivity is also used to mean **reduction factor**: fraction of tuples in a table retrieved by the access path or selection condition.
- Eg. Consider the selection:
 $\text{day} < 8/9/94$ **AND** $\text{bid} = 5$ **AND** $\text{sid} = 3$
 - Tree Index(day)
 - Hash index (bid,sid)

Query Execution Plans

- A tree of database operators: each operator is a RA operator with specific implementation
- Selection σ : Index Scan or Table Scan
- Projection π :
 - Without DISTINCT : Table Scan
 - With DISTINCT : requires sorting or index scan
- Join \bowtie :
 - Nested loop joins (naïve)
 - Index nested loop joins
 - Sort merge joins

Nested Loop Join

S1	<u>sid</u>	sname	rating	age
	22	Dustin	7	45.0
	31	Lubber	8	55.5
	58	Rusty	10	35.0

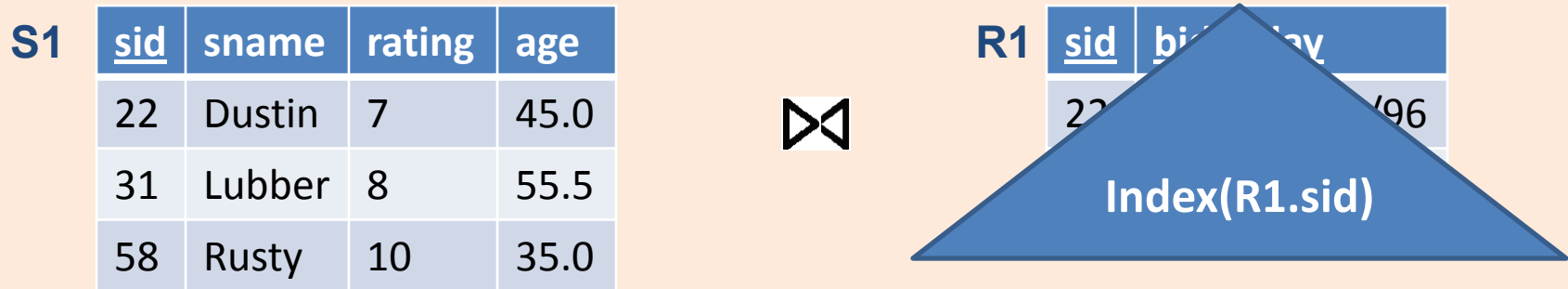


R1	<u>sid</u>	<u>bid</u>	<u>day</u>
	22	101	10/10/96
	58	103	11/12/96

For each data page P_{S1} of S1
 For each tuple s in P_{S1}
 For each data page P_{R1} of R1
 For each tuple r in P_{R1}
 if ($s.sid == r.sid$)
 then output s, r

