### The Ising Partition Function: Zeros and Deterministic Approximation

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### Spin systems (aka MRFs or graphical models)

We focus on two-state systems only:

- An undirected graph (or hypergraph) G = (V, E). Let n = |V|, m = |E|.
- Configuration:  $\sigma \in \Sigma$ , where  $\Sigma := \{+, -\}^V$ .
- Edge potentials:  $\varphi_e: \Sigma \times \Sigma \times \cdots \times \Sigma \to \mathbb{C}$ . W.l.o.g.  $\varphi_e(-, \dots, -) = 1$ .
- Vertex potentials:  $\psi_{\nu}: \Sigma \to \mathbb{C}$ . W.l.o.g.  $\psi_{\nu}(+) = \lambda, \psi_{\nu}(-) = 1$ .

#### Definition (Partition function)

$$\begin{split} Z_{G}^{\boldsymbol{\varphi}}(\lambda) &= \sum_{\boldsymbol{\sigma}: V \to \{+, -\}} \underbrace{\prod_{e \in E} \varphi_{e}(\boldsymbol{\sigma}\big|_{e}) \prod_{v \in V} \psi(\boldsymbol{\sigma}(v))}_{\text{weight of configuration } \boldsymbol{\sigma}} \\ &= \sum_{\boldsymbol{\sigma}: V \to \{+, -\}} \prod_{e \in E} \varphi_{e}(\boldsymbol{\sigma}\big|_{e}) \, \lambda^{|\{v: \boldsymbol{\sigma}(v) = +\}|} \end{split}$$

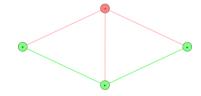
### Example: Ising model on graphs

For  $\beta, \lambda \in \mathbb{R}_+$ ,

- Configuration:  $\sigma \in \{+, -\}^V$
- Edge potentials:  $f_e = \begin{pmatrix} 1 & \beta \\ \beta & 1 \end{pmatrix}$  related to the "temperature"
- Vertex potentials:  $f_{
  m v} = inom{\lambda}{1}$  "external field"

### Ising model

$$Z_G^{\beta}(\lambda) = \sum_{S \subseteq V} \beta^{|E(S,V \setminus S)|} \lambda^{|S|}$$



A spin configuration with weight  $\beta^3 \lambda^3$ 

- $\beta$  < 1: Ferromagnetic; the model favors small cuts
- $\beta >$  1: Anti-ferromagnetic; the model favors large cuts

### Approximating the partition fucntion

We will be interested in multiplicative approximation of Z

### Definition (Fully polynomial-time approximation scheme)

An FPTAS for a function  $f(\cdot)$  is an algorithm with:

- Input:  $\varepsilon > 0, \mathbf{x}$
- Output:  $\widetilde{f(x)}$  such that  $\left|f(x) \widetilde{f(x)}\right| \le \varepsilon |f(x)|$
- Running time:  $poly(|\mathbf{x}|, 1/\varepsilon)$

© For self-reducible problems, this notion of approximability is robust.

### Antiferromagnetic Ising model: fully understood

# Theorem (Weitz, Sinclair-Srivastava-Thurley, Li-Lu-Yin, Sly-Sun, Galanis-Stefankovic-Vigoda)

For any  $\beta > 1$ ,  $\lambda > 0$ , there is a threshold  $\beta_c(\lambda, d)$  s.t.

- If  $\beta < \beta_c(\lambda, d)$ , then there is an FPTAS to approximate Z on graphs of maximum degree d; (Weitz's algorithm)
- If  $\beta > \beta_c(\lambda, d)$ , then it is NP-hard to approximate Z on d-regular graphs.

#### Remark

This threshold  $\beta_c(\lambda, d)$  coincides with the threshold for uniqueness of the Gibbs measure on the infinite d-regular tree.

### Ferromagnetic Ising model

There is also a uniqueness phase transition in the ferromagnetic regime, but there is no approximability transition:

#### Theorem (Jerrum-Sinclair 1993)

For  $0 < \beta < 1$  and  $\lambda > 0$ , there exists a **randomized** MCMC algorithm (FPRAS) for approximating the partition function of the ferromagnetic Ising model on graphs.

