

IT5005 Artificial Intelligence

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2. Uninformed Search

Learning Objectives of the Session

- Understand workflow of goal-based agents with atomic representation
- Modeling a problem into an implicit state space graph

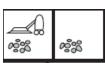
- Learn how to search for a solution on implicit state space graphs
 - Uninformed search

Recap

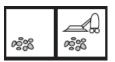
Atomic Representation

Two-Room Vacuum World: Eight States

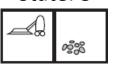
State: 1



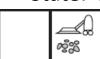
State: 2



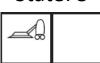
State: 3



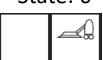
State: 4



State: 5



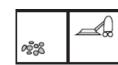
State: 6



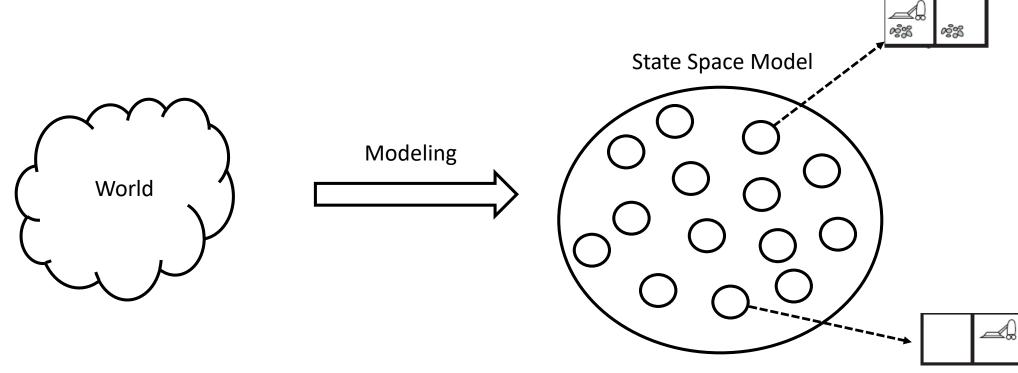
State: 7



State: 8

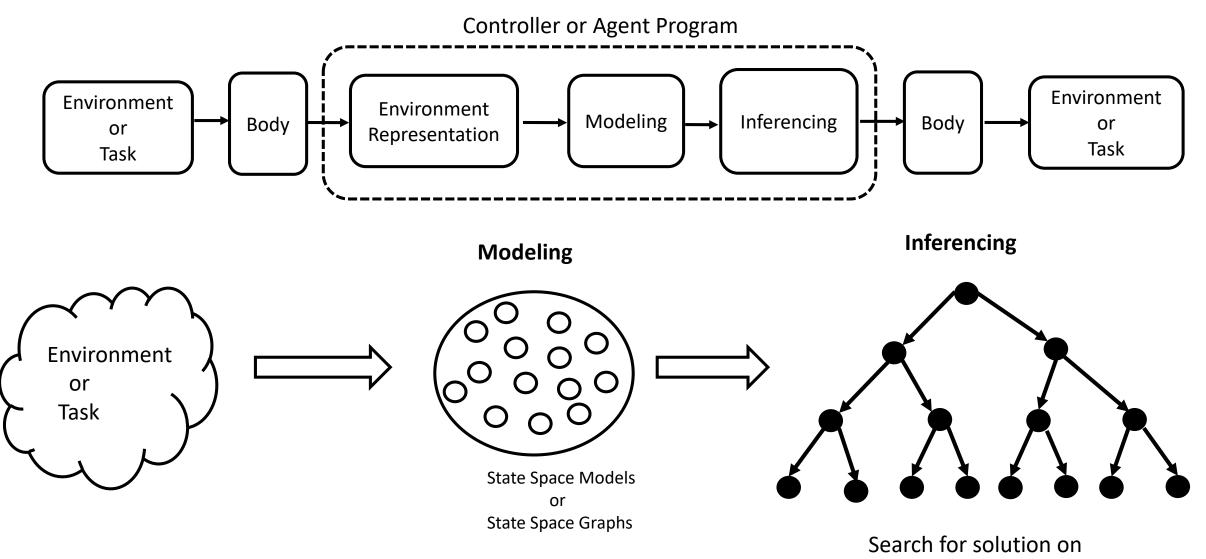


Modeling with Atomic Representation



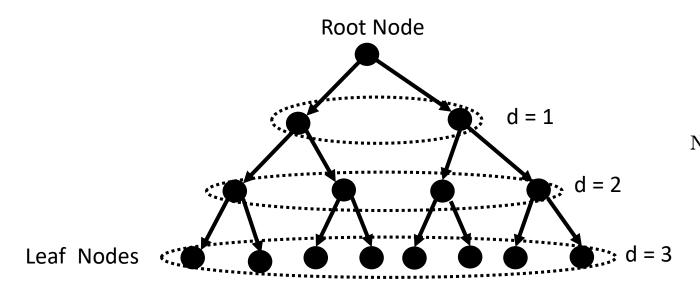
- Each node (vertex) is related to a state of the world
- Goal should also be represented by state (node/vertex)

Problem Solving by Search: Workflow



State Space Graph

Preliminaries: Trees



Branching Factor = bMaximum Depth = mNumber of Nodes at depth $d = b^d$ Number of Leaf Nodes $= b^m$ Number of Nodes $= 1 + b + b^2 + \dots b^d + \dots + b^m$ m: maximum depth

Preliminaries: Performance Measures

- Completeness
 - Complete if algorithm can find a solution (achieve the goal)
- Optimality (aka rationality)
 - Optimal if algorithm finds an optimal (lowest cost) path to solution
- Time Complexity
 - Time taken to find the solution
 - Measured in terms of number of nodes generated
- Space Complexity
 - Memory needed to find the solution
 - Measured in terms of number of nodes stored while building the graph

Uninformed Search

Applications

- Puzzles
 - Slide Puzzles
 - Missionaries and Cannibals Problem, etc.
- Games:
 - Pacman, etc.
 - Maze Navigation
- Real-World Applications
 - Route Planning
 - Robot Motion Planning
 - VLSI Layout Planning, etc.

Agenda

Workflow

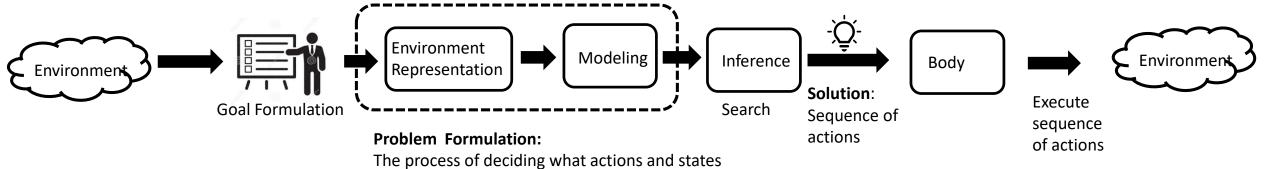
Goal Formulation

Problem Formulation or Modeling

Building Blocks of Search Tree

• Inference

Workflow



to consider given a problem

Agenda

Workflow

Goal Formulation

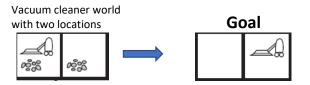
Problem Formulation or Modeling

Building Blocks of Search Tree

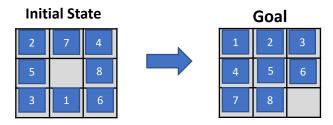
• Inference

Goal Formulation: Examples

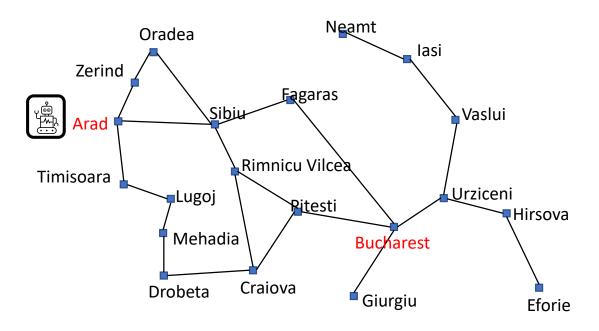
Two-Room Vacuum World:



