INTRODUCTION TO NETWORK STORAGE -

| (1) | TBOD, DAS, NAS, SAN & CAS evolution and companion |
|----------|--|
| | Refer Unit I Evolution of various storages te mologies . CAS is |
| | the lastest technology after IP SAN. CAS is discussed in further retion |
| | in this unit. CAS - Content-Adherable Storage. |
| ė | |
| | Din-1-Stongs |
| | Driet-Attached Storage (DAS) - |
| | It is an auchitecture where storage connects directly to surver. |
| | Applications occurs date from DAS using block-burl access protocols. |
| - | Applications of DAS - Internal HDD of a hort, take libraries and directly |
| | connected enternal HDD packs |
| - , → | Phynial Elements of DAS - |
| | |
| | CPU Connectionity STORAGE |
| | Motherboard Internal Hard dish (s) CD-RON drive, |
| | Cluster group of processors Enternal Optical drive, Removal media |
| y .= . | Processor cards Tape library, RAID |
| | Complète system Intelligent anay(s), Portable medie denies |
| → | Types of DAS- |
| | (1) Internal DAS - In this architecture, the storage deince is internating |
| | connected to the host by a seval or parallel his. |
| | (2) Enternal DAS - In this architecture, the sever cornects duitty to the |
| | Enternal storage devices |
| <u>_</u> | Connecturity - |
| , | Thtunas → serial or farallel his , Entimal → whis having SCSI or FC fustoral |
| → | DAS Management - |
| | Internal - |
| | · Most provides din partitioning (volume management), file rystem dayout |
| | and data addressing. |
| | DAS managed individually through the server and the OS. |
| | |
| | Mycompanion |

| | SAS - Ferral Attached SCSI |
|---------------|---|
| | Enternal - |
| <u> </u> | Avay based management, Availability - multi-fath I/O |
| | how Total Cost of Ownership (TCO) for managing date and storage infrastrution |
| \rightarrow | Security - |
| | (i) harge role DAS not very secure because of the distributed nature of the sew |
| | (2) DAS hosted on Windows sewer can be made seeme by uning group policies |
| - | (3) DAS does not provide security that is amounted with a SAN or NAS configure |
| - | Standard - (Other one IDE/ATA, SATA, SAS) |
| | (1) SCSI -> Small Computer hystem Interface. |
| | (2) FC - File Connectivity (usually point-to-point) |
| | himited no of ports knows and |
| | (4) DAS does not reale well (NO xcalability) himited handwidth & derhams |
| | (2) Unused Regauses example to early vallocated. |
| <u>-</u> | (3) Date not accomble ly deviser unes good groups. |
| | (4) Allows only one use at a Time |
| | (5) High administrature wits. |
| * ** <u> </u> | (6) DAS performance is effected by dish utilization, throughput, cache |
| | memory of storage device, virtual memory of host, RAID livel configurations, |
| | stage storage controller protocol and the efficiency of the his - |
| \rightarrow | Benefits - |
| | (i) how cost |
| | (2) Configuration is simple and can be deployed covily and rapidly. |
| | (3) Storage management tasks is confir small and medicine enterprises. |
| | Cu Regimes box hardware and reftunce eliments. |
| - - | (5) Reliabliz |
| | (6) how confunity. |
| | (1) Abrence of storage interconnects and network laterry. |
| | Q . |
| | |
| | |
| | |

