A SECOND LOOK AT COMPLEX EVENT PROCESSING

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http://siddhi.sourceforge.net/





OUTLINE



- Introduction to CEP
- Second look at CEP Implementations
- Siddhi Architecture
- Siddhi Performance
- Conclusion and Future Topics





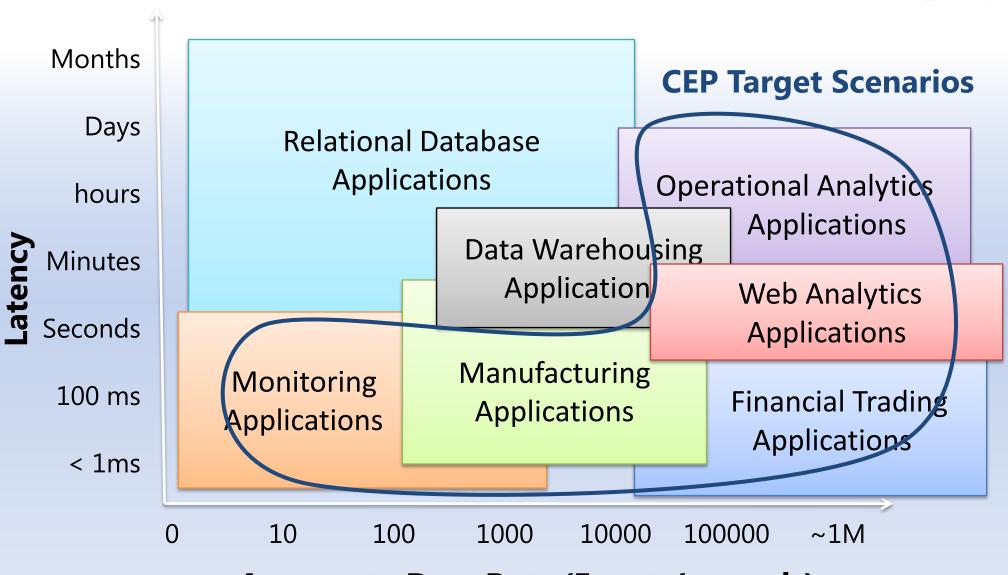
WHAT IS CEP?

Complex Event Processing (CEP) is identifying meaningful patterns, relationships and data abstractions among unrelated events and fire an immediate response.

	Database Applications	Event-driven Applications
Query Paradigm	Ad-hoc queries or requests	Continuous standing queries
Latency	Seconds, hours, days	Milliseconds or less
Data Rate	Hundreds of events/sec	Tens of thousands of events/sec or more
	request	output stream stream



SCENARIOS OF EVENT PROCESSING



Aggregate Data Rate (Events/seconds)





- E-Science often deals with National and Global scale usecases while we try to understand the world around us better.
- Following is an important class of E-science applications
 - Receives data about the world around us from sensors deployed across the country/world
 - Try to make sense, react to, predict, and/or control the world around us
 - Examples: Weather, Traffic Data, Surveillance, Smart Grid etc...
- CEP is a powerful enabling technologies for these usecases