- Worst case number of disk reads
 = $Npages(S1) + |S1| * Npages(R1)$

Index Nested Loop Join



For each data page P_{S1} of S1

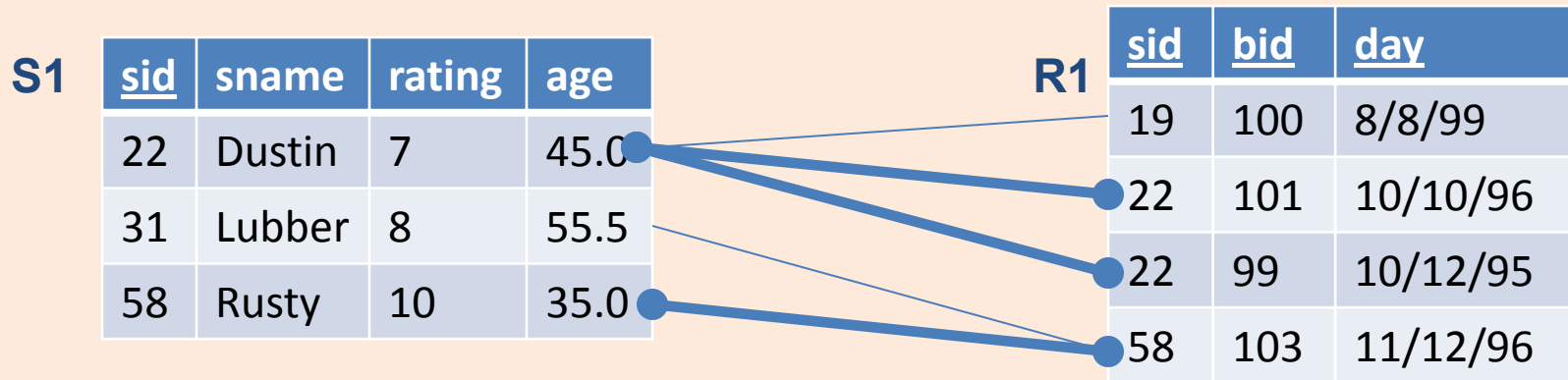
For each tuple s in P_{S1}

if ($s.sid \in \text{Index}(R1.sid)$)

then fetch r & output $\langle s, r \rangle$

- Worst case number of disk reads with tree index
 $= \text{Npages}(S1) + |S1| * (1 + \log_f \text{Npages}(R1))$
- Worst case number of disk reads with hash index
 $= \text{Npages}(S1) + |S1| * 2$

Sort Merge Join



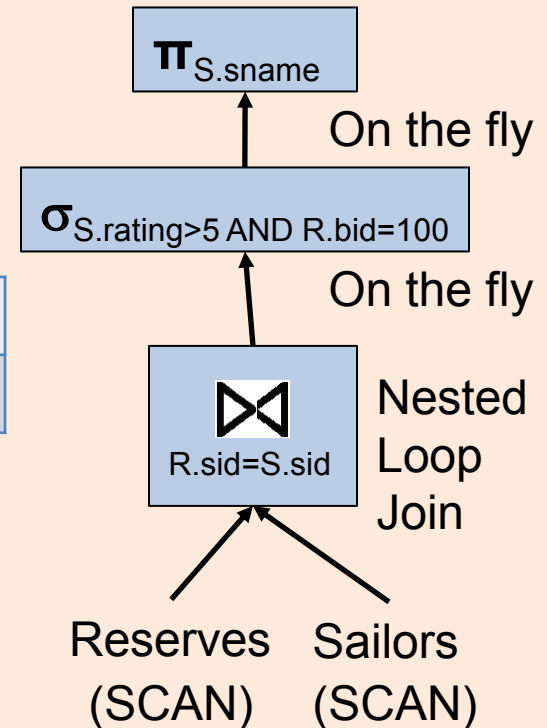
1. Sort S1 on SID
 2. Sort R1 on SID
 3. Compute join on SID using Merging algorithm
- If join attributes are relatively unique, the number of disk pages
= $N_{\text{pages}}(S1) \log N_{\text{pages}}(S1)$
+ $N_{\text{pages}}(R1) \log N_{\text{pages}}(R1)$
+ $N_{\text{pages}}(S1) + N_{\text{pages}}(R1)$
 - What if the number of duplicates is large?
 - the number of disk pages approaches that of nested loop join.

Example

```
SELECT S.sname
FROM Reserves R, Sailors S
WHERE R.sid=S.sid AND R.bid=100 AND S.rating>5
```

Reserves	40 bytes/tuple	100 tuples/page	1000 pages
Sailors	50 bytes/tuple	80 tuples/page	500 pages

- Nested Loop Join cost 1K+ 100K*500
- On the fly selection and project does not incur any disk access.
- Total disk access = 500001K (worst case)



Example: Predicate Pushdown

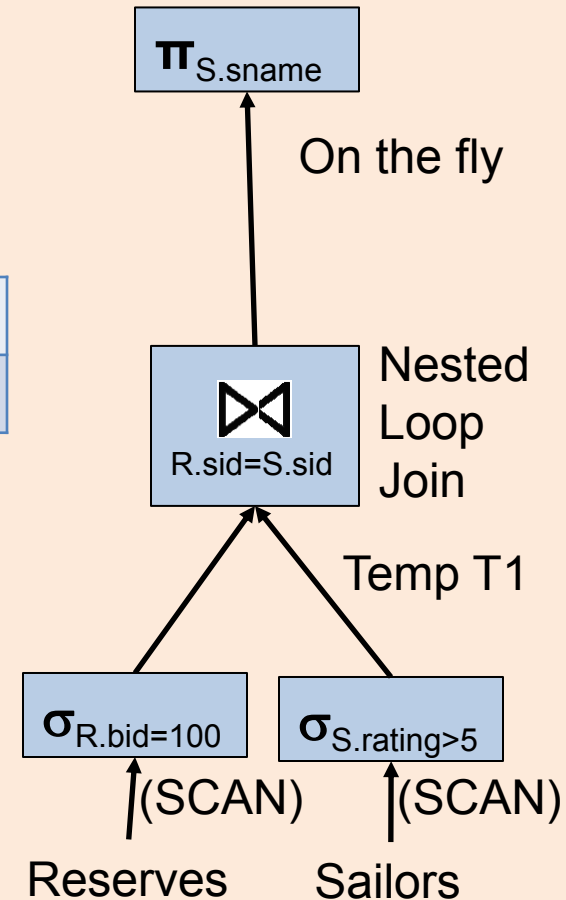
SELECT S.sname
FROM Reserves R, Sailors S
WHERE R.sid=S.sid **AND** R.bid=100 **AND** S.rating>5