IDE/ATA -> Integrated Device Electronics / Advantaged technology attachment

JATA - Serial ATA

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| 3 | Network Attacked Store (hine) | | |
|-------------|--|---|------------------------|
| | The state of the s | | and the second second |
| | - Darch file - Marina | device attached to | a local once network |
| | | | |
| | 1 what | convolichation throu | igh file-level data |
| | accers and sharing | | |
| | Applications - NAS chables both UN. | IX and Microsoft We | ndows were to |
| | shone the same data scamlendy. | V | |
| | It was NES for UNIX, CIES for W Physical Elements of NAS- | Indows and FTP an | d other brotons |
| | | | |
| : | [WINDOWS] CIFS [IP] | NAS HEAD | NETWORK INTERFACE |
| | [WINDOWS] CIFS CLUBS | THIS READ | NFS CIFS |
| | And the second s | | NAS DEVICE OS |
| | Δ. | STORAGE ARRAY | STORAGE INTERFACE |
| | Components of NAS one - | and the second | |
| 4 | () NAS Head (CPU and Memory) | | |
| 3 | 2) One or more network interface could (NJ) | (s) which he said | 0 |
| | networks. Eg-Gigatit Ethernet, Fast Ethe | unet | onnerlink to the |
| 3 | 3) An optimized OS for managing NAS for | m chimalt. | |
| 4 | 1) NFS and CIFS protocob for file tha | nin. | |
| 5 | 5) Industry standard storage protocols to a | Omas of Gard many | |
| | resources such as ATA, SESI & FC. | end menuge | shyneal disk |
| → 1 | VAS implementations - | | and provide the second |
| I | integrated NAS - 9t has all the component | and NOS | |
| H | trace in a ringle enclosure in frame. This make | gras, wer as the | NAS head and |
| en | unonment. | Structured NAS | a relf-contained |
| | Simple management functions. Scalar | 11th and 1 | |
| G | ateury NAS - It connists of an independent | The similed | |
| 0 | mays. | MAS read and or | e or more storage |
| | Compler management functions. Scalar | الاستان ما الاستان ال | |
| → <u>Ce</u> | muelting - | my as per required | |
| 11 | Integrated NAS - Standard IP network | | |
| 6 | nycompanion | | |

Gateway NAS - Front-end connectivity is done by IP Network het communication between the NAS and Storage yetern is done by FC SAN. UNIX SERVER INTEGRATED WINDOWS SERVER CLIENT 1 UNIX SERVER WINDOWS SERVER FC SAN CLIENT 2 NAS GATEWAY CLIENT 3 NAS Management ->NAS Fila I/O -NAS uns file-level occess for all of its I/O operations. File I/O is a high-level request that specific the file to be accepted, but does not specify its Inguis Work address 1) File hystem and Remote file thaning - A file system is a thurtured way of Tring and organizing data files. Many file systems maintain a file accustable to springlify the process of funding and occerning file 2) Accuming a Filwhysten - A file yetern must be mounted before it conferred When mounting a file system, the OS organizes files and directories in a tree-like shurting and grants the user the frivilege of accorning this should Files one located at the leaf rudes, and directories and new-directories one licated of intermediate node 3) File Moning - A user who exotes the file (owner) determines the type of access to be given to other uses (read, write, execute, append, debite, and list). and controls changes to the file.

File Transfer Protocol (FTP) is used in remote file shawing between

the severs and was clinto

Mycompanion

| NASI | 10 | Operations | 400 |
|------|----|------------|-----|
| | | | |

The NFS and CIFS protocols brandle file I/O requests to a remote file system, which is managed by the NAS device. The process of NAS I/O is as follows -

- (1) The requestor packages and I/O request into TCP/IP and forwards it though the network stack. The NAS device receives this request from the network.
- (2) The NAS device converts the I/O request into an appropriate physical storage request shock level I/O, and the performs the operation against the physical storage foot.

 (3) When the data is returned from the physical storage pool, the NAS device presents and

repackages the data into an appropriate file protocol response

(4) The NAS device packages this response into TCP [IP again and forward it to the clint

- Chwigh the network.

| | | | | Name and Address of the Owner, where the Owner, which is the Owner, where the Owner, which is the Owner, where the Owner, which is the Ow |
|---|-------------------|---|-------------------|--|
| - | APPLICATION | | STORAGE INTERFACE | < |
| | 1 0S | | NETWORK PROTOCOL | |
| 2 | I/O REDIRECT | | NAS OS | • 13 |
| | NES / CIFS | | NFS/CIFS | |
| | TEP IP STACK | | TEP/JP STACK | |
| | NETWORK INTERFACE | 3 | NETWORK INTERFACE | |
| | CLIENT | 4 | NAS DEVISE | |

CLIENT U NAS DEVICE

Histing and accessing files on NAS- (Stepwie pures)

- (1) Create storage array volumes
- (2) heate NAS volumes
- (3) heate NAS file rystems
- (4) Mount file systems
- (5) Acus the file systems

> Standards -

NAS file shaving protocob-

(1) NFS (Network file yearn) -

It is a chint/ sewer protocol. It was a machine independent model to represent uses data.