Deterministic approximation is currently known only up to the uniqueness threshold:

#### Theorem (Zhang, Liang and Bai 2011

For  $\frac{\Delta-1}{\Delta+1} < \beta < 1$  and  $\lambda > 0$ , there exists an FPTAS for approximating the partition function of the ferromagnetic Ising model on graphs of maximum degree  $\Delta$ .

The presence of the uniqueness phase transition is an obstacle to *decay of correlations*, but not an obstacle to approximability

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### Zeros of partition functions

Instead of making use of the uniqueness property, we appeal to the classical notion of phase transition, as zeros of the partition function:

#### Theorem (Lee-Yang 1952)

For  $0 < \beta \le 1$ , the zeros of  $Z_G^{\beta}(\lambda)$  (viewed as a polynomial in  $\lambda$ ) satisfy  $|\lambda| = 1$ .



 $Z^{\beta}(\lambda)$  is zero-free except on the unit circle in complex plane

#### Our results

#### Theorem

Fix any  $\Delta > 0$ . There is a FPTAS for the Ising partition function  $Z_G^{\beta}(\lambda)$  in all graphs G of maximum degree  $\Delta$  for all edge activities  $-1 \leq \beta \leq 1$  and all (possibly complex) vertex activities  $\lambda$  with  $|\lambda| \neq 1$ .

#### Remark

This is the first deterministic FPTAS for (almost) the whole range of  $\beta$ ,  $\lambda$ . We can also allow edge-dependent activities  $\beta_e$  provided all of them lie in [-1,1].

### Our results (cont'd)

#### Definition (Ising Model on Hypergraphs)

$$Z_H^{\beta}(\lambda) = \sum_{S \subseteq V} \beta^{|E(S,V \setminus S)|} \lambda^{|S|}.$$

#### Theorem (Lee-Yang Theorem for Hypergraphs)

Let H=(V,E) be a hypergraph with maximum hyperedge size  $k \geq 3$ . Then all the zeros of the Ising model partition function  $Z_H^{\beta}(\lambda)$  lie on the unit circle if and only if the edge activity  $\beta$  lies in the range

$$-\frac{1}{2^{k-1}-1} \le \beta \le \frac{1}{2^{k-1}\cos^{k-1}\left(\frac{\pi}{k-1}\right)+1}.$$

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### Our results (cont'd)

In combination with our Lee-Yang theorem for hypergraphs:

#### Theorem

Fix any  $\Delta>0$  and  $k\geq 3$ . There is an FPTAS for the Ising partition function  $Z_H^\beta(\lambda)$  in all hypergraphs H of maximum degree  $\Delta$  and maximum edge size k, for all edge activities  $\beta$  such that

$$-\frac{1}{2^{k-1}-1} \le \beta \le \frac{1}{2^{k-1}\cos^{k-1}\left(\frac{\pi}{k-1}\right)+1}$$

and all vertex activities  $|\lambda| \neq 1$ .

### Our results (cont'd)

Recall

$$Z_G^{oldsymbol{arphi}}(\lambda) = \sum_{oldsymbol{\sigma}: V 
ightarrow \{+,-\}} \prod_{e \in E} arphi_e ig(oldsymbol{\sigma}ig|_eig) \, \lambda^{|\{v:oldsymbol{\sigma}(v)=+\}|}$$

Together with Suzuki-Fisher 1971 (Lee-Yang theorem for general ferromagnetic 2-spin models):

#### Theorem

Fix any  $\Delta>0$  and  $k\geq 2$  and a family of edge activities  $\boldsymbol{\varphi}=\{\varphi_e\}$  satisfying

- symmetry:  $\varphi_e(\sigma) = \overline{\varphi_e(-\sigma)}$ ;
- "ferromagnetism":  $|\varphi_e(+,\cdots,+)| \geq \frac{1}{4} \sum_{\sigma \in \{+,-\}^V} |\varphi_e(\sigma)|$ .

Then there exists an FPTAS for the partition function  $Z_H^{\varphi}(\lambda)$  in all hypergraphs H of maximum degree  $\Delta$  and maximum edge size k for all vertex activities  $\lambda \in \mathbb{C}$  such that  $|\lambda| \neq 1$ .

