

A SIMPLE FPTAS FOR COUNTING EDGE COVERS

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Overview

1 INTRODUCTION

2 OUR RESULT

Edge cover

Definition

For an undirected input graph $G = (V, E)$, an **edge cover** of G is a set of edges C covering all vertices.

Example

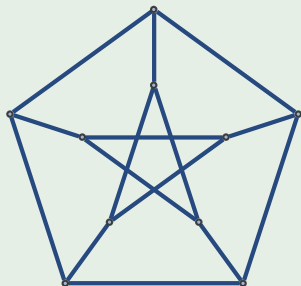


Figure : An edge cover for Petersen graph

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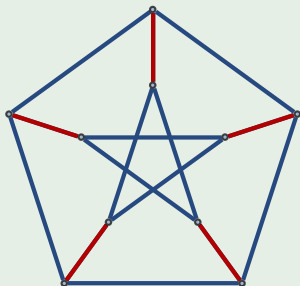


Figure : An edge cover for Petersen graph, with edges chosen being highlighted in red. Note that this is also a perfect matching.

Edge cover is related to many other problems such as:

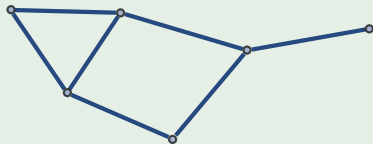
- Matching problem.
- Rtw-Mon-CNF. (read twice monotone CNF)
- Holant problem.
-

Relation to Matching

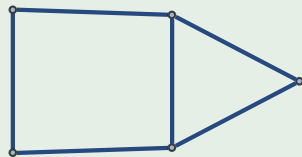
The minimal edge cover could be found via a greedy algorithm based on a maximum matching.

Example

Find edge covers by maximal matching?



(a) G has a perfect matching.



(b) G doesn't have a perfect matching.

Remark

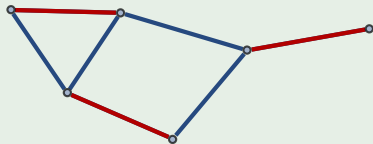
For a graph with a perfect matching, enumerating (sampling) perfect matchings is equivalent to enumerating (sampling) minimum edge covers.

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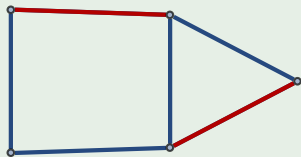
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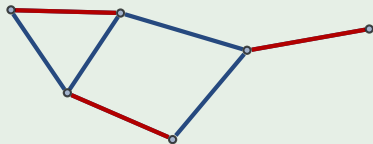
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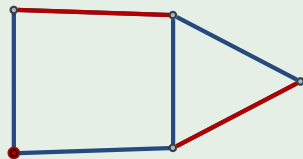
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Relation to Rtw-Mon-CNF

Consider the following graph, its edge covers are exactly satisfying assignments to the following CNF formula, if we treat edges as variables, and vertices as clauses.

$$\phi = (e_1 \vee e_2 \vee e_3) \wedge (e_1 \vee e_4) \wedge (e_4 \vee e_5 \vee e_2) \wedge (e_3 \vee e_5).$$

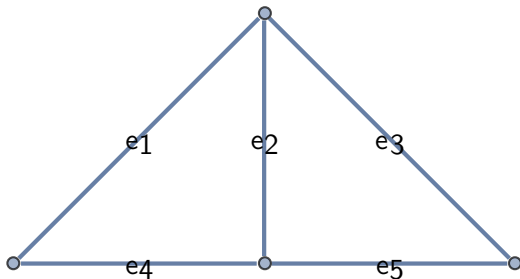


Figure : Graph representation for ϕ .

Counting Problems

A list of problems in their search, optimization, and counting versions.

Search problems:

- SAT.
- Find a (perfect) matching.
- Find an edge cover.
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Optimizations:

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Why counting?

Besides theoretical computer science, counting problems are also related to many problems from other discipline such as:

- Partition function of Statistical physics.
- Graph polynomials.
- Sampling, learning and inference.
- Query evaluations of probabilistic database.
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Approximate Counting

Many interesting problems in the exact counting regimes, including counting edge cover, is hard ($\#P$ -complete). Instead we look for these two types of polynomial time approximation scheme:

Definition (FPTAS)

For given parameter $\varepsilon > 0$ and an instance of a particular problem class, if the algorithm outputs a number \hat{N} such that $(1 - \varepsilon)N \leq \hat{N} \leq (1 + \varepsilon)N$, where N is the accurate answer of the problem instance, and the running time is bounded by $\text{poly}(n, 1/\varepsilon)$ with n being the size of instance, this is called the **FPTAS (fully polynomial time approximation scheme)**.

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Main Result

Previous work:

- Only an MCMC-based FPRAS is known for counting edge covers in graphs with maximum degree 3.
- The correlation decay based FPTAS for anti-ferromagnetic 2-spins systems (e.g. counting independent sets) goes beyond the best known MCMC based FPRAS and achieves the boundary of approximability.

Our main result is a correlation based FPTAS for counting edge covers for general graphs, which provides another example where the tractable range for correlation decay based FPTAS exceeds the sampling based FPRAS.

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Dangling instance

Definition

A **dangling edge** $e = (u, -)$ of a graph is such singleton edge with exactly one end-point vertex u , as shown in the Figure 4a.

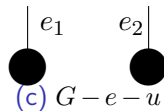
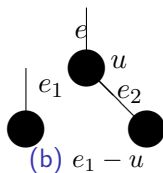
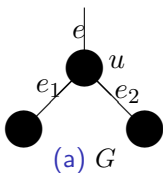


Figure : Dangling edges examples.

Counting v.s. Marginal Probability

Recall our counting problem:

Problem

Let $EC(G)$ be the set of edge covers.

Goal: estimate $|EC(G)|$.

Let X be an edge cover sampled uniformly from $EC(G)$, consider the following marginal probability:

for an edge e , we write $P(G, e) \triangleq \Pr(e \notin X)$.

Solution: estimate $P(G, e)$.

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Counting from marginal probability

Recall that the set of all edges E is an edge cover. For a randomly sampled edge cover X , what is $\Pr(X = E)$?

$$\Pr(X = E) = \frac{1}{|EC(G)|}$$

Therefore,

$$\frac{1}{|EC(G)|} = \prod_i (1 - P(G_i, e_i)).$$

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Estimating Marginal Probability

We focus on e is dangling.

$$P(G, e) = \frac{1 - \prod_{i=1}^d P(G_i, e_i)}{2 - \prod_{i=1}^d P(G_i, e_i)},$$

where $G_1 \triangleq G - e - u$, and $G_{i+1} \triangleq G_i - e_i$.

Truncate with a modified recursion depth:

$$P(G, e, L) = \begin{cases} \frac{1}{2}, & \text{if } L \leq 0; \\ \frac{1 - \prod_{i=1}^d P(G_i, e_i, L - \lceil \log_6(d+1) \rceil)}{2 - \prod_{i=1}^d P(G_i, e_i, L - \lceil \log_6(d+1) \rceil)}, & \text{otherwise.} \end{cases}$$

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Proposition

Given graph G , edge e and depth L ,

$$|P(G, e, L) - P(G, e)| \leq 3 \cdot \left(\frac{1}{2}\right)^{L+1}$$

THANK YOU!