

## Memory Structures

(1) System Global Area (SGA)  
- Shared SQL Area

(2) Program Global Area (PGA)  
- Private SQL Area

## SQL statement processing

- (1) parsing ←
- (2) binding
- (3) exec
- (4) fetching

Plan

- ★ (1) Library Cache Latch
- ★ (2) Shared Pool Latch
- (3) Transform
- (4) Estimate
- (5) Plan Generator

## CBO

### (1) Query Transformer

- to rewrite the original SQL statements to find out more possible access paths

{ OR --> UNION ALL  
  Subquery with IN --> Join form  
  etc



# CBO Estimator



## (1) Optimizer statistics

- rows
- index
- B\*Tree or Bitmap
- NDV
- etc

1
2
3
4
5
6
7
8
9
10

5  
.  
5

Plan 1

Plan 2

Plan 3

Histogram

(1) selectivity: % of data to be filtered (0 ~ 1)

- where col = 5 10% --> 0.1

=  $1/\text{NDV}$  (assumption: even distribution)

=  $1/10$

(2) cardinality: number of rows to be processed

- 1 row

= rows x selectivity

=  $10 \times 1/10$

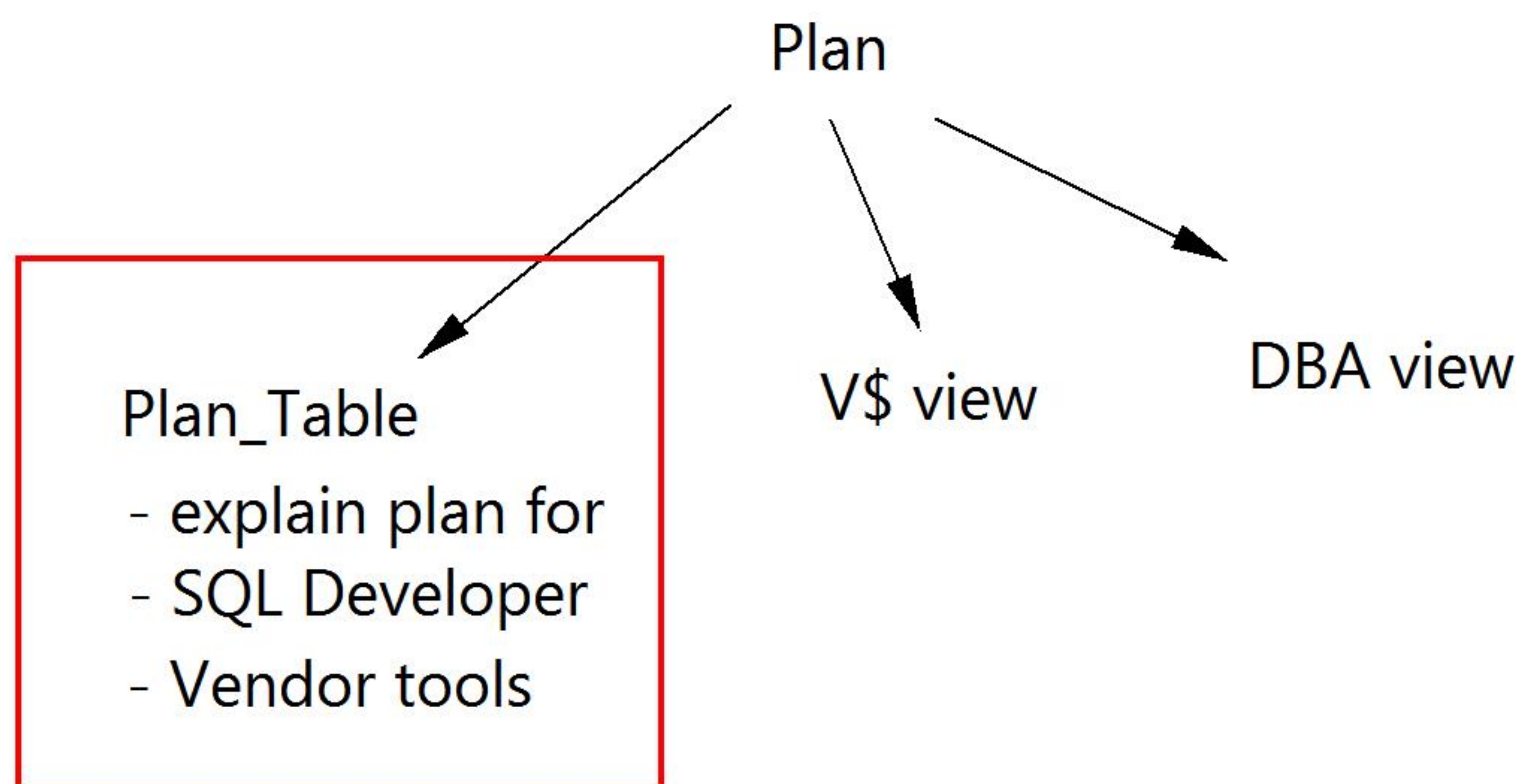
= 1 row

Cost (System Statistics)



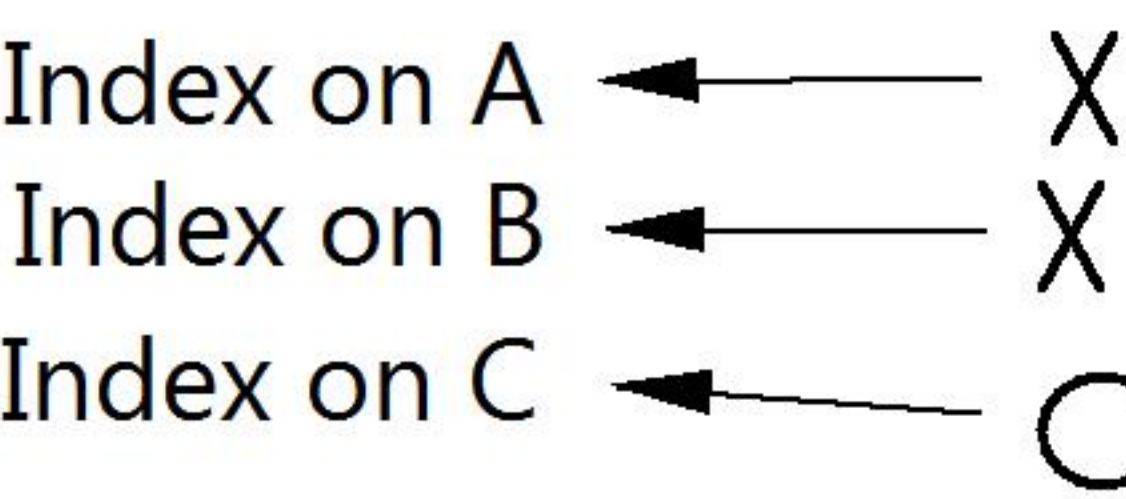
Plan Cost Unit = Estimated Time to Read a Single Block

dbms\_xplan  
- display  
- display\_cursor  
- display\_awr





A, B, C



Index on (A, B, C)

## Single Column Index vs Composite Column Index

- Which one is better? Single Column Index

If multiple columns appear in the WHERE clause, then you can possibly take advantage of Composite Key Index

How should I create a composite key index?

- Index on (A, B, C) frequently
- What should be in A position?
  - (1) The most selective column should be in A.
  - (2) If all three columns are equally selective, then choose the column with the best selectivity.