10%

50%

Reserves	40 bytes/tuple	100 tuples/page	1000 pages
Sailors	50 bytes/tuple	80 tuples/page	500 pages

- Nested Loop Join requires materializing the inner table as T1.
- With 50% selectivity, T1 has 250 pages
- With 10% selectivity, outer “table” in join has 10K tuples
- Disk accesses for scans = 1000 + 500
- Writing T1 = 250
- NLJoin = 10K * 250
- Total disk access = 2500.175 K (worst case)



What happens if we make the left leg the inner table of the join ?

Example: Sort Merge Join

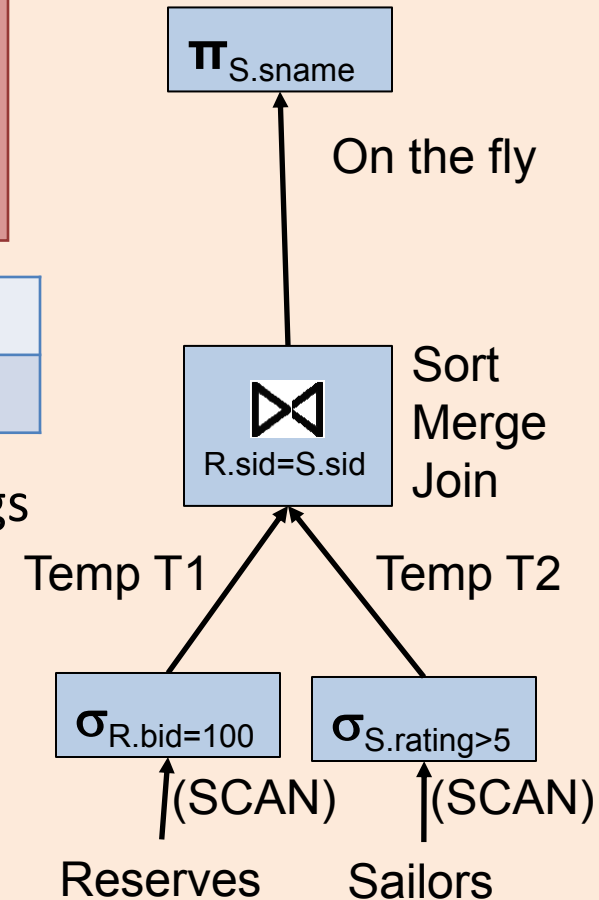
SELECT S.sname
FROM Reserves R, Sailors S
WHERE R.sid=S.sid **AND** R.bid=100 **AND** S.rating>5

10%

50%

Reserves	40 bytes/tuple	100 tuples/page	1000 pages
Sailors	50 bytes/tuple	80 tuples/page	500 pages

- Sort Merge Join requires materializing both legs for sorting.
- With 10% selectivity, T1 has 100 pages
- With 50% selectivity, T2 has 250 pages
- Disk accesses for scans = 1000 + 500
- Writing T1 & T2 = 100 + 250
- Sort Merge Join = $100 \log 100 + 250 \log 250 + 100 + 250$ (assume 10 way merge sort)
- Total disk access = 52.8 K



What happens if we make the left leg the inner table of the join ?