It was RPC muhanimo over PCP purtous my companion

Mount point grant occess to remote hierarchical file structures for local file system thucking. Access to the mound can be controlled by permissions Three versions of NFS are-NFS version 2 (NFS v 2) - Uses GETUDP to provide stateless network connection Features - locking are handled outside the probably NFS vernion 3 (NFSV3) - Une TCPOTUDP to provide statelan network connection Features - anynchmous writes, 64 bit file size NFS remon 4 (NFSV4) - Un TEP to provide stateful network connection Feature - Enhanced sicurity (2) CIFS (Common Internet File System) -It is a chief reme application protocol that enables client programs to make regisests for files and sensite on remote computers over TEP/IP. It is a fulling or open, variation of server Menage Rlock (SMB) protocol. It is a statiful protoco To ennue data intiguity -(1) It uses file and record locking in It runs over TCP (iii) It supports fault tohence and can automatically restore connections and reopen files that were open fruir to interruption. (i) F Embanual receiving due to the use of NFSV4 protocol (2) Provide Use tuller authentication and authorization (1) how huformance as compand to SAN. R) Communio more bandwidth (3) IP natural night be congested: (4) Factors that affect NAS performance one number of hops, authentication with the directory service, retransmission, overutilized soutes and switches, file deserting lookup and metadata requests, Overutilized NAS devices and Overutilized chints

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| | | e dag marka a sa |
|---------------------|--|---|
| → | Benefits of WAS- | |
| | | 5) Simplified management |
| | | |
| | 9 | 6) Scalability |
| | | 1) high availability |
| | (4) Centralized storage | s) Security. |
| | | |
| (4) | Content-Addressed Storage (CAS) - | |
| | It is an object-based system that has been | n purposely built for storing |
| | fined content data. It is designed for secure online & | wage and retrival of fixed |
| | content. The stoud object is arrighed a globally and | |
| | content address (CA). | |
| | It provides an oftimized and centrally managed | Storage volution that can |
| | support ringle-instance storage (SiS) to eliminate me | |
| > | Applications - Health cone rotation (Storing Patient | |
| | (storing financial records) | Appendix of the second of the |
| > | Physical Elements of CAS - | |
| | CLIENT SERVER API | Physical elements of CAS |
| <u>, </u> | ARLESS NODES STORAGE NODE | 9 |
| | | CAS land storage |
| | CAS SYSTEM | La hogical elements of CAS |
| | Private LAH | are API, metadata any |
| | CAS ARCHITECTURE | Object - level access proturals |
| | A client accesses the (AS-land storage over a L | , |
| | runs the KLAS API which is responsible for perform | |
| | application to store and whieve the data. | |
| · Service · · · · · | | male + a - (PATN) |
| | CAS architecture is a redundant away of indepe | Mas (M.Z.) |
| | CAS management - | |
| • | > Object storage and retrieval in CAS - | |
| | some trimologies- | , , , , , , , , , , , , , , , , , , , |
| +2 | C-Clib - A virtual beckers that contains dates as | lit amounted CDF |

C-Clip Descriptor file (CDF) - XML file that a system custes while making a C-Clip

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| | (A is calculated) Recalculate (A for validation |
|---------------|--|
| | Storing data object on CAS - JE |
| 1 | Ameldate CA+mulate > C-tilly OBJECT) |
| | CLIENT CONTECT D API (C-Clip) |
| | with C-CILLED CAS SYSTEM |
| | EGF CON TO |
| | INTERNAL STORME & CONTRACT OF AND AND HOLE |
| | Data object exterior from CAS - |
| | PROUNT SERVER DESECT DESECT |
| 1 | CITEME C-CLIP |
| | Object CAS SYSTEM |
| | C-Up) |
| | INTERNAL INTERNAL |
| | Mandard STORAGE |
| \rightarrow | Connectinty - Use IP protocol (LAN for inside the CAS system), |
| | Security - Data recently is provided by constantly communication between |
| | the occase nodes and the storage nodes |
| \rightarrow | himilation - |
| | (i) Com be shower than SAN, NAS on DAS |
| ٠. | |
| | (2) Afflication intigration |
| | (3) Still cot of ownership is higher even though T(O is raprificantly lower. |
| -> | Broutis |
| | (4) Content a Martin (5) Retention enforcement |
| · . · | (2) Control Transition and duportion |
| | (3) horodin Indiference (1) Technology Indefendence |
| | (4) Single-without stage (Sis). (6) Fast second returned. |
| | |
| (B) | Storage Aug Netume (CAN) - |
| | |
| | It is a high speed, dedicated network of severs and showed strage devices. |
| | Traditionally connected over Files Channel (FC) networks, as AN forms a more |
| | storage part and facilitates data containation and comolidation. |
| > | File Channel (FC)- |
| | 9k is a high speed network technology that num on high - speed oftered |
| | files (frequed for front-end SAN connectivity) and revial cables (frequenced fra |
| | <i>my</i> companion |