unique or non unique

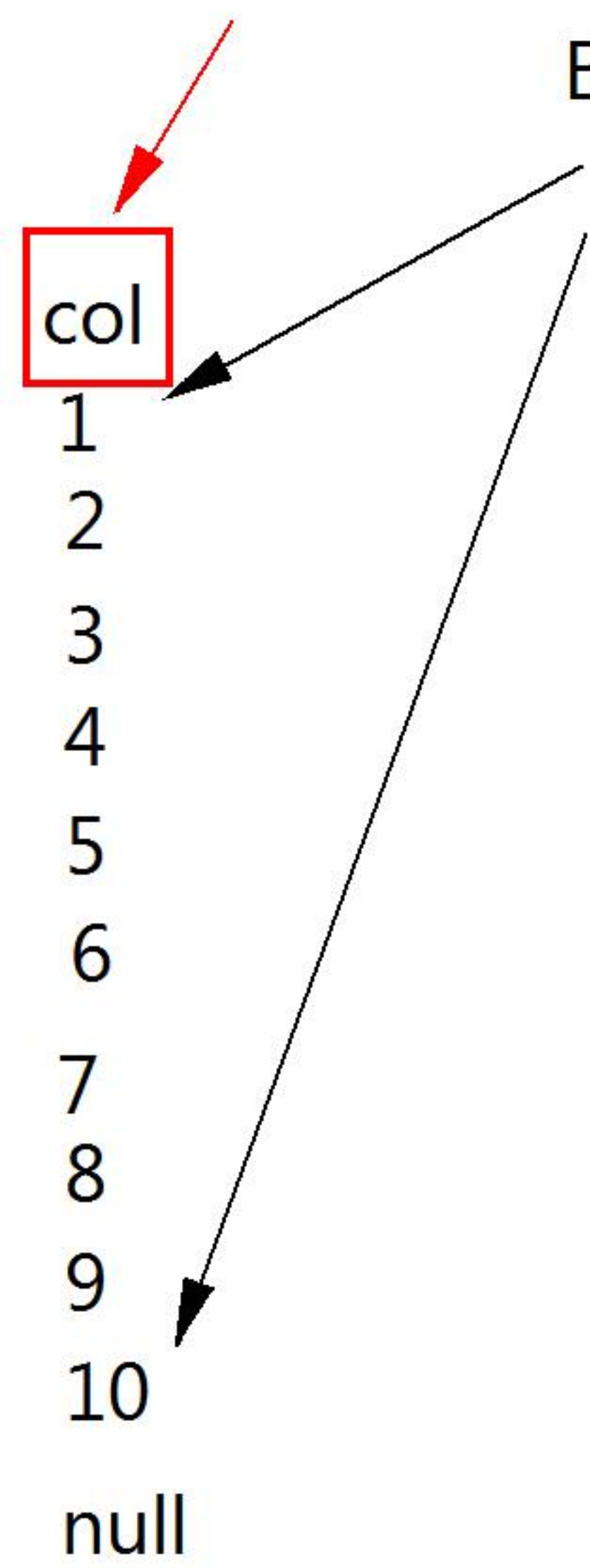
gender

postal\_code

Index Skip Scan



## B\*Tree Index



select count(col)  