Example: Index Nested Loop Join

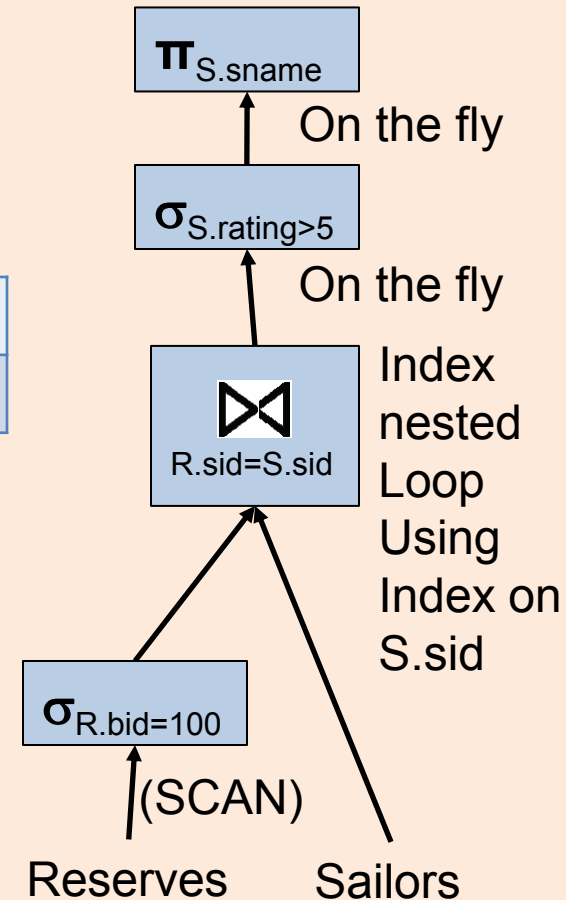
SELECT S.sname
FROM Reserves R, Sailors S
WHERE R.sid=S.sid **AND** R.bid=100 **AND** S.rating>5

10%

50%

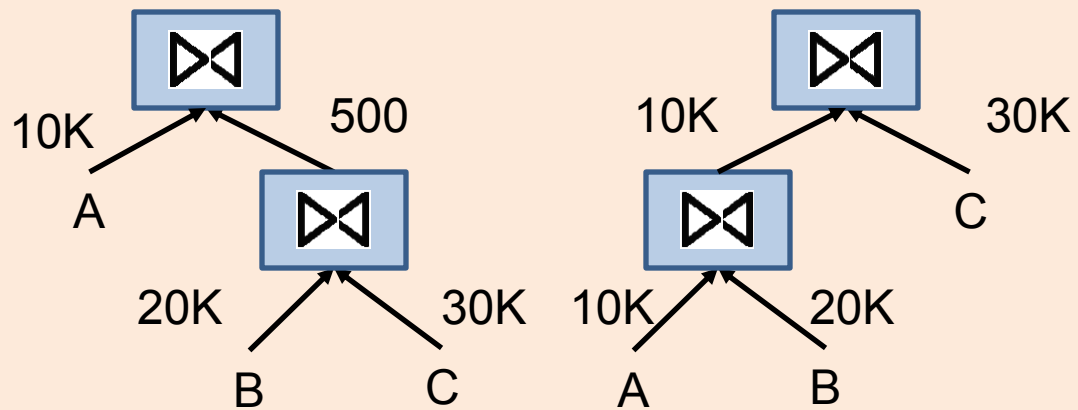
Reserves	40 bytes/tuple	100 tuples/page	1000 pages
Sailors	50 bytes/tuple	80 tuples/page	500 pages

- With 10% selectivity, selection on R has 10K tuples
- Disk accesses for scan = 1000
- Index Nested Loop Join = $10K * (1 + \log_{10} 500) = 37K$
- Total disk access = 38 K



What happens if we make the left leg the inner table of the join ?

Join Ordering



Relations	Tuples	Pages
A	10K	1000
B	20K	2000
C	30K	3000
A join B	10K	1000
B join C	1K	100

- Independent of what join algorithm is chosen, the order in which joins are performed affects the performance.
- Rule of thumb: do the most “selective” join first
- In practice, left deep trees (eg. the right one above) are preferred --- why ?

Statistics & Cost Estimation

- Page size
- Data Statistics:
 - Record size -> number of records per data page
 - Cardinality of relations (including temporary tables)
 - Selectivity of selection operator on different columns of a relation
- (Tree) Index Statistics
 - number of leaf pages, index entries
 - Height
- Statistics collection is user triggered
 - DB2: RUNSTATS ON TABLE mytable AND INDEXES ALL