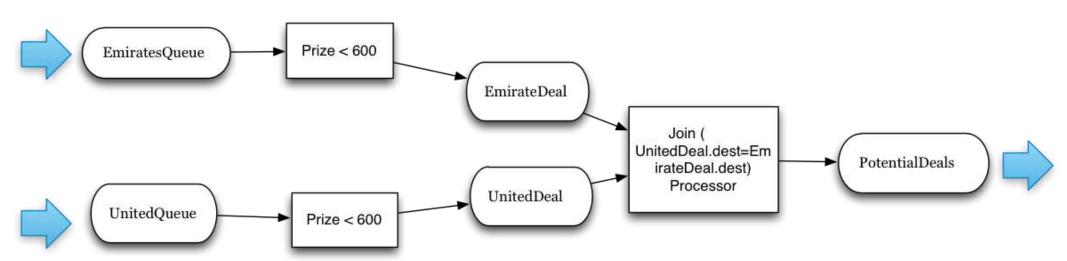
CEP QUERIES

- Each Event consists of properties (name value pairs)
- We separate different events to streams
- Use a SQL like query language, but queries are evaluated on continuous event streams
- Types of queries are following
 - Selection (filtering) and projection (e.g. like select in SQL)
 - Windows events are processed within a window (e.g. for aggregation and joins). Time window vs. length window.
 - Ordering sequences and patterns (before, followed by conditions e.g. new location followed by small and a large purchase might suggest a fraud)
 - Join and split
 - Aggregation



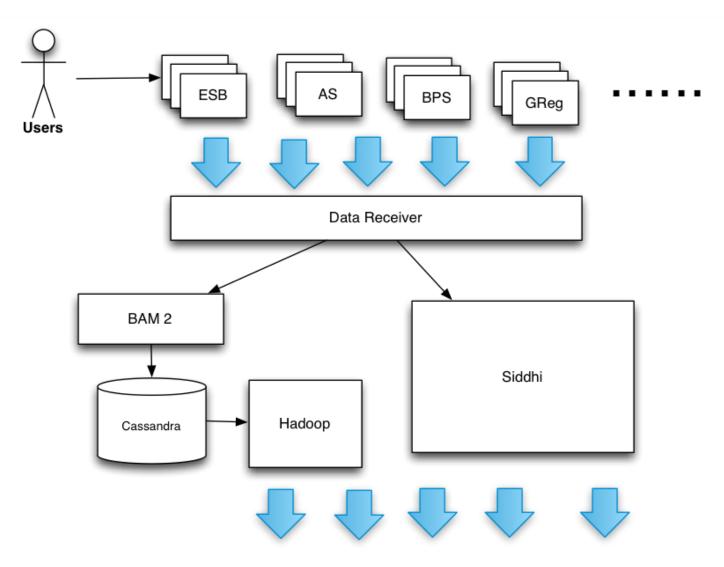
EXAMPLE QUERY

EmirateDeal= Select * from Emirate where prize < 600
UnitedDeal = Select * from United where prize < 600
PotentialDeal = Select * from EmirateDeal join UnitedDeal
where UnitedDeal .dest=EmirateDeal .dest</pre>





CEP ROLE WITHIN WSO2 PLATFORM



React to Data collected

MOTIVATING USE CASE: LOS ANGELES SMART GRID DEMONSTRATION PROJECT

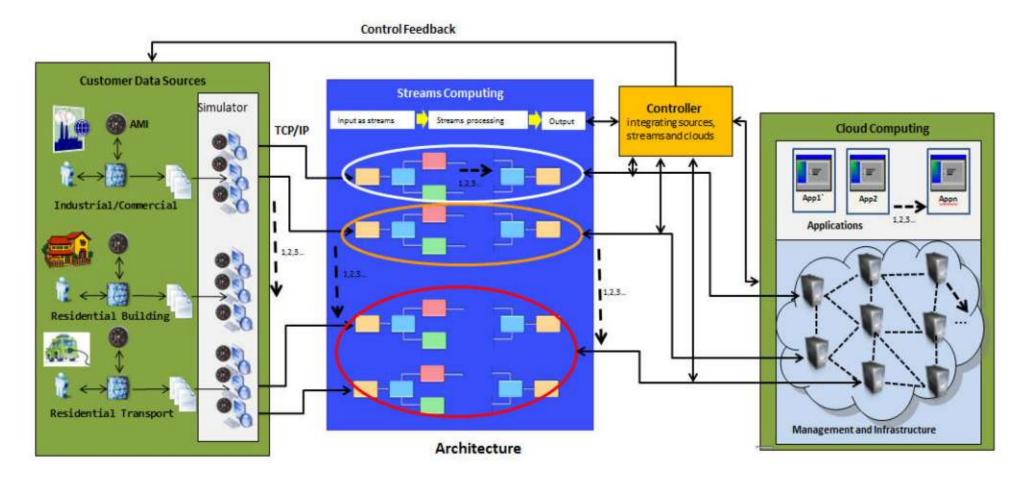


Figure 2: Design for adaptive rate control feedback from stream processing system to AMIs to optimize bandwidth into Cloud platform