• 8-Puzzle



Goal Formulation: Examples



"Map of Romania"

Route Planning

Goal: Find a path from Arad to Bucharest

Agenda

Workflow

Goal Formulation

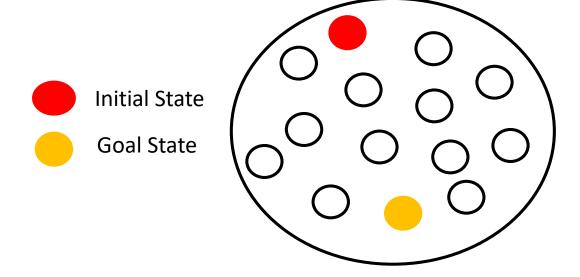
Problem Formulation or Modeling

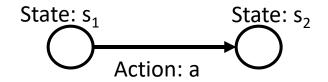
Building Blocks of Search Tree

• Inference

Modeling or Problem Formulation

- 1. States
- Initial State
- 3. Actions: Actions(s)
 - Legal (applicable) actions for an agent at a state s
- 4. Transition Model: RESULT(s, a)
 - Defines the result state for an action at a given state
 - Ex: $s_2 = RESULT(s_1, action)$
- 5. Goal State:
 - Test goal state using IS GOAL(s)
- 6. Action Cost Function: ACTION COST(s, a, s')
 - Cost of an action at a state

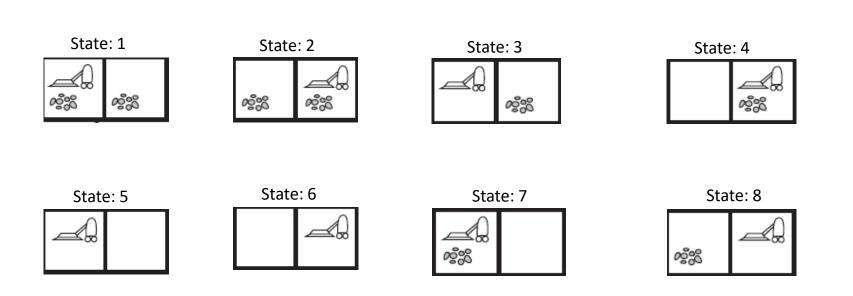




Action Cost: c(s₁,a,s₂)

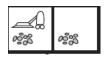
Two-Room Vacuum World: Modeling

1. States



Two-Room Vacuum World: Modeling

2. Initial State



3. Actions:

Move Left (L) Move Right (R) Suck dirt (S) No-op (N)

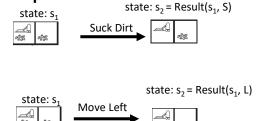
6. Action Cost

Suck Dirt: 1 Move Left: 2 Move Right: 2 No-Op: 0

4. Transition Model:

If room is dirty, *Suck Dirt* make it clean
If room is clean, *Suck Dirt* does nothing
If agent is in left room, move Right takes it to right room
If agent is in left room, move left does nothing, etc

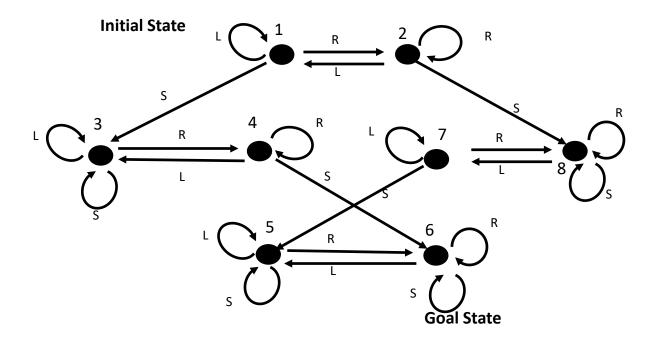
Example:



5. Goal State



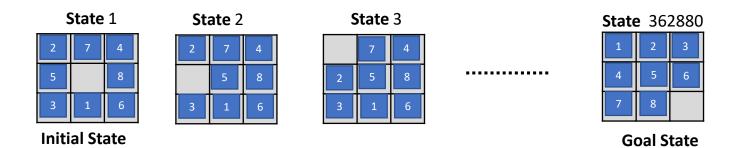
Two-Room Vacuum World: Modeling



8-Puzzle: Modeling

1. State Representation

Number of states = 9! = 362,880



8-Puzzle: Modeling

2. Initial State

2	7	4
5		8
3	1	6

3. Actions:

Movement of blank space Slide Left (L) Slide Right (R) Slide Below (B) Slide Above (A)

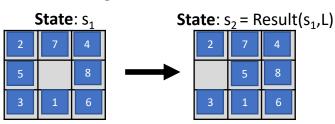
6. Action Cost

Each action costs 1

4. Transition Model:

Given a state and action return the resultant state

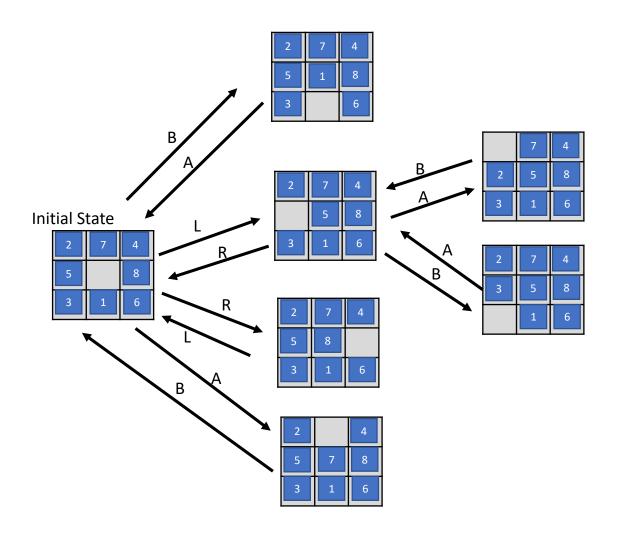
Eg:



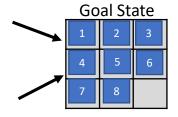
5. Goal State

1	2	3
4	5	6
7	8	

8-Puzzle: State Space Model



State Space Graph for 8-Puzzle (impossible to show complete state space graph)