from table;  
where col is null

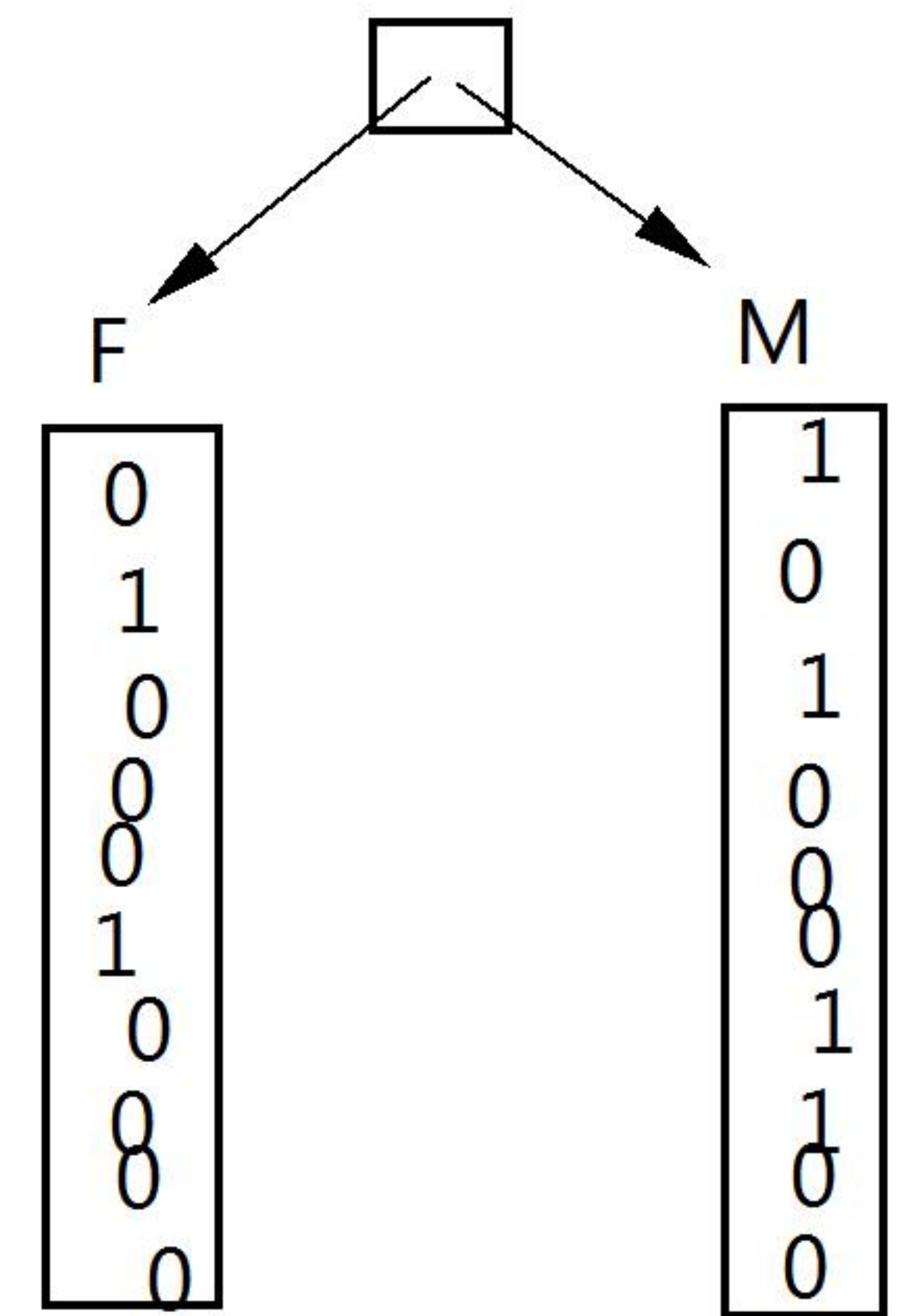
When using index --> 10  
11

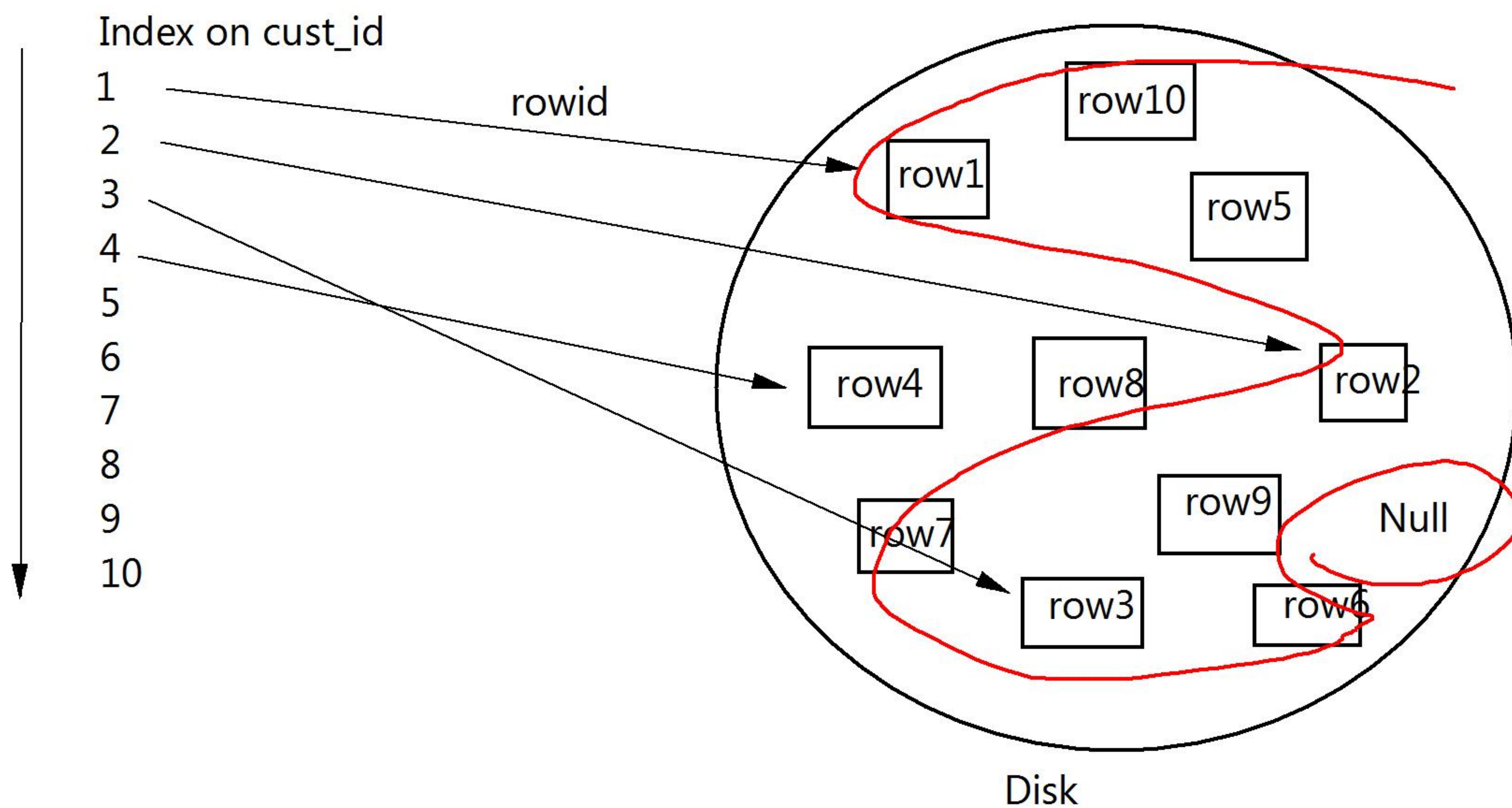
