A Simple FPTAS for Counting Edge Covers

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Overview

Edge cover

Definition

For an undirected input graph $G = (V, E)$, an **edge cover** of G is a set of edges C covering all vertices.

Example

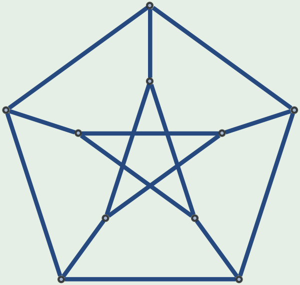

Figure : An edge cover for Petersen graph

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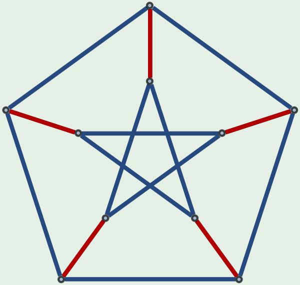


Figure : An edge cover for Petersen graph, with edges chosen being highlighted in red. Note that this is also a perfect matching.

Edge cover

Edge cover is related to many other problems such as:

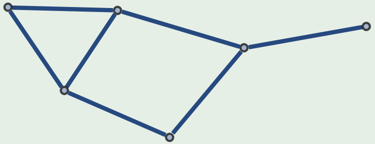
- Matching problem.
- Rtw-Mon-CNF. (read twice monotone CNF)
- Holant problem.
- ....

Relation to Matching

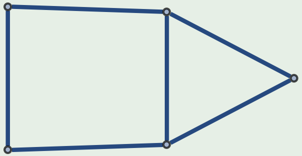
The minimal edge cover could be found via a greedy algorithm based on a maximum matching.

Example

Find edge covers by maximal matching?



(a) G has a perfect matching.



(b) G doesn't have a perfect matching.

Remark

For a graph with a perfect matching, enumerating (sampling) perfect matchings is equivalent to enumerating (sampling) minimum edge covers.

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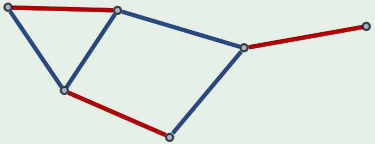

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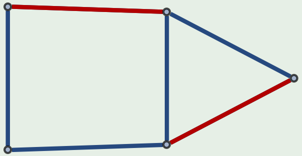

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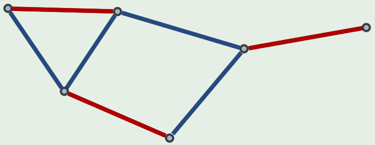

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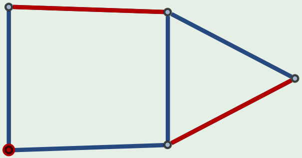

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Relation to Rtw-Mon-CNF

TBA.

Counting Problems

A list of problems in their search, optimization, and counting versions.

Search problems:

- SAT.
- Find a (perfect) matching.
- Find an edge cover.
- ....

Optimizations:

- MAX-SAT.
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Counting problems:

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Besides theoretical computer science, counting problems are also related to many problems from other discipline such as:

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Approximate Counting

Many interesting problems in the exact counting regimes, including counting edge cover, is hard (#P-complete). Instead we look for these two types of polynomial time approximation scheme:

Definition (FPTAS)

For given parameter $\varepsilon > 0$ and an instance of a particular problem class, if the algorithm outputs a number $\hat{N}$ such that $(1-\varepsilon)N \leq \hat{N} \leq (1+\varepsilon)N$, where N is the accurate answer of the problem instance, and the running time is bounded by $poly(n, 1/\varepsilon)$ with n being the size of instance, this is called the **FPTAS (fully polynomial time approximation scheme)**.

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Main Result

Previous work:

- Only an MCMC-based FPRAS is known for counting edge covers in graphs with maximum degree 3.
- The correlation decay based FPTAS for anti-ferromagnetic 2-spins systems (e.g. counting independent sets) goes beyond the best known MCMC based FPRAS and achieves the boundary of approximability.

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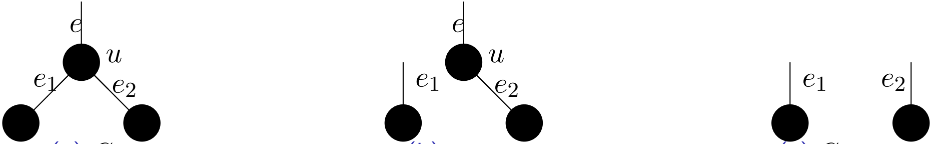
Definition

A **dangling edge** $e = (u, \_)$ of a graph is such singleton edge with exactly one end-point vertex u, as shown in the Figure 3a.



(a) G (b) $e_1 - u$ (c) $G - e - u$

Figure : Dangling edges examples.

Counting v.s. Marginal Probability

Recall our counting problem:

Problem

Let $EC(G)$ be the set of edge covers.
Goal: estimate $|EC(G)|$.

Let X be an edge cover sampled uniformly from $EC(G)$, consider the following marginal probability:
for an edge e, we write $P(G,e) \triangleq \Pr(e \notin X)$.

Solution: estimate $P(G,e)$.

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Counting from marginal probability

Recall that the set of all edges E is an edge cover. For a randomly sampled edge cover X, what is $\Pr(X = E)$?

$$\Pr(X = E) = \frac{1}{|EC(G)|}$$

where $G_1 = G, G_{i+1} = G_i - e_i - u_i - v_i$.
Therefore,

$$\frac{1}{|EC(G)|} = \prod_i (1 - P(G_i, e_i)) .$$

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Estimating Marginal Probability

We focus on e is dangling.

$$P(G,e) = \frac{1 - \prod_{i=1}^{d} P(G_i, e_i)}{2 - \prod_{i=1}^{d} P(G_i, e_i)},$$

where $G_1 \triangleq G - e - u$, and $G_{i+1} \triangleq G_i - e_i$.

Truncate with modified recursion depth:

$$P(G,e,L) = \frac{1 - \prod_{i=1}^{d} P(G_i, e_i, L - \lceil \log_6 (d+1) \rceil)}{2 - \prod_{i=1}^{d} P(G_i, e_i, L - \lceil \log_6 (d+1) \rceil)},$$

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Correlation Decay

Proposition

Given graph G, edge e and depth L,

$$|P(G,e,L) - P(G,e)| \leq 3 \cdot (\frac{1}{2})^{L+1}$$

Thank you!