A: Above

B: Below

L: Left

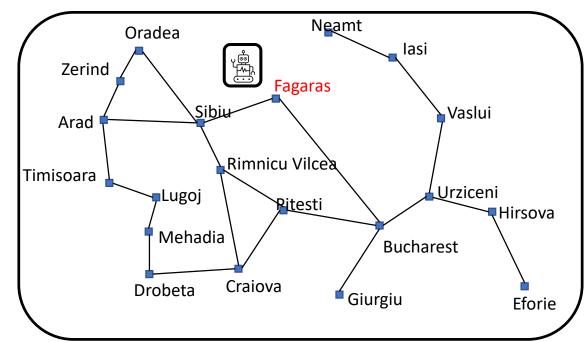
R: Right

Romania Map (Route Planning): Modeling

1. States

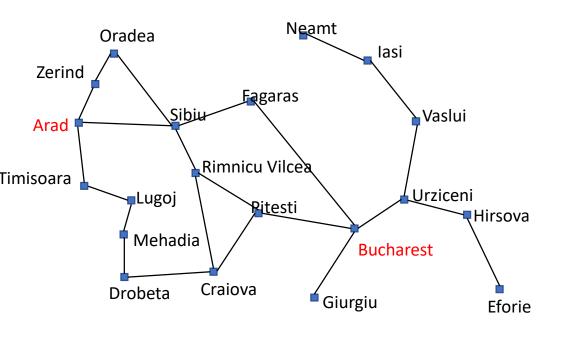
• State: location of agent

State: Arad Neamt Oradea lasi Zerind **Fagaras** Sibju Vaslui Arad Rimnicu Vilcea Timisoara Urziceni Lugoj **Ritesti** Hirsova Mehadia **Bucharest** Craiova Drobeta Giurgiu **Eforie** State: Fagaras



Romania Map (Route Planning): Modeling

Problem Formulation



1. States:

Arad, Oradea, etc.

2. Initial State:

Arad

3. Actions:

Eg: for *Arad*:

{ToZerind, ToSibiu, ToTimisoara}

4. Transition Model:

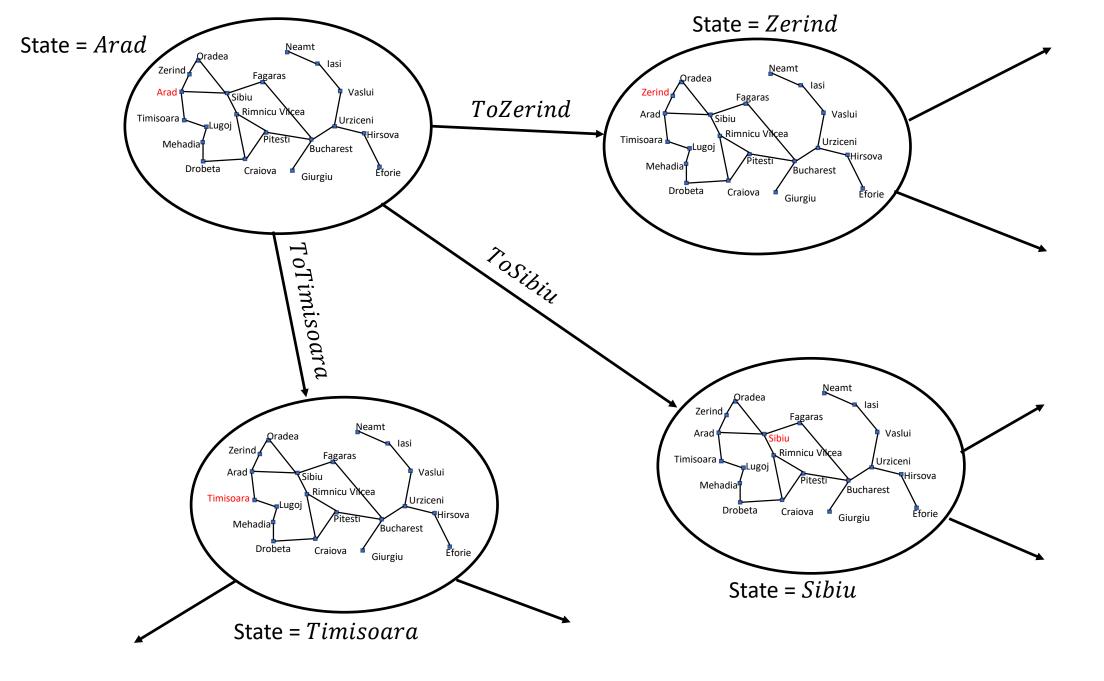
RESULT(Arad, ToZerind) = Zerind

5. Goal-Test:

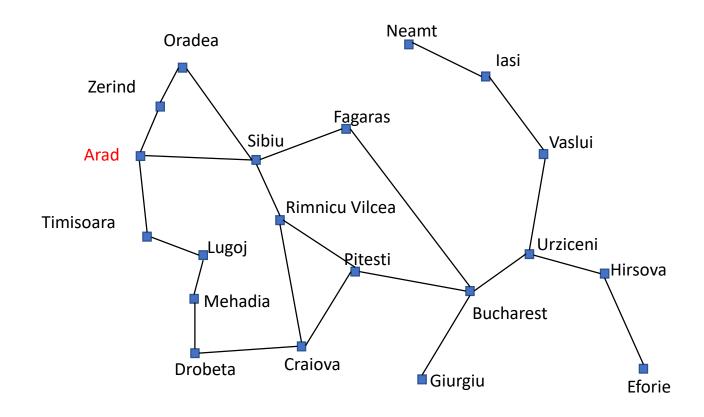
IS - GOAL(Bucharest)

6. Action-Cost:

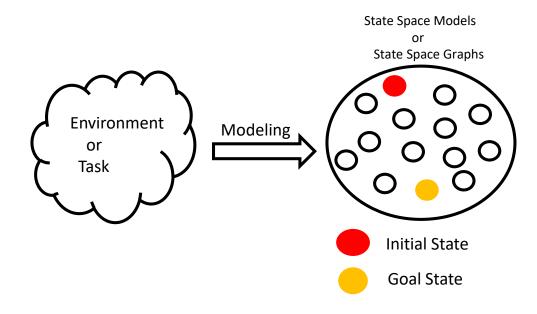
Distance between cities



Romania Map (Route Planning): Modeling



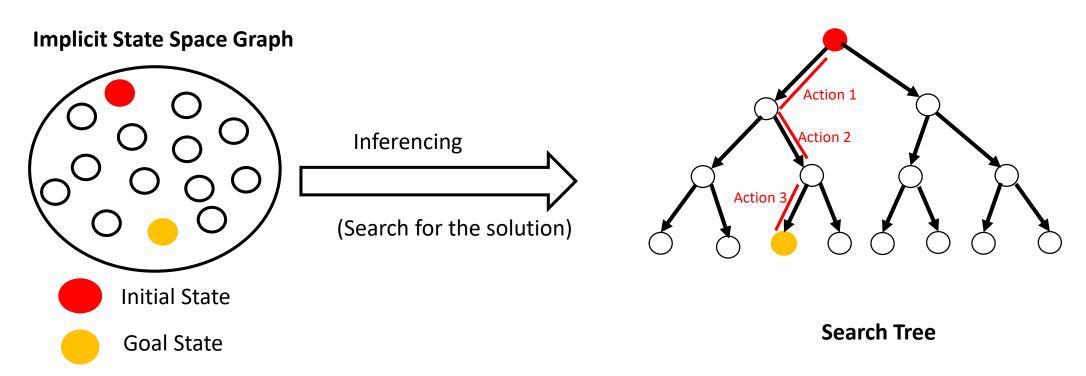
Until Now: Modeling



Note:

- Implicit Graph.
- Only provided a framework (model/problem formulation) to generate this graph

Next: Inference (Search for the solution)



Inferencing:

Finding the path from Initial State to Goal State

Agenda

Workflow

Goal Formulation

Problem Formulation or Modeling

Building Blocks of Search Tree

• Inference

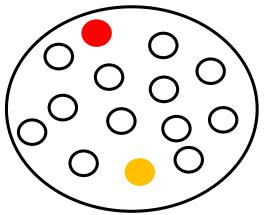
Building Blocks of Search Tree

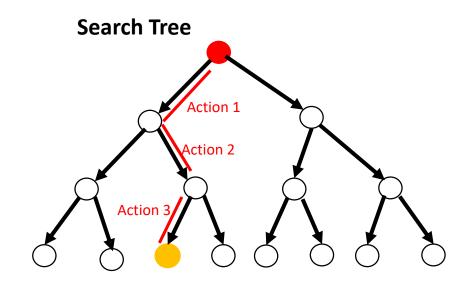
Need to create the search tree dynamically

- Require data structures
 - to generate nodes in search tree
 - for navigating around the search tree
- Let's start with node

Illustration with Romania map

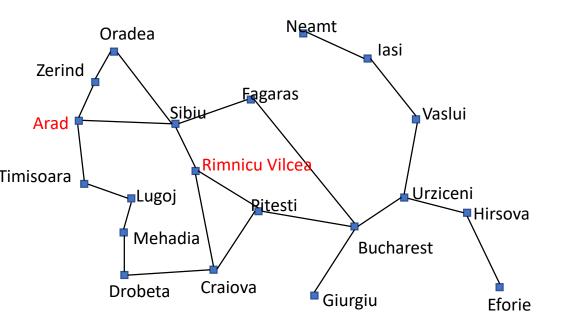






Romania Map: Route Planning

Problem Formulation



1. States:

Arad, Oradea, etc.

2. Initial State:

Arad

3. Actions:

Eg: for Arad:

{ToZerind, ToSibiu, ToTimisoara}

4. Transition Model:

RESULT(Arad, ToZerind) = Zerind

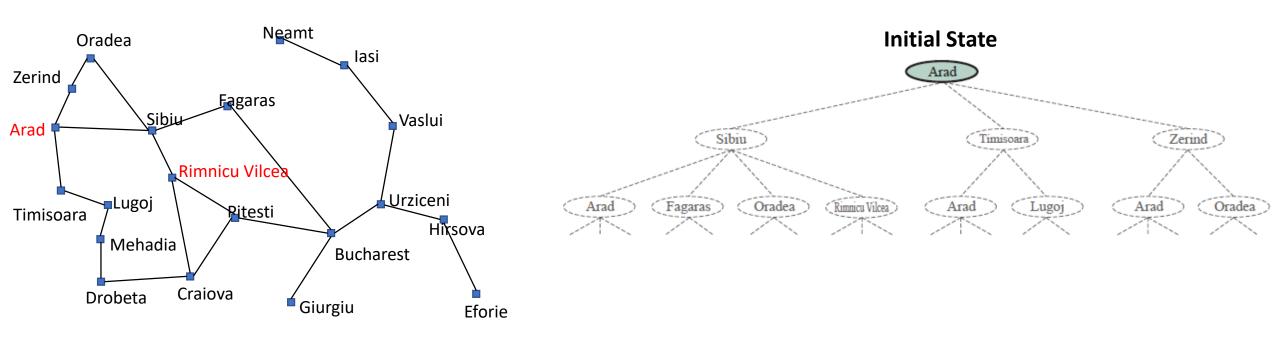
5. Goal-State: Rimnicu Vilcea

 $IS - GOAL(Rimnicu \ Vilcea)$

6. Action-Cost:

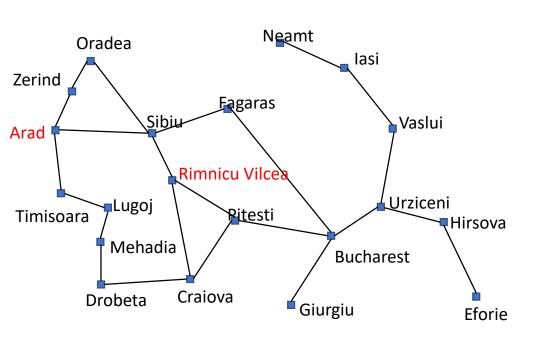
Distance between cities

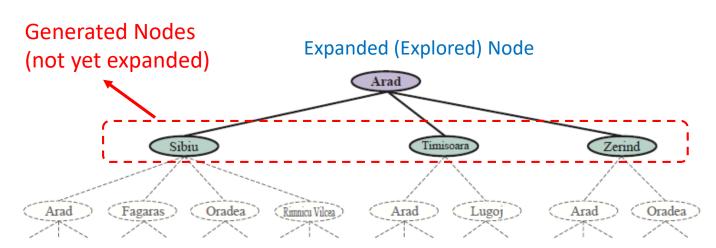
Search Tree: Arad to Rimnicu Vilcea



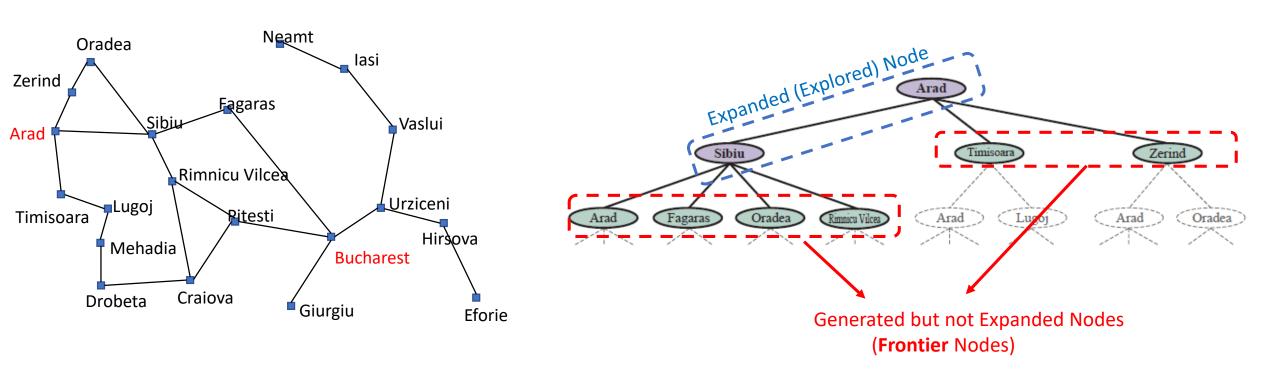
Actions at InitialState: {ToZerind, ToSibiu, ToTimisoara}