RELATED WORK ON CEP IMPLEMENTATIONS

- A very good reference: "G. Cugola and A. Margara.
 Processing flows of information: From data stream to complex event" processing. ACM Computing Surveys, 2011.
- Initial work from Databases, active databases. But the "store and process" model was too slow.
- Stream processing uses a pub/sub model, very scalable but lacks temporal support
 - (Aurora [5, 2], PIPES [23, 7], STREAM [30], Borealis [6,1] S4 [25], and Storm.)
- CEP engines
 - Mainly uses a NFA based model
 - SASE (data flow and NFA)
 - Esper NFA and Delta networks
 - Cayuga NFA and re-subscriptions (single thread)

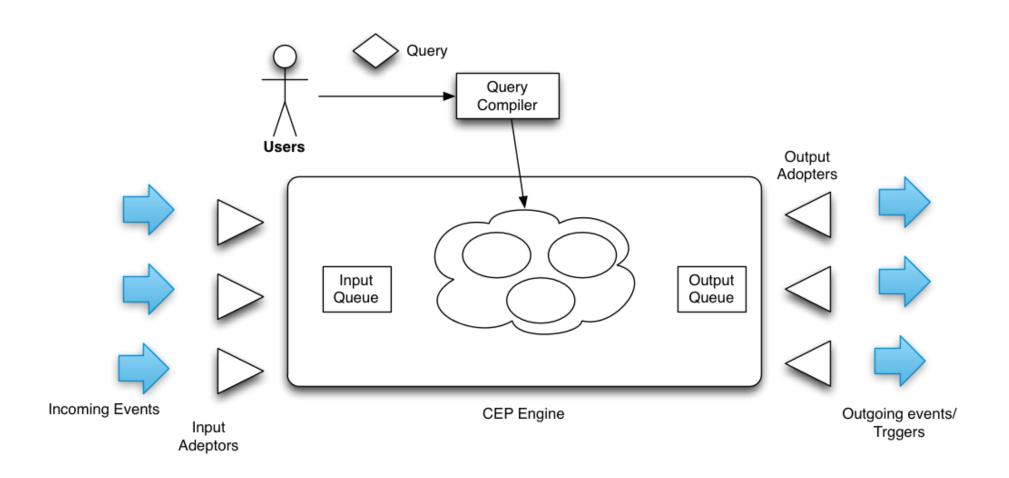


GOAL: IDEAS FOR IMPROVEMENT

- Borrowing ideas from Stream processing: Pipeline based architecture with a pub/sub model inspired by Stream processing.
- Breaking single thread evaluations: Improving parallelism in processing through a pipeline based model
- A New State machine implementation that avoids checking a series of conditions
- Improved support for query chaining with complex queries