#### Overview

- Approximate counting, sampling and motivations
- Our results
- Proof sketch of our FPTAS
- 4 Lee-Yang theorem

### Approximation via the log-partition function

#### Theorem (Barvinok, Barvinok and Soberon)

For a zero free region,  $\log Z$  can be approximated to within  $\pm \varepsilon$  by its k-th order Taylor series, for  $k = O(\log(n/\varepsilon))$ .

### To make use of the analyticity of $\log Z$

- Taylor expansion of  $\log Z$  around  $\lambda = 0$
- ullet By Lee-Yang theorem,  $|\lambda| < 1$  is zero free
- The first k terms of the Taylor series require the first k + 1 coefficients of Z

$$Z_G^{\beta}(\lambda) = \sum_{i=0}^n \left( \sum_{\substack{S \subseteq V \\ |S|=i}} \beta^{|E(S,\overline{S})|} \right) \lambda^i,$$

$$\log Z = \sum_{i=0}^{k-1} \frac{\lambda^i}{i!} \left( \frac{\mathrm{d}^i}{\mathrm{d}\lambda^i} \log Z \Big|_{\lambda=0} \right) + \cdots$$

Naively, computing the first k coefficients of Z takes time  $O(n^k) \implies$  quasi-polynomial time algorithm

### Computing coefficients of Z

#### Theorem (Patel and Regts)

If the first k coefficients can be represented as a sum over induced subgraphs, one can compute them in time  $\Delta^{O(k)} = \text{poly}(n/\varepsilon)$  for graphs of bounded degree  $\Delta$ . In particular, for the Ising model, if  $\lambda = 1$  and  $|\beta - 1| < 0.34/\Delta$ , there is an FPTAS.

For graphs of maximum degree  $\Delta$ :

- the number of labeled *induced subgraphs* of size k is  $O(n^k)$
- the number of labeled *connected* induced subgraphs of size k is at most  $n(e\Delta)^k$

Main idea: reduce a sum over all induced subgraphs to sum over connected induced subgraphs.

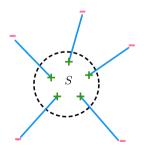
 $\odot$  The Ising model, when viewed as a polynomial in  $\lambda$ , is **not** a sum over induced subgraphs.

### Insects in graphs

Main idea: Generalize the notion of *induced subgraphs* to *induced sub-insects*. Recall the Ising partition function:

$$Z_G^{\beta}(\lambda) = \sum_{S \subseteq V} \beta^{|E(S,V \setminus S)|} \lambda^{|S|}.$$

Given a configuration  $\sigma$ , let S be the set of vertices assigned +-spins:

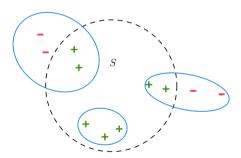


The coefficient  $\beta^{|E(S,V\setminus S)|}$  of a configuration depends only on the *induced sub-insect*  $G^+[S]$ 

### Insects in hypergraphs

Recall that w.l.o.g  $\varphi_e(-,\cdots,-)=1$ .

$$Z_{H}^{\varphi}(\lambda) = \sum_{\sigma: V \to \{+, -\}} \prod_{e \in E} \varphi_{e}(\sigma|_{e}) \lambda^{|\{v: \sigma(v) = +\}|} = \sum_{S \subseteq V} \prod_{e: e \cap S \neq \emptyset} \varphi_{e}(S) \lambda^{|S|}.$$



The coefficient  $\prod_{e:e\cap S\neq\emptyset} \varphi_e(\sigma|_e)$  of a configuration depends only on the induced sub-insect  $H^+[S]$ 

Note: the number of labeled *connected* sub-insects of size t is at most  $n(e\Delta k)^t$ .