Search Tree: Arad to Rimnicu Vilcea





Search Tree: Arad to Rimnicu Vilcea

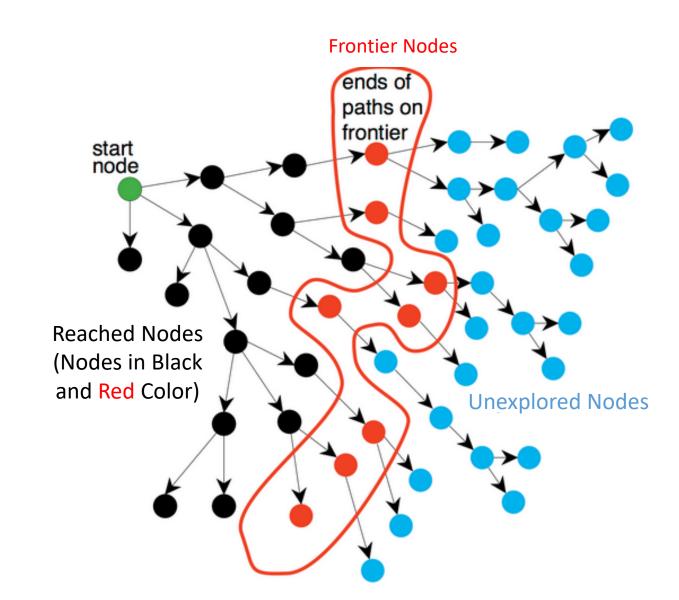


Reached Nodes:

Expanded Nodes and Frontier Nodes

Search Tree Terminology

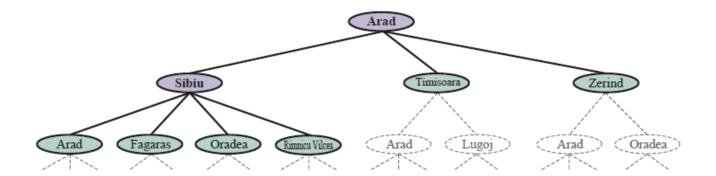
- Reached Nodes
 - Nodes in black and red color
 - Includes both expanded and not expanded nodes
- Frontier Nodes
 - Nodes in red color
 - Generated but not yet expanded
- Unexplored Nodes
 - Nodes in blue color
 - Not yet generated



Data Structures for Search Tree

Generation of search tree

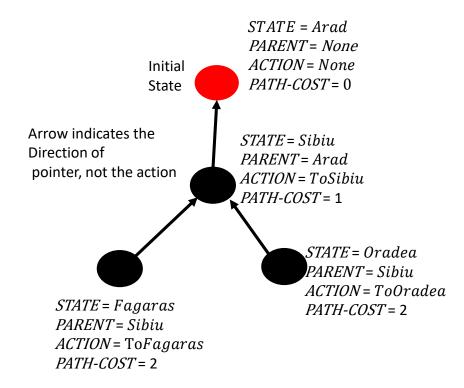
- Representation of
 - Reached nodes
 - Frontier nodes



Generation of Search Tree

- Node (node) contains four components:
 - State Information (node. STATE)
 - Eg: Sibiu
 - Parent Information (node. PARENT)
 - Pointer from child to parent node
 - Need this for backtracking
 - Action (*node*. *ACTION*)
 - Action at parent node that leads to this node
 - Need this for final solution
 - Path-Cost (node.PATH COST)
 - Cost of reaching this node from initial state
 - For checking optimality of a path

Node Data Structure



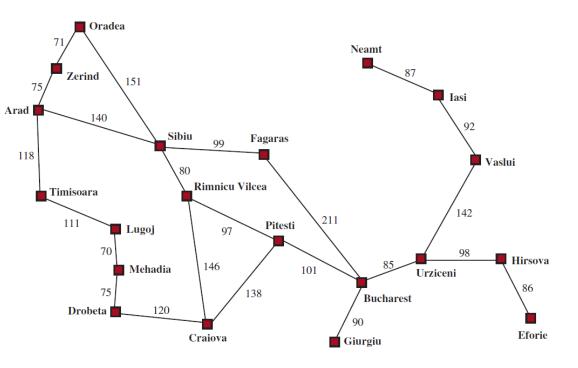
Expanding a node

Node generator function

```
function EXPAND(problem, node) yields nodes
  s \leftarrow node.STATE
  for each action in problem.ACTIONS(s) do
    s' \leftarrow problem.RESULT(s, action)
     cost \leftarrow node.PATH-COST + problem.ACTION-COST(s, action, s')
    yield Node(State=s', Parent=node, Action=action, Path-Cost=cost)
```

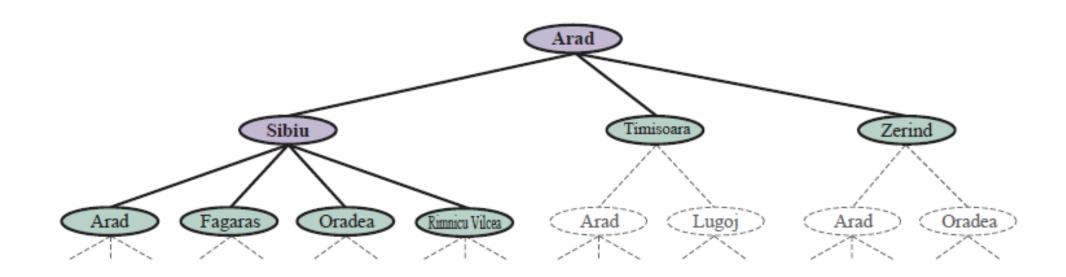
Yields a child node of the given node with state *node*. *STATE* for each call

Expanding a node: Example

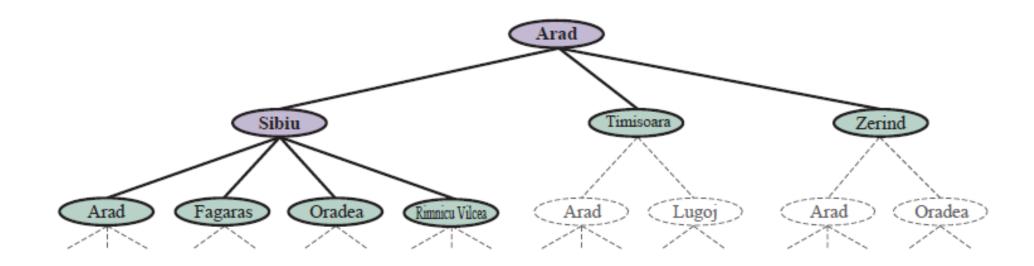


```
function Expand(problem, node) yields nodes s \leftarrow node. State for each action in problem. Actions(s) do s' \leftarrow problem. Result(s, action) cost \leftarrow node. Path-Cost + problem. Action-Cost(s, action, s') yield Node(State=s', Parent=node, Action=action, Path-Cost=cost)
```

Expanding a Node



Node Vs State

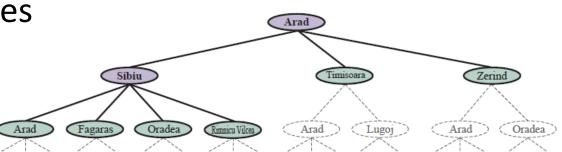


Representation of *reached* Nodes

Need a data structure to store reached nodes

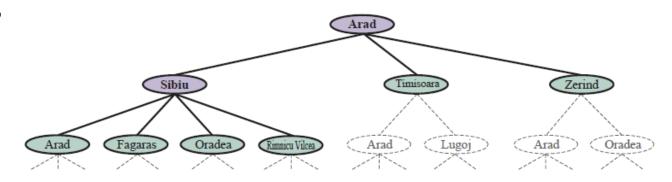
To check whether a node with a state is already generated