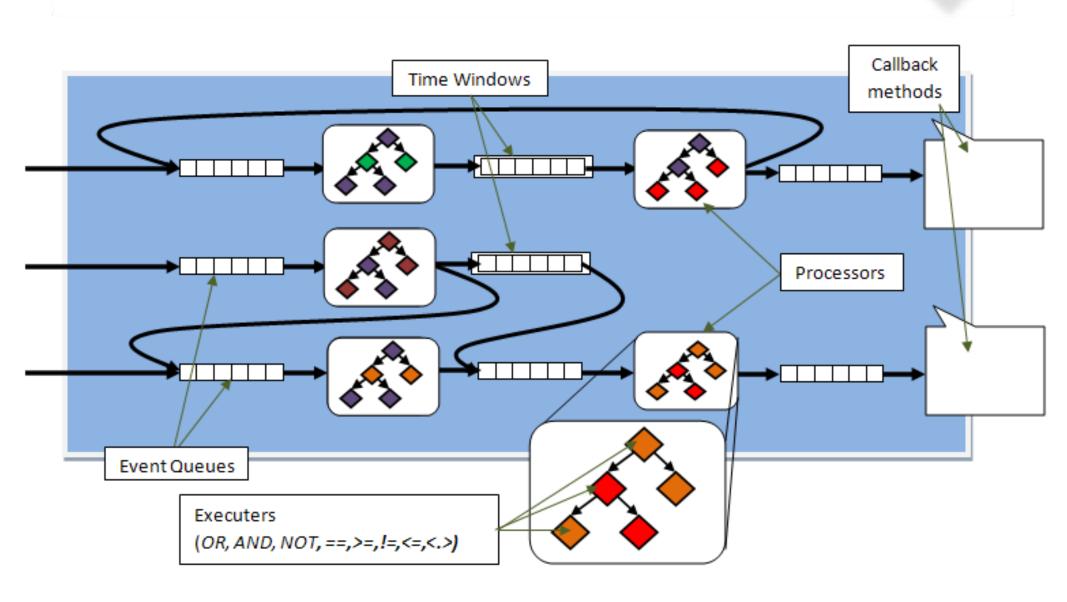


SIDDHI HIGH LEVEL ARCHITECTURE





SIDDHI PROCESSING IMPLEMENTATION





SIDDHI PROCESSING IMPLEMENTATION

- Uses a pipeline architecture
- Consist of a pipeline of processors (support all other operators except windows)
- Processors are connected by queues (they implement windows)
- Processors take inputs from a queue by subscribing to them and placing results in a output queue
- Many threads are assigned to processors, where they take data from input queues, process them, and place them in output queues.





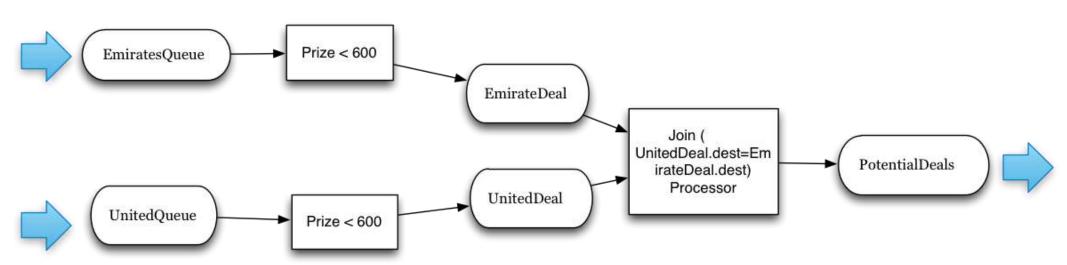
API TO WRITE QUERIES

```
//Instantiate SiddhiManager
SiddhiManager siddhiManager = new SiddhiManager();
//Get the QueryFactory to create gueries
QueryFactory qf = siddhiManager.getQueryFactory();
InputEventStream stockStream = new InputEventStream(
    "StockStream",
    new String[]{"symbol", "price", "volume"},
    new Class[]{String.class, Float.class, Long.class}
//Setting a time window of 1 hour
stockStream.setTimeWindow(3600000);
//Assign input stream
siddhiManager.addInputEventStream(stockStream);
//Create Query
Query query = qf.createQuery(
    "StockQuote",
    qf.output("symbol=StockStream.symbol","price=avg(StockStream.price)"),
    qf.inputStream(stockStream),
    qf.and(
            qf.condition("StockStream.symbol", EQUAL, "IBH"),
            qf.condition("StockStream.volume", GREATERTHAN, "10000")
//Assign query
siddhiManager.addQuery(guery);
```



