

## 作业——1

0. 请证明补码加法公式： $[x]_{\text{补}} + [y]_{\text{补}} = [x + y]_{\text{补}} \pmod{2^w}$ 。 $[*]_{\text{补}}$ 表示整型数据\*的补码表示，机器字长为  $w$ 。

1. 将 8 位无符号数 130 转换为 8 位浮点数（exp 域宽度为 4 bits, frac 域宽度为 3bits）

Exp = ?

Frac = ?

2. We are running programs on a machine with the following characteristics:

- Values of type int are 32 bits. They are represented in two's complement（补码）, and they are right shifted arithmetically. Values of type unsigned are 32 bits.
- Values of type float are represented using the 32-bit IEEE floating point format, while values of type double use the 64-bit IEEE floating point format.

We generate arbitrary values  $x$ ,  $y$ , and  $z$ , and convert them to other forms as follows:

```
/* Create some arbitrary values */
int x = random();
int y = random();
int z = random();
/* Convert to other forms */
unsigned ux = (unsigned) x;
unsigned uy = (unsigned) y;
double dx = (double) x;
double dy = (double) y;
double dz = (double) z;
```

For each of the following C expressions, you are to indicate whether or not the expression always yields 1.

Expression	Always True?
$(x < y) == (-x > -y)$	Y N
$((x+y) << 4) + y - x == 17*y + 15*x$	Y N
$\sim x + \sim y + 1 == \sim(x+y)$	Y N
$ux - uy == -(y - x)$	Y N
$(x >= 0) \parallel (x < ux)$	Y N
$((x >> 1) << 1) <= x$	Y N
$(\text{double})(\text{float}) x == (\text{double}) x$	Y N
$dx + dy == (\text{double})(y+x)$	Y N
$dx + dy + dz == dz + dy + dx$	Y N

3. In the following questions assume the variables **a** and **b** are **signed integers** and that the machine uses two's complement representation. Also assume that MAX\_INT is the maximum integer, MIN\_INT is the minimum integer, and W is one less than the word length (e.g., W = 31 for 32-bit integers). Match each of the descriptions on the left with a line of code on the right (write in the letter).

//1's Complement: 反码, 即按位取反

//2's Complement: 补码

1. One's complement of a

\_\_\_\_\_

a.  $\sim(\sim a \mid (b \wedge (\text{MIN\_INT} + \text{MAX\_INT})))$

2. a.

\_\_\_\_\_

b.  $((a \wedge b) \& \sim b) \mid (\sim(a \wedge b) \& b)$

3. a & b.

\_\_\_\_\_

d.  $(a \ll 4) + (a \ll 2) + (a \ll 1)$

4. a \* 7.

\_\_\_\_\_

e.  $((a < 0) ? (a + 3) : a) \gg 2$

f.  $a \wedge (\text{MIN\_INT} + \text{MAX\_INT})$

5. a / 4 .

\_\_\_\_\_

g.  $\sim((a \mid (\sim a + 1)) \gg W) \& 1$

h.  $\sim((a \gg W) \ll 1)$

6. (a < 0) ? 1 : -1 .

\_\_\_\_\_

i.  $a \gg 2$

4. Match each of the assembler routines on the left with the equivalent C function on the right.

```

foo1:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%eax
    sall $4,%eax
    subl 8(%ebp),%eax
    movl %ebp,%esp
    popl %ebp
    ret

foo2:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%eax
    testl %eax,%eax
    jge .L4
    addl $15,%eax
.L4:
    sarl $4,%eax
    movl %ebp,%esp
    popl %ebp
    ret

foo3:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%eax
    shr1 $31,%eax
    movl %ebp,%esp
    popl %ebp
    ret

```

```

int choice1(int x)
{
    return (x < 0);
}

int choice2(int x)
{
    return (x << 31) & 1;
}

int choice3(int x)
{
    return 15 * x;
}

int choice4(int x)
{
    return (x + 15) / 4;
}

int choice5(int x)
{
    return x / 16;
}

int choice6(int x)
{
    return (x >> 31);
}

```