L35: Sorting and searching And linked lists

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Version: release

Announcements

Lab 12 is due on April 21 at 12 pm (noon)

Today:

Sorting and searching

Next week:

Special topics

Searching through a large array is MUCH faster if the array is sorted

Searching through a large array is MUCH faster if the array is sorted

Unsorted array

- Start at beginning of list
- ► Look through the list one element at a time until you find a match or reach the end of the list
- On average, look at half of the list
- \triangleright Complexity: $\mathcal{O}(n)$

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Searching through a large array is MUCH faster if the array is sorted

Unsorted array

- Start at beginning of list
- Look through the list one element at a time until you find a match or reach the end of the list
- On average, look at half of the list
- ▶ Complexity: $\mathcal{O}(n)$

Sorted array

- ► Start at the middle of list
- ► Bisection method: Eliminate half of the list (all elements either before or after the current element) at each step
 - ightarrow Divide and Conquer
- ► At each step, divide the size of the problem by 2
- ▶ Complexity: $\mathcal{O}(\log_2(n))$

n: size of the list to search. 3/15

Average number of comparisons needed to search unsorted and sorted versions of the same list:

	n/2	$\log_2(n)$
Search E7 roaster for an SID	≈ 210	pprox 9
Search DMV records for a California driver's license or ID number ¹	pprox 16.5 millions	≈ 25
Search IRS records for a U.S. 2016 individual tax return ²	pprox 75 millions	≈ 27

n: size of the list to search

^{1:} based on numbers from State of California DEPARTMENT OF MOTOR VEHICLES STATISTICS FOR PUBLICATION JANUARY THROUGH DECEMBER 2015, https://www.dmv.ca.gov/portal/dmv/detail/about/dmvinfo, retrieved April 19th 2017

² based on numbers from https://www.irs.gov/uac/soi-tax-stats-tax-stats-at-a-glance, retrieved April 19th 2017.

Sorting struct arrays

```
>> scientists
scientists =
  1×3 struct array with fields:
    ln
    dob
    pob
>> scientists(1)
ans =
  struct with fields:
     fn: 'Pierre'
     In: 'Currie'
    dob: 1859
    pob: 'Paris'
>> scientists(2)
ans =
  struct with fields:
     fn: 'Harie'
     In: 'Currie'
    dob: 1867
    pob: 'Warsaw'
>> scientists(3)
ans =
  struct with fields:
     fn: 'Isaac'
     In: 'Newton'
    dob: 1643
    pob: 'Woolsthorpe-by-Colsterworth'
```

```
>> [~, indices] = sort([scientists(:).dob]);
>> for i = indices
       scientists(i)
   end
ans =
  struct with fields:
     fn: 'Isaac'
     In: 'Newton'
    dob: 1643
    pob: 'Woolsthorpe-by-Colsterworth'
ans =
  struct with fields:
     fn: 'Pierre'
     ln: 'Currie'
    dob: 1859
    pob: 'Paris'
ans =
  struct with fields:
     fn: 'Marie'
     In: 'Currie'
    dob: 1867
    pob: 'Warsaw'
>>
```

Sort vector of dates of birth (dob), which are doubles

Sorting struct arrays

```
>> scientists
scientists =
  1×3 struct array with fields:
    fn
    ln
    dob
    pob
>> scientists(1)
ans =
  struct with fields:
     fn: 'Pierre'
     In: 'Currie'
    dob: 1859
    pob: 'Paris'
>> scientists(2)
ans =
  struct with fields:
     fn: 'Harie'
     In: 'Currie'
    dob: 1867
    pob: 'Warsaw'
>> scientists(3)
ans =
  struct with fields:
     fn: 'Isaac'
     In: 'Newton'
    dob: 1643
    pob: 'Woolsthorpe-by-Colsterworth'
```

```
>> [~, indices] = sort({scientists(:).pob});
>> for i = indices
       scientists(i)
   end
ans =
  struct with fields:
     fn: 'Pierre'
     In: 'Currie'
    dob: 1859
    pob: 'Paris'
ans =
  struct with fields:
     fn: 'Marie'
     ln: 'Currie'
    dob: 1867
    pob: 'Warsaw'
ans =
  struct with fields:
     fn: 'Isaac'
     ln: 'Newton'
    dob: 1643
    pob: 'Woolsthorpe-by-Colsterworth'
>>
```