| Voik-end dish commentarity) |
|--|
| (FC-AL) |
| SAN Islands FC Arbitrated hoof -> Interconnected SANS FC switched falmic |
| Enterprise SANSS FC sixtehed facts fabric 4 |
| SERVER SERVERT SERVERT SERVERT SERVERT SERVERT |
| EC SWITCH FC SWITCH |
| FC Switch FC Switch |
| FC Hub FC Switch FC SWITCH FC SWITCH |
| STORAGE STORAGE STORAGE STORAGE ARRAY 1 ARRAY 2 ARRAY 1 ARRAY 2 |
| |
| SAY Implementation - |
| SERVER 1 STORAHE ARRAY 1 |
| SERVER 2 FC SAN } |
| SERVER3 STORAGE ARRAY 2 |
| |
| Components of SAN - (Physical Elements) |
| Three vani components - Servers, Network Infrastructure and storage. |
| These components can be further holen dozun into five key clements - |
| (1) Node Pat - |
| Each node regions one or more furto to provide a physical interface for |
| communication with of her nodes. These historia integral components of an HRA and |
| communicating with other nodes. These parts are integral components of an HBA and the storage front-end adapters. |
| |
| Port 0 1x |
| NODE LINK |
| Porta |
| |
| (2) <u>Cabring</u> - |
| SAN implementation unes ofteral file coulding Copper can be used for |
| shorter distances for lock-end connectivity, as it provides a better signal-to-noise |
| The state of the s |
| |

Optical file carles carry data in form of light. There are time types of ofitical

| H | |
|---|---|
| | calles & multimode file (MMF) -> kany multiple beams of light and single- |
| | mode film (SMF) -) cang a ringle team of light |
| | MMFs are generally used within data centers for shorter distance runs, |
| | certilo SMFs one used for longer distance ruins. MMF tromseccives on less enpensive |
| | as compared to SMF transceives |
| | Connection - Standard Connection (SC) -> 1 hb/s transmission speed |
| | herent connector (LC) -> 466/s transmin speed, |
| | Straight Tip (ST) connector has a plug and a rocket that is wiked inth a hal |
| | twited bayonet lock |
| - | Small Form-foctor Pluggable (SFP) is an oftical transcewer used in |
| | optical communication. (106/16/s data rates) |
| | (3) Interconnecting Derries - |
| | ful 6 Pas - Bus topology, Provident (No longer used), Bandwilt Ma |
| | Snorth - Stone topology point to point, No Bandwidth sharing |
| - | Directors -> hauges than ninteres and one deployed for dates centres implementate |
| | The function of directors is similar to that of FC ninteles, but directors have |
| | high pat went and fault tohence capability. |
| | (4) Stange anous - |
| | I promet high a wile litty high as and redundancy, improved perform |
| | business continuity and multiple had connectually. |
| | (5) SAN Management Miles |
| | It manages the interfaces between host, interconnect devices and storage |
| | anoys. It provides key managements functions - mapping of storage dencis, |
| | surtates, and sever monitoring and generating alests for discovered devices a |
| | logical partitioning of the SAN, called zoning. |
| | |
| _ | FC Connectivity - |
| | There love interpretent Alterna - |
| _ | Three boni interestants oftens - |
| | (1) Point-to-Point- |
| _ | Simplest FC configuration - two duries one connected directly to |

Offers limited connecturity and no realability

each other.