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### Reducing to a connected sub-insect count

Let  $r_1, \dots, r_n$  be the complex zeros of  $Z_G^{\varphi}(\lambda)$ :

$$Z_G^{\varphi}(\lambda) = \prod_{i=1}^n (1 - \lambda/r_i) = \sum_{i=0}^n (-1)^i e_i(G) \lambda^i,$$

#### Review of our goal:

- To compute the Taylor series of  $\log Z$ , we need the coefficients  $e_i(G)$
- $e_i(G)$  is the elementary symmetric polynomial evaluated at  $(\frac{1}{r_1}, \cdots, \frac{1}{r_n})$
- From the definition of *Z*.

$$e_i(G) = (-1)^i \sum_{\substack{S \subseteq V \\ |S|=i}} \prod_{e: e \cap S \neq \emptyset} \varphi_e(S)$$

• Notice that  $e_i(G)$  is a weighted sub-insect count, but not necessarily connected

Instead, we consider a related quantity: the t-th power sum given by  $p_t = \sum_{i=1}^n 1/r_i^t$ 

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### Reducing to a connected sub-insect count (cont'd)

Let  $p_t = \sum_{i=1}^n 1/r_i^t$  be the *t*-th power sum. By Newton's identities:

$$p_t = \sum_{i=1}^{t-1} (-1)^{i-1} p_{t-i} e_i + (-1)^{t-1} t e_t.$$

#### Proof sketch

- Recall that  $e_i$  is a weighted sub-insect count
- Lemma: product of weighted sub-insect counts is also a weighted sub-insect count
- Thus  $p_t$  is also a weighted sub-insect count
- ullet Notice that  $p_t$  is additive in the sense that  $p_t(G_1 \cup G_2) = p_t(G_1) + p_t(G_2)$
- Lemma: a weighted sub-insect count is additive iff it is a connected sub-insect count
- Thus  $p_t$  is supported only on *connected* sub-insects up to size t.

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# Summary of FPTAS for $Z_G^{\beta}(\lambda)$

#### Taylor approximation:

- Since  $Z_G^{\beta}(\lambda) = \lambda^n \cdot Z_G^{\beta}(1/\lambda)$ , w.l.o.g.  $|\lambda| < 1$
- To get a  $(1 \pm \varepsilon)$  multiplicative approximation of Z, it suffices to get a  $\pm \frac{\varepsilon}{4}$  additive approximation of  $\log Z$  (by standard complex analysis)
- By Barvinok et. al., the *t*-th order Taylor series of  $\log Z$  around  $\lambda=0$  is a  $\pm \varepsilon$  approximation for  $t=O(\log(n/\varepsilon))$  at any point  $\lambda$  such that  $B(0,|\lambda|)$  is a zero-free region
- ullet By the Lee-Yang theorem, there are no zeros of Z in  $|\lambda| < 1$

# Summary of FPTAS for $Z_G^{\beta}(\lambda)$

#### Computing coefficients by reducing to a connected sum:

- The *t*-th order Taylor series of  $\log Z$  depends only on the first t+1 coefficients
- Recall that  $Z_G^{\beta}(\lambda) = \sum_{i=0}^n (-1)^i e_i(G) \lambda^i$ ,  $e_i$  is the *i*-th coefficient
- $\bullet$   $e_t$  can be computed using Newton's identities given  $p_t$
- $p_t$  can be computed efficiently by enumerating over *connected* sub-insects for  $t = O(\log(n/\varepsilon))$

#### Overview

- Approximate counting, sampling and motivations
- Proof sketch of our FPTAS
- 4 Lee-Yang theorem

### Lee-Yang theorem

#### Definition (Lee-Yang property)

A multilinear polynomial P is said to have the *Lee-Yang property*, denoted by  $P \in LY$ , if  $P(\lambda_1, \dots, \lambda_n) \neq 0$  for any  $\lambda_1, \dots, \lambda_n$  such that  $|\lambda_i| \geq 1$  for all i, and  $|\lambda_i| > 1$  for some i.

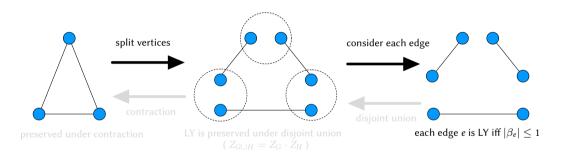
#### Definition (Multivariate Ising model)

$$Z_G^{\vec{\beta}}(\lambda_1, \cdots, \lambda_n) = \sum_{S \subseteq V} \prod_{e \in E(S, \overline{S})} \beta_e \prod_{i \in S} \lambda_i$$

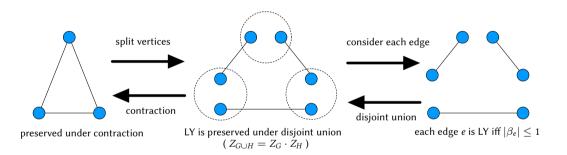
#### Theorem (Lee-Yang Theorem 1952)

If  $0 < \beta_e < 1$ , then  $Z_G^{\vec{\beta}}(\lambda_1, \dots, \lambda_n)$  has the Lee-Yang property.