Sort cell array of places of birth (pob), which are character strings

Sorting struct arrays

```
>> scientists
scientists =
  1×3 struct array with fields:
    fn
    ln
    dob
    pob
>> scientists(1)
ans =
  struct with fields:
     fn: 'Pierre'
     In: 'Currie'
    dob: 1859
    pob: 'Paris'
>> scientists(2)
ans =
  struct with fields:
     fn: 'Harie'
     In: 'Currie'
    dob: 1867
    pob: 'Warsaw'
>> scientists(3)
ans =
  struct with fields:
     fn: 'Isaac'
     In: 'Newton'
    dob: 1643
    pob: 'Woolsthorpe-by-Colsterworth'
```

```
>> names = transpose({scientists(:).ln; scientists(:).fn})
 3×2 cell array
    'Currie'
                 'Pierre'
    'Currie'
                 'Harie'
    'Newton'
                'Isaac'
>> [~, indices] = sortrows(names, [1, 2]);
>> for i = transpose(indices)
       scientists(i)
   end
ans =
 struct with fields:
     fn: 'Harie'
     In: 'Currie'
    dob: 1867
    pob: 'Warsaw'
 struct with fields:
     fn: 'Pierre'
     ln: 'Currie'
    doh: 1859
    pob: 'Paris'
ans =
 struct with fields:
     fn: 'Isaac'
     In: 'Newton'
    dob: 1643
    pob: 'Woolsthorpe-by-Colsterworth'
```

Use sortrows on the cell array of last and first names

Searching sorted vectors

Approach 1: Use Matlab's built-in find function e.g.,

```
>> rng(0); a = sort(randi([1, 1e5], [1, 1e5]));
>> indices = find(a==102)
indices =
    96    97    98    99
>> indices = find(a==104)
indices =
    1x0 empty double row vector
```

Searching sorted vectors

Approach 1: Use Matlab's built-in find function *e.g.*,

```
>> rng(0); a = sort(randi([1, 1e5], [1, 1e5]));
>> indices = find(a==102)
indices =
   96   97   98   99
>> indices = find(a==104)
indices =
   1x0 empty double row vector
```

Approach 2: Use user-defined function (see my_find_in_sorted.m) *e.g.*,

```
>> rng(0); a = sort(randi([1, 1e5], [1, 1e5]));
>> [~, indices] = my_find_in_sorted(a, 102)
indices =
    96    97    98    99
>> [~, indices] = my_find_in_sorted(a, 104)
indices =
    []
```

Note: this approach uses an algorithm similar to bisection

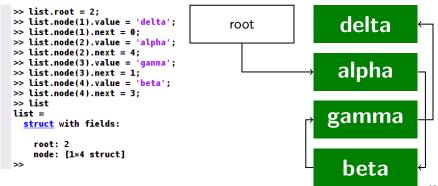
Searching sorted vectors: efficiency

```
>> rng(0)
>> a = sort(randi([1, 1e5], [1, 1e5]));
>> a(93:102)
ans =
   100 101 101 102 102 102 102 103 105
                                                          105
>> % Without using the fact that "a" is sorted
>> % Matlab looks at all of the elements in the array!
>> % Complexity: O(n)
>> tic(); for i = 1:1000; find(a==102); end; toc()
Elapsed time is 0.220739 seconds.
>> tic(); for i = 1:1000; find(a==104); end; toc()
Elapsed time is 0.167810 seconds.
>> % Using the fact that "a" is sorted. Complexity: O(log 2(n))
\Rightarrow tic(); for i = 1:1000; my find in sorted(a, 102); end; toc()
Elapsed time is 0.004417 seconds.
\Rightarrow tic(); for i = 1:1000; my find in sorted(a, 104); end; toc()
Elapsed time is 0.001637 seconds.
```