| | SERVER 1 |
|----|---|
| | SERVER2 STORAGE ARRAY |
| | SERVE 23 |
| | (2) Filre Channel Stituted hoop (FC-AL) - |
| | Derrie are attached to a shoup loop. It has characteristics of a token |
| | ring topology and a physical stare topology. |
| | At a given time, only one device can perform I/O operations on the loop. |
| | Devices on the loop must "arbitrate" to gain control of the loop. |
| - | himilted scalability due to bandwidth shaving, was 8-hit addressing (127 dense), |
| _ | and adding or removing a device results in loop re-initialization. |
| | SERVER 1 |
| ķ. | SERVER 2 STORAGE ARRAY |
| | SERVER 3 FC HOB |
| | (3) File Channel Switched Fabric (FC-SW or fatric connect) - |
| _ | 9t purioles interconnected deines, declicated landwidth and |
| | reability |
| | A fahie is a logical space in which all nodes communicate with one |
| | another in a network which is created using a switch or a network of rwitchs. |
| | Each whiteh in a fabric contains a unique domain identifier, which is part of |
| | the fabrus addrening scheme. |
| | SERVER 1 STURME ARRAY |
| | SERVER 3 FC SWITCH |
| | |
| | Filre Channel Ports - |
| | (1) the N-port (Node Port) - An end point in the facili |
| | (2) NL-port (Node hoof Port) -> Node port that repeats the arbitrated with topology. |
| | (3) E-port (Enpennion Port) → FC port that forms the connection between two FC norther. |
| | (4) F_port (Falsi Port) -> A port on a ninter that connects on N-port |
| | (5) FL-port \$1 -> A fabric part that parthupoles in FC-A2. |
| | (6) G-port → Agenein port that can ofunte as an E-port or an E-port. |
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| Filre | Channel | Architecture - | (Standards) |
|-------|---------|----------------|--------------|
| | | | |

The FC architecture represents true channel network integration with standard interconnecting devices.

Channel technologies provide high levels of performance with low protocol overheads.

File Channel Partonal (FCP) is the implementation of social SCSI-3 over an FC network. In the FCP architecture, all enternal and remote storage devices attached to the SAN appear as local devices to the Hort OS. The key advantages of FIP ene as follows-

- (1) Sustained transmission boundwidth over long distances,
- (2) hepport for a large number of addressable devices over a network (15 million)
- (3) Exhibits the characteristics of channel transport and provides speed upto 8.5 Gb/s

- Filre Channel Protocol Stack -

| | | APPLI | CATION | | |
|------|------------------------|--------|--------|------|----|
| FC-4 | SCSI | HIPPI | ESCON | ATM | ΙP |
| FC-2 | FRAMING FLOW CONTROL | | | | |
| FC-1 | | ENCODE | DECODE | | |
| F4-0 | عرباد | 240/s | 40615 | Sab/ | S |

FC-9 (1866 for fortical) - Defines the application interfaces and the way Uffer by from (IIIR) as maffed to the lower FC layers.

fair of the feetbook on SCST, His Reference Bralled Interfere (MRPP)

framing funtional , Partitions Strong Come Koy (ESCON), ATM and IP.

and destination hat, and link control information.

It also defines fatine tensio, closes of some, flow control and control of FC-1 (Transmission Pratical) - Includes sound encoding and decoding value throat character used, and enor control.

FC-0 (Physical Interfore) - defines the physical introduce, media, and transmission of naw hits.