EXAMPLE QUERY IMPLEMENTATION

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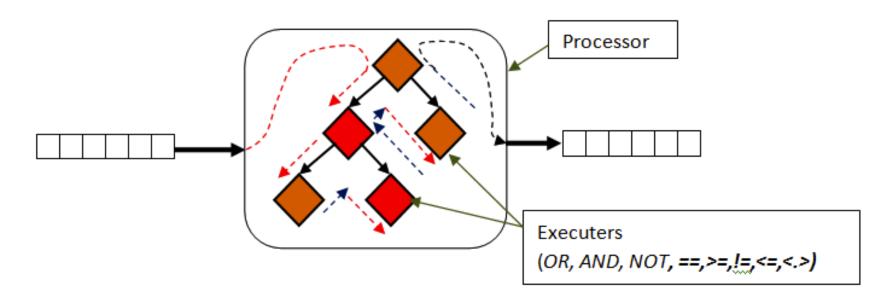






IMPLEMENTING FILTERS

 Evaluate the tree, and optimizations to stop the evaluations as soon as possible

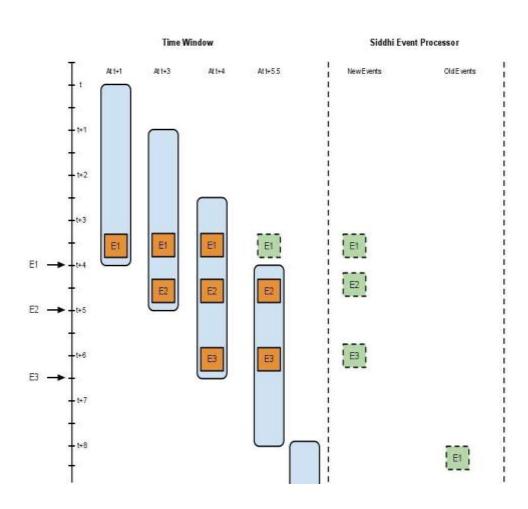






IMPLEMENTING WINDOWS

- Implemented as a queue
- Use events to notify the windows event addition and removals using a pub/sub pattern to processors





AGGREGATION

Siddhi supports avg, sum, count, max, min

```
Query query = qf.createQuery(

"StockQuote",
    qf.output("symbol=CSEStream.symbol", "avgPrice=avg(CSEStream.price)"),
    qf.inputStream(cseEventStream),
    qf.condition("CSEStream.price", GREATERTHAN_EQUAL, "20")
);
query.groupBy("CSEStream.symbol");
query.having(qf.condition("StockQuote.avgPrice", GREATERTHAN, "50"));

Notify Processor when
events are added or removed

Aggregation
Processor

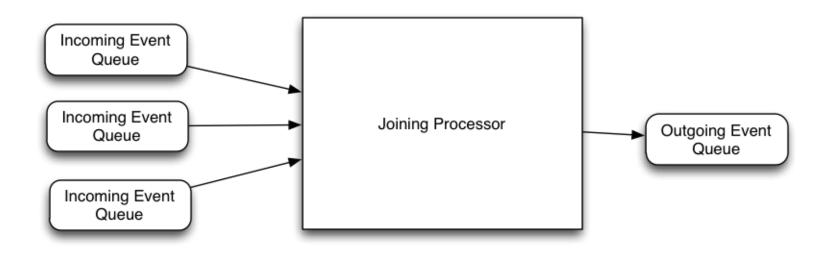
EmirateDeal
```

• These are implemented within processor and each addition or removal of a event updates the value (if it is a window based, queues notify subscribers when thing change).





IMPLEMENTING JOINS



- Implemented as a processor that moves data from several queues to a different queue after joining
- When a event arrives, if it matches the join condition, we send them over to the next queue, else keeps them to match with the future events.
- We keep them as long as the window condition allows (this is done by keeping references to queues, and checking within them for matching)