where col is not null

## gender

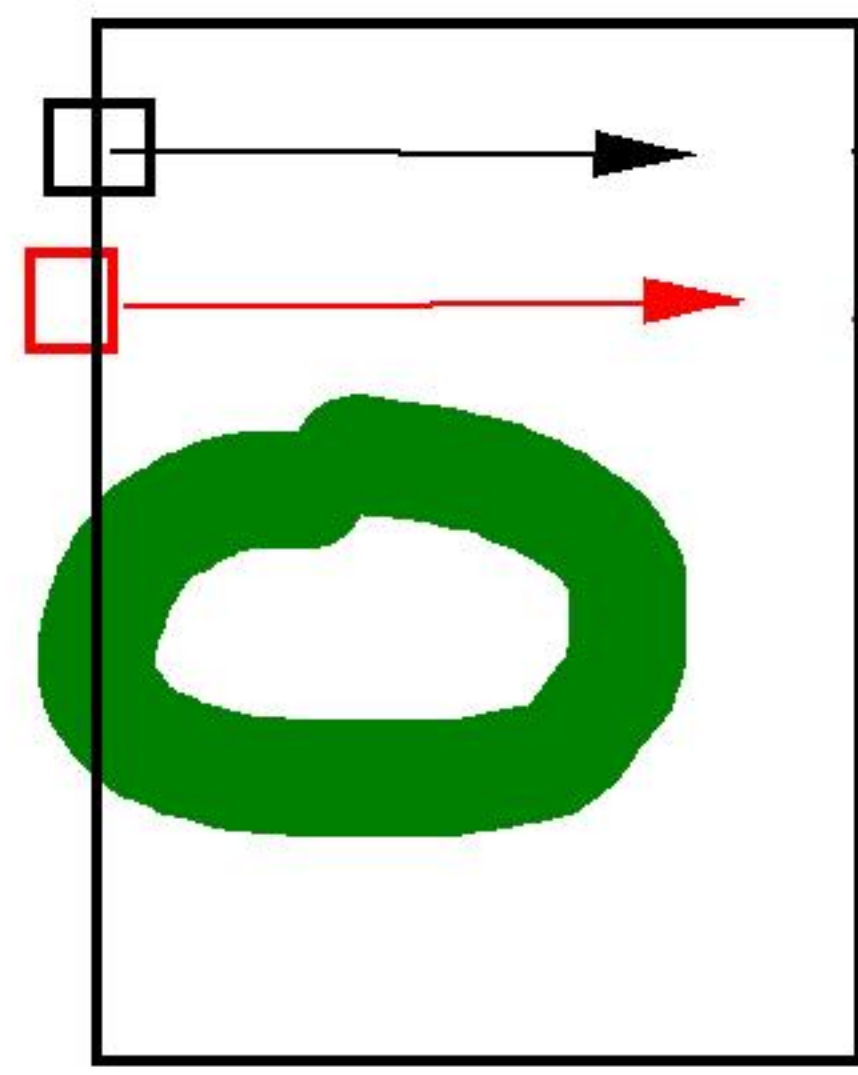
1 M  
2 F  
3 M  
4  
5  
6 F  
7 M  
8 M  
9 M  
10