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| - | File Chan | mel addressing - | | | | <u> </u> | |
|------------------------|---|--------------------------|--------------|------------------------|-----------------------|-------------|---------------------------|
| | | aclohis is dy | | | | | |
| | 24 bit FC | address of N-por | <u>x</u> - d | 0 | | <i>J</i> | V |
| | , | | AREA . | T.D ROI | QI TD |] | - |
| | 8 ti 17 address | ts (239) ed are reverved | 8-bits | | n'ts (256) | = 15, | 663,104 adduns |
| - | Area | -ID is used to | identify | a group of E-h | nt, = | | |
| | 24 bit FC | address of NL-p | ort- | | | · . | |
| | | LOUP ID / U | | AL-PA I | 0 | | |
| | LOOP ID V | n cose of public lo | oh (6.61% | s) 8-bits (| 127) | | |
| | Uhnud | in case of mirate | woh (16-61 | - / | | | |
| | rava | loop > when a | | | | | - |
| | | e Names (WWH | | similar | | | |
| | Each | device in the F | e envio | ment is arrigi | ned 0 641 | it unqu | e identifer |
| | called the | WWN. 9tisa | static na | me for each d | ence on a | in Fe he | turk. |
| _ * | | this - World W | | | | | |
| | II | · | | Note Town I will serve | | | · · |
| | | | | | | | |
| <u> </u> | FC Frame | | | | | Endel H | lane |
| delings > | SOF | FRAME HEADER | | TA FIELD 2112 BYTES | CR K 4BYTES | EOF | edeliniters |
| delimiters | Start of frame | | - | | Cyclic Red | underry Che | k |
| | 6 | RUUTING CONTROL | (R-CTL) | DESTINAT | TON ID (| 0-10) | |
| | | ECIFIC CONTROL (C | | | D (S_IC | | |
| 8.7 | TYPE FRAME CONTRO | | | | | | |
| | SEQUE | NCE ID (SEQ_I | o) | DATA FIELD (ON | ITROL (DE_ | CTL) SEC | (SEQ_(NT) QUENCE (OUNT |
| 8 . E + S S | 1 | ORIGINATING EX | CUANGE I | 0 (0x-10) | | | ONDER EXCHANGE (RX_ID |
| | ₹ | 1BYTE | OFFSE | * 1 | BYTE - | | 2 BYTES |
| | Tuho - | deniles the 1111. | | | | +1 | |
| | Route Co | denites the Upper | hatta d | Harrist A. | in and | ou frame | |
| | llan de | ntul-denotes | TO I TO | une frome is a l | ink control | frome or | a data frame. |
| | Data Field control - Should any optioned headers present at the beginning or not. | | | | | | |
| | vara fuld | untiel - Indu | atto any | optional headers | present a | the begin | migg a not |

Frame control - Contain control information arlated to frame content

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Yes.

Moduata

| Structure and Organi | whom of FC date - | <u>.</u> | A Company | | | | | | |
|--|--|--|---|--|--|--|--|--|--|
| | | | | | | | | | |
| two heaple, wherehe a frame represents a word, a requerce represents a sentence, | | | | | | | | | |
| | | | | | | | | | |
| | | | | | | | | | |
| • | | | | | | | | | |
| Frame - fund | amount unit of etala | namper | | | | | | | |
| U | A STATE OF THE STA | | | | | | | | |
| Flow Control - | y kamanan ara-gata ya ta | | | | | | | | |
| | have of the flow of | data frames during d | ater transfermon | | | | | | |
| | | | | | | | | | |
| | | | | | | | | | |
| | | | | | | | | | |
| | | | | | | | | | |
| heless and continuous to send haves it the count is execute them zero | | | | | | | | | |
| 2) Ford-to-End Credit (CE Credit) - Similar to BR Credit | | | | | | | | | |
| 11 | | | | | | | | | |
| | | | | | | | | | |
| | | | | | | | | | |
| | U | UP | đ | | | | | | |
| Cland Service - | | | | | | | | | |
| | CAASS 1 | CLASS 2 | CLASS 3 | | | | | | |
| Companication Type | | Non deducated connection | Non dedicated connect | | | | | | |
| | | | BB_credit | | | | | | |
| · · | | 3 | Order not guarant | | | | | | |
| • | 4 | Acknowledged | Not teknowledge | | | | | | |
| | In an FC nete two people, whereby and an enchangenely Enchange operation Legionie - Conti Frame - fund Flow Control - 9t defines the FC technology was to 1) Briffer to - Briffer (21 is used frames. I helpes and continue 2) Ford-to-End Credit with each other, they | In an FC returns, data transport two people, whereby a frame represents a conservation. Enchange of returns a conservation. Enchange of returns a conservation. Enchange of received frame Frame — fundamental unit of etada Flow Control — 3t define the frame of the flow of FC technology was two flow-control me 1) Briffer-to-Briffer Credit (BB Credit) The firm and continues to send frames if 2) End-to-End Credit (FE Credit) — with each other, they enchange the EE C 9t differ to the flow control for class Class of terrice— CLASS 1 Companication Type Dedicated Commention Flow Control EE-credit Frame Delairey 3n order Delivery | 9t define the paire of the flow of data frames during de FC technology was two flow-control mechanism — 1) Briffy -to-Briffer Gedit (BB_Credit)— 3t is used for hardware-based flow control. It controls to make the part of frames. The transmitting part maintains a count of helpers and continuous to said frames if the count is greater the 2) Ford-to-End Gredit (FE_Credit)———————————————————————————————————— | | | | | | |