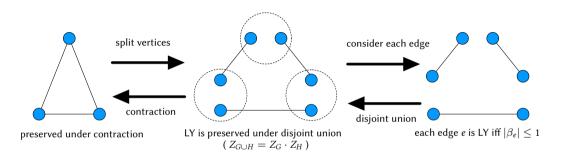
### Asano's proof of the Lee-Yang theorem



### Asano's proof of the Lee-Yang theorem



### Asano's proof of the Lee-Yang theorem



What about hypergraphs? The above scheme is very general except:

- LY holds for each hyperedge
- LY is preserved under contraction

### Characterizing Lee-Yang theorems for symmetric polynomials

#### Lemma (Criterion for Lee-Yang property)

Given a multilinear polynomial  $P(z_1, z_2, ..., z_n)$ , define multilinear polynomials  $A_j$  and  $B_j$  in the variables  $z_1, ..., z_{j-1}, z_{j+1}, ..., z_n$  such that

$$P = A_j z_j + B_j$$

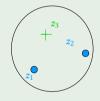
If P is symmetric, i.e.,  $P(z) = \prod_{i=1}^n z_i \cdot \overline{P(1/z)}$ , then  $P \in LY$  if and only if  $A_j \in LY$  for all j.

- For a single hyperedge: LY for one of the leading coefficients  $A_j$  implies LY for a hyperedge.
- For contraction: LY for the disjoint union implies that LY for all the leading coefficients.

### Lee-Yang theorem on a single hyperedge

### The leading coefficients $A_i$

Recall that  $Z = A_j z_j + B_j$ ,



 $A_3 = z_1 z_2 + \beta z_1 + \beta z_2 + \beta$  in a hyperedge of size 3.

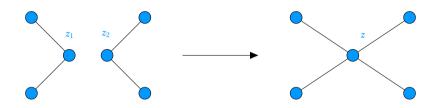
More generally,  $A_j \in \mathrm{LY}$  (that is,  $A_j = 0$  for  $|z_i| \geq 1$ ) is equivalent to

$$\frac{1}{eta} = 1 - \prod_{\substack{i=1 \ i 
eq j}}^k \left(1 + \frac{1}{z_i}\right).$$

This characterizes the range in our theorem:

$$-\frac{1}{2^{k-1}-1} \le \beta \le \frac{1}{2^{k-1}\cos^{k-1}\left(\frac{\pi}{k-1}\right)+1}.$$

#### Asano contraction



As an contraction: suppose that  $Az_1z_2 + Bz_1 + Cz_2 + D \in LY$ , need to show  $Az + D \in LY$ .

- $Az_1z_2 + Bz_1 + Cz_2 + D \neq 0$  for  $|z_1|, |z_2| > 1$
- $Az^2 + (B+C)z + D = 0$  only if  $|z| \le 1$
- $\left|\frac{D}{A}\right| \le 1$ , using Vieta's formula for product of zeros
- By our lemma,  $A \in LY$ , so  $A \neq 0$ . Thus Az + D = 0 only if  $|z| = \left| \frac{D}{A} \right| \leq 1$
- $Az + D \in LY$

### Discussions and Open problems

#### Open problem

What about  $\lambda = 1$ ?

- There are zeros arbitrarily close to  $\lambda = 1$  at low temperature
- Our algorithm works for all  $|\beta| \le 1$ . FPTAS for  $\lambda = 1$  and  $-1 < \beta < 0$  would give FPTAS for counting perfect matchings in general (non-bipartite) graphs

#### Open problem

Connections of locations of zeros, and algorithms such as MCMC and the correlation decay approach?

- Jerrum-Sinclair's MCMC works in subgraphs world instead of the spins world, which by Lee-Yang theorem is real-rooted
- Analog of Griffiths inequality for the self-avoiding walk tree

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