## Bitmap



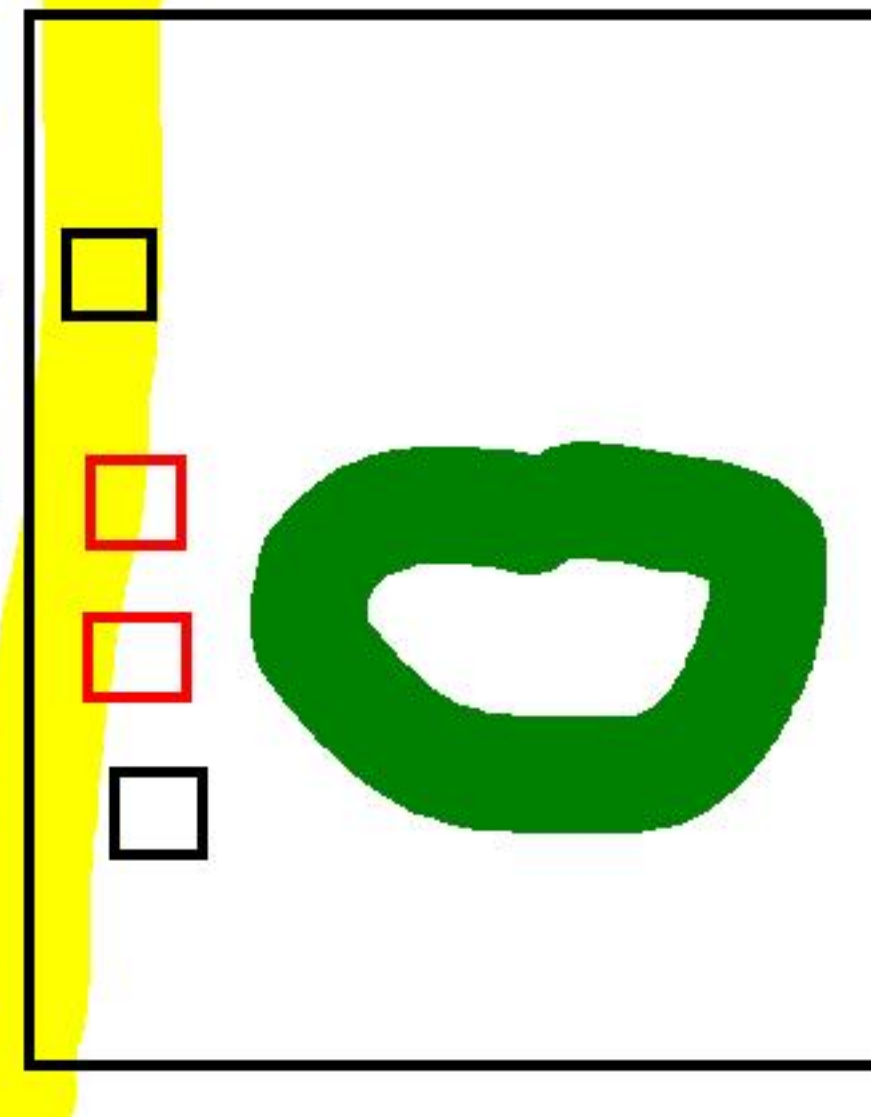






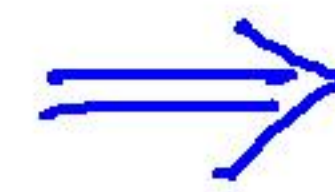
Driving Table

10  
100  
100B



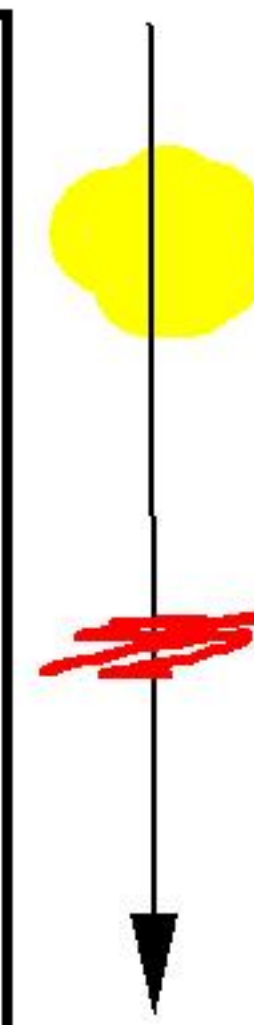
Driven Table

100  
1000  
1B



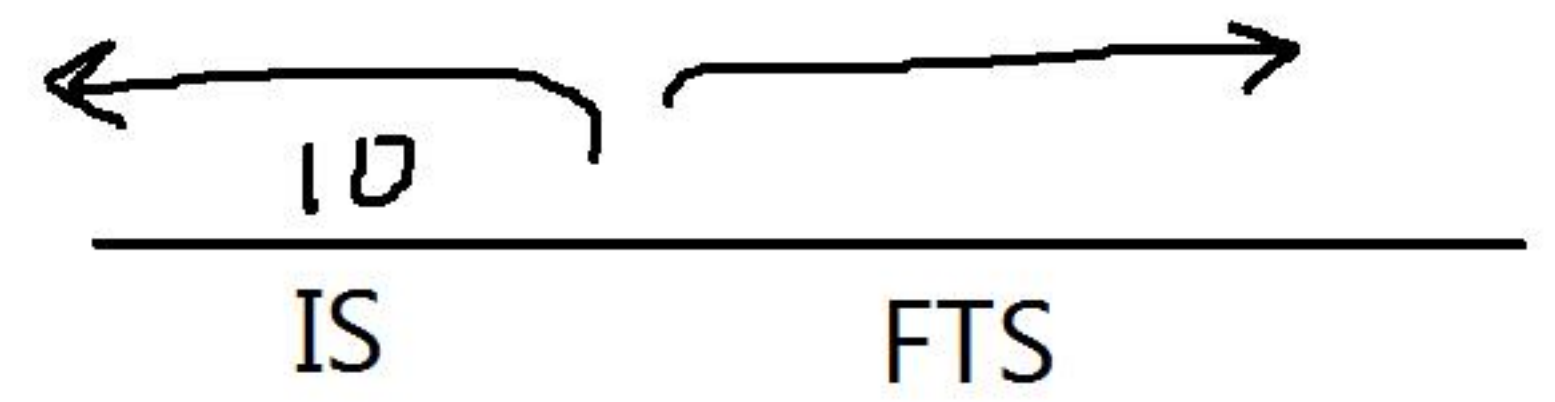