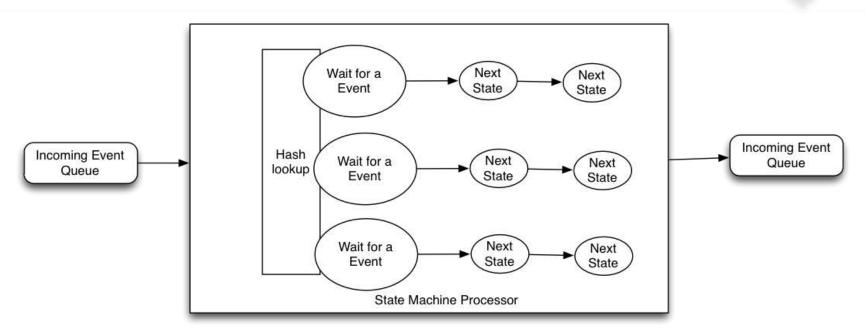


PATTERNS AND SEQUENCES

- Patterns and sequences handle series of events like A->B->C, there are two classes patterns and sequences
- Patterns matches with events ignoring events in between while sequences matches exactly. You can use star (*) with sequences to mimic patterns
- Supported through a state machine (NFA)



IMPLEMENTING PATTERNS AND SEQUENCES



- Use a event model
- Say we need A->B->C, then we break down it to list of listeners, which when triggered remove itself and registers the next in the list.
- We apply different variations of this to support *, and sequences.





PERFORMANCE RESULTS

- We compared Siddhi with Esper, the widely used opensource CEP engine
- For evaluation, we did setup different queries using both systems, push events in to the system, and measure the time till all of them are processed.
- We used Intel(R) Xeon(R) X3440 @2.53GHz, 4 cores 8M cache 8GB RAM running Debian 2.6.32-5-amd64 Kernel

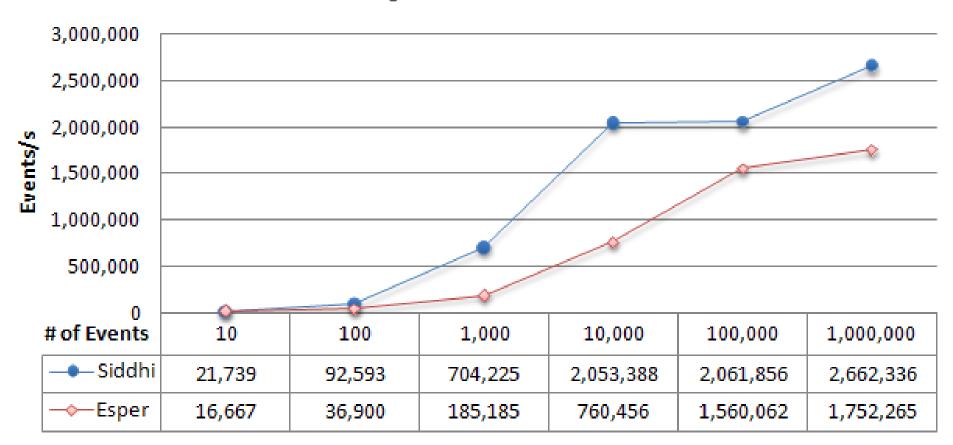




PERFORMANCE COMPARISON WITH ESPER

Simple filter without window

select symbol, price
from StockTick(price>6)