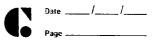
Multipliani

Bandwidth Hilyation

Zoning -It is a FC nirtch function that enables nodes within the foodric to be logically segmented into groups that can communicate with each other SEREVER I STORAGE FC SAH ARRAY SERVER2 SERVER 3 ZONING (also called Zone Configurations)
Zone Sets (hibrary) Only one zone set can be centre cettire at a time RONE SET ZONEL ZONE2 Zone (hibrary) MEMBER3 MEMBER4 MEMBER 2 Members WNN'S > Types of Zoning - Three types -(1) Port zoning - It was the FC address of the physical parts to define zones It is also called Hard zoning. Although this method is secure, it requires updating of zoning configuration information in the event of fabric reconfiguration (2) WWN Zoning - It was World Wide Names to define zones. It is also called soft zoning - It is flerible because does not require zones abdating of zoning configuration information in the event of fabric reconfiguration (3) Mined zoning - It combines the quality of both WWN zoning and port zoning Using mind zoning enables a specific port to be tied to the WWN of a node. - Fabri Channel hogin Types -Fahir Kurus define three light types -(1) Favri hogin (FLOGI) - between N-port and F-port FhOhI frame contains WWNN LWWPN franchetes kind by N- part & F-port kind Accept (ACC) from to N- port. N- ports und the frame at FC address FFFFFE

(2) Port hogin (PLU h) - between N-port and another N-port to enchange (3) Process hogin (PRLI) - between N-port and another N-port to enchange the information about the FC type in use, the SCSI initiates and or the target, my companion

| 1 | |
|---|--|
| • | FC topologies - |
| | (1) Core-Edge Falmi - |
| | In this topology, there are two types of ninter trees in the false. |
| | Edge ties - Comprises neithers and offers an invententive approach to adding more |
| | hoste in a fatir. The traction of the edge forward from the The at the war. The |
| | nodes on the edge can communicate with each other. |
| | Conetice - Compains enterprise directors that ensure high fahir waitability |
| | Additionally all ties traffic has to either havens through a terminate at this ties. |
| | This topology increases connectinty within the SAN while consuming overall |
| | port utilization. Two variations one - |
| | Single Core topology - E04E TIER |
| | FC SWITCHI FC CWITCH FC SWITCH 3 |
| | SERVER STORAGE ARRAY |
| | CORE TIER DIRECTOR |
| | Dual are tohology - |
| | FC SWITCHI FC SWITCHE FC SWITCHE |
| | SERVER STORAGE ARRAY |
| | DIRECTOR DIRECTOR |
| | CORE TIEL |
| | Beautis - Promise one hop storage accen to all storage in the system |
| | Offerailes comes colombian of ISL looking and traffic patterns |
| | (3) Scalable to lance engreenests |
| | himletion - (1) bear hop lead to toanmining delay |
| | QI Preference is officed by weener in Isls and domain count in the fatition |
| | G |
| | (2) Mesh Topelogy - |
| | Each minted to directly connected to other puritches by using ISLS. |
| , | It enhanced commentanty. Two type - Partial Mesh and Full Mesh |
| | Survey make the bound how is agained |
| | SERVER SERVER SERVER |
| | STORAGE STORAGE ARRAY |
| | mycompanion Partial Mesh Full Mesh |
| Ш | |



| > | Benefits of SAN - |
|----------------|---|
| | (1) High Bondwidth |
| | (2) SCSI Entensión |
| | (3) Renounce Comolidation |
| | (4) Scalability |
| | (5) Secure Acces |
| | |
| _ → | Security - 9 robotions and filtering features one present in SAN. |
| - | Management - |
| | (1) Infrastruture prolection. |
| | (2) Fabric Management |
| | 3) Storage Allocation |
| | II . , , |
| | (3) Performence Management |
| | (5) Performance Management |
| -> | Applications - |
| | Und in disk anays, take libraries and officed Takelones to enhace stronge |
| | dencis |
| | - ANNUS |
| | himitations - |
| | |
| | (1) Very Expensive |
| | (2) Management of SAH is difficult (3) Complexity of SAH |
| | Co companity of JAN |
| | |
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| | <i>My</i> COMPANION |
| ļ | Mycorrigation |