Nested Loop

- (1) small data sources
- (2) good filter conditions  
- use indexes
- (3) OLTP
- (4) SGA



Index

selectivity < 0.1

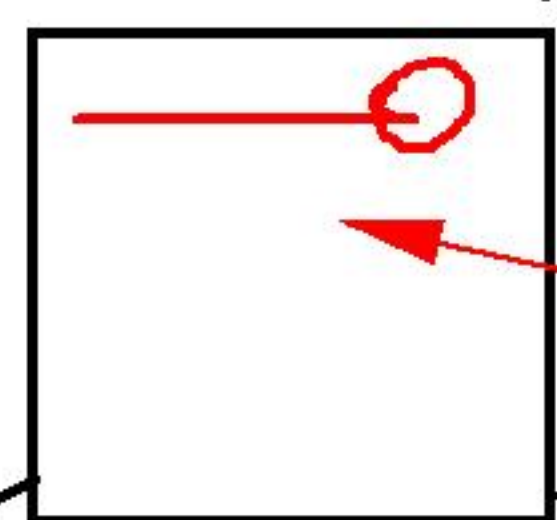
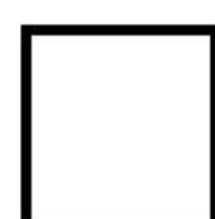