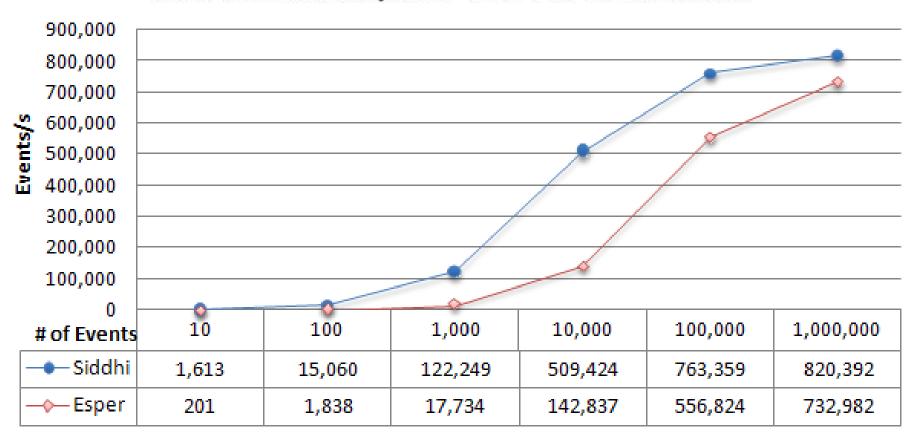




PERFORMANCE COMPARISON WITH ESPER

Average calculation of a given symbol within time window

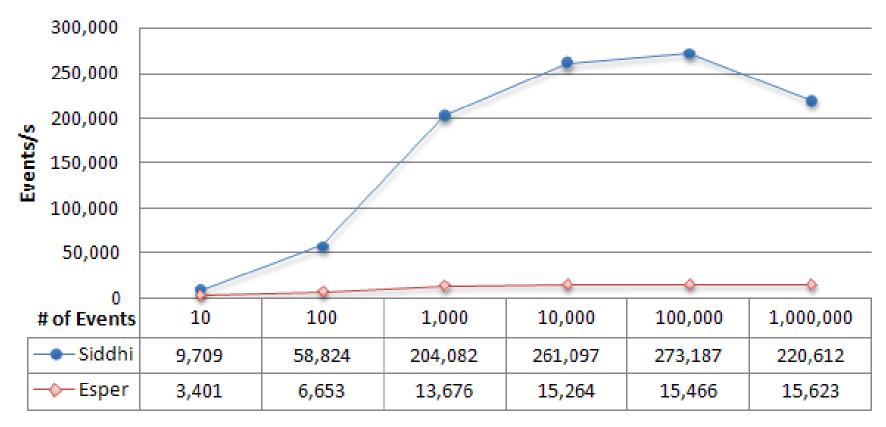
select irstream symbol, price, avg(price)
from StockTick(symbol='IBM').win:time(.005)





PERFORMANCE COMPARISON WITH ESPER

State machine query for pattern matching





SIDDHI vs. ESPER (FUNCTIONALITY)

- Following describe Siddhi in contrast to Esper using the comparisons defined in Cugola et al.
- In terms of Functional and Processing Models and Data, Time, and Rule Models (see Table I and III of Cugola et al.), Esper and Siddhi behave the same.
- In terms of Supported language model (table IV),
 - Esper supports Pane and Tumble windows, User Defined operators and Parameterization while Siddhi does not support them
 - Siddhi supports removal of duplicates that is not supported by Siddhi.
- Both support all other language constructs.



FEATURES WE DID NOT DISCUSSED

- Improvements in the Selection Operator
 - using MVEL
- Support conditions on the aggregation functions
 - having, group-by
- Event quantifications in sequential patterns
 - kleene closure (*)
- Unique function



CURRENT USERS OF SIDDHI

- "Los Angeles Smart Grid Demonstration Project"
 - It forecasts electricity demand, respond to peak load events, and improves sustainable use of energy by consumers.
 - http://ceng.usc.edu/~simmhan/pubs/simmhan-usctr2011-smartgridinformatics.pdf
- Open MRS NCD module <u>_</u> idea is to detect and notify patient when certain conditions have occurredhttps://wiki.openmrs.org/display/docs/Notifiable+Condition+Detector+%28NCD%29+Module
- WSO2 CEP Server (Perspective)





FUTURE WORK



- SQL like query language for Siddhi
- Building scalable processing network using Siddhi nodes
 - Siddhi is currently a java library, build a Thrift based server for Siddhi
 - Supporting a network of Siddhi processors (clustering)
 - Query partition and optimization for scale
- Support for Pane and Tumble windows, User Defined operators and Parameterization
- Available under Apache License

http://www.flickr.com/photos/garryknight/3650151941/



IMMEDIATE PLANS



- Have resource allocations and development efforts to continue next year
- Looking forward to build a opensource community around it
- Looking for usecases, real data feeds, and of course Guinea pigs



QUESTIONS?