Helps in pruning the tree



- Should support quick insert and quick access
- Can be implemented using Hash Table
 - Key-value pairs
 - Key is state
 - Value is node
 - Dictionary in Python

- Need a data structure to store frontier nodes
- Should support following operations
 - Is Empty(frontier)
 - *POP*(frontier)
 - *TOP*(*frontier*)
 - *ADD*(node, frontier)

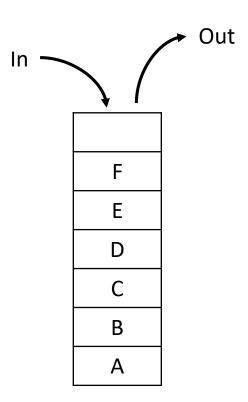


- Implementation depends on search strategy
 - Last-in-First-Out Queue (stack)
 - First-in-First-Out Queue (queue)
 - Priority Queue

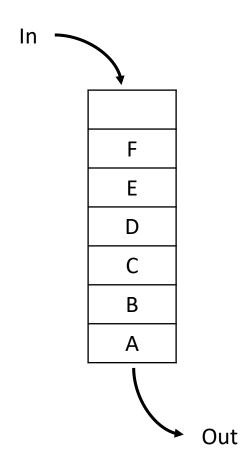
Last-In-First-Out Queue (Stack)



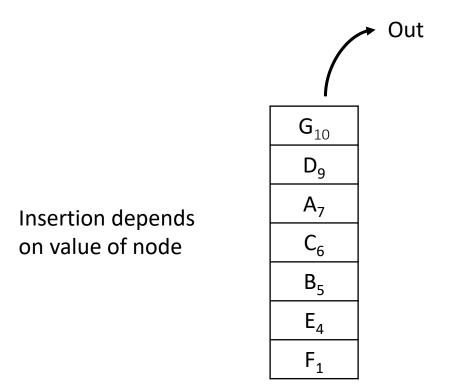
Similar to a stack of plates, adding or removing is only possible at the top.



• First-In-First-Out Queue



- Priority Queue
 - Nodes are ordered based on evaluation of nodes



Building Blocks of Search Tree

Problem:

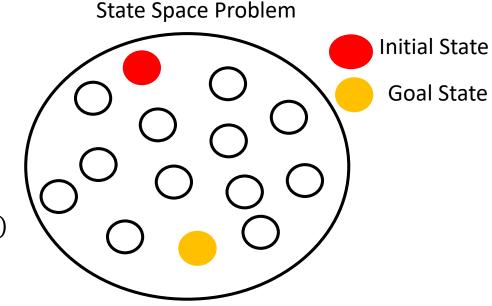
Initial State: problem.INITIAL - STATE

Goal: problem. Is - GOAL(s)

Actions: problem.ACTIONS(s)

Transition Model: *problem*. *RESULT*(*s*, *action*)

Action-Cost Function: problem. ACTION - COST(s, action, s')



Node:

State: node. STATE

Parent: node. PARENT

Path-Cost: node. PATH - COST

Action: node. ACTION

```
\begin{aligned} & \textbf{function EXPAND}(\textit{problem}, node) \textbf{ yields} \text{ nodes} \\ & s \leftarrow node. \text{STATE} \\ & \textbf{for each } \textit{action in problem}. \text{ACTIONS}(s) \textbf{ do} \\ & s' \leftarrow \textit{problem}. \text{RESULT}(s, action) \\ & \textit{cost} \leftarrow node. \text{PATH-Cost} + \textit{problem}. \text{ACTION-Cost}(s, action, s') \\ & \textbf{yield Node}(\text{STATE} = s', \text{Parent} = node, \text{ACTION} = action, \text{Path-Cost} = \textit{cost}) \end{aligned}
```

Agenda

Workflow

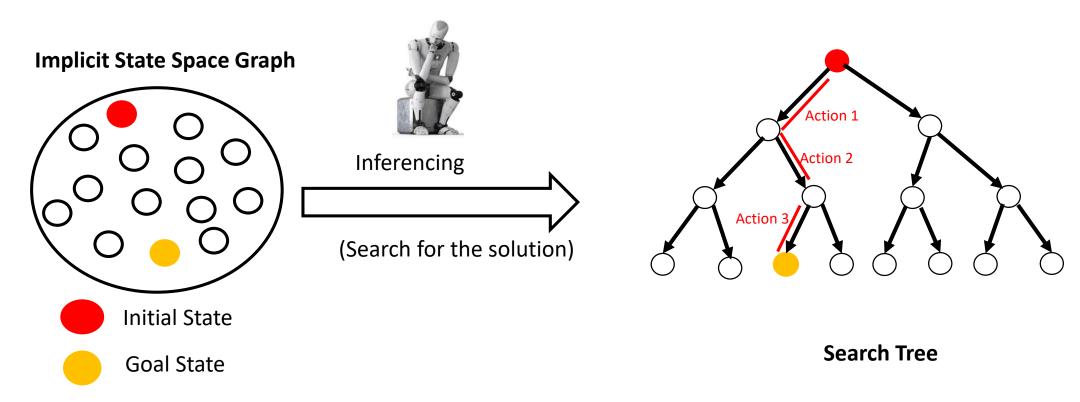
Goal Formulation

Problem Formulation or Modeling

Building Blocks of Search Tree

• Inference

Inference (Search for the solution)



Inferencing:

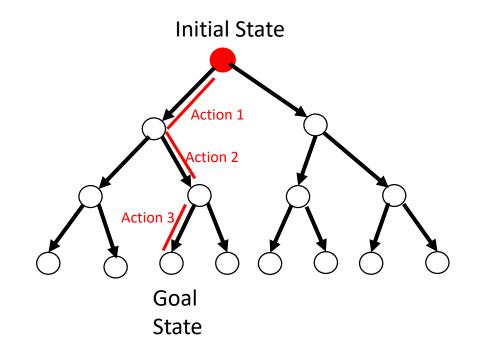
Finding the path from Initial State to Goal State

Solution: Sequence of Actions

[Action 1, Action 2, Action 3]

Uninformed Search Algorithms

- Based on storage of nodes
 - Graph Search
 - Tree Search
- Based on Traversal Strategy
 - Breadth-First Search
 - Depth-First Search
 - Depth-Limited Search
 - Iterative Deepening Search
 - Bidirectional Search
 - Uniform Cost Search



Graph Search vs Tree Search

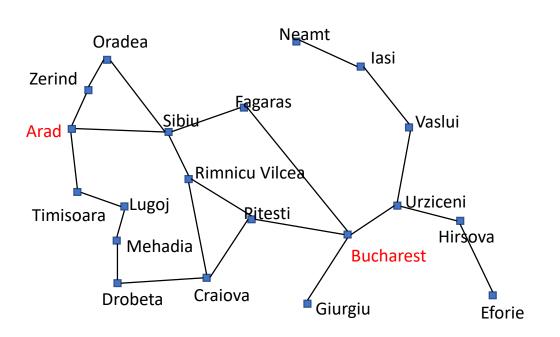
- Graph Search
 - Stores reached nodes
- Tree Search
 - Does not store reached nodes

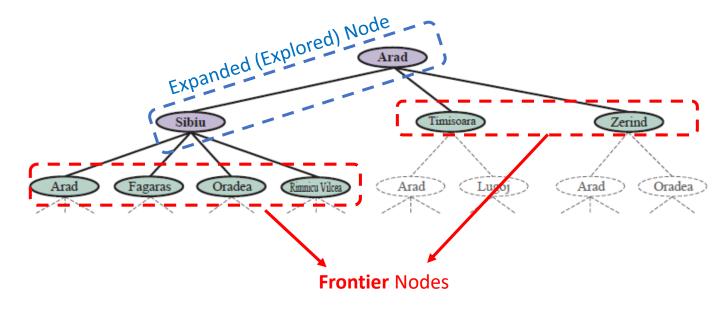
Storing reached nodes?

If model contains

Cycle or Loopy Paths Redundant Paths

Storing reached nodes helps in pruning search tree





Reached Nodes:

Breadth-First Search (BFS)

```
function Breadth-First-Search(problem) returns a solution node or failure
  node \leftarrow Node(problem.INITIAL)
  if problem.Is-GOAL(node.STATE) then return node
  frontier \leftarrow a FIFO queue, with node as an element
  reached \leftarrow \{problem.INITIAL\}
   while not IS-EMPTY(frontier) do
     node \leftarrow Pop(frontier)
    for each child in EXPAND(problem, node) do
       s \leftarrow child.STATE
       if problem.IS-GOAL(s) then return child
       if s is not in reached then
          add s to reached
          add child to frontier
  return failure
```

Trace of BFS

Node	S	Is-Goal(s)	Frontier	Reached

Search Tree of BFS

Performance of BFS

- Complete:
 - 3
- Optimal
 - 5
- Time Complexity
 - [
- Space Complexity
 - [

Depth-First Graph Search (DFGS)

```
function DEPTH-FIRST-SEARCH(problem):
          node \leftarrow NODE(problem.INITIAL-STATE)
          if problem.IS-GOAL(node.STATE) then return node
          frontier ←a LIFO queue (stack) with node as element
          reached \leftarrow {problem.INITIAL-STATE}
          while not IS-EMPTY(frontier) do
                     node ← POP(frontier)
                     for each child in EXPAND(problem, node) do
                                s \leftarrow child.STATE
                                if problem.IS-GOAL(s) then return child
                                if s is not in reached then
                                           add s to reached
                                           add child to frontier
          return failure
```

Trace of DFGS

Node	S	Is-Goal(s)	Frontier	Reached

Search Tree of DFGS

Performance of DFGS

- Complete:
 - 7
- Optimal
 - [
- Time Complexity
 - ?
- Space Complexity
 - [

Depth-First Tree Search (DFTS)

```
function DEPTH-FIRST-TREE-SEARCH(problem):
         node \leftarrow NODE(problem.INITIAL-STATE)
         if problem.IS-GOAL(node.STATE) then return node
         frontier ←a LIFO queue (stack) with node as element
         while not IS-EMPTY(frontier) do
                   node \leftarrow POP(frontier)
                   if problem.IS-GOAL(node.STATE) then return node
                   if not IS - CYCLE (node) do
                             for each child in EXPAND(problem, node) do
                                       add child to frontier
         return failure
```

How to check cycles?

Trace of DFTS

Node	s	Is-Goal(s)	Is-Cycle	Frontier

Search Tree of DFTS

Performance of DFTS

- Complete
 - 7
- Optimal
 - ?
- Time Complexity
 - [
- Space Complexity
 - [

Depth-limited Search (DLS)

```
function DEPTH-LIMITED-SEARCH(problem, \ell) returns a node or failure or cutoff
  frontier \leftarrow a LIFO queue (stack) with NODE(problem.INITIAL) as an element
  result \leftarrow failure
  while not IS-EMPTY(frontier) do
     node \leftarrow Pop(frontier)
    if problem.IS-GOAL(node.STATE) then return node
    if DEPTH(node) > \ell then
       result \leftarrow cutoff
     else if not Is-CYCLE(node) do
       for each child in EXPAND(problem, node) do
         add child to frontier
  return result
```

Trace of DLS

Node	Is-Goal(s)	Depth(Node) > l	Is-Cycle(Node)	Frontier	Result

Search Tree of DLS

Performance of DLS

- Complete
 - 3
- Optimal
 - [
- Time Complexity
 - ;
- Space Complexity
 - [

Iterative Deepening Search (IDS)

```
function Iterative-Deepening-Search(problem) returns a solution node or failure for depth = 0 to \infty do result \leftarrow Depth-Limited-Search(problem, depth) if result \neq cutoff then return result
```

Trace of IDS

Trace of DLS

Node	Is-Goal(s)	Depth(Node) > l	Is-Cycle(Node)	Frontier	Result

Search Tree of IDS

Overhead in IDS

- Number of nodes generated with solution at depth-d:
 - Breadth-First Search

•
$$N(BFS) = 1 + b^1 + b^2 + \dots + b^d$$

- Iterative Depth-Limited Search
 - $N(IDLS) = (d)b + (d-1)b^2 + (d-2)b^3 + \dots + (1)b^d$

Overhead in IDS

- Let b = 10 and d = 5
 - N(BFS) = 111,110
 - N(IDLS) = 123,450

• Overhead =
$$\frac{N(IDLS)-N(BFS)}{N(BFS)}$$
 = 11%

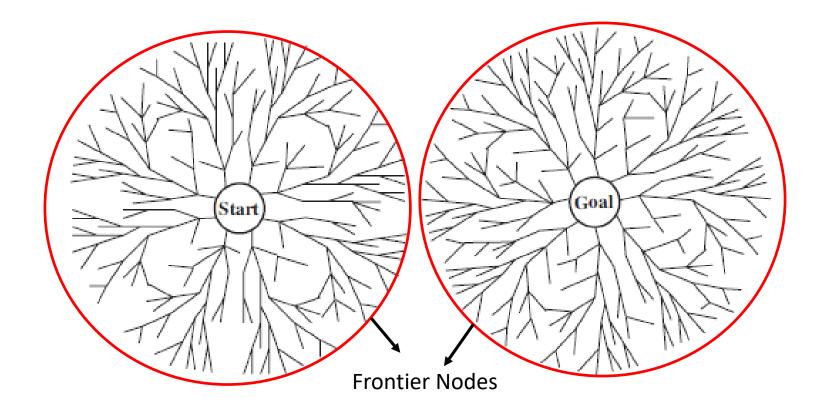
 IDLS: Huge savings in memory with little overhead in number of nodes generated compared to BFS

Performance of IDS

- Complete
 - 3
- Optimal
 - 3
- Time Complexity
 - [
- Space Complexity
 - [

Bidirectional Search

- Idea:
 - Search from initial and goal states
 - If frontiers of both meet, solution is found