## Hash Join

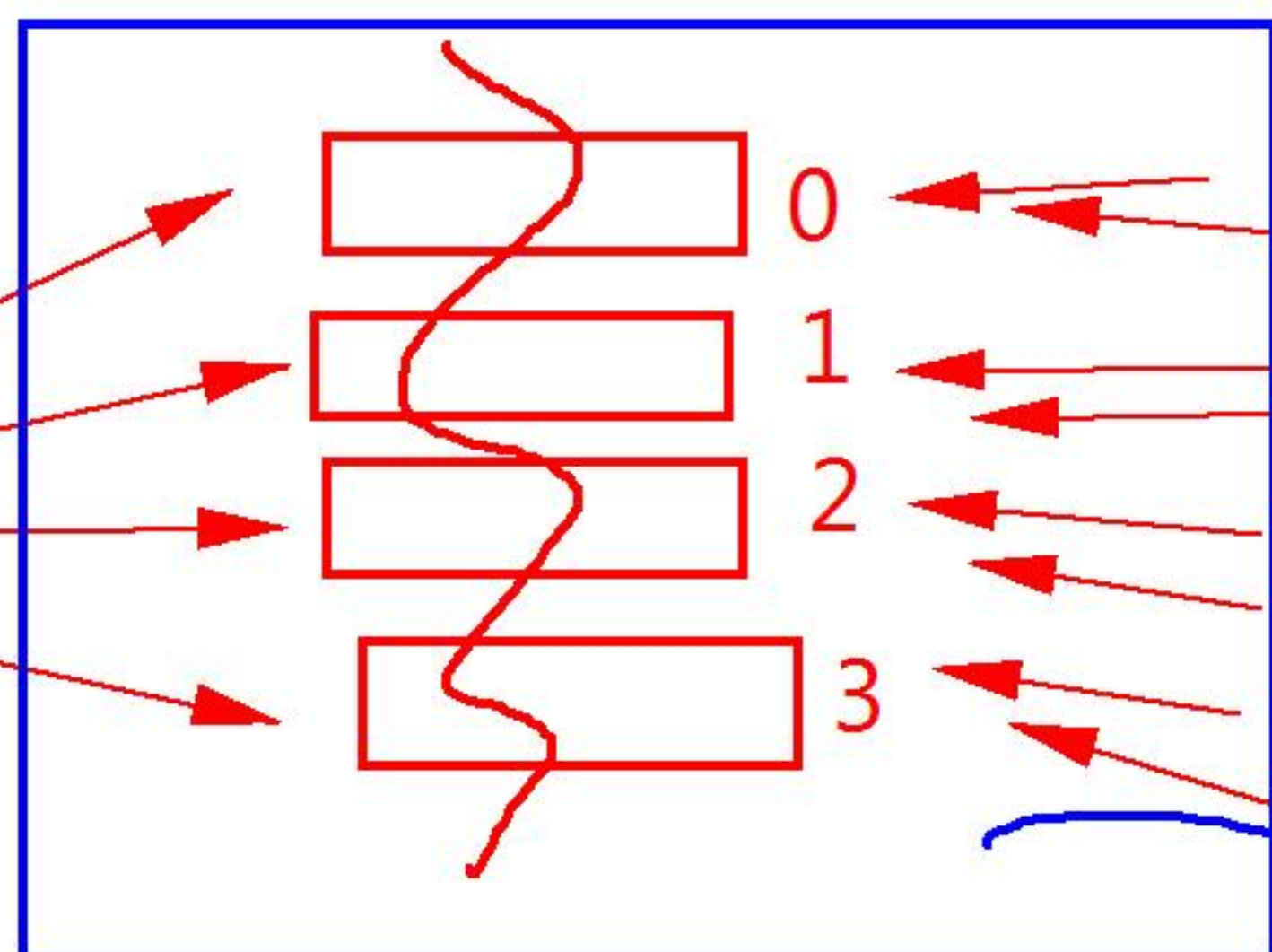
- small-large
- large-large
- uses Hash function
- uses PGA



DWH



sales



$$\underline{f(\text{key})} = \frac{\text{key}}{1} \% 4$$

1

Step 18 in Practice 3



## Merge Join

- Large Data Sources

- case where HJ is NOT possible ( $<>$  between)
- pre-sorting