A linked list is a data structure with a root element and nodes

- ▶ The root is a reference to the first node in the list
- Nodes contain data and a reference to the next node in the list

The logical order of the data (as defined by the references in the nodes) may be different from the physical order of the data (here: the order of the nodes in the **struct** array)



Traversing a linked list

To traverse a linked list, we do not look at the elements in the order in which they appear in the struct array. Rather, we follow the order from one element to the next, as indicated by the reference "next" of each node. Here, we use a zero to indicate that there is no "next" element (other choices are possible *e.g.*, -1)

```
function [output] = my ll traverse(linked list)
% Traverses the linked list, and gathers the corresponding
% values (in the logical order of the list) in the cell
% array "output".
output = {};
current = linked list.root;
while current ~= 0
    output{end+1} = linked list.node(current).value;
    current = linked list.node(current).next;
end
end
```

Traversing a linked list

```
>> list.root = 2;
>> list.node = struct('value', 'delta', 'next', 0);
>> list.node(2) = struct('value', 'alpha', 'next', 4);
>> list.node(3) = struct('value', 'gamma', 'next', 1);
>> list.node(4) = struct('value', 'beta', 'next', 3);
>> % Physical order of the data
>> {list.node(:).value}
ans =
 1x4 cell array
   'delta' 'alpha' 'gamma' 'beta'
>> % Logical order of the data
>> values = my ll traverse(list)
values =
 1x4 cell array
   'alpha' 'beta' 'gamma' 'delta'
```

Deleting nodes from a linked list

We can "delete" nodes without physically removing them from memory (watch out for memory leaks!), by re-setting the "next" values of the appropriate nodes to skip the "deleted" nodes

```
function [linked_list] = my_ll_delete(linked_list, value)
% Delete all nodes whose value is "value" from the linked list.
previous = 0:
current = linked_list.root:
while current "= 0
    next = linked_list.node(current).next;
    if strcmp(linked_list.node(current).value, value)
        linked_list.node(current).value = '';
        linked_list.node(current).next = 0:
        if previous = 0
            % We just deleted the first node
            linked_list.root = next:
        else
            linked_list.node(previous).next = next;
        end
    else
        previous = current;
    end
    current = next:
end
end
```

Deleting nodes from a linked list

```
>> list.root = 2;
>> list.node = struct('value', 'delta', 'next', 0);
>> list.node(2) = struct('value', 'alpha', 'next', 4);
>> list.node(3) = struct('value', 'gamma', 'next', 1);
>> list.node(4) = struct('value', 'beta', 'next', 3);
>> [list] = my ll delete(list, 'gamma');
>> % Physical order of the data
>> {list.node(:).value}
ans =
 1x4 cell array
   'delta' 'alpha' '' 'beta'
>> % Logical order of the data
>> values = my ll traverse(list)
values =
 1x3 cell array
   'alpha' 'beta' 'delta'
```

Advantages:

- We can sort, insert, and remove elements from a linked list without having to move data around in memory (as opposed to, for example, inserting an element in the middle of a large vector of class double)
- ► Thus, it makes it easier to keep the list sorted as nodes are added and/or removed

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 Need to come up with alternate ways to search the list; we cannot easily apply the bisection method used earlier for sorted vectors of class double

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Disadvantages:

Need to come up with alternate ways to search the list; we cannot easily apply the bisection method used earlier for sorted vectors of class double

Examples of additional functions for maintaining linked lists:

- Add an element to the list (in correct order)
- Compress linked list to get rid of unused nodes