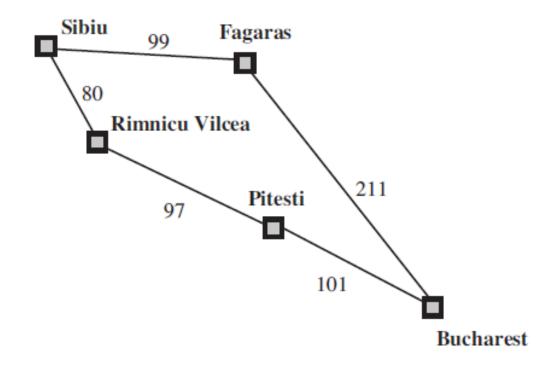
Motivation:

$$b^{\frac{d}{2}} + b^{\frac{d}{2}} < b^d$$

Performance of Bidirectional Search

- Complete
 - ?
- Optimal
 - 3
- Time Complexity
 - }
- Space Complexity
 - [

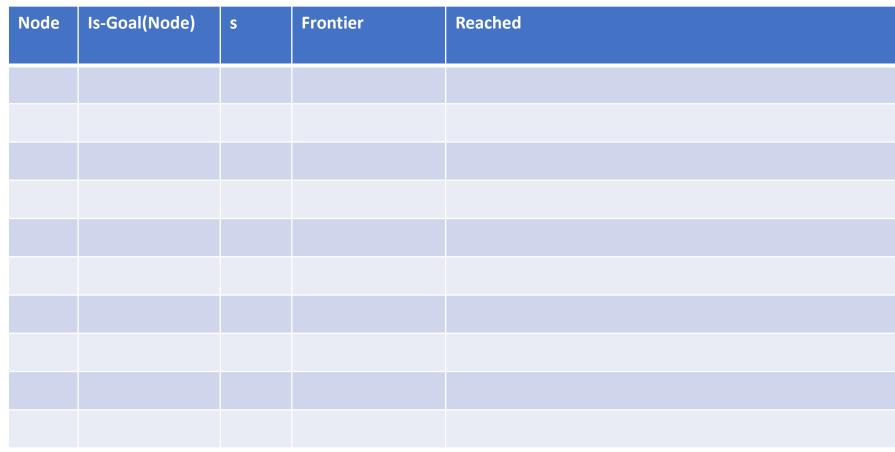
Uniform-Cost Search: Romania Map

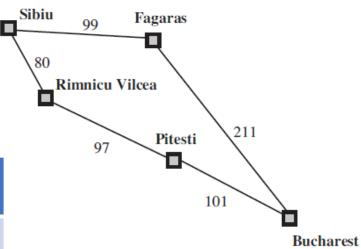


Uniform-Cost Search

```
function UNIFORM-COST-SEARCH(problem) returns a solution node, or failure
  return BEST-FIRST-SEARCH(problem, PATH-COST)
function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure
  node \leftarrow Node(State = problem.Initial)
  frontier \leftarrow a priority queue ordered by f, with node as an element
  reached \leftarrow a lookup table, with one entry with key problem. INITIAL and value node
  while not Is-EMPTY(frontier) do
     node \leftarrow \mathsf{POP}(frontier)
     if problem.Is-Goal(node.State) then return node
     for each child in EXPAND(problem, node) do
       s \leftarrow child.STATE
       if s is not in reached or child.PATH-COST < reached[s].PATH-COST then
          reached[s] \leftarrow child
         add child to frontier
  return failure
```

Trace of UCS

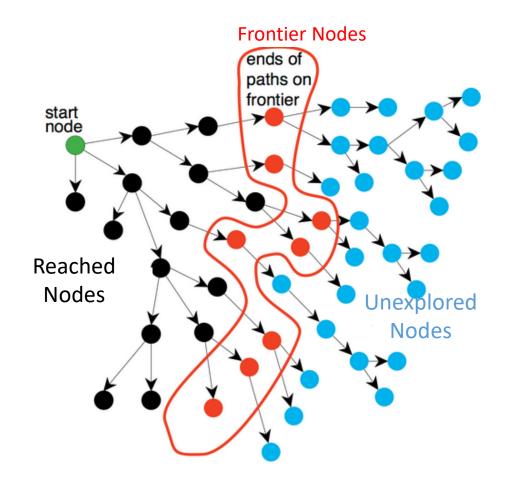




UCS Search Tree

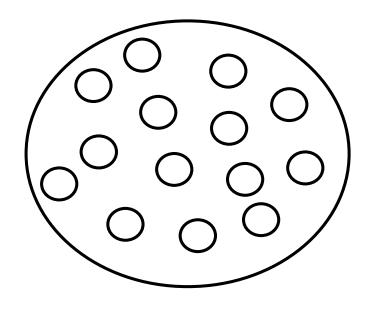
Summary of Uninformed Search

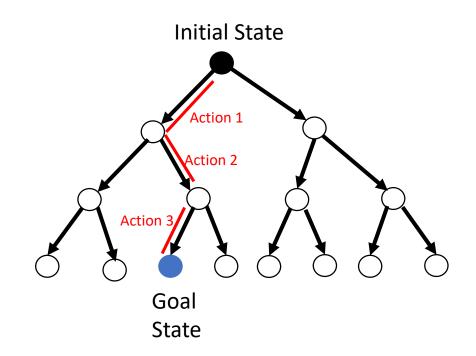
- Classification based on traversal methods
 - Breadth-First Search: Frontier is a Queue (FIFO)
 - Depth-First Search: Frontier is a stack (LIFO)
 - Uniform-Cost Search: Frontier is a priority queue
- Tree-Search Vs Graph-Search
 - Tree-Search
 - Doesn't store reached nodes
 - Leads to cycles and redundant paths
 - Cycles can be avoided with Cycle-Check
 - Graph-Search
 - Maintains a set of reached states
 - Avoids cycles and redundant paths

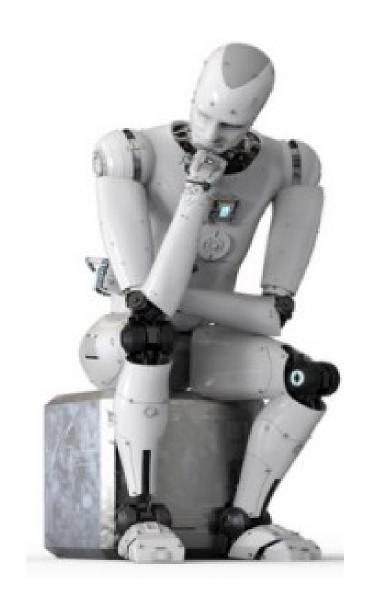


Why is it called uninformed search?

Also called Blind Search or Brute Force Search







When can an agent use uninformed search algorithms?

Design Space

Dimension	Values
Environment	Static, Dynamic
Representation Scheme	States, Features, Relations
Observability	Fully Observable, Partially observable
Parameter Types	Discrete, Continuous
Uncertainty	Deterministic, Stochastic
Learning	Knowledge is given (known), knowledge is learned (unknown)
Number of Agents	Single Agent, Multiple Agent

Need well-defined goal Interaction is offline No limits on resources (memory and time)

Conclusion

Atomic Representation of World

• Implicit representation of state space graphs

Inferencing using search algorithm

Issues related to cycles in state space graphs