

# Quantum state discrimination

## I. CLASSICAL INFORMATION OVER QUANTUM CHANNELS

$\Phi \in \text{Chan}(\mathcal{X}, \mathcal{Y})$  and  $\Psi \in \text{Chan}(\mathcal{Z})$ . The channel  $\Phi$  emulates  $\Psi$  if

$$\exists \Xi_E \in \text{Chan}(\mathcal{Z}, \mathcal{X}) \ \Xi_D \in \text{Chan}(\mathcal{Y}, \mathcal{Z}) \ \Psi = \Xi_D \Phi \Xi_E \quad (1)$$

The channel  $\Xi_E$  is called an encoding channel and  $\Xi_D$  is called a decoding channel.

**Definition 1**  $\Psi_0, \Psi_1 \in \text{Chan}(\mathcal{Z})$ .  $\Psi_0$  is  $\epsilon$ -approximation to  $\Psi_1$  if

$$\|\Psi_0 - \Psi_1\| < \epsilon \quad (2)$$

**Definition 2**  $\Phi \in \text{Chan}(\mathcal{X}, \mathcal{Y})$ .  $\mathcal{Z} = \mathbb{C}^{\{0,1\}}$ .  $\Delta \in \text{Chan}(\mathcal{Z})$  is the completely dephasing channel.

A value  $\alpha > 0$  is achievable if  $\forall \epsilon > 0$ , for all but finitely many  $n \in \mathbb{N}$ , the channel  $\Phi^{\otimes n}$  emulates an  $\epsilon$ -approximation to the channel  $\Delta^{\lfloor \alpha n \rfloor}$ .

The classical capacity of  $\Phi$ , denoted  $C(\Phi)$ , is the supremum value of all achievable rates for classical information transmission through  $\Phi$ .

**Remark 1** When considering an emulation of the  $m$ -fold tensor product  $\Delta^{\otimes m}$  of this ideal classical channel by the channel  $\Phi^{\otimes n}$ , no generality is lost in restricting one's attention to classical-to-quantum encoding channels  $\Xi_E$  and quantum-to-classical decoding channels  $\Xi_D$ .

$$\Xi_E = \Xi_E \Delta^{\otimes m} \quad \Xi_D = \Delta^{\otimes m} \Xi_D \quad (3)$$

**Theorem 1**  $\Phi \in \text{Chan}(\mathcal{X}, \mathcal{Y})$

$$C(\Phi^{\otimes k}) = k C(\Phi) \quad (4)$$

**Definition 3** Entanglement-assisted classical capacity

$$Z \mapsto (\Xi_D(\Phi \Xi_E \otimes \mathbb{1}_{\text{Lin}(\mathcal{W})}))(Z \otimes \xi) \quad (5)$$

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The entanglement-assisted classical capacity of  $\Phi$ , denoted  $C_E(\Phi)$ , is the supremum value of all achievable rates for classical information transmission through  $\Phi$ .

**Theorem 2**  $\Phi \in \text{Chan}(\mathcal{X}, \mathcal{Y})$

$$C_E(\Phi^{\otimes k}) = k C_E(\Phi) \quad (6)$$

**Definition 4**  $\Phi \in \text{Chan}(\mathcal{X}, \mathcal{Y})$ . The Holevo capacity of  $\Phi$  is defined as

$$\chi(\Phi) = \sup_{\eta} \chi(\Phi(\eta)) \quad (7)$$

**Theorem 3** It is enough to maximize Holevo information using pure states. And what's more, the size of the alphabet of the ensemble can be restricted to  $\dim(\mathcal{X})